

# Report of ACE CW2

## 1 Maze

This Java program is about finding the shortest time from the start cell to the destination cell in a maze using Dijkstra's Algorithm.

### 1.1 Selection of data structures

- 2D Array: Using a 2D array 'maze' to represent the maze.
- 2D Array Access: Efficiently accessing elements in the 2D array maze using coordinates. e.g.: `maze[s2][s1]`
- Constants: Used to represent propositional logic formulas and to manipulate substrings during calculations. e.g.: 2147483647
- Scanner: Used to read input from the user.
- Array: Used to store values extracted from lists for specific purposes. e.g.: `arr0`, `arr1`
- Lists: Using `List<Integer>` to dynamically store integers (the parsed input values, vertex coordinates, and distances in Dijkstra's algorithm).

### 1.2 Algorithm design

The program is designed to scan and collect the input of the grid, start cell and destination cell, using Dijkstra's algorithm to find the shortest time to walk from the start cell to the destination cell.

- Input scanning and checking: The 'Scanner' can read input by user and the method 'isLegal' and 'isLegalTwoPairs' can check if the input string is in right format.
- Dijkstra algorithm: the core of the whole algorithm to find the shortest time
  - Initialization: Initialize arrays and lists for distance (`dist`), paths (`path1`, `path2`), and vertex coordinates (`vertex1`, `vertex2`) and set initial distances for neighboring cells of the starting point.
  - Loop: First select the unprocessed cell which has the minimum `dist` and add it to the set of vertex. Then update `dist[ ][ ]` if a shorter path is found. Repeat until all vertices are processed or the destination is reached.
  - Backtrack: Once the destination is reached, backtrack the path from the destination to the starting point using the stored paths.
- Output: The final result 'time' is printed to the console.

### 1.3 Algorithm correctness justification and efficiency analysis

The algorithm can run correctly after several different tests, as it successfully reads and analyses the input maze, start cell and destination cell, and outputs the shortest time rightly.

```
PS C:\Users\86135\Desktop\ACE\CW 2\20411212> javac Maze.java
PS C:\Users\86135\Desktop\ACE\CW 2\20411212> java Maze
6 3
(1, 1) (6, 3)
1 1 1 1 1 1
1 1 1 1 1 1
1 1 1 1 1 1
7
```

Figure 1: Maze

Time complexity:

The input parsing has a time complexity of  $O(n + m)$ , where  $n$  is the length of the first input line and  $m$  is the length of the second input line. The dominant factor is the time complexity of Dijkstra's algorithm:  $O(V^2)$  where  $V = m * n$ .

The overall time complexity is  $O(n + m + (m * n)^2)$ .

Space complexity:

The maze is represented as a 2D array 'maze', contributing  $O(m * n)$  to the space complexity.

The space complexity related to Dijkstra's algorithm is dominated by the 2D distance array 'dist', which is  $O(m * n)$ , where  $m$  is the number of rows and  $n$  is the number of columns.

The overall space complexity is  $O(m * n)$ .

## 2 Walk

This java program is about finding the minimal number of steps taken from the start cell to the end cell in an area using DFS (depth-first search) and backtrack.

### 2.1 Selection of data structures

- `ArrayList<ArrayList<Integer>>`: Used to represent a 2D `ArrayList` of integers, e.g.: 'all'. It can record all possible routes.
- `int[][]` array: Used to represent a two-dimensional array for the grid where each cell can be either 0 or 1.
- `ArrayList<Integer>`: Use 'step' to record the steps in the current path during the DFS process.
- `List<Integer>`: Used to store integers obtained from user input and parsing.

### 2.2 Algorithm design

The program is designed to scan and check the area of  $n * m$  grid, non-negative int k, the start cell and the end cell from users' input, and use DFS to get and output the minimum number of step it takes from start cell to end cell.

- Read and check input values: Read users' input for grid dimensions ( $n$  and  $m$ ), starting and ending coordinates ( $(s1, s2)$  and  $(e1, e2)$ ), the number of block cells that can be ignored ( $k$ ), and the grid itself. Then check if the format is right using 'isLegal', 'isLegalTwoPairs', 'isNonNegaInt', 'isLegalGrid'.
- DFS: Define a recursive DFS function 'dfsFindWay' to explore possible paths from the starting point to the destination. The DFS function considers different directions: up, down, left, right, upper and lower diagonals, and ignores block cells up to the specified limit 'k'. The current path 'step' is updated during the DFS traversal.
- Backtrack: When the destination is reached, record the current path in the all ArrayList, and then backtrack to explore other paths.
- Find minimum steps and output: After exploring all paths, find the path with the minimum number of steps by iterating through the recorded paths in the 'all' ArrayList, then print it out.

### 2.3 Algorithm correctness justification and efficiency analysis

The algorithm can run correctly after several different tests, as it successfully output the shortest number of jumping steps.

```
PS C:\Users\86135\Desktop\ACE\CW 2\20411212> javac Walk.java
PS C:\Users\86135\Desktop\ACE\CW 2\20411212> java Walk
6 3
(3, 1) (5, 1)
0
0 1 0 1 0 1
0 0 0 0 0 1
0 1 0 0 0 0
2
```

Figure 2: Walk

Time complexity:

Grid validation includes iterating over the input grid and coordinates. Since the grid size and coordinates are bounded by  $n$  and  $m$ , the complexity of grid validation is  $O(n * m)$ . The DFS search explores different paths in the grid. In the worst case, it might explore all possible paths, leading to a time complexity of  $O(3^{n*m})$ , where 3 represents the number of directions (up, down, left, right, diagonals) excluding the one that led to the current cell.

The dominant factor of the overall time complexity is the DFS search, so the overall time complexity is  $O(3^{n*m})$ .

Space complexity:

The grid is stored as a 2D array, contributing  $O(n * m)$  to the space complexity.

The ArrayList 'all' stores all recorded paths, and each path is an ArrayList storing coordinates. The space complexity for recorded paths is  $O(k * avg)$ , where  $avg$  is the average length of the recorded paths. The overall space complexity is  $O(n * m + k * avg)$ .