



COC Berlin Code of Conduct





CATEGORY THEORY FOR PROGRAMMERS



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Category Theory for

Programmers

Main Concepts	Supporting Concepts	
Universal Construction (technique)		
A) Objects	Naturality Condition	
B) Morphisms Also known as "arrows"	Isomorphism	
C) Category (A + B)	Hom-Set	
D) Functor (C + B) Bifunctor, Profunctor, Contravariant Functor Product, Coproduct Maybe, List, Reader	Hom-Functor	
E) Natural Transformation (D + B)	Yoneda Lemma	
	Yoneda Embedding	

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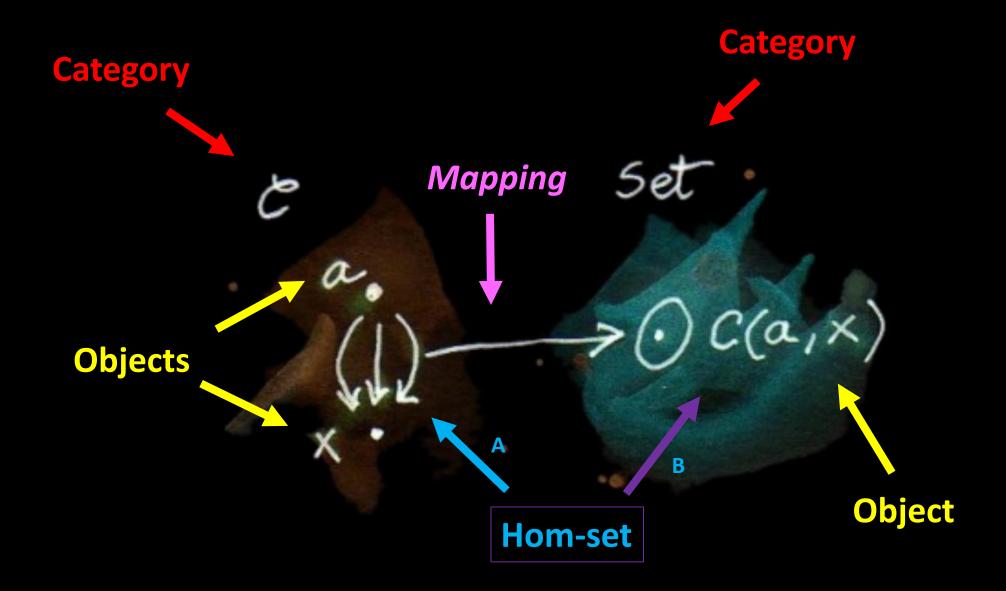
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14.1 The Hom Functor

Every category comes equipped with a canonical family of mappings to **Set**. Those mappings are in fact functors, so they preserve the structure of the category. Let's build one such mapping.

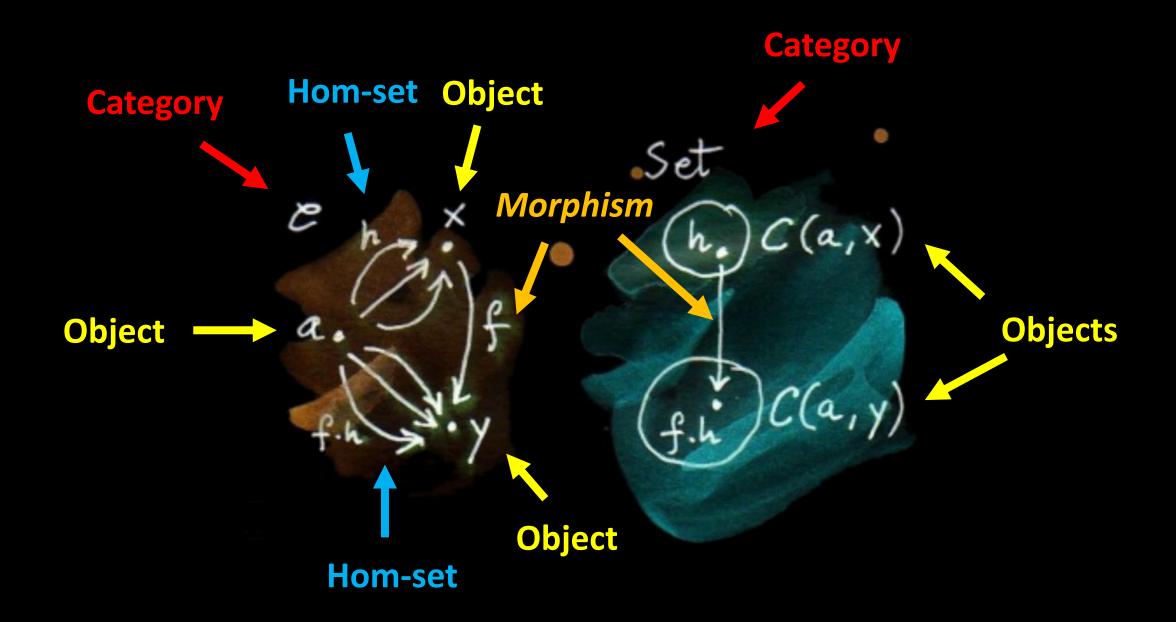
Let's fix one object a in C and pick another object x also in C. The hom-set C(a, x) is a set, an object in **Set**. When we vary x, keeping a fixed, C(a, x) will also vary in **Set**. Thus we have a mapping from x to **Set**.





If we want to stress the fact that we are considering the hom-set as a mapping in its second argument, we use the notation C(a, -) with the dash serving as the placeholder for the argument.

This mapping of objects is easily extended to the mapping of morphisms. Let's take a morphism f in C between two arbitrary objects x and y. The object x is mapped to the set C(a, x), and the object y is mapped to C(a, y), under the mapping we have just defined. If this mapping is to be a functor, f must be mapped to a function between the two sets: $C(a, x) \rightarrow C(a, y)$



There two subsections across chapters called **The Hom-Functor**Ch 8 Functors Section 8 Ch 14 Representable Functors Section 1

The above examples are the reflection of a more general statement that the mapping that takes a pair of objects a and b and assigns to it the set of morphisms between them, the hom-set C(a, b), is a functor. It is a functor from the product category $C^{op} \times C$ to the category of sets, **Set**.

Functors

T THE RISK OF SOUNDING like a broken record, I will say this about functors: A functor is a very simple but powerful idea. Category theory is just full of those simple but powerful ideas. A functor is a mapping between categories. Given two categories, C and D, a functor

