

# Pseudo-random Function

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We devise a pseudo-random function from pseudo-random generator  $G$  as follows. Let

$G : \{0, 1\}^n \rightarrow \{0, 1\}^{2n}$  be a PRG. Let

$r = r_0 r_1 \dots r_{n-1}$  be a random sequence. Define  $G_0$  and  $G_1$  as the left and right parts of the output of  $G(x)$ .

$$G_0(x) = \text{Left}(G(x))$$

$$G_1(x) = \text{Right}(G(x))$$

$$G(x) = G_0(x) || G_1(x)$$

Let  $F_k : \{0, 1\}^n \rightarrow \{0, 1\}^n$  be defined as follows,

$$F_k(r) = G_{r_{n-1}}(G_{r_{n-2}}(\dots(G_{r_0}(k))))$$

where  $k$  is given key of length  $n$ .

If  $F_k(r)$  were distinguishable from a truly random function, the PRG  $G$  used would be distinguishable from a truly random generator because a truly random generator must give a truly random function following the aforesaid process. Hence,  $F_k(r)$  is a pseudo random function.