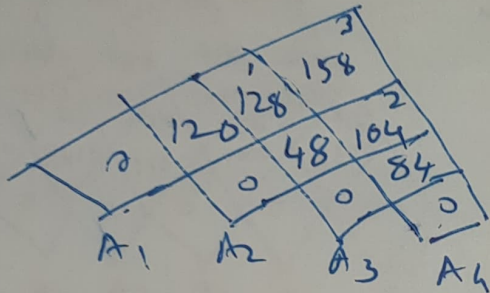


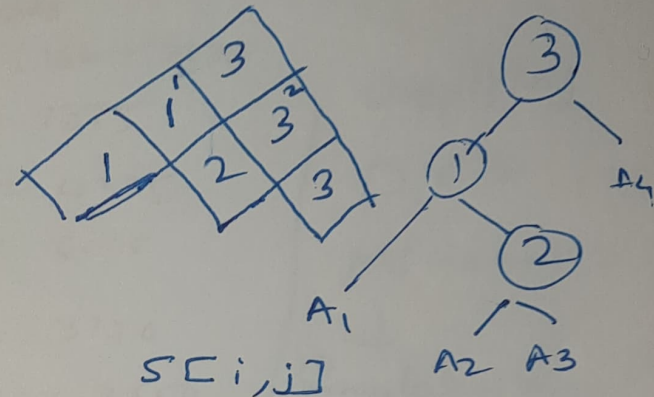
MCM - solution-order

initial set of dimensions $\langle 5 \ 4 \ 6 \ 2 \ 7 \rangle$

A_1	A_2	A_3	A_4	A_5
5×4	4×6	6×2	2×7	



$m[i, j]$



order $(A_1, (A_2 A_3)) A_4$

$$\begin{aligned} \textcircled{1} \quad & \begin{matrix} 5 \times 4 & 4 \times 2 \\ (A_1) & (A_2 A_3) \end{matrix} = 40 + 48 \\ & \begin{matrix} 5 \times 6 & 6 \times 2 \\ (A_1 A_2) & (A_3) \end{matrix} = 60 + 120 \end{aligned}$$

$$\begin{aligned} \textcircled{2} \quad & \begin{matrix} 4 \times 6 & 6 \times 7 \\ A_2 & (A_3 A_4) \end{matrix} = 168 + 84 \\ & \begin{matrix} 4 \times 2 & 2 \times 7 \\ (A_2 A_3) & A_4 \end{matrix} = 56 + 48 \end{aligned}$$

$$\begin{aligned} \textcircled{3} \quad & \begin{matrix} 5 \times 4 & 4 \times 7 \\ A_1 & (A_2 A_3 A_4) \end{matrix} = 140 + 104 = 244 \\ & \begin{matrix} 5 \times 6 & 6 \times 7 \\ (A_1 A_2) & (A_3 A_4) \end{matrix} = 120 + 84 + 290 \\ & \begin{matrix} 5 \times 4 & 4 \times 2 & 2 \times 7 \\ A_1 & (A_2 A_3) & A_4 \end{matrix} = 48 + 40 + 70 = 158 \\ & \begin{matrix} 5 \times 4 & 4 \times 2 & 2 \times 7 \\ (A_1 A_2 A_3) & A_4 \end{matrix} = 70 + 128 \end{aligned}$$

```

n ← length LPF - 1
for i = 1 to n
  do m[i, i] = 0
for L = 2 to n do
  (L is length of subchain)
  for i = 1 to n - L + 1
    do j = i + L - 1
      m[i, j] = ∞
      for k = i to j - 1
        do
          q ← m[i, k] + m[k + 1, j]
            + p[i - 1] p[k] p[j]
          if q < m[i, j]
            then
              m[i, j] ← q
              S[i, j] ← k
  return m, S
  
```

matrix chain multiplication

A_1 A_2 A_3 A_4 A_5 A_6
 30×35 35×15 15×5 5×10 10×20 20×25

~~2625~~
~~1500~~
~~7875~~

35
 $\times 150$ 7

 5250

① $(A_1 A_2) A_3 = \text{7875}$

min

\checkmark $A_1 (A_2 A_3) = \text{4375}$ 7875
 $(30 \times 35 \times 5) + 2625 = 5250 + 2625 = 7875$

② $\checkmark (A_2 A_3) A_4 = 4375$
 min $A_2 (A_3 A_4) = 6000$

③ $(A_3 A_4) A_5 = 3750$
 $A_3 (A_4 A_5) = 2500$

④ $(A_4 A_5) A_6 = 3500$
 $A_4 (A_5 A_6) = 5750$

applying

D P to

KLP - class \rightarrow

\downarrow

Formulations may
 not be efficient
 as no. of subprbs
 is large
 (should be poly. number)

Dyna. -

substructure
 table structure
 bottom up comput.

D & C

divide sub problems

independent subprb.
 (subprb share
 sub-sub prb)

guarantee
 optimal soln

choose best at each
 current step

Dyna

D & C

Greedy

y

y

n

~~y~~

y

n

y

y

n

y