KOOBE: Towards Facilitating Exploit Generation of Kernel Out-Of-Bounds Write Vulnerabilities

W. Chen et al., USENIX Security '20

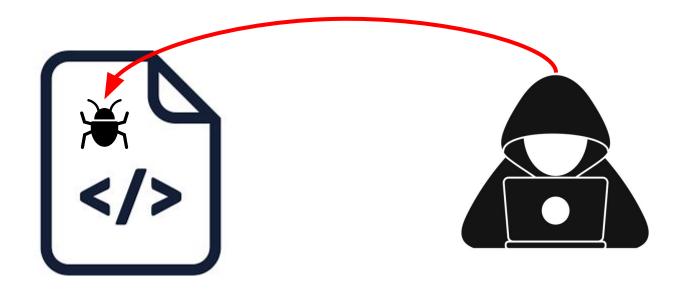
Presenter: Hyunsu Kim



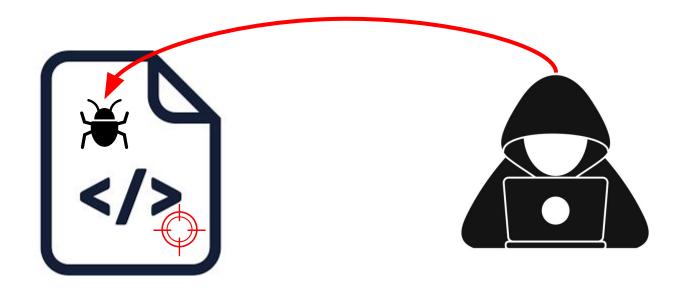
Vulnerability



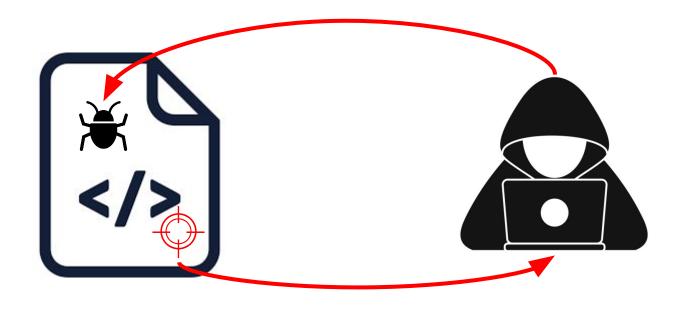
Exploitability



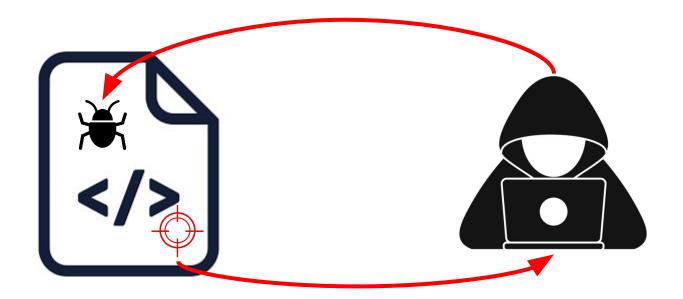
Vulnerability



Vulnerability



Vulnerability

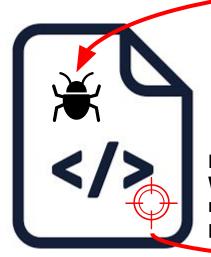


Vulnerability

- Use-After-Free (UAF)
- Out-of-Bound Write (OOB)

- Control-flow hijacking
- Privilege escalation

About 1,000 bugs fixed annually



Missing: What really matters in prioritization



Vulnerability

- Use-After-Free (UAF)
- Out-of-Bound Write (OOB)

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Missing: What really matters in prioritization



Vulnerability

- Use-After-Free (UAF)
- Out-of-Bound Write (OOB)

- Control-flow hijacking
- Privilege escalation

Example

```
1.
     struct Type1 { ..; };
     struct Type2 { Type1 sk; uint64 t option; ..; };
     struct Type3 { int (*ptr)(); ..; };
     Type1 * vul = NULL; Type3 * tgt = NULL;
     void sys socket() //sizeof(Type1) == sizeof(Type3)
 7.
         vul = kmalloc(sizeof(Type1));
 8.
     void sys accept()
         vul = (Type2*)vul;
                                   //type confusion
         vul->option = gsock.option; //vulnerability pt
13.
     void sys setsockopt(val) //not invoked in given PoC
14.
         if (val == -1) return;
15.
         gsock.option = val;
16.
17.
     void sys create tgt()
         tgt = kmalloc(sizeof(Type3));
18.
19.
         tgt->ptr = NULL;
                            //init ptr
20.
21.
     void sys deref() { if (tgt->ptr) tgt->ptr(); }
```

Struct definition Initialization

System calls

Possible exploit

sys socket();

sys create tgt();

for (*) { sys create tgt(); } // cache exhaustion

// vuln obj

// target obj

Example

```
sys setsockopt(0xdeadbeef);
                                         sys accept();  // tgt->ptr = 0xdeadbeef
     struct Type1 { ..; };
                                     6.
                                         sys deref();
 2. struct Type2 { Type1 sk; uint64 t
     struct Type3 { int (*ptr)(); ..; };
     Type1 * vul = NULL; Type3 * tqt = NULL;
     void sys socket() //sizeof(Type1) == sizeof(Type3)
 7.
       vul = kmalloc(sizeof(Type1));
 8.
     void sys accept()
      vul = (Type2*)vul; //type confusion
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Possible exploit

sys socket();

sys create tgt();

for (*) { sys create tgt(); }

Example

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sys setsockopt(0xdeadbeef);
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```

```
// tgt->ptr = 0xdeadbeef
                         Vulnerable Object
Memory Object
                         Target Object
          Overwritten Data
```

6 Critical data

Candidate:

6 Target Offset

6 Desired Payload

► **①** Size

// cache exhaustion

// vuln obj

// target obj

Heap feng shui

Capability: 0

OOB Offset ___

3 OOB Length &

OOB Value

O Size

Design & Implementation

- Capability summarization & exploration
 - Capability = (Vulnerability point * OOB write) list
 - One vulnerability may result in several capabilities
 - From one PoC, populate more capabilities (Def-Use manner)
- 2. Symbolic tracing
 - Together with Kernel addresssanitzer (KASAN)
 - Avoid known crash in while fuzzing (performance)
 - Exploitability evaluation w/ approximation
- 3. Exploit synthesis
 - Heap feng shui

Evaluation

Among 28 Out-of-bound Write vulnerabilities, KOOBE was able to run against 17 (10) of them. Finally, was able to generate working exploit of 11 (5) vulnerabilities.

(*): # of Non-CVEs (found by syzkaller)

Questions