Control-Flow Integrity
principles, implementations, and Applications

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Introduction

- Control Flow Integrity(CFI) tries to prevent attacks from arbitrarily controlling program behavior
- Adopts binary instrumentation to enforce CFI on Windows x86
- Compatible with existing software and simple to enforce with low overhead.

Mitigations

- StackGuard(USENIX `98)
  - Buffer Overflow Detector by inserting random value
- CRED(NDSS `04)
  - Runtime Elimination of Buffer Overflows
- Secure Program Execution via Dynamic Information Flow Tracking(ASPLOS `04)
  - Tainting of suspect Data

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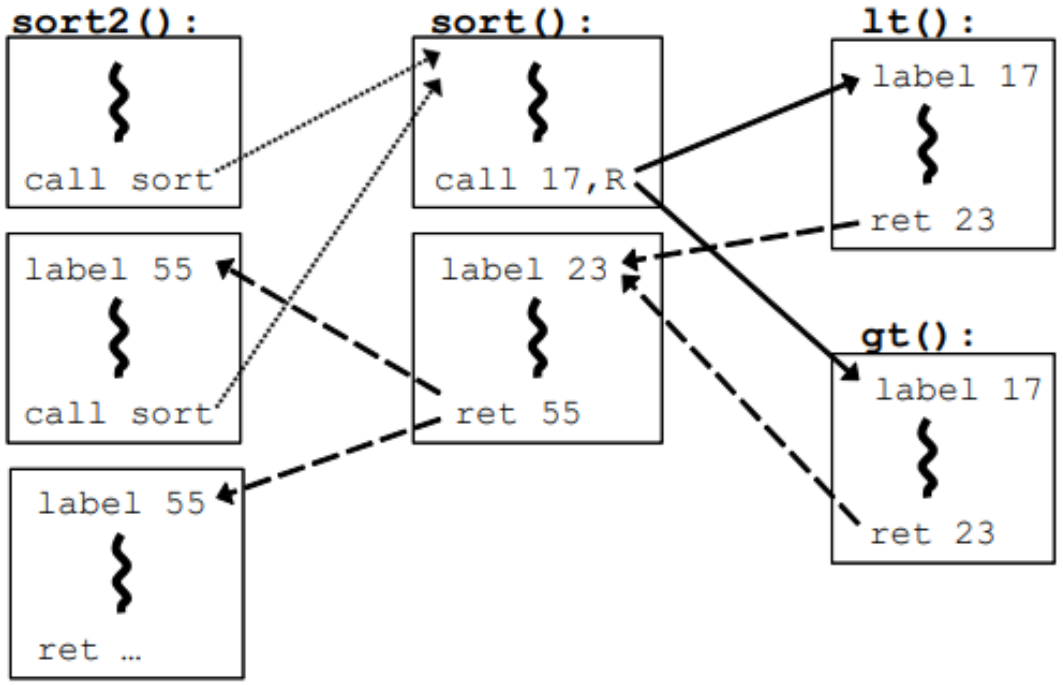
Hard to catch practicalness & performance

Is call/ret targets a valid destination?

```
bool lt(int x, int y) {
    return x < y;
}

bool gt(int x, int y) {
    return x > y;
}

sort2(int a[], int b[], int len)
{
    sort( a, len, lt );
    sort( b, len, gt );
}
```



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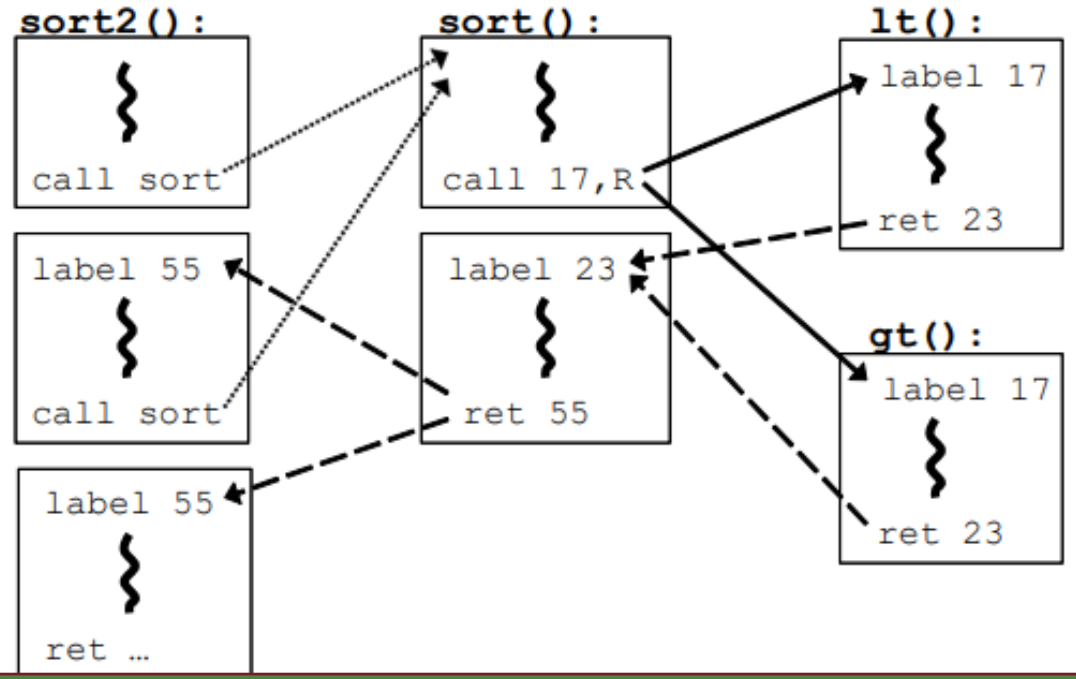

Can be determined by a Control Flow Graph(CFG) ☺

Control Flow Integrity

- Software execution must follow a path of a CFG determined ahead of time.

- The CFG can be defined by static (source code, binary) analysis.

CFI Enforcement

- At each destination, instrumentation inserts a bit pattern, or ID
  - Use same bit pattern for equivalent destination

- Also Insert check routine to ensure the runtime destination has the ID or bit pattern.

Assumptions

- Unique IDs
  - After CFI instrumentation, the bit pattern must not be present anywhere in the code memory except in ID's and ID-checks
- Non-Writable Code
  - Modifying code memory at runtime is not allowed
- Non-Executable Data
  - It's not possible to execute data as if it were code.

CFI Instrumentation of ‘call’ and ‘ret’

Function Call

Opcode bytes	Instructions	
FF 53 08	call [ebx+8]	; call fptr

Function Return

Opcode bytes	Instructions	
C2 10 00	ret 10h	; return

are instrumented using `prefetchnta` destination IDs, to become

8B 43 08	mov eax, [ebx+8]	; load fptr
3E 81 78 04 78 56 34 12	cmp [eax+4], 12345678h	; comp w/ID
75 13	jne error_label	; if != fail
FF D0	call eax	; call fptr
3E 0F 18 05 DD CC BB AA	prefetchnta [AABBCCDDh]	; label ID

8B 0C 24	mov	ecx, [esp]	; load ret
83 C4 14	add	esp, 14h	; pop 20
3E 81 79 04	cmp	[ecx+4],	; compare
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FF E1	jmp	ecx	; jump ret

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Load Function Pointer and
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Load Return Pointer and
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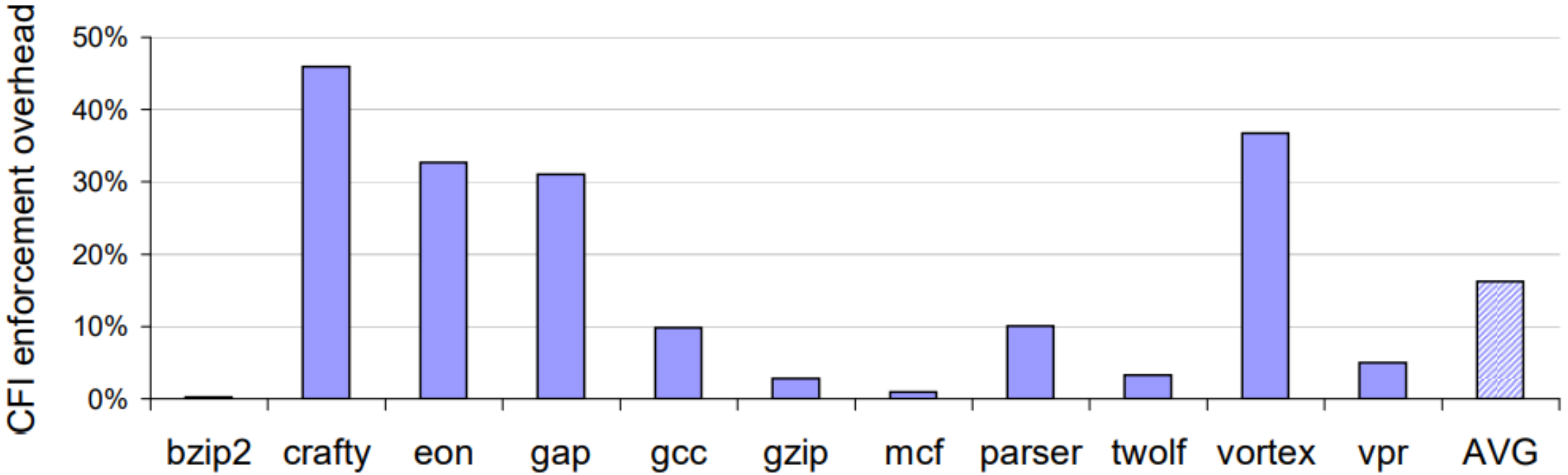
Evaluation Setup

- Windows XP SP2 in “Safe Mode”
  - Most daemons and kernel modules are disabled
- Pentium 4 x86 processor with 512 MB RAM
- Target binaries were compiled with MS Visual C++ 7.1 using full optimizations

Execution overhead of inlined CFI



Measurements

- CFG construction + CFI instrumentation = 10 sec
- Binary increasing = 8%
- Overhead took 0~45%

- This is competitive with the cost of most comparable technique.
  - CRED: up to 130%
  - PointGuard: up to 20%
  - etc.

Function Pointer Overwrite

```
int median( int* data, int len, void* cmp )
{
        // must have 0 < len <= MAX_LEN
        int tmp[MAX_LEN];
        memcpy( tmp, data, len*sizeof(int) );
        qsort( tmp, len, sizeof(int), cmp );
        return tmp[len/2];
}
```

Function Pointer Overwrite

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Input Data passed

Function Pointer Overwrite

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Stack-based Buffer Overflow

Function Pointer Overwrite

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```

Function Call ‘cmp’ Overwritten!

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Vtable Overwrite :D

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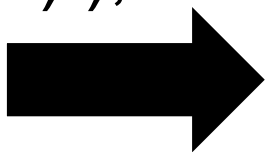

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mov  eax, [ebx+8]          ;
cmp  [eax+4], 12345678h    ;
jne  error_label           ;
call eax                   ;
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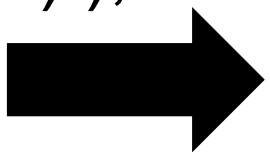

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call eax                   ;
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```

Fails :p

Security-Related Experiments

- Prevented
  - Jump to libc
  - Virtual Table Overwrite
  - etc.

- Not prevented
  - Incorrect parsing of input strings

Critique

+ Simple yet effective mitigation for many real world programs
+ Low overhead, Practical solution

- We can still use valid CFG as exploit method in some cases.

Questions? ☺