

Reducing the Storage Overhead of Main-Memory OLTP Databases with

Hybrid Indexes

SIGMOD'16

Huanchen Zhang

Co-authors:

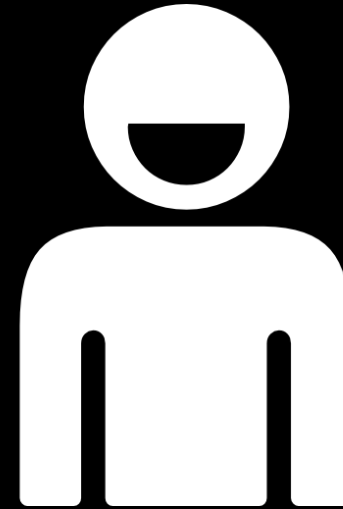
David G. Andersen

Andrew Pavlo

Michael Kaminsky

Lin Ma

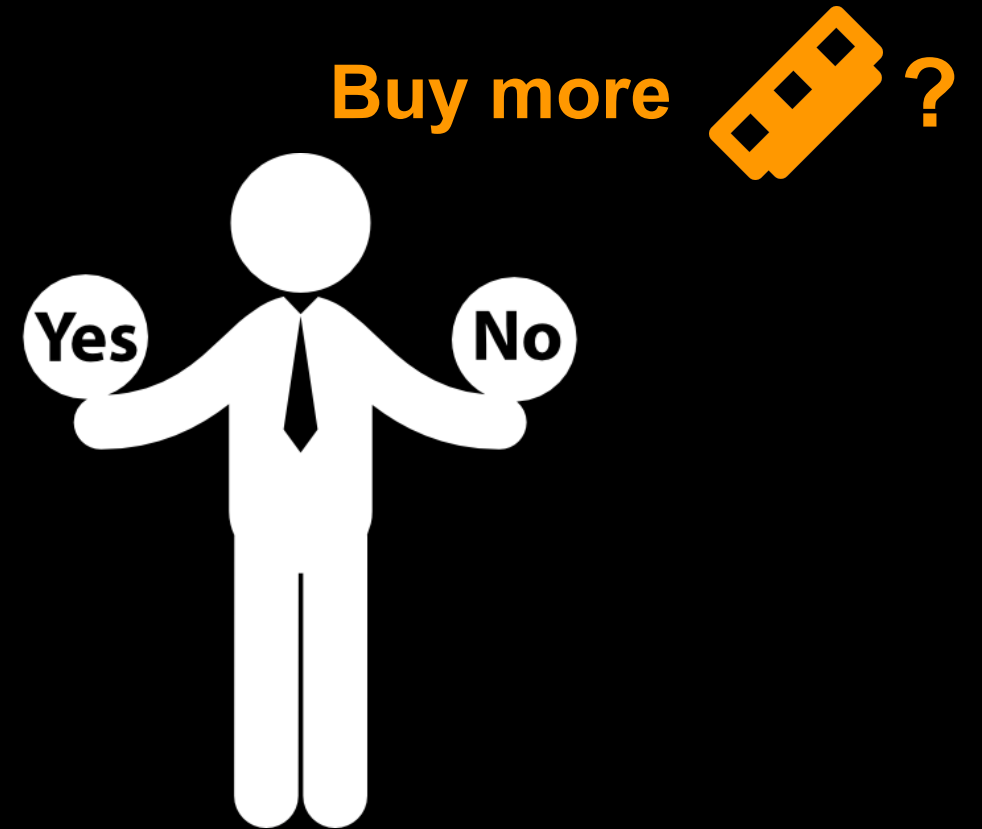
Rui Shen



You are running out of memory



You are running out of memory

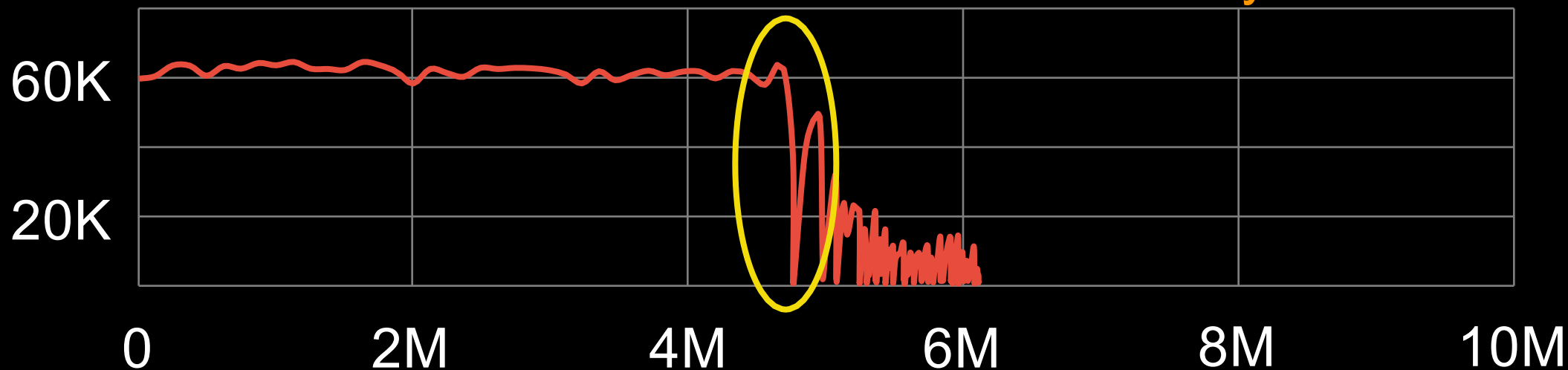


 **You are running out of memory**

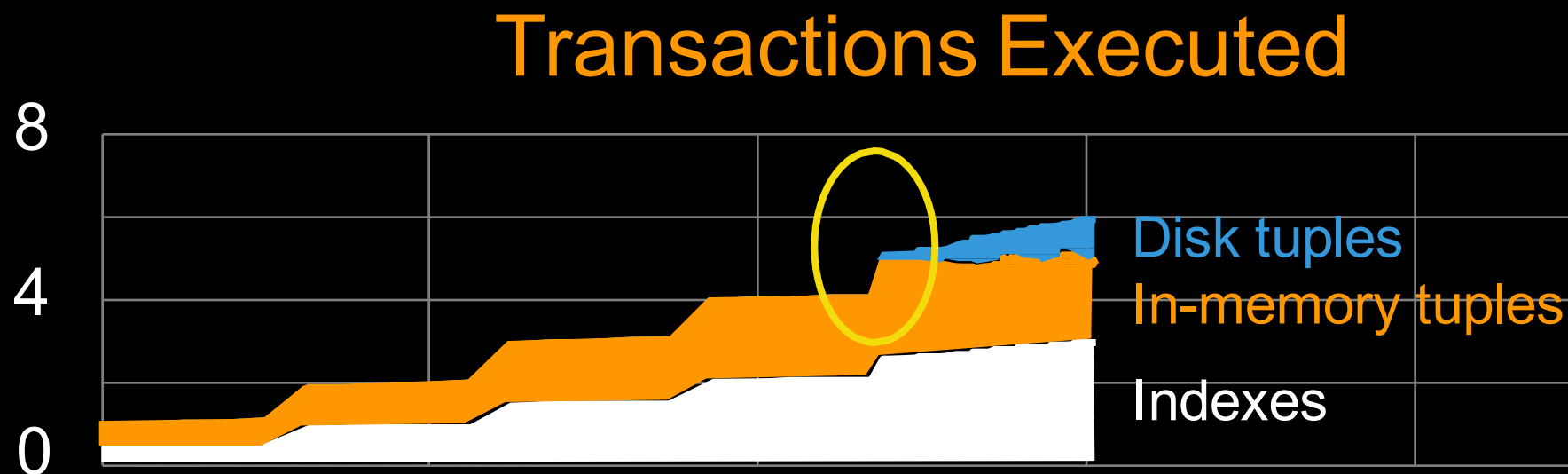
TPC-C on **H**-Store

Memory Limit = 5GB

Throughput



Memory (GB)



I GOT STUCK



SO I WENT TO SLEEP

The better way:
Use memory more efficiently



Indexes are **LARGE**

Benchmark	% space for index		Hybrid Index
TPC-C	58%	→	34%
Voter	55%	→	41%
Articles	34%	→	18%

Our Contributions

- ① The hybrid index architecture
- ② The Dual-Stage Transformation
- ③ Applied to 4 index structures
 - B+tree - Skip List
 - Masstree - Adaptive Radix Tree (ART)

Performance



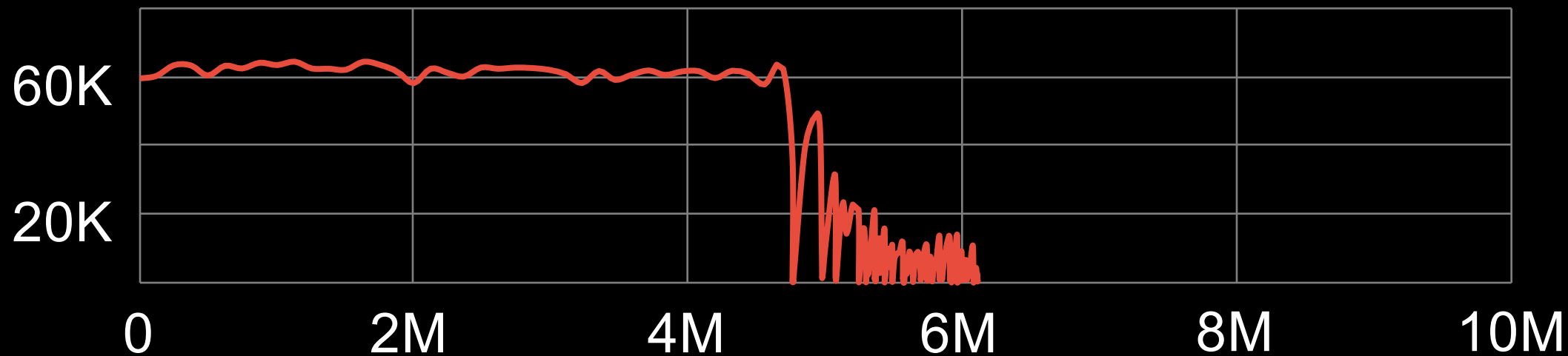
Space

30 – 70% ↓

TPC-C on **H**-Store

Throughput (txn/s)

Did we solve this problem?



Transactions Executed

How do hybrid indexes achieve
memory savings ?

0th Static

Hybrid Index: a dual-stage architecture

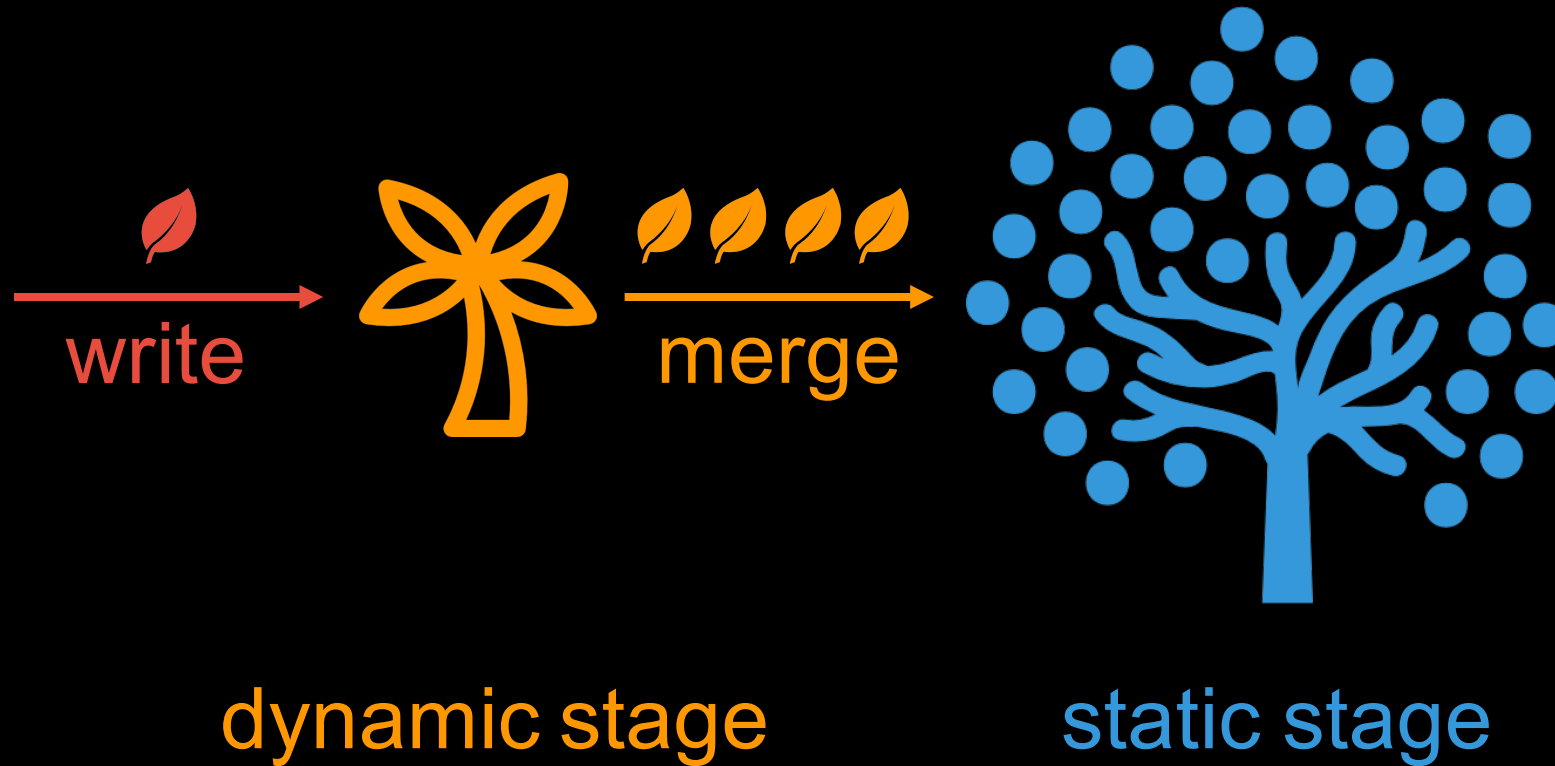


dynamic stage

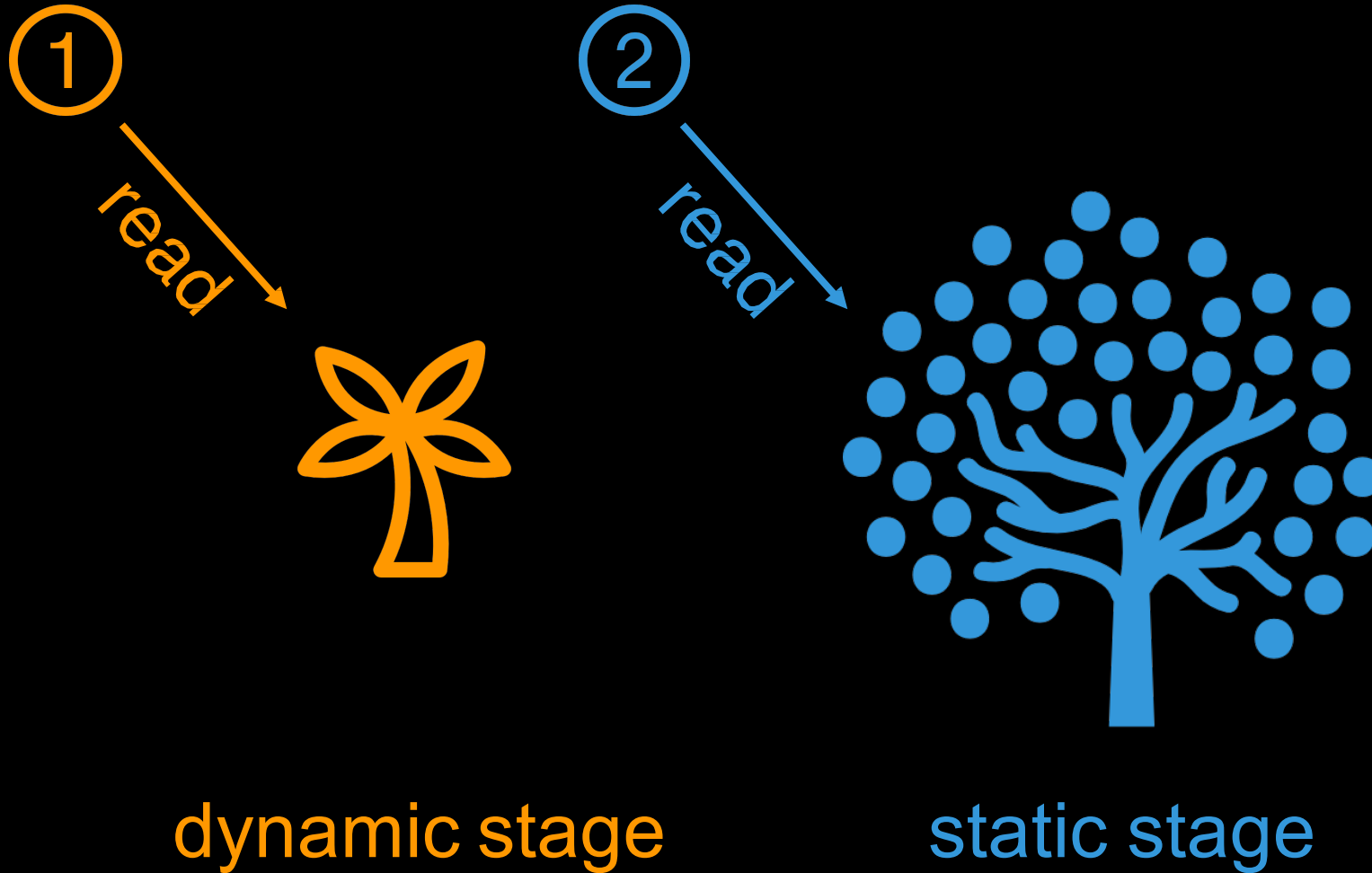


static stage

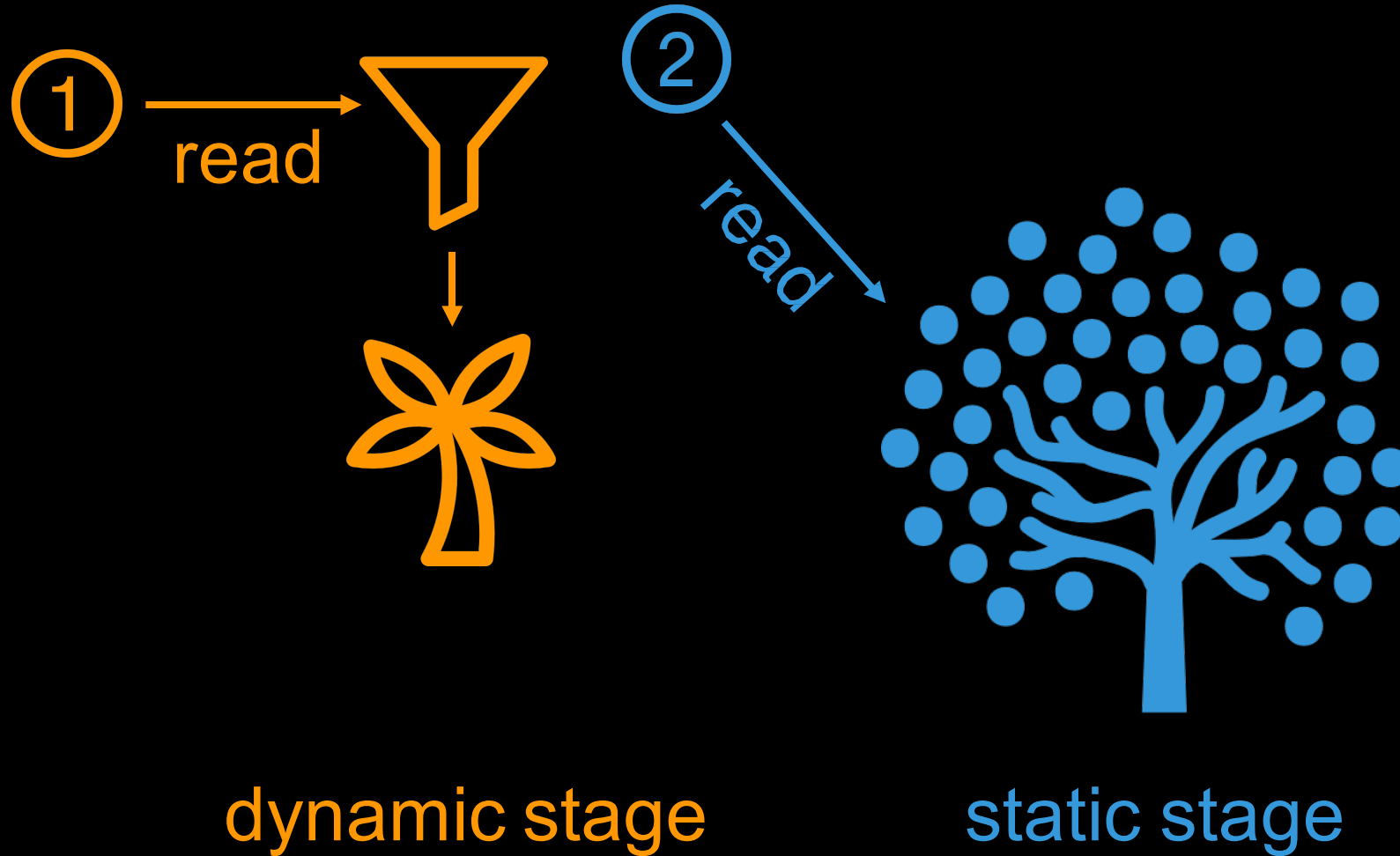
Inserts are batched in the dynamic stage

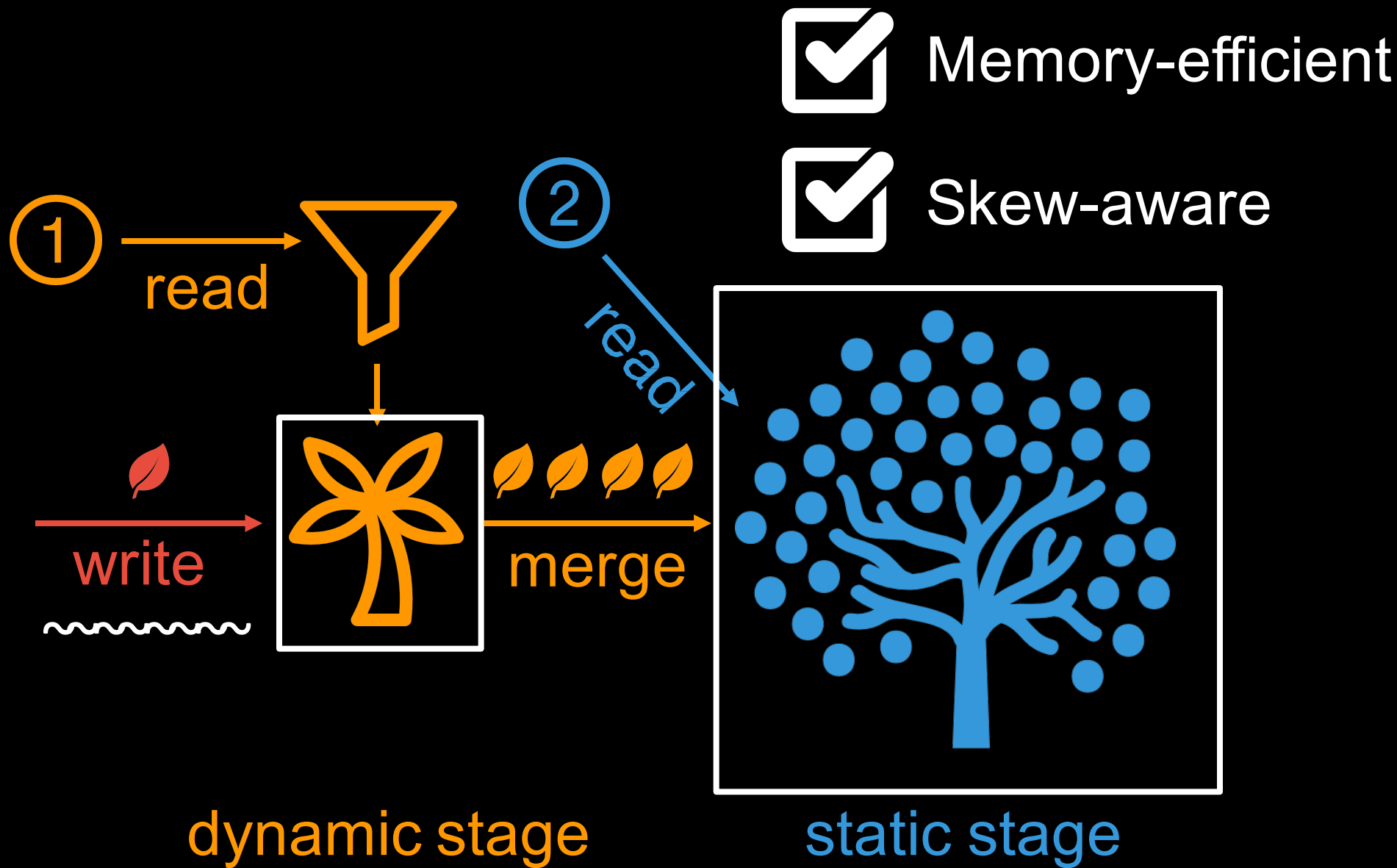


Reads search the stages in order

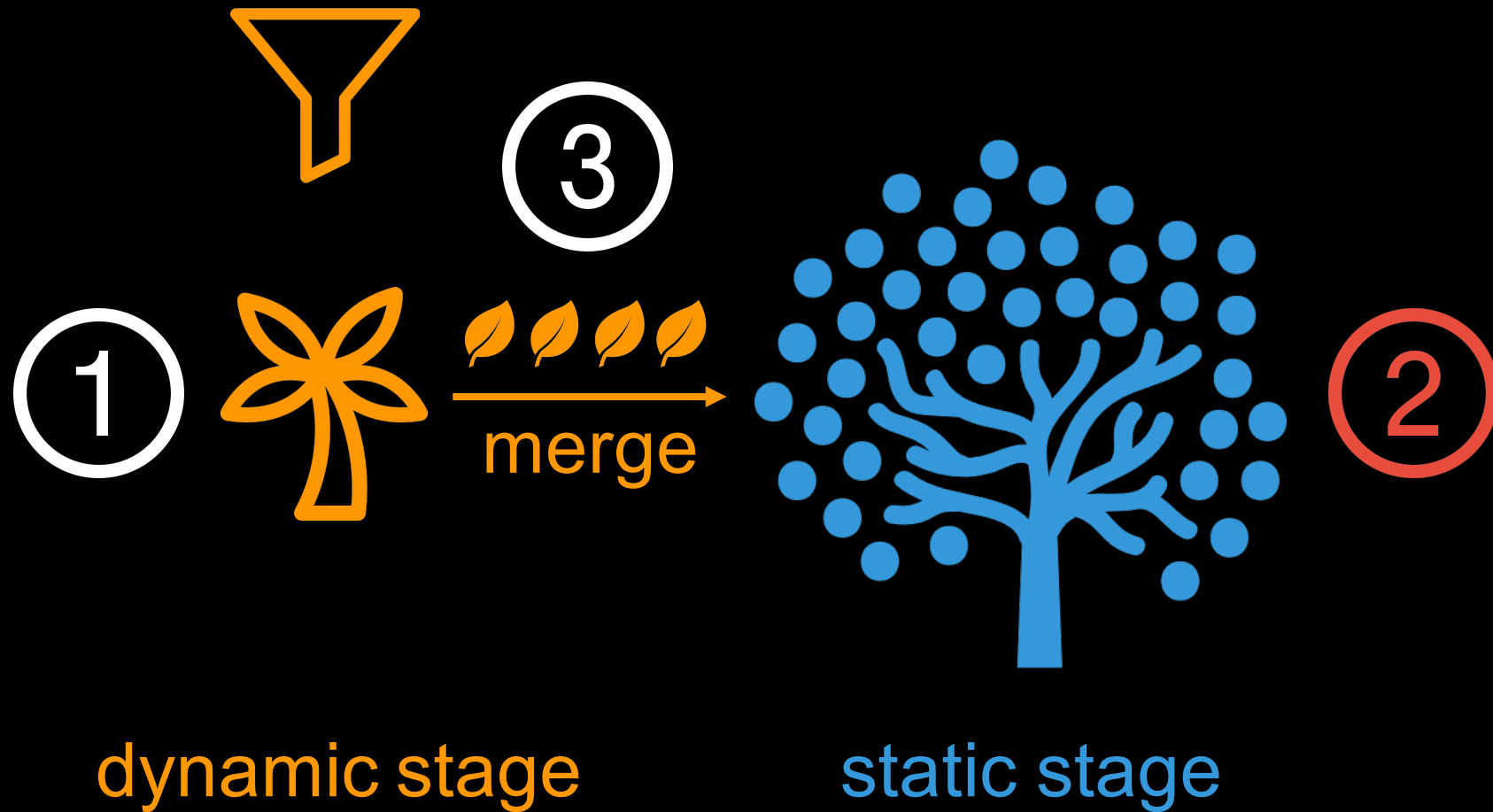


A Bloom filter improves read performance





The Dual-Stage Transformation



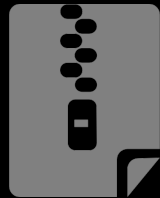
The Dynamic-to-Static Rules



Compaction

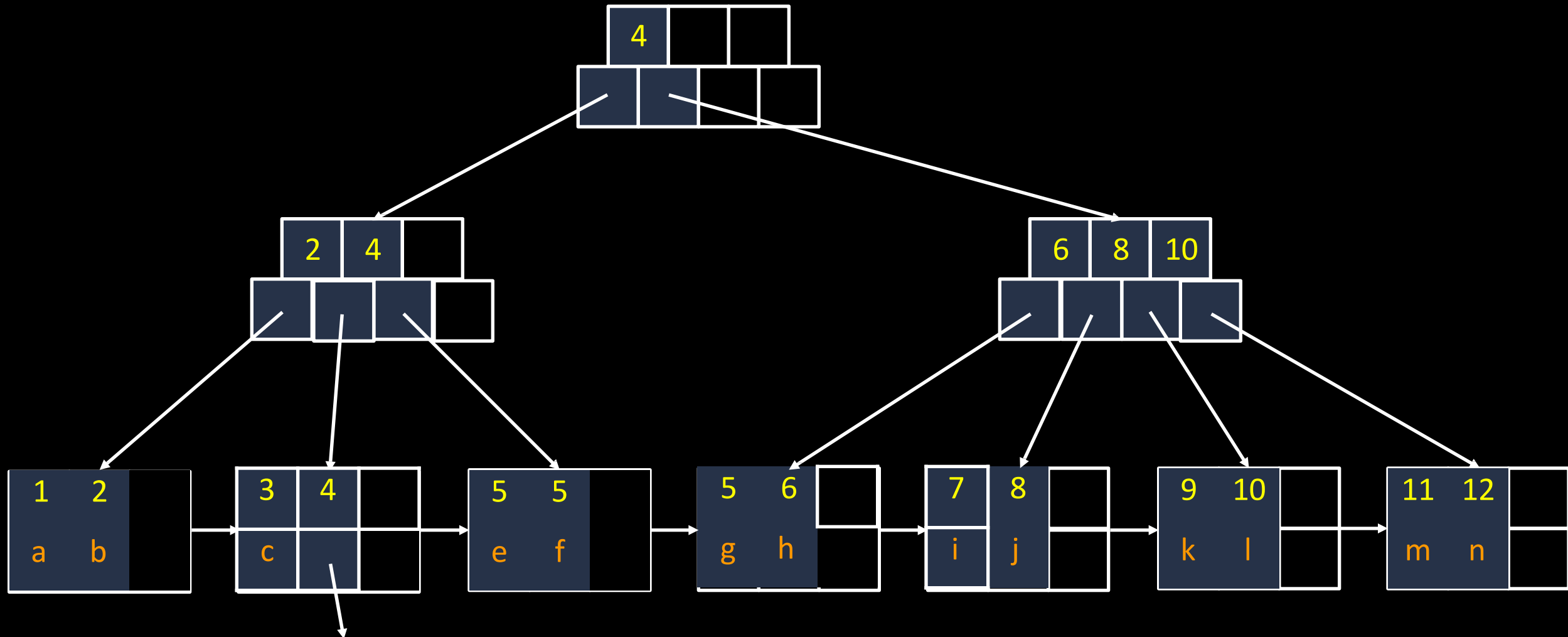


Reduction

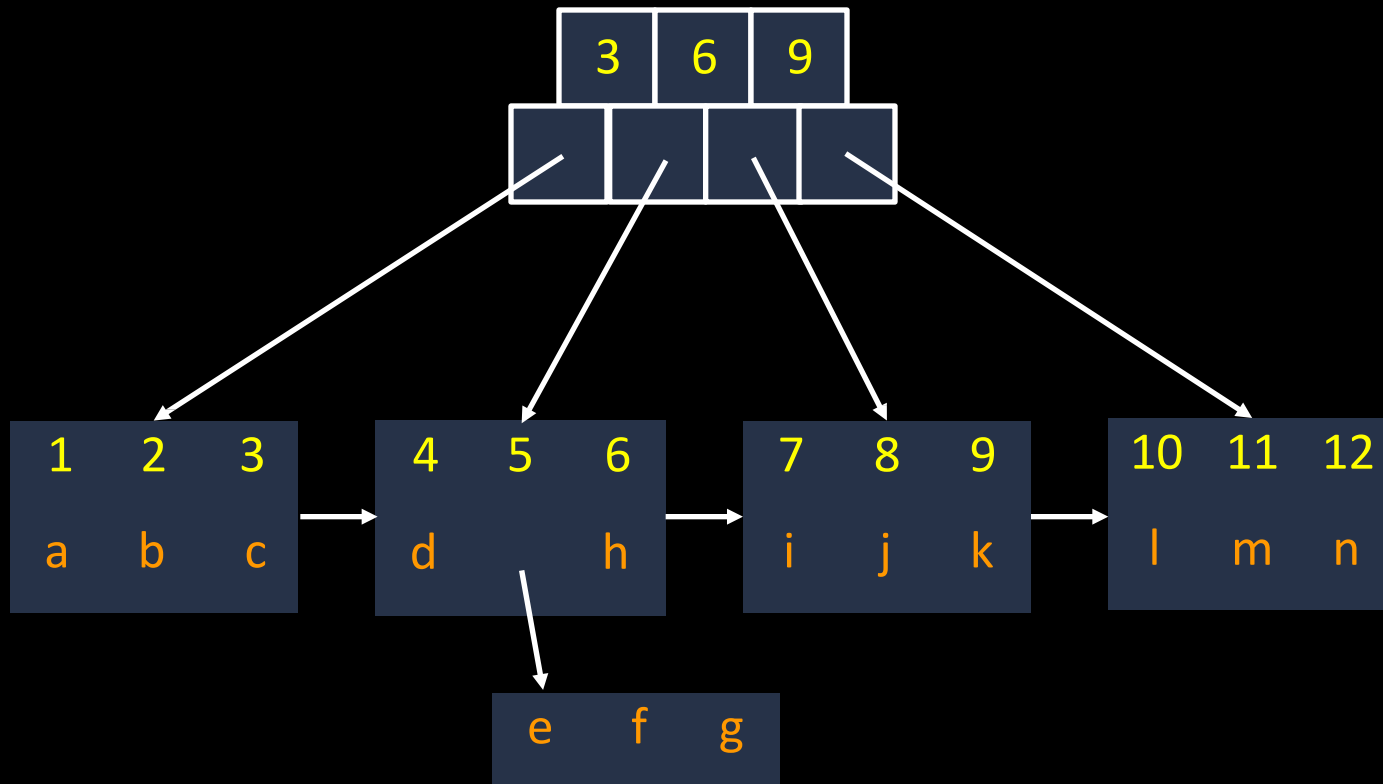


Compression

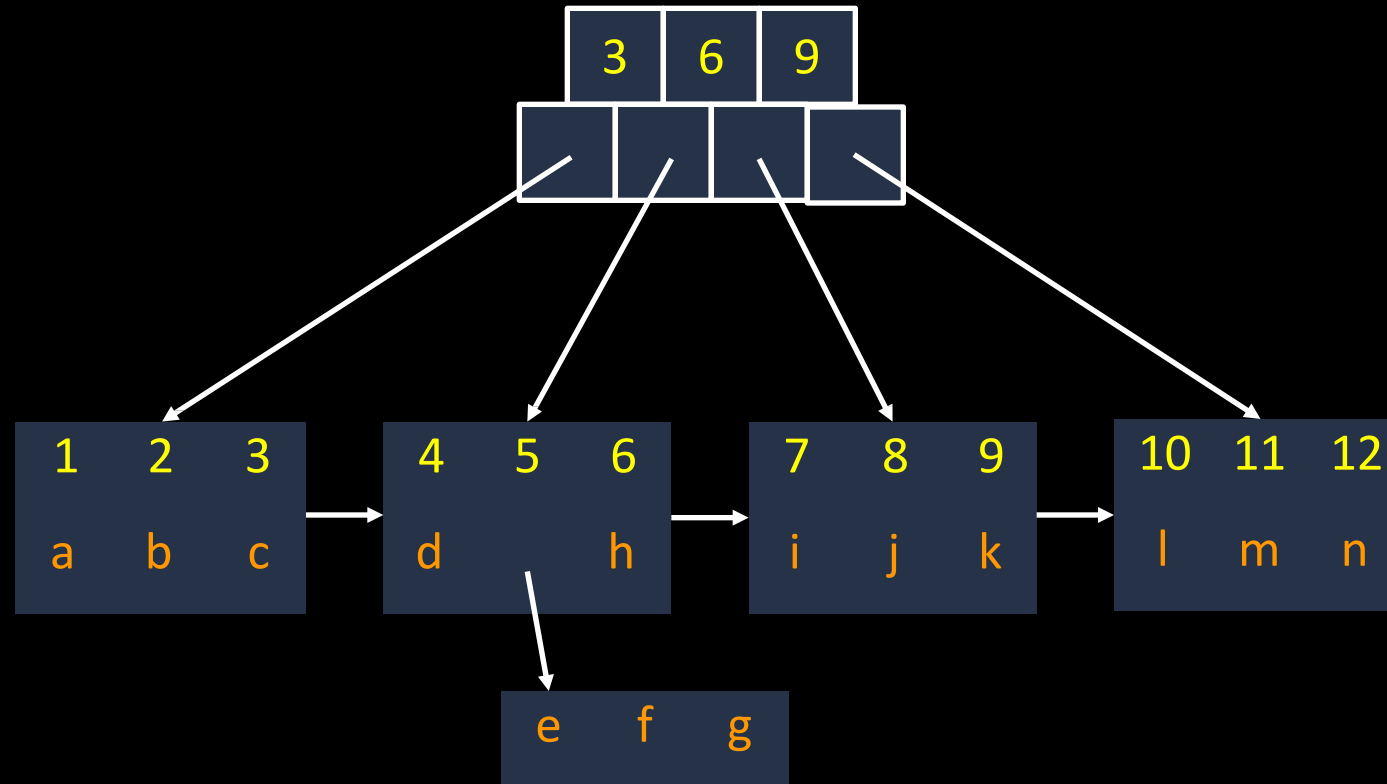
Compaction: minimize # of memory blocks



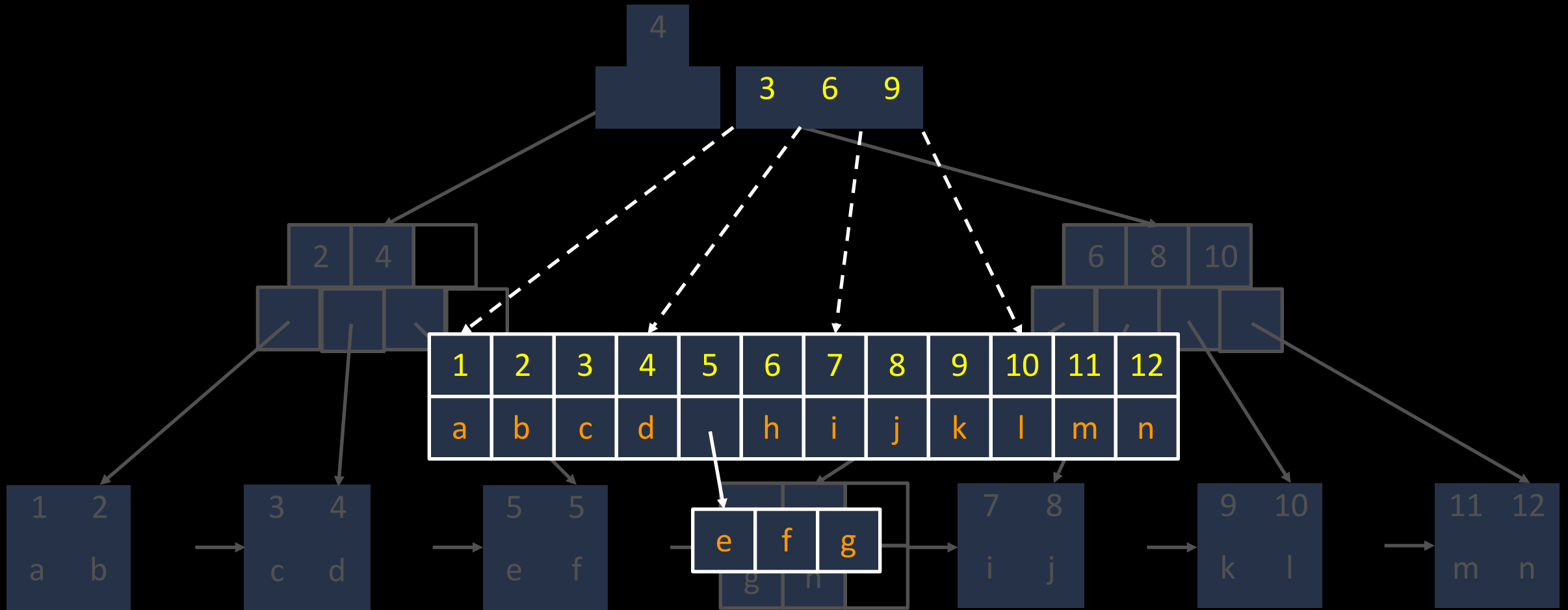
⚡ Compaction: minimize # of memory blocks



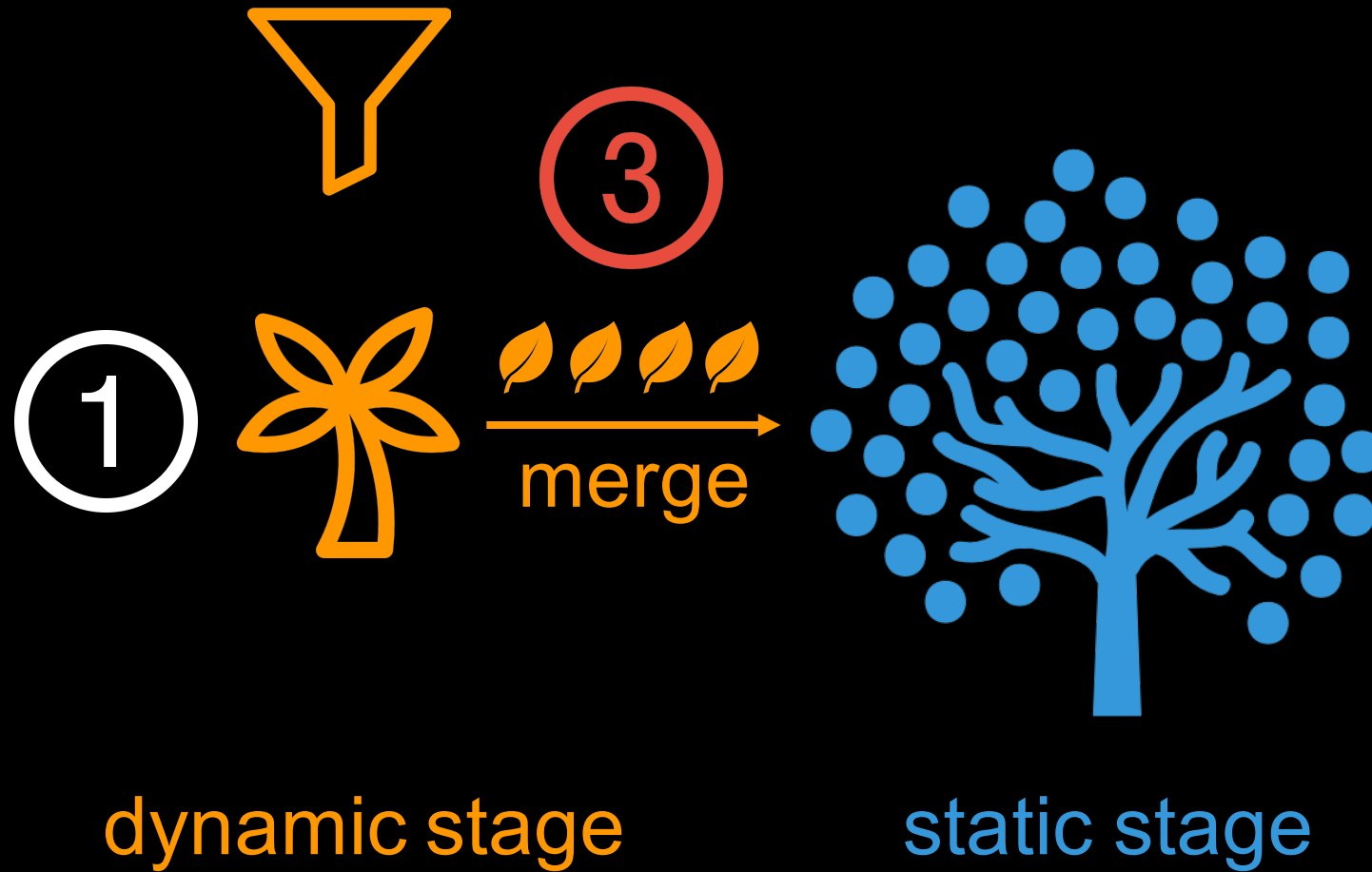
✂ Reduction: minimize structural overhead



✂ Reduction: minimize structural overhead



The Dual-Stage Transformation



The Dual-Stage Transformation



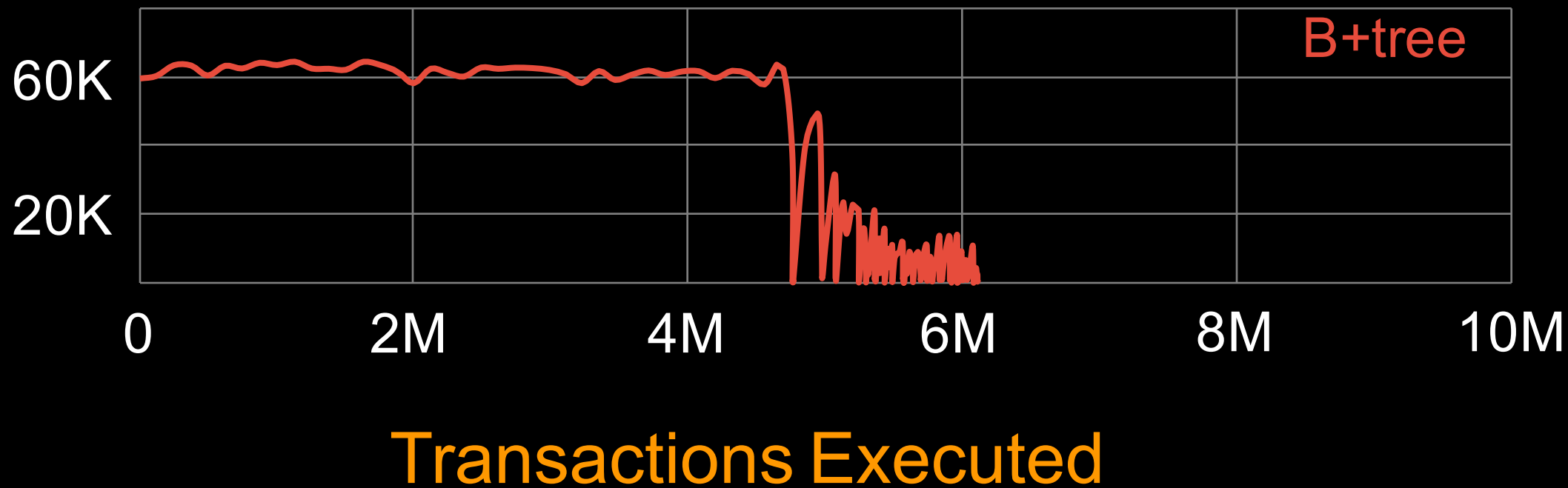
dynamic stage

Merge Questions:

1. Partial? 🍃 🍃 ?
2. When? ⌚ ?
3. Blocking? 🔒 ?

static stage

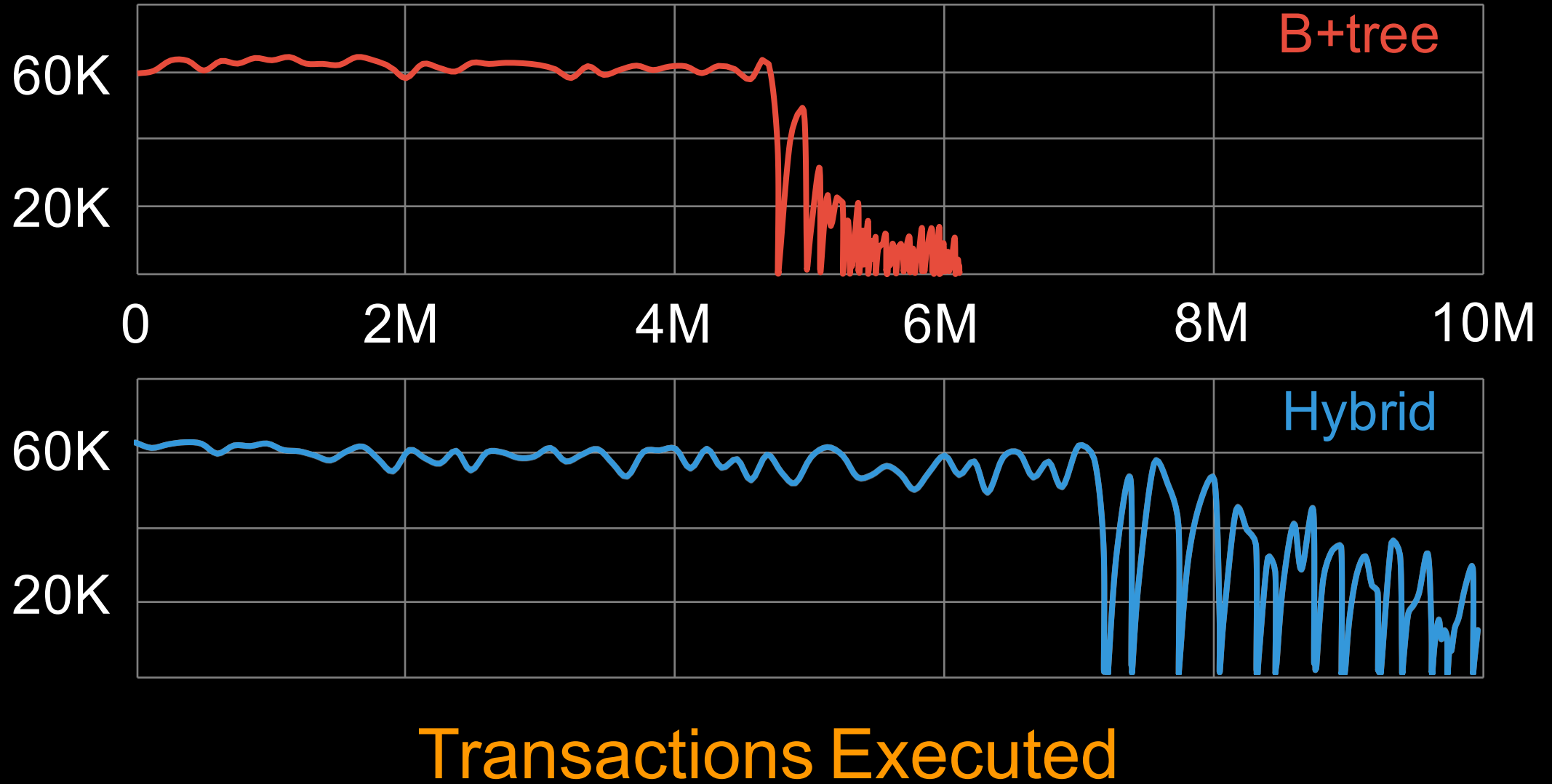
Did we solve this problem?



Yes, we improved the DBMS's capacity!

TPC-C on **H**-Store

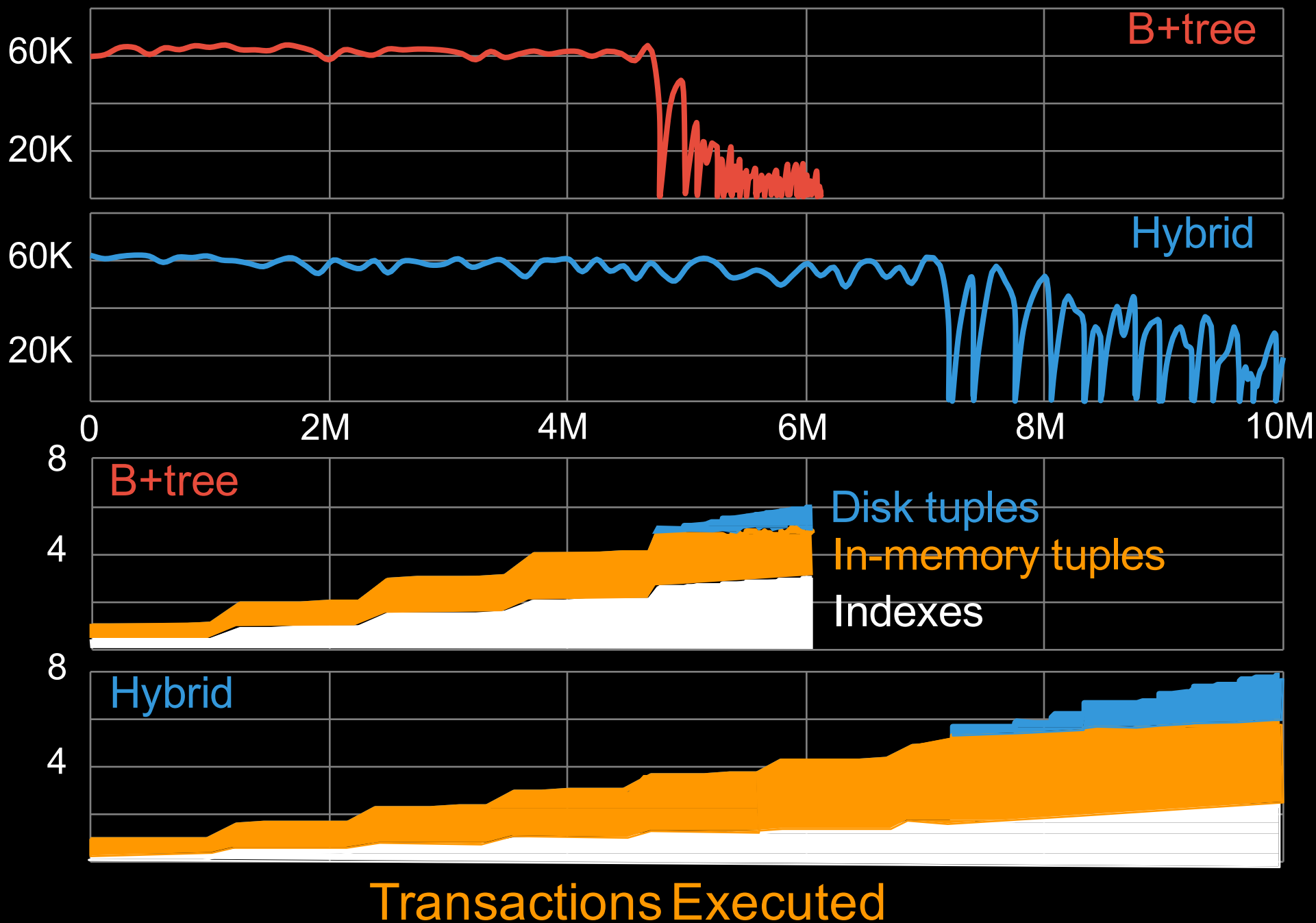
Throughput (txn/s)



TPC-C on H-Store

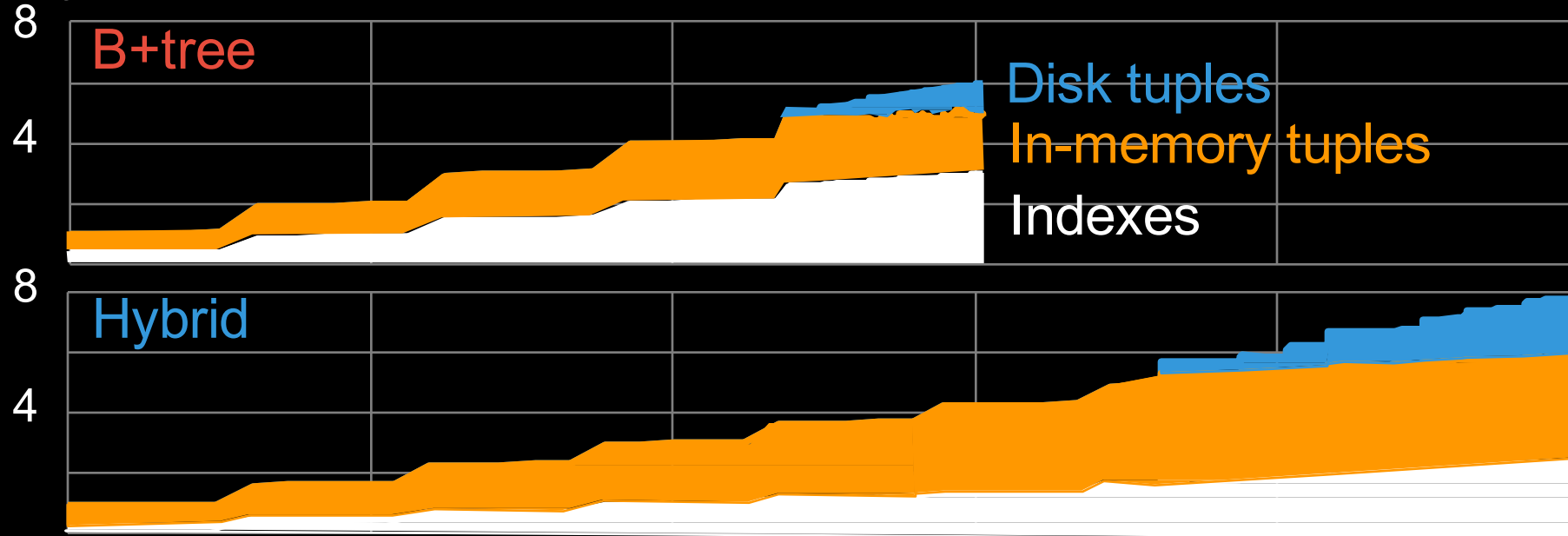
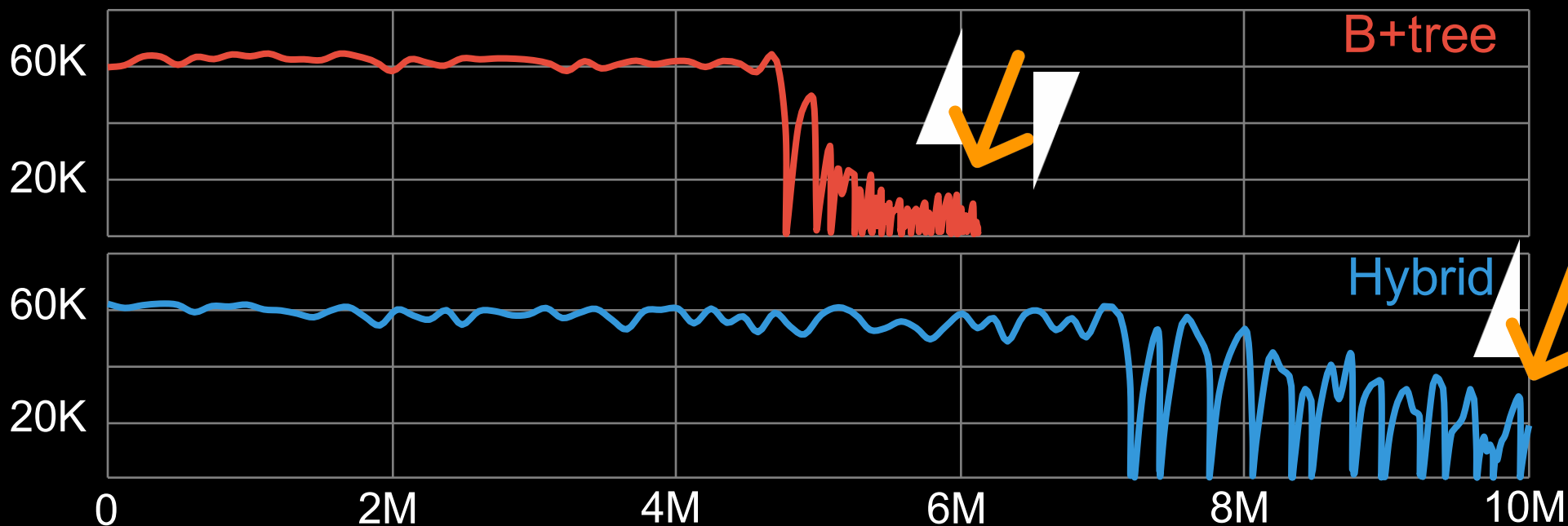
Throughput (txn/s)

Memory (GB)



TPC-C on H-Store

Throughput (txn/s)

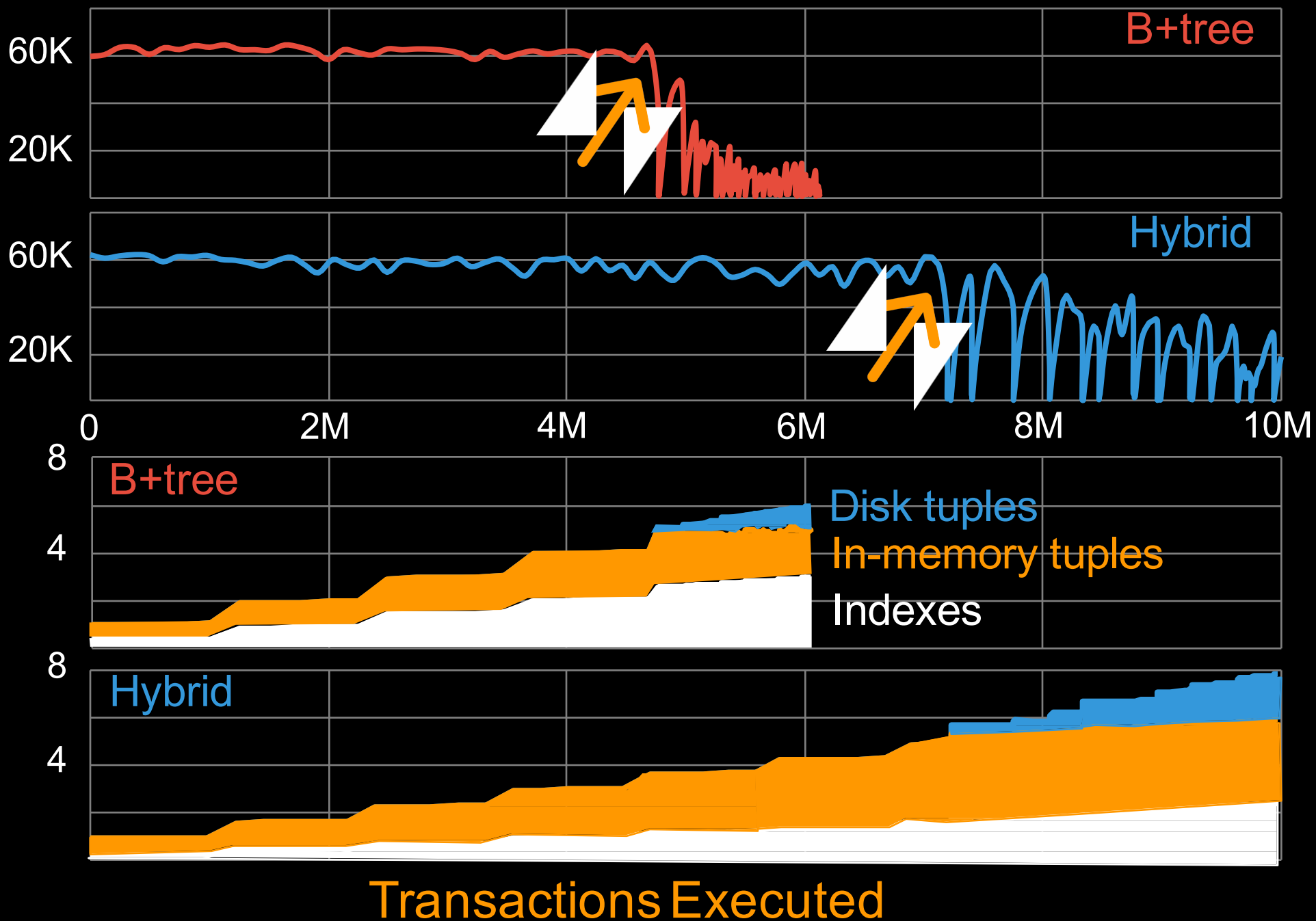


Transactions Executed

TPC-C on H-Store

Throughput (txn/s)

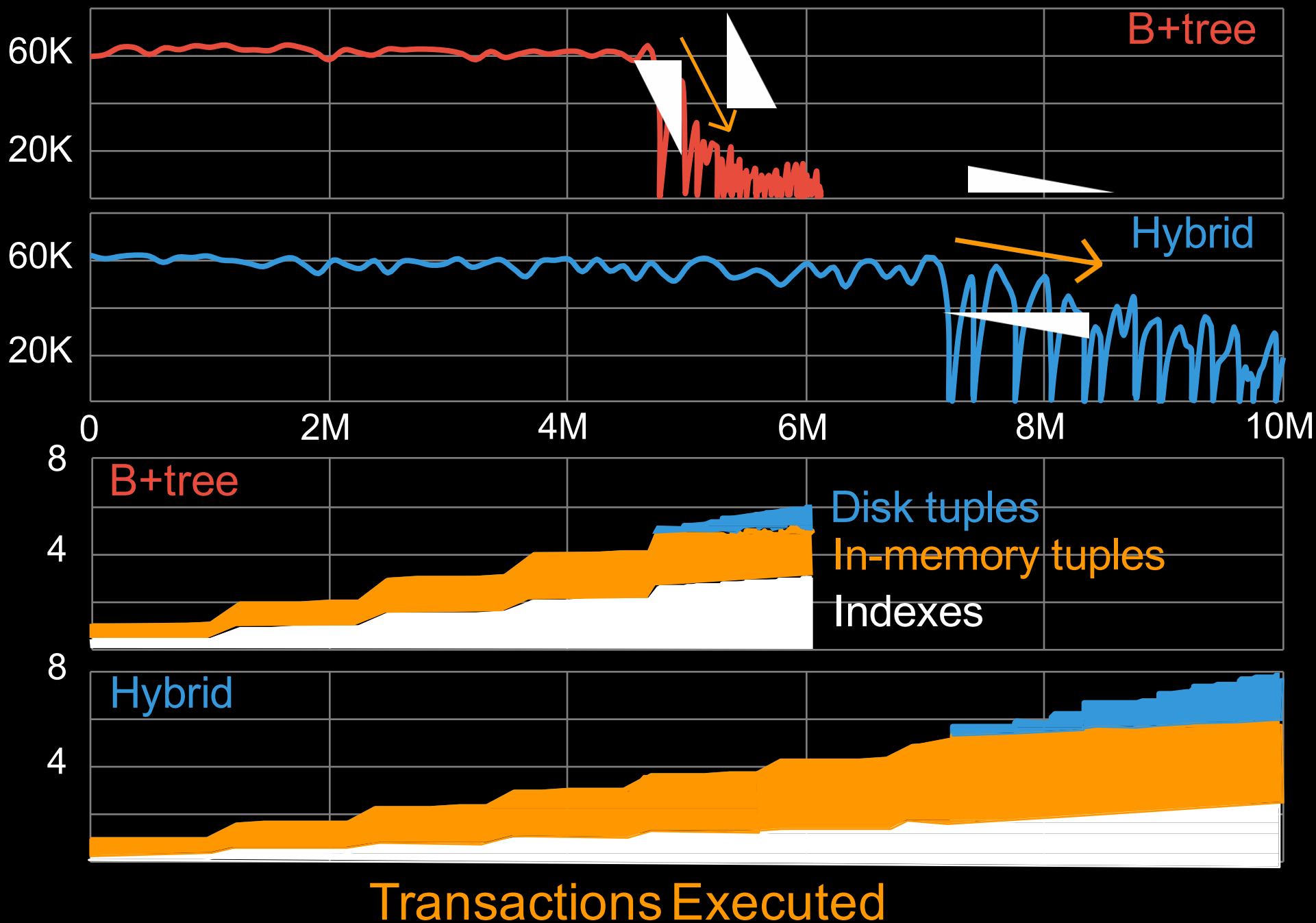
Memory (GB)



TPC-C on H-Store

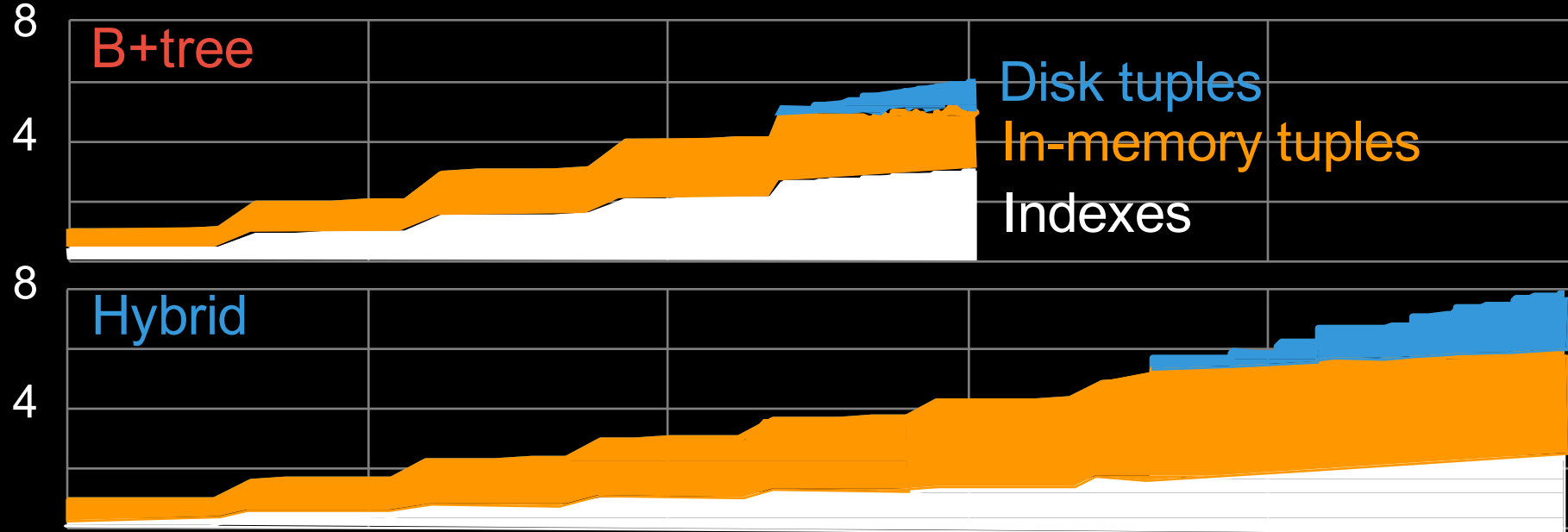
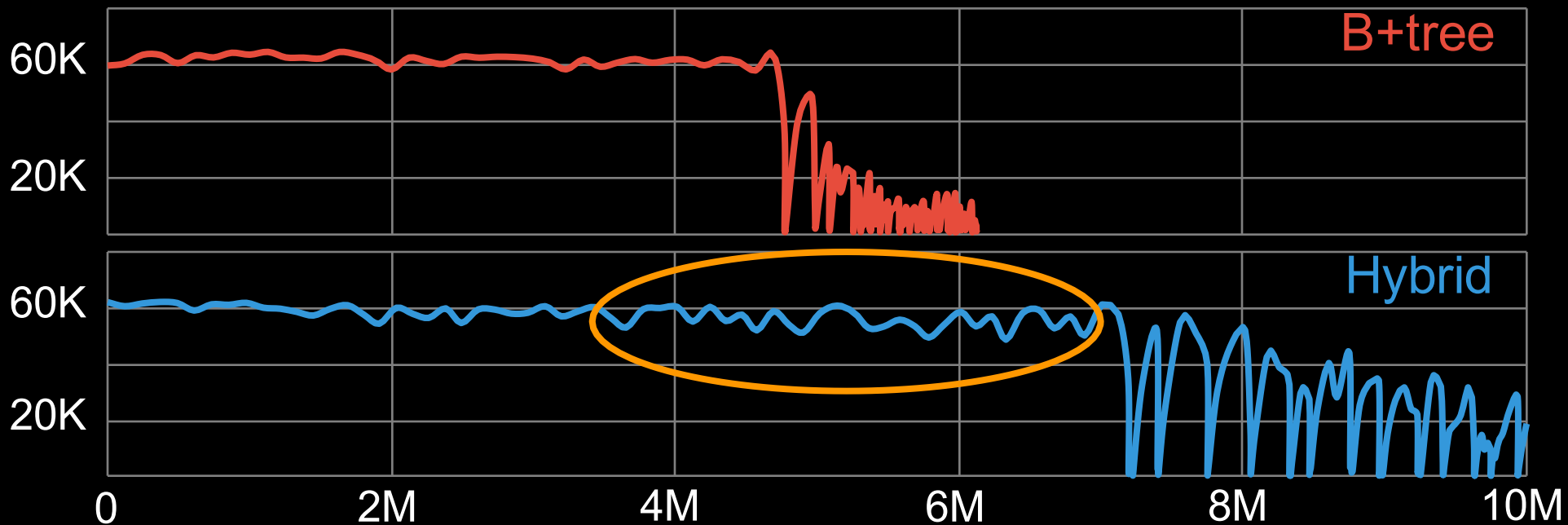
Throughput (txn/s)

Memory (GB)



TPC-C on H-Store

Throughput (txn/s)



Transactions Executed

Take Away:

Memory saved
by indexes



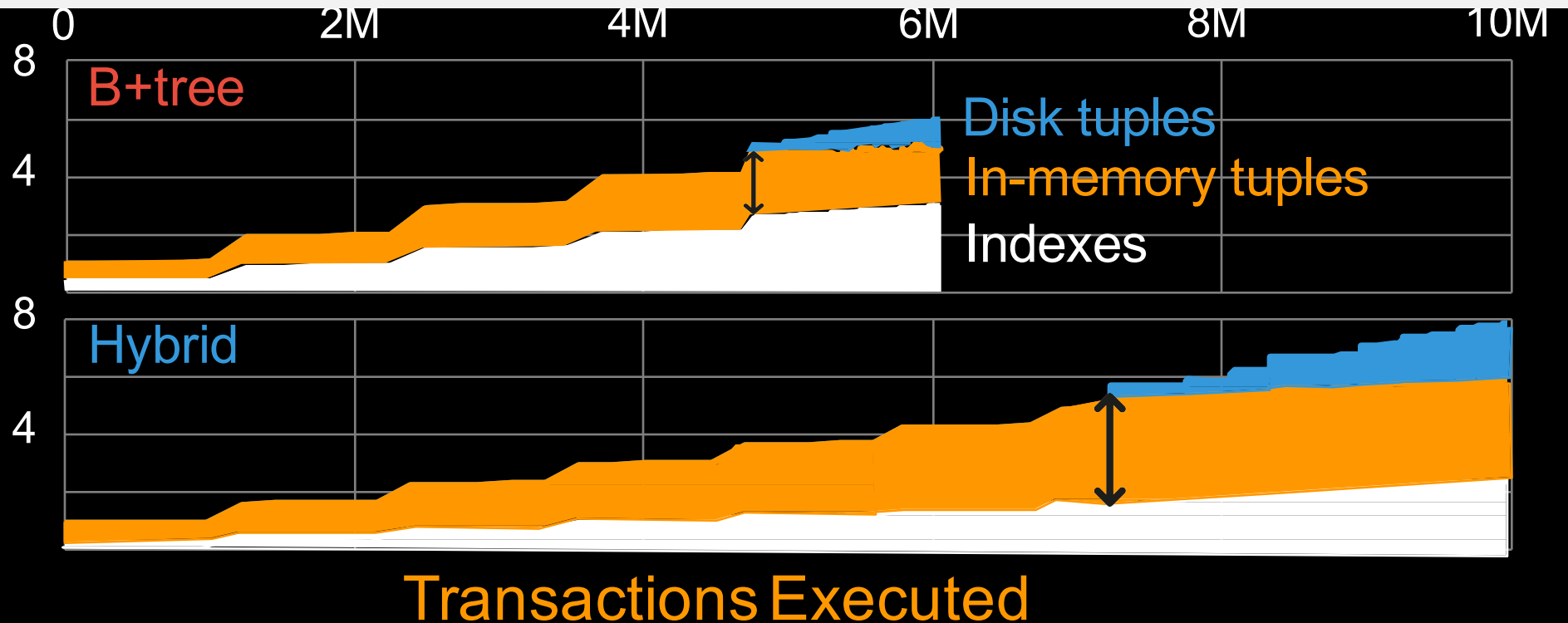
Larger working
set in memory



Higher
throughput

TPC-C on

Memory (GB)



This is just the **BEGINNING**

Conclusions

- ① The hybrid index architecture **GENERAL**
- ② The Dual-Stage Transformation **PRACTICAL**
- ③ Applied to 4 index structures **USEFUL**
 - B+tree
 - Skip List
 - Masstree
 - Adaptive Radix Tree (ART)

Hybrid Index Inspirations:

1. A Tradeoff Research among Data Tuples, Indexes and Evicted Tables in Memory Consumption
2. Dual-stage Architecture of Indexes in Hybrid Memory
3. A Memory-efficient Hash Table Index with Range Query Optimization (DHT might be possible)
4. Non-blocking Merging for Hybrid Indexes with COW

Anti-caching Inspirations:

- 1. Revisit Anti-caching Mechanism for OLTP Workloads in Hybrid Memory (Evict cold tuples to NVM)**
- 2. Non-blocking Eviction in NVM-Optimized Anti-caching**

Non-volatile Memory Inspirations:

- 1. DiRedis: A NVM-Optimized KV-Store Based on Redis**
- 2. Rethinking Program Scheme in Persistent Memory Era**

Thanks !

Note website:

<http://kaixinhuang.com/Research/hybrid-index-db/>

The Art of Research Presentation