



# Insert Interval (medium)

We'll cover the following ^

- Problem Statement
- Try it yourself
- Solution
  - Code
  - Time complexity
  - Space complexity

## Problem Statement #

Given a list of non-overlapping intervals sorted by their start time, **insert a given interval at the correct position** and merge all necessary intervals to produce a list that has only mutually exclusive intervals.

### Example 1:

Input: Intervals=[[1,3], [5,7], [8,12]], New Interval=[4,6]

Output: [[1,3], [4,7], [8,12]]

Explanation: After insertion, since [4,6] overlaps with [5,7], we merge d them into one [4,7].

### Example 2:

Input: Intervals=[[1,3], [5,7], [8,12]], New Interval=[4,10]

Output: [[1,3], [4,12]]



Explanation: After insertion, since [4,10] overlaps with [2,8] & [8,12], we merged them into [4,12].



### Example 3:

Input: Intervals=[[2,3],[5,7]], New Interval=[1,4]

Output: [[1,4], [5,7]]

Explanation: After insertion, since [1,4] overlaps with [2,3], we merge d them into one [1,4].

## Try it yourself #

Try solving this question here:



```
1 using namespace std;
2
3 #include <iostream>
4 #include <vector>
5
6 class Interval {
7 public:
8     int start = 0;
9     int end = 0;
10
11     Interval(int start, int end) {
12         this->start = start;
13         this->end = end;
14     }
15 };
16
17 class InsertInterval {
18 public:
19     static vector<Interval> insert(const vector<Interval> &intervals, Interval newInte
20         vector<Interval> mergedIntervals;
21         // TODO: Write your code here
22         return mergedIntervals;
23     }
24 };
25
26 int main(int argc, char *argv[]) {
27     vector<Interval> input = {{1, 3}, {5, 7}, {8, 12}};
```



```
28     cout << "Intervals after inserting the new interval: ";
29     for (auto interval : InsertInterval::insert(input, {4, 6}))
30         cout << "[" << interval.start << "," << interval.end << "] ";
31 }
```



## Solution

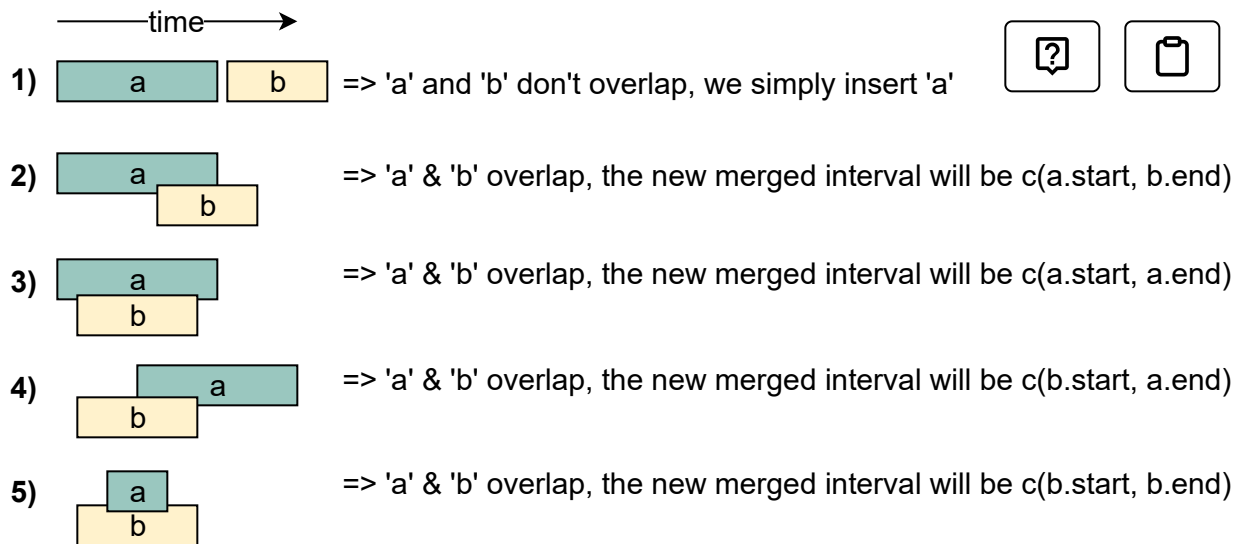
If the given list was not sorted, we could have simply appended the new interval to it and used the `merge()` function from [Merge Intervals](#). But since the given list is sorted, we should try to come up with a solution better than  $O(N * \log N)$

When inserting a new interval in a sorted list, we need to first find the correct index where the new interval can be placed. In other words, we need to skip all the intervals which end before the start of the new interval. So we can iterate through the given sorted list of intervals and skip all the intervals with the following condition:

```
intervals[i].end < newInterval.start
```

Once we have found the correct place, we can follow an approach similar to [Merge Intervals](#) to insert and/or merge the new interval. Let's call the new interval 'a' and the first interval with the above condition 'b'. There are five possibilities:





The diagram above clearly shows the merging approach. To handle all four merging scenarios, we need to do something like this:

```
c.start = min(a.start, b.start)
c.end = max(a.end, b.end)
```

Our overall algorithm will look like this:

1. Skip all intervals which end before the start of the new interval, i.e., skip all **intervals** with the following condition:

```
intervals[i].end < newInterval.start
```

2. Let's call the last interval 'b' that does not satisfy the above condition. If 'b' overlaps with the new interval (a) (i.e.  $b.start \leq a.end$ ), we need to merge them into a new interval 'c':

```
c.start = min(a.start, b.start)
c.end = max(a.end, b.end)
```

3. We will repeat the above two steps to merge 'c' with the next overlapping interval.

## Code #



Here is what our algorithm will look like:



C++



```
1 using namespace std;
2
3 #include <iostream>
4 #include <vector>
5
6 class Interval {
7 public:
8     int start = 0;
9     int end = 0;
10
11     Interval(int start, int end) {
12         this->start = start;
13         this->end = end;
14     }
15 };
16
17 class InsertInterval {
18 public:
19     static vector<Interval> insert(const vector<Interval> &intervals, Interval newInte
20         if (intervals.empty()) {
21             return vector<Interval>{newInterval};
22         }
23
24         vector<Interval> mergedIntervals;
25
26         int i = 0;
27         // skip (and add to output) all intervals that come before the 'newInterval'
28         while (i < intervals.size() && intervals[i].end < newInterval.start) {
29             mergedIntervals.push_back(intervals[i++]);
30         }
31
```



## Time complexity

As we are iterating through all the intervals only once, the time complexity of the above algorithm is  $O(N)$ , where 'N' is the total number of intervals.



# Space complexity



The space complexity of the above algorithm will be  $O(N)$  as we need to return a list containing all the merged intervals.

[← Back](#)[Merge Intervals \(medium\)](#)[Next →](#)[Intervals Intersection \(medium\)](#)☒ Completed[⚠ Report an Issue](#)