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# BLG335E Analysis of Algorithm

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## Project 1



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## A - Asymptotic Upper Bound

Merge Sort is run on  $O(n \lg(n))$ .

My code for Merge Sort:

```
if (lowerBound < upperBound)
{
    int median = (lowerBound + upperBound) / 2;
    sort(unsorted, lowerBound, median);
    sort(unsorted, median + 1, upperBound);
    merge(unsorted, lowerBound, median, upperBound);
}
```

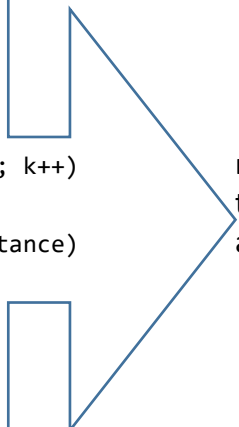


$O(1)$

$T(n) = 2T(n/2) + \text{merge part}$

Merge Part:

```
for (int i = 0; i < leftLength; i++)
{
    Left[i] = unsorted[lowerBound + i];
}
for (int j = 0; j < rightLength; j++)
{
    Right[j] = unsorted[median + 1 + j];
}
for (int k = lowerBound; k <= upperBound; k++)
{
    if (Left[i].distance <= Right[j].distance)
    {
        unsorted[k] = Left[i];
        i++;
    }
    else
    {
        unsorted[k] = Right[j];
        j++;
    }
}
```



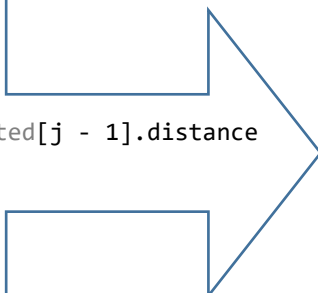
run (upperBound-lowerBound) times maximum so it depends on  $n$  and runs on  $O(n)$

These fors run on  $O(n)$ . So it is  $T(n) = 2T(n/2) + O(n)$ . It becomes finally 1 due to binary division so the result is  $O(n) + O(n) + \dots + O(n)$ . There are  $\lg(n)$  times  $O(n)$ . It shows that it is  $O(n \lg(n))$ .

Insertion Sort is run on  $O(n^2)$ .

My code for Insertion Sort:

```
int j;
Point temp;
for (int i = 1; i < length; i++)
{
    j = i;
    while (j > 0 && unsorted[j].distance < unsorted[j - 1].distance)
    {
        temp = unsorted[j];
        unsorted[j] = unsorted[j - 1];
        unsorted[j - 1] = temp;
        j--;
    }
}
```



for loop  $n$  times while loop  $n$  times; therefore,  $O(n^2)$

It shows that it is  $O(n^2)$ .

Linear Search is run on  $O(n^2)$ .

My code for Linear Search:

```
while (j != total)
{
    for (int i = 0; i < length; i++)
    {
        if (unsorted[i].distance > higherDistance)
        {
            higherDistance = unsorted[i].distance;
            higherIndex = i;
        }
    }
    for (j ; j < total; j++)
    {
        if (unsorted[j].distance < higherDistance)
        {
            higherDistance = unsorted[j].distance;
            temp = unsorted[j];
            unsorted[j] = unsorted[higherIndex];
            unsorted[higherIndex] = temp;
            break;
        }
    }
}
```

$(n-k)*k$  complexity

Best case  $O(1)$

Worst case  $O(n^2/4)=O(n^2)$

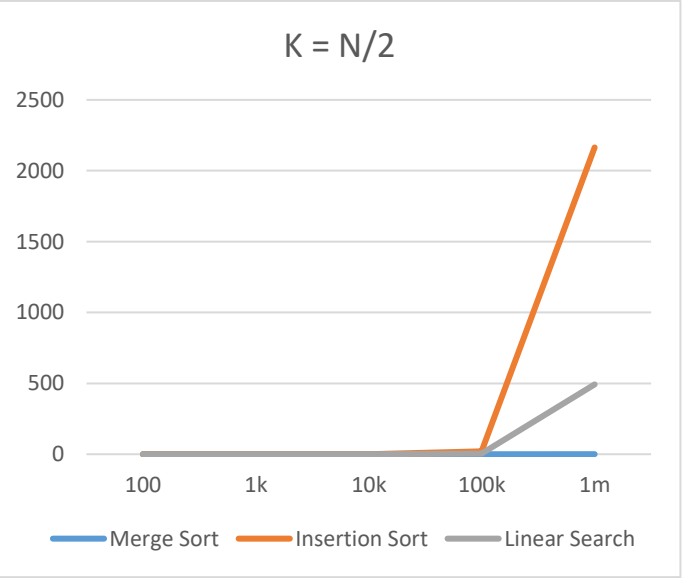
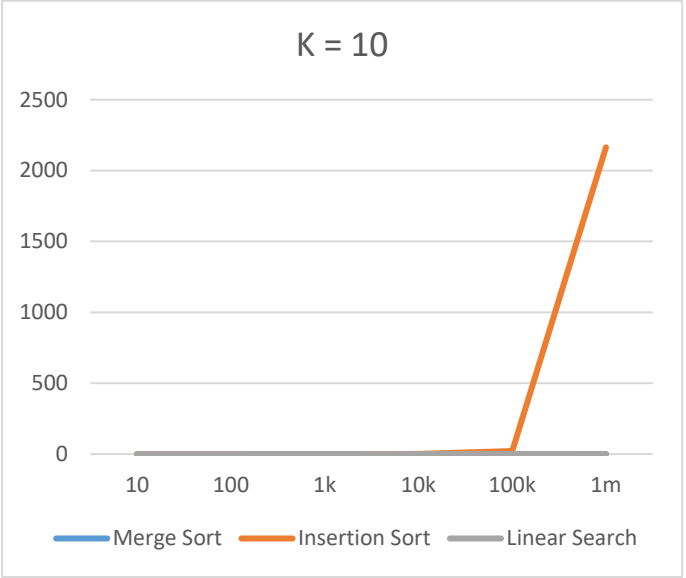
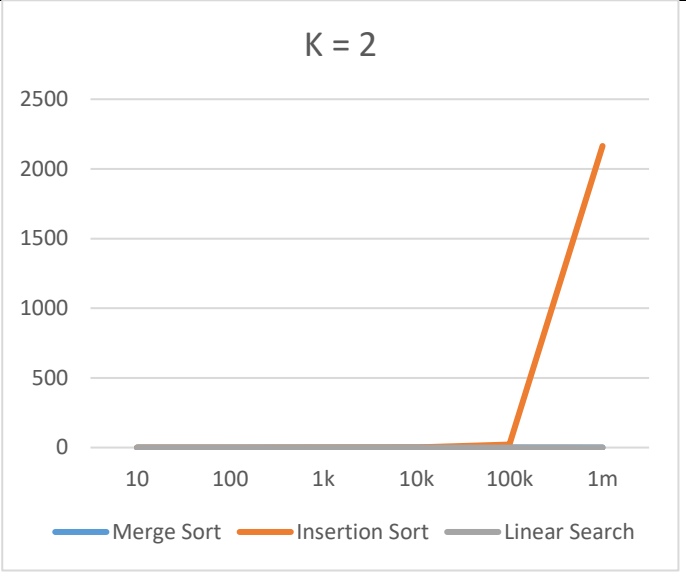
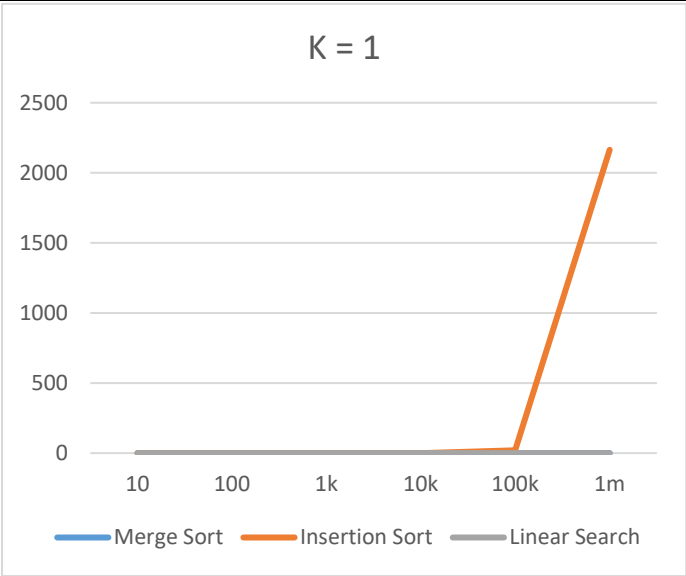
Worst case =>  $N = 1000000$ ,  $K = 500000$

## B – Calculation Times

N	K	Merge Sort	Insertion Sort	Linear Search
10	1	0	0	0
10	2	0	0	0
10	5	0	0	0
100	1	0	0	0
100	2	0	0	0
100	10	0	0	0
100	50	0	0	0
1k	1	0	0.002	0
1k	2	0	0.002	0
1k	10	0	0.002	0
1k	500	0	0.002	0
10k	1	0.005	0.208	0
10k	2	0.005	0.208	0
10k	10	0.005	0.208	0
10k	5k	0.005	0.208	0.047
100k	1	0.049	20.851	0
100k	2	0.049	20.851	0
100k	10	0.049	20.851	0
100k	50k	0.049	20.851	4.696
1m	1	0.511	2164.27	0.003
1m	2	0.511	2164.27	0.003
1m	10	0.511	2164.27	0.003
1m	500k	0.511	2164.27	492.604

\*The unit is second

C – Graph



		ALL																						
		10	100	1k	10k	100k	1m	10	100	1k	10k	100k	1m	10	100	1k	10k	100k	1m	100	1k	10k	100k	1m
Merge Sort		0	0	0	0.005	0.049	0.511	0	0	0	0.005	0.049	0.511	0	0	0	0.005	0.049	0.511	0	0	0.005	0.049	0.511
Insertion Sort		0	0	0.002	0.208	20.85	2164	0	0	0.002	0.208	20.85	2164	0	0	0.002	0.208	20.85	2164	0	0.002	0.208	20.85	2164
Linear Search		0	0	0	0	0	0.003	0	0	0	0	0	0.003	0	0	0	0	0	0.003	0	0	0.047	4.696	492.6

It is clearly seen that Merge Sort is better in these algorithms. Merge Sort is  $O(n \lg(n))$  because it has lower asymptotic function view on the graph.

Until  $N = 10000$  all algorithms almost have same time.  $K$  factor affects only Linear Search algorithm because Linear Search algorithm complexity depends on  $K$  factor.

If  $(n-k)*k$  is small number I select Linear Search Algorithm. Otherwise Merge Sort is better than Linear Search. Insertion Sort is worst algorithm in these algorithms.