# Advanced Functional Programming

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Assignment 1-4 are *individual* assignments and Assignment 5 and 6 are group assignments, where a group consists of two students. You will receive the assignment descriptions for Assignment 4-6 according to the schedule. We expect timely submission of the assignments, late assignments count as failed assignments. We expect genuine solutions and plagiarism will lead directly to fail. Please read the submission guide for assignments 4-6.

#### 1 Assignment

Perform  $\beta$  reduction of the following terms and write the  $\beta$ -contractum: (Lambda calculus slides, p. 21.)

- 1.  $\lambda x.x$
- 2.  $(\lambda x.x)y$
- 3.  $(\lambda x.xx)y$
- 4.  $(\lambda x.z)y$
- 5.  $(\lambda x \lambda y.x)z$

## 2 Assignment

Does the following terms have a  $\beta$  Normalform? Please briefly explain why.  $\beta$  Normalforms are presented in (Lambda calculus slides, p. 21.)

- 1.  $\mathbf{I} = \lambda x.x$
- 2.  $\Omega = \omega \omega$  with  $\omega = \lambda x.xx$
- 3. **KI** $\Omega$  with **K** =  $\lambda x \ y.x$
- 4.  $(\lambda x.\mathbf{KI}(x\ x)\lambda y.\mathbf{KI}(y\ y))$
- 5.  $(\lambda x.z(xx))\lambda y.z(yy)$

## 3 Assignment

A reduction path of a  $\lambda$  term M is a finite or infinite sequence of the following form:

$$M \to_{\beta} M_1 \to_{\beta} M_2 \to_{\beta} \dots$$

We call a term weak normalizing, if it has a Normalform. A term is called strong normalizing, if all of its reduction paths end in a Normalform.

- 1. Which of the previous five terms are weak normalizing and which of them are strong normalizing?
- 2. In which of these cases does the term end in different Normalforms?

### 4 Assignment

Our calculator is missing the functionality to operate with predefined constants, e.g. pi. Please extend the calculator from the *Calculator* example, by these functionalities.

- 1. Extend the discriminated union type Expr by defining a type called Var which is a string.
- 2. Implement the function memory, which should return a Map<string,float>. Please include the keys and values: "pi", 3.14152976, "x", 10.0.
- 3. Implement the function lookup, which takes a key x and a storage stor of type Map<string, ConstType> as arguments and returns the value of key x from storage stor. If the key it not contained in stor then it should fail with an error message.
- 4. Extend the function eval to evaluate variables of type Var, such that it performs a lookup on a key in the storage.
- 5. Run the program and test that your output is correct.

### 5 Assignment

This assignment is about implementing an interpreter for terms of untyped  $\lambda$  calculus as presented in (Lambda calculus slides). We will use various simplifications:

- We ignore all the  $\alpha$  equivalence, as presented (Lambda calculus slides, p. 16).
- We ignore handling free variables and potential collisions (Lambda calculus slides, p. 19).

In order to implement the interpreter follow the steps below.

- 1. Implement the Embedded Domain Specific Language, Term, for the grammar *pre-terms* presented in (Lambda calculus slides, p. 8.).
- 2. Implement the function substitute with the signature: let rec substitute (term: Term) (var: VarType) (sub: Term). The function substitute should perform the substitution that is defined in (Lambda calculus slides, p. 20). Ignore the free variable side condition in case 4.
- 3. Implement the function reduce with the signature: let rec reduce ((term,red): (Term \* bool)). The function reduce should perform the  $\beta$  reduction that is defined in (Lambda calculus slides, p. 21).
  - **Hints:** (1) the application case should be handled by recursive descend in both arguments. (2) Use the second member of the tuple, **red**, to indicate that substitution has been performed.
- 4. Implement the function eval with the signature: let rec eval (term : Term). The function eval should perform a  $\beta$  reduction, implemented in reduce, until a Normalform is reached. The  $\beta$  Normalform is defined in (Lambda calculus slides, p. 21).
  - **Hint:** a Normalform is reached if no further substitution can be performed. Use the boolean variable red for this check.
- 5. Run the program and test that your output is correct.