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Full Stack Software Development



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In the previous class, we covered...



Homework Discussion

1. Remove characters from the first String which are present in the second String.

Sample Input:

education

rate

Sample Output:

ducion

('e', 'a', 't' are present in 2nd string too, so remove them from 1st string)

1. Print first non-repeated character from String.

Sample Input:

triggered

Sample Output:

t

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Today's Agenda

1 Trees





Poll 1 (15 Sec.)

Which of the below is a linear data structure?

- 1. Array
- 2. Linked List
- 3. Stacks
- 4. Queues



Poll 1 (Answer)

Which of the below is a linear data structure?

- 1. Array
- 2. Linked List
- 3. Stacks
- 4. Queues

We explored linear data structures such as arrays and linked lists. However, in the real world, data modelling is not always possible in a linear fashion.

Can you think of some real-world example where data is stored in a non-linear fashion?

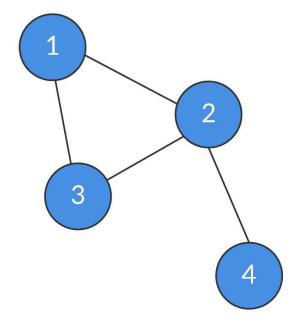
How do you think Facebook suggests friends when we visit someone's profile? Or get connected with someone new?

Non-linear data structures give us a powerful mechanism to store and manipulate data, making a lot of algorithms possible, which would not be easy to do if we used linear data structures.

They allow us to associate or connect one data element to several other data elements, or even model them in a <u>hierarchical</u> way, if we want to.

However, for the <u>manipulation of data</u>, be it insertion, searching or deletion, BSTs (Binary Search Trees) provide better running time.

Non-Linear Data Structures



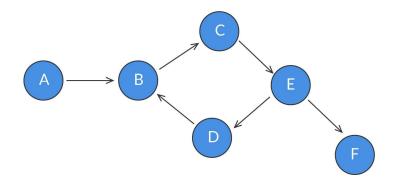
A graph 'G' can be represented by two attributes:

- Its set of vertices (V), and
- Its set of edges connecting these vertices (E).

Edges are represented as **pairs** of vertices. In the graph shown in this slide, the edges have <u>no</u> 'direction'.

This type of graph is called an **undirected graph**.

Non-Linear Data Structures



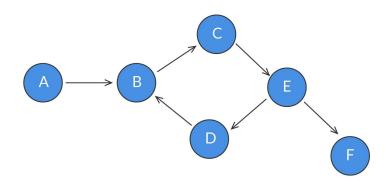
We can also have <u>graphs</u> where the <u>edges</u> have a 'direction'.

In the adjoining graph,

- The edge (A, B) is <u>not</u> the <u>same</u> as the edge (B, A).
- A graph might have both the edges (A,B) and (B,A), if need be.

These graphs are called **directed graphs**.

In directed graphs, the edges are represented in the format (u, v), where 'u' is the <u>source vertex</u> and 'v' is the <u>destination vertex</u>.



Can you think of a real-world model where a directed graph would be useful?

How about one where undirected graphs would make more sense?

What kind of graph is implemented by Facebook? How about Instagram?



Poll 2 (15 Sec.)

Graph in which edges have no direction is called:

- 1. Directed graph
- 2. Undirected graph



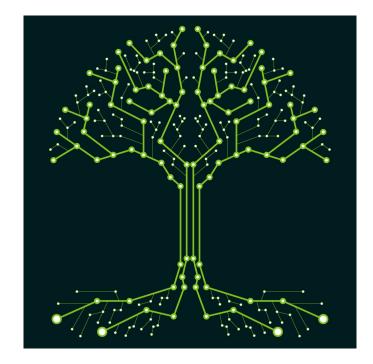
Poll 2 (Answer)

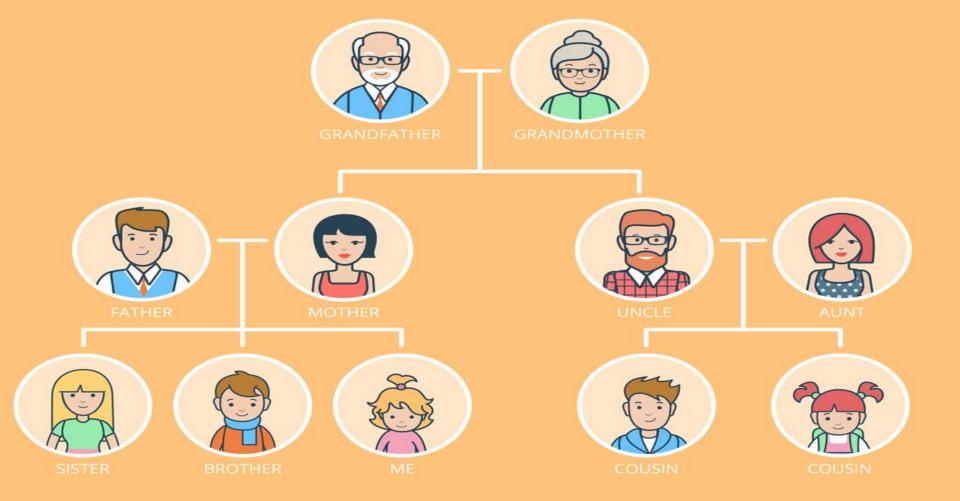
Graph in which edges have no direction is called:

- 1. Directed graph
- 2. Undirected graph

Trees up**Grad**

Tree data structure is a subclass of graph that deserves individual focus, simply because of the vast amount of power that it gives us in terms of efficient algorithm design. This subclass of graphs is called **Trees**.

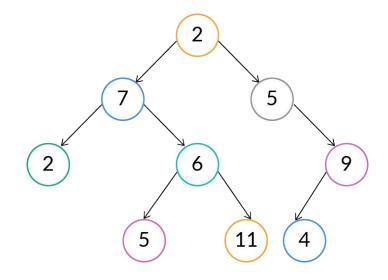




Trees up**Grad**

A tree is a graph that:

- is **simple**,
- connected,
- undirected,
- has no cycles
- and has *n*-1 edges if the number of vertices is *n*.



Trees

A tree also has a few other features:

- The graph has no cycles, but a cycle will be formed even if one edge is added to
 it.
- The graph is connected, but will not remain connected if any vertex is removed from it.
- Any two nodes/vertices in the graph can be connected by one simple unique path.

Trees

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Now, let us take a look at a few tree terms and definitions:

- A tree consists of one 'root' node, zero or more 'internal' nodes and zero or more 'leaf' nodes.
- Each node can have zero or more children.
- A node can have at most one parent.
- A binary tree is a special kind of tree where a node can have zero to two nodes only.



Poll 3 (15 Sec.)

State true or false:

A tree consists of one 'root' node, zero or more 'internal' nodes and zero or more 'leaf' nodes.

- 1. True
- 2. False



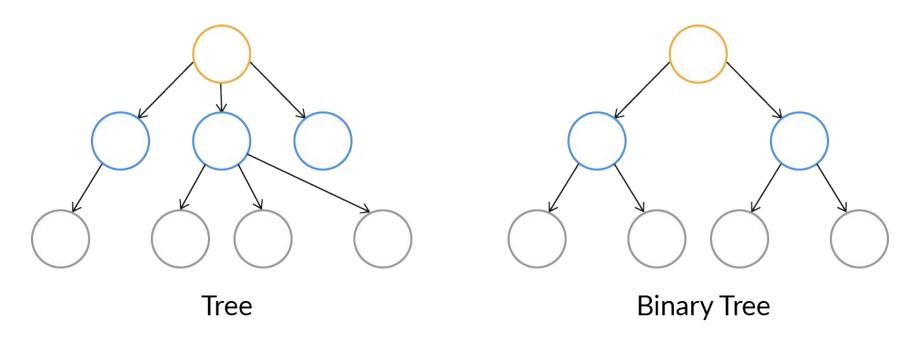
Poll 3 (Answer)

State true or false:

A tree consists of one 'root' node, zero or more 'internal' nodes and zero or more 'leaf' nodes.

- 1. True
- 2. False

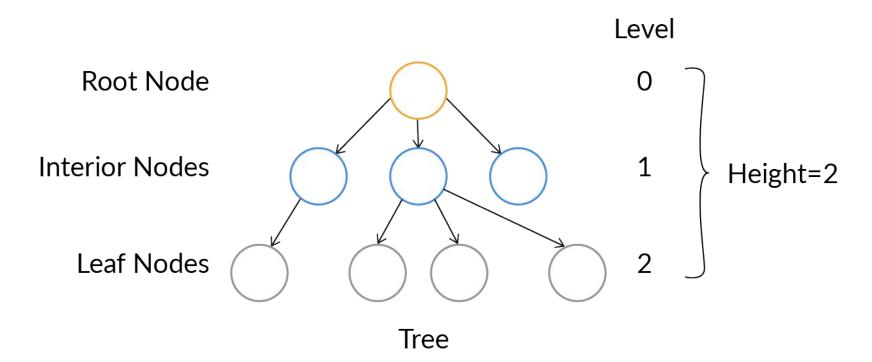
Here are diagrammatic representations of a tree and a binary tree:



Trees

Now, let's take a look at a few more terms and definitions related to a tree:

- The 'root' of a tree has **no parent**. (Trees are similar to traditional family trees.)
- The 'leaves' of a tree have no children.
- All other nodes are interior nodes.
- The height of a tree is the distance from the root of the tree to the farthest leaf node.
- The depth of any node is the distance from the root to the node.
- The 'child' nodes of all the non-leaf nodes can themselves be considered as trees in their own right, which are rooted at these nodes.





Poll 4 (15 Sec.)

State true or false:

The height of a tree is the distance from the root of the tree to the farthest leaf node.

- 1. True
- 2. False



Poll 4 (Answer)

State true or false:

The height of a tree is the distance from the root of the tree to the farthest leaf node.

- 1. True
- 2. False

Trees



Now, let us learn about a few tree terms and definitions:

- The 'degree' of a node is the number of sub-trees or children the node has. As we can understand, the leaves are nodes with a degree of '0'.
- Children of the same parent node are known as siblings.
- All the nodes in the path starting from the root of the tree right up to the parent of a node are known as its *ancestors* or *predecessors*.

Binary trees are the most popular application of trees, and they help us in a large number of algorithms in reducing the time complexity. Binary trees, as we saw earlier, are a special type of trees, where the maximum degree of each node is 2. Binary trees, in turn, can be of a few types as shown below:

- Degenerate trees, where all the nodes have one child.
- Fully binary or strictly binary trees, where all the non-leaf nodes have two children.
- **Balanced** binary trees, where most of the nodes have two children and the sub-trees are mostly of equal height.
- **Complete** trees, where all the tree levels are complete, except possibly the last, which in itself must be filled starting from the left-most node.



Poll 5 (15 Sec.)

State true or false:

In fully binary trees, most of the nodes have two children and the sub-trees are mostly of equal height.

- 1. True
- 2. False

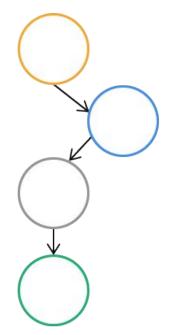


Poll 5 (Answer)

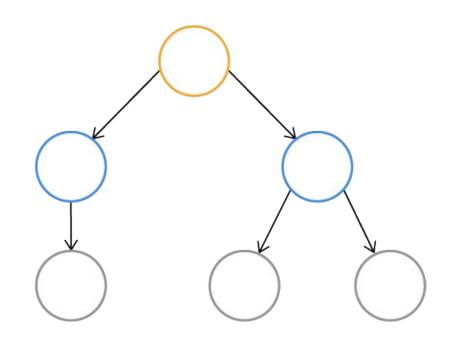
State true or false:

In fully binary trees, most of the nodes have two children and the sub-trees are mostly of equal height.

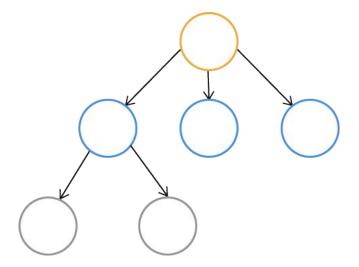
- 1. True
- 2. False



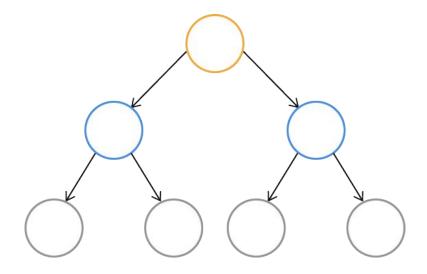




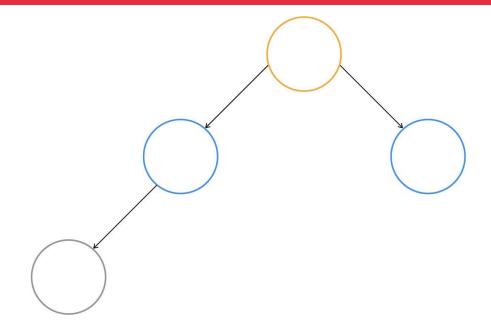
Balanced binary tree



Complete binary tree



Full and Complete binary tree

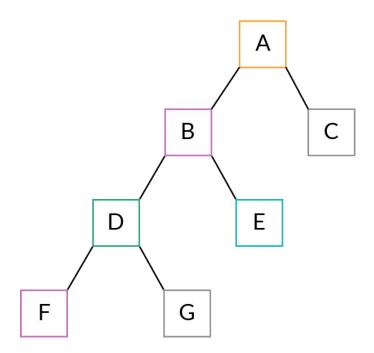


Almost complete binary tree

It's your turn now. Are these binary trees:

Complete?

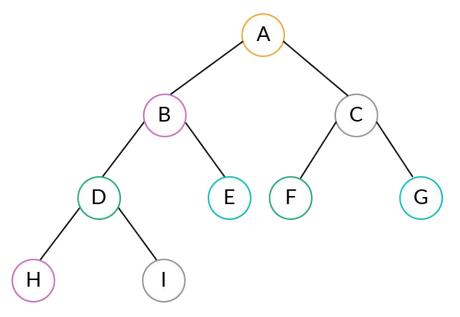
Strictly binary?



It's your turn now. Are these binary trees:

Complete?

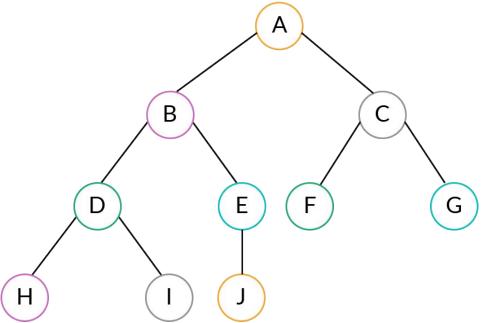
Strictly binary?



It's your turn now. Are these binary trees:

Complete?

Strictly binary?



Now, let's try to figure out a bit of mathematics:

- In a binary tree, what is the maximum number of nodes at the 'k'th level?
- What is the maximum number of nodes in a binary tree of height 'h'?
- What is the minimum number of nodes that can be present in a complete binary tree of height 'h'?

Now that we know what a binary tree is, let us see how we can represent the same in our data structures.

Binary trees can be traditionally represented in two ways:

- 1. With **arrays**, which is also known as the sequential representation.
- 2. A **linked** representation, using nodes and pointers, similar to the ones we used while discussing Linked Lists.

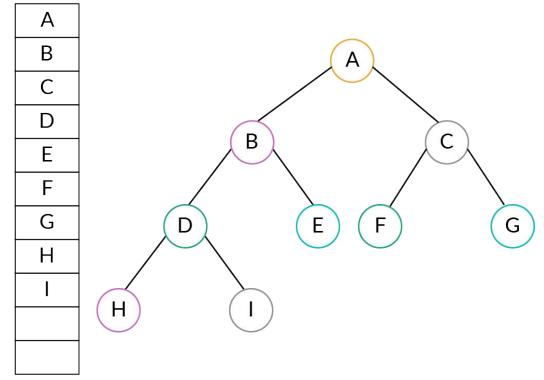
Binary Tree

Array representation:

In this representation, we define an array of size 2^h-1, if 'h' is the height of the binary tree. The root is placed at the first array position. Subsequently, for any node stored at the 'x'th index:

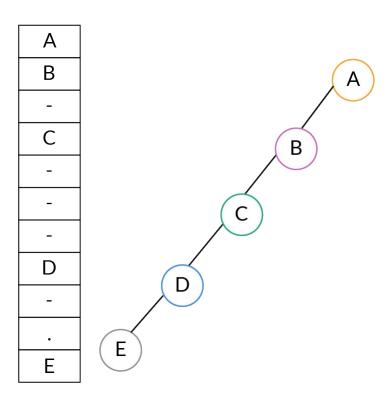
- The left child of the node will be at an array index of $2^*x + 1$.
- The right child of the node will be at an array index of $2^*x + 2$.
- The parent node will be at an array index position of (x-1)/2 if x is odd and at an array index position of (x-2)/2 if x is even.

Array representation:



Array representation:

We can see, as a tree trends towards being degenerate, there is tremendous wastage of storage space!





Poll 6 (15 Sec.)

The left child of the binary tree node will be at an array index of 2*x + 1.

- 1. True
- 2. False



Poll 6 (Answer)

The left child of the binary tree node will be at an array index of 2*x + 1.

- 1. True
- 2. False

Binary Tree

The problem of storage wastage posed by an array representation of a binary tree is solved by making use of a linked representation. In this representation, we define an abstract data type (Nodes) like the ones we did while creating Linked Lists.

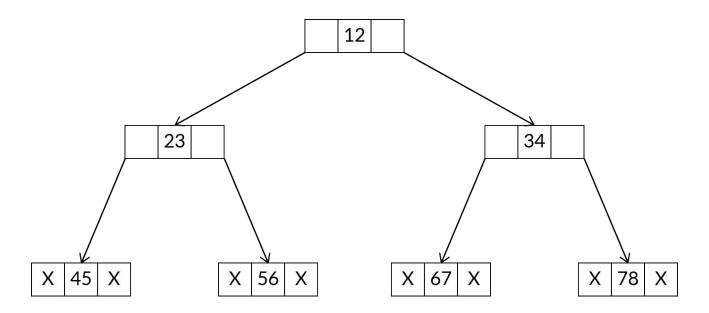
However, there is a slight difference now. Instead of having a pointer to the 'next' sequential node, we will now have pointers to the left and the right child of the binary tree.

Left	Data	Right
------	------	-------

The left and the right child pointers of these nodes will be **null**.

Binary Tree

As you can see, there will be no wastage of storage space as we create nodes that are equal in number to the data elements that we have.



Now that you know how to represent a binary tree, you will now learn how to go about traversing them. Conventionally, we are interested in three type of binary tree traversals:

- Inorder traversal
- Preorder traversal
- Postorder traversal

Binary Tree upGrad

Inorder Traversal

Inorder traversal follows a very simple algorithm. Whenever we encounter a node, we:

- Perform an inorder traversal of the left sub-tree of the node.
- 'Traverse' or 'process' the node in consideration.
- Perform an inorder traversal of the right sub-tree of the node.

Now, can anyone tell which algorithmic construct we shall be revisiting here?

Inorder Traversal

Let's take a look at a high-level algorithm for inorder traversal.

```
void inorder (Node root)
{
    if (root != null) {
        inorder (root -> left);
        process (root);
        inorder (root -> right);
    }
}
```

Preorder and Postorder Traversals

Preorder and Postorder traversals work very similarly to Inorder traversal, with the same subtle difference that prefix and postfix expressions had from infix expressions.

So, when you are required to perform a **Preorder** traversal, and you come across a node:

- 'Traverse' or 'process' the node in consideration.
- Perform a preorder traversal of the left sub-tree of the node.
- Perform an preorder traversal of the right sub-tree of the node.

Similarly, when you are required to perform **Postorder** traversal, and you come across a node:

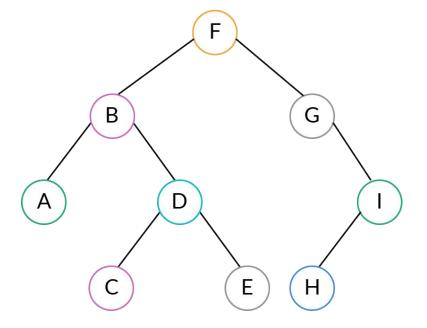
- Perform an postorder traversal of the left sub-tree of the node.
- Perform an postorder traversal of the right sub-tree of the node.
- 'Traverse' or 'process' the node in consideration.

Now, let us practice traversals a bit with a few examples:

Preorder: FBADCEGIH

Inorder: ABCDEFGHI

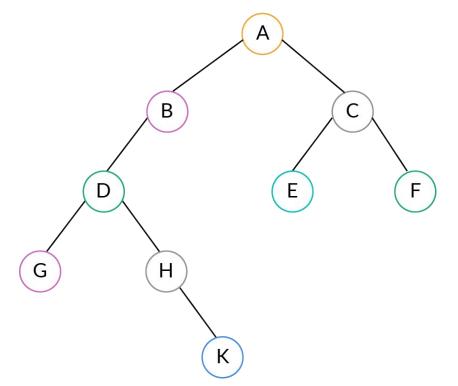
Postorder: A C E D B H I G F



Preorder:

Inorder:

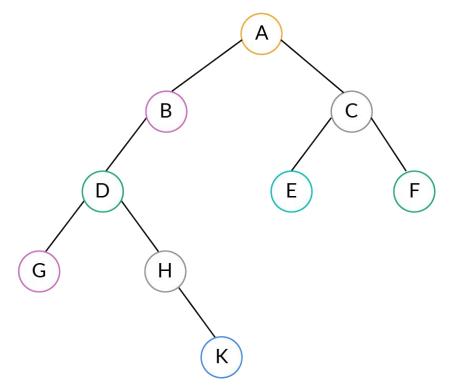
Postorder:



Preorder: ABDGHKCEF

Inorder: GDHKBAECF

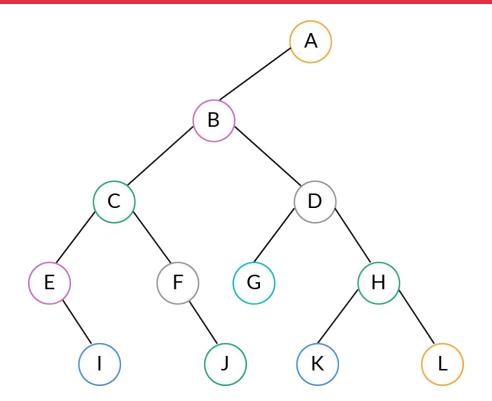
Postorder: GKHDBEFCA



Preorder:

Inorder:

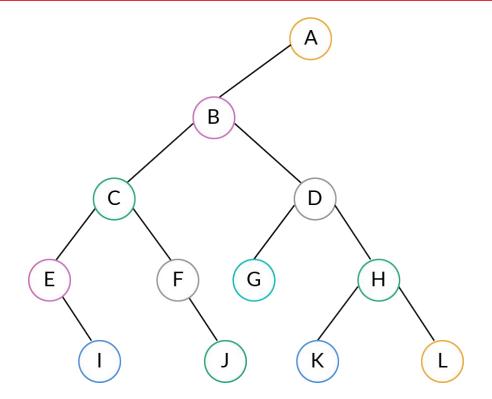
Postorder:



Preorder: ABCEIFJDGHKL

Inorder: EICFJBGDKHLA

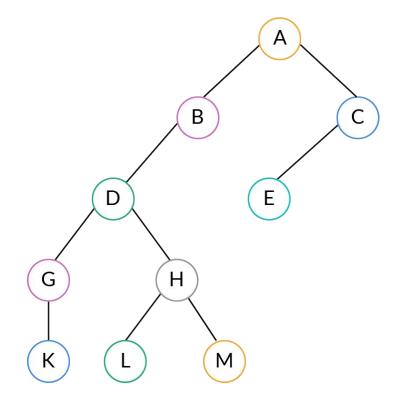
Postorder: IEJFCGKLHDBA



Preorder:

Inorder:

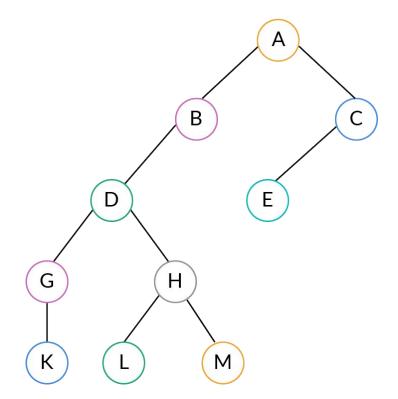
Postorder:



Preorder: ABDGKHLMCE

Inorder: K G D L H M B A E C

Postorder: K G L M H D B E C A



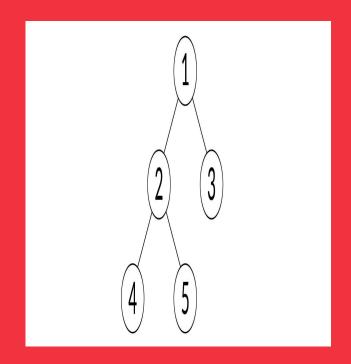
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Homework

You are given a binary tree. You need to print the maximum height of the binary tree. If the tree is NULL (empty tree), print the height of the tree as 0.

Output: 3

Here, in the given tree, the maximum height of the binary tree is 3.





Tasks to complete after the session

Homework

MCQs

Coding Questions

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In the next class...

BST





Thank You!