

Vertex Covers

Alexander Golovnev

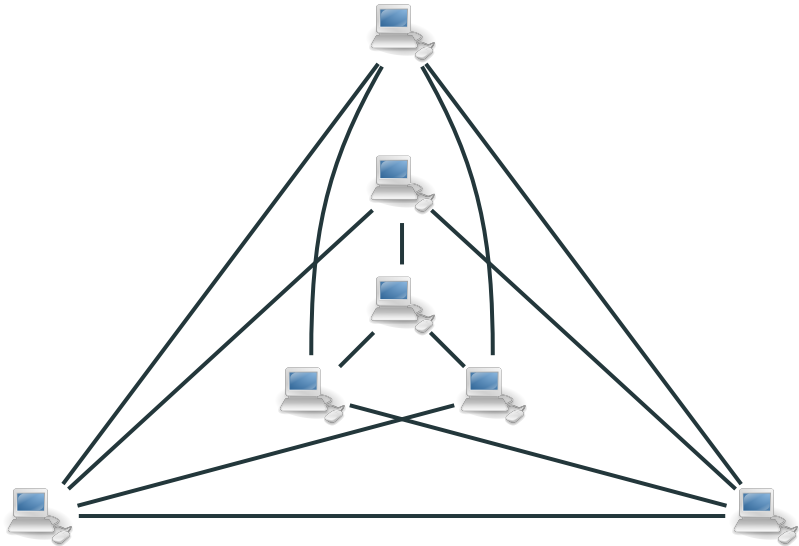
Outline

Antivirus System

Vertex Covers

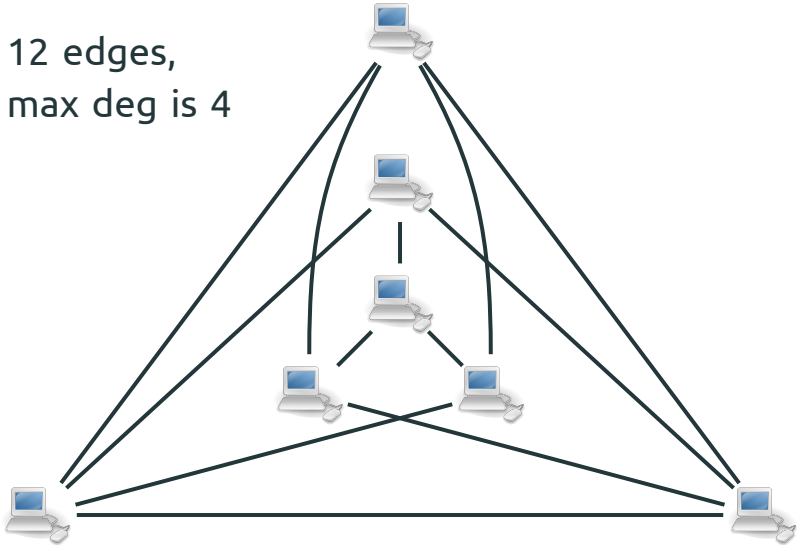
König's Theorem

Antivirus System



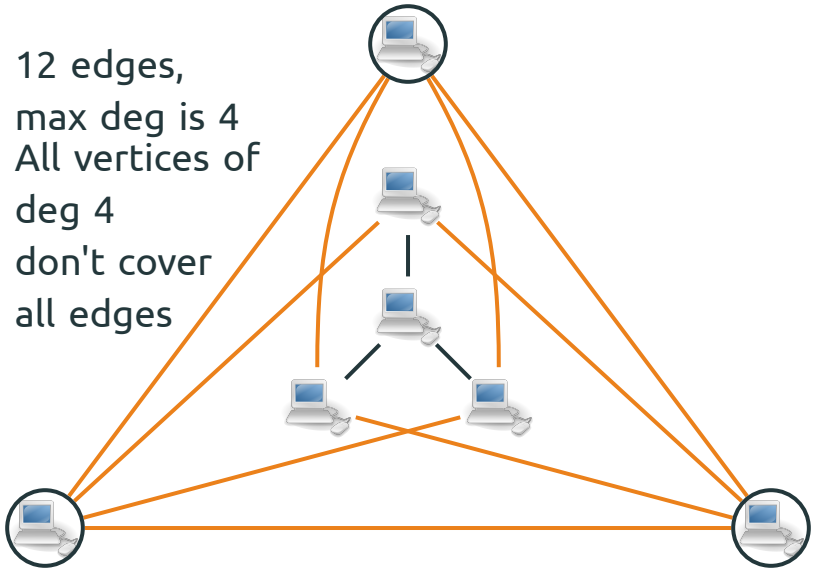
Antivirus System

12 edges,
max deg is 4



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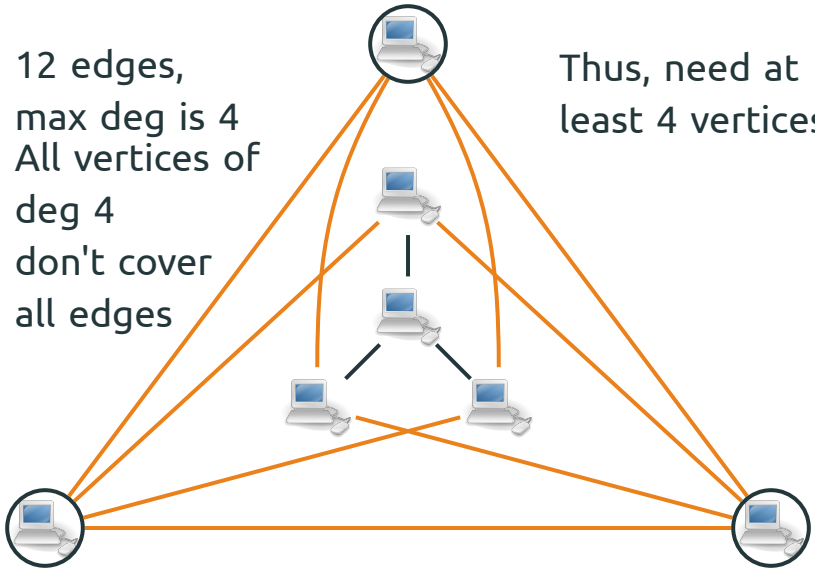
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All vertices of
deg 4
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Antivirus System

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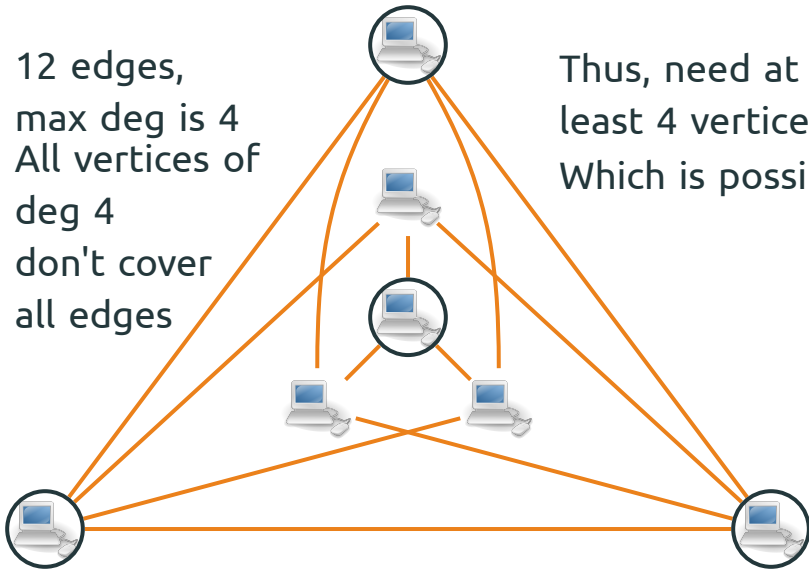
Thus, need at
least 4 vertices



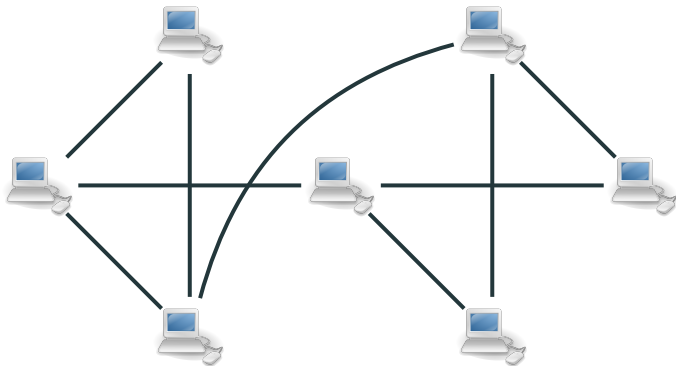
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Thus, need at
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Which is possible

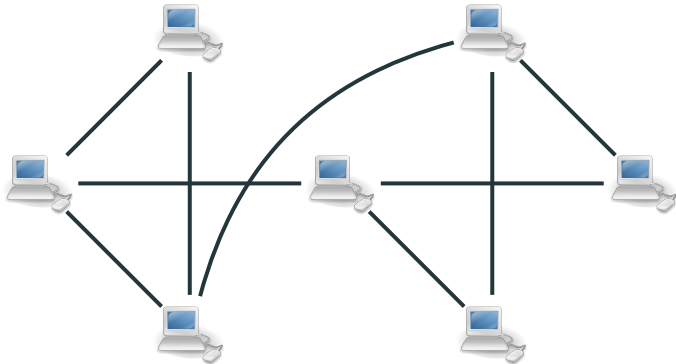


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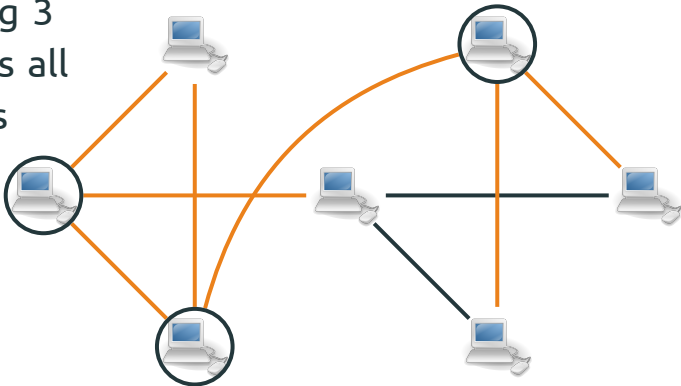
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9 edges,
max deg is 3



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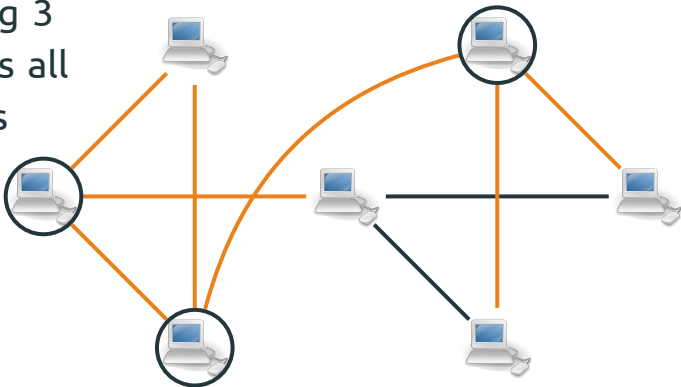
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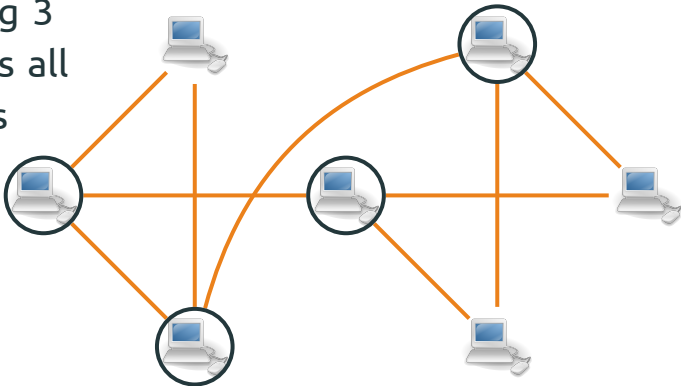
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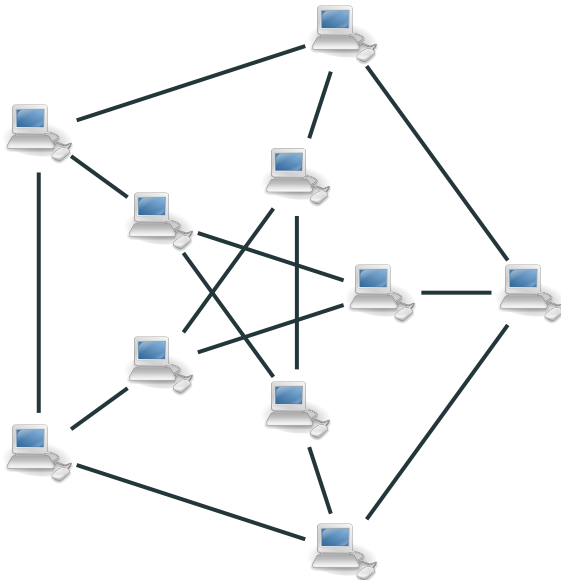
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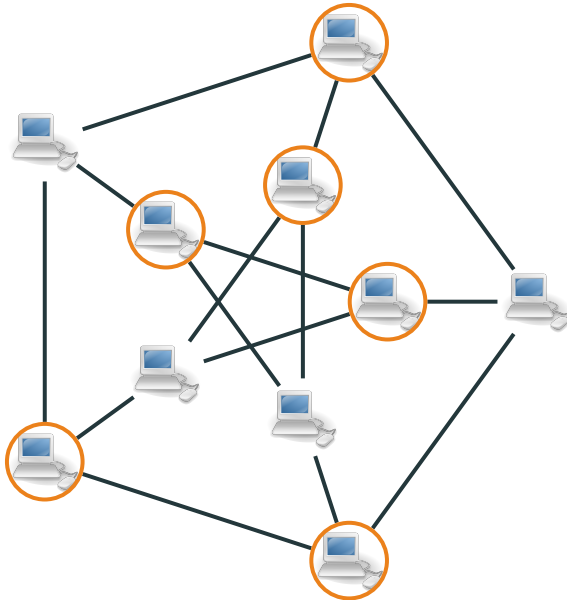
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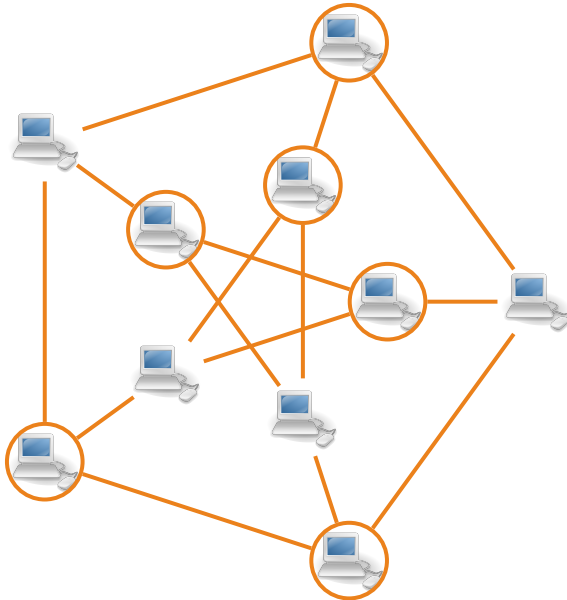
Antivirus System



Antivirus System



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Outline

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Vertex Covers

König's Theorem

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- A **Vertex Cover** of a graph G is a set of vertices C such that every edge of G is connected to some vertex in C .

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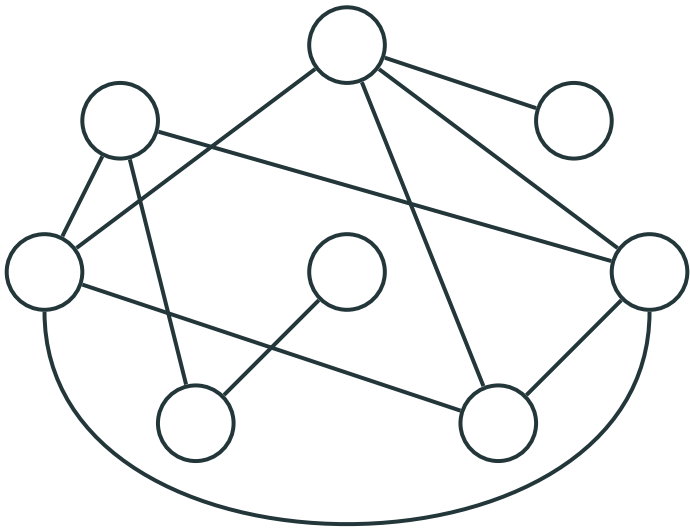
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Vertex Covers

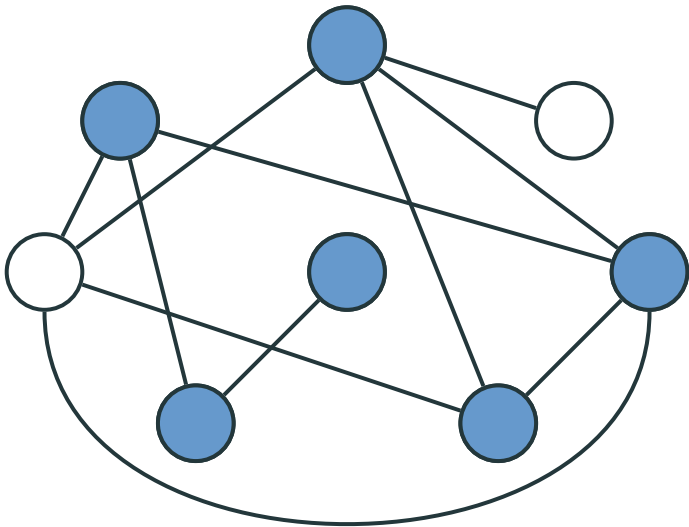
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- The Size of a Minimum Vertex Cover is denoted by $\beta(G)$.

Vertex Covers: Examples



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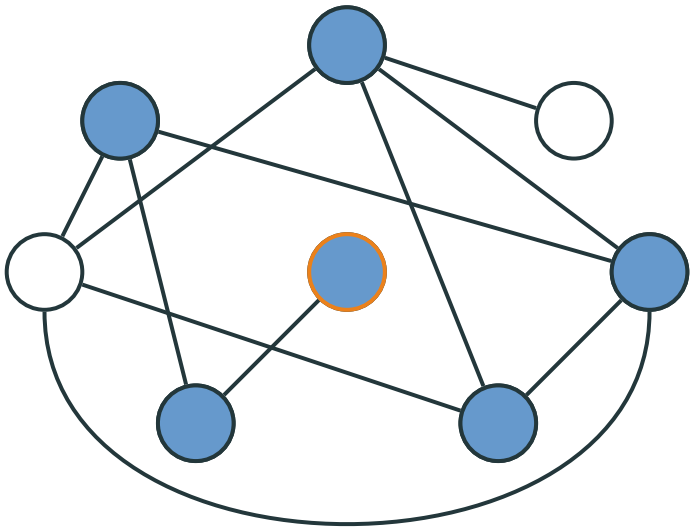
A Vertex
Cover



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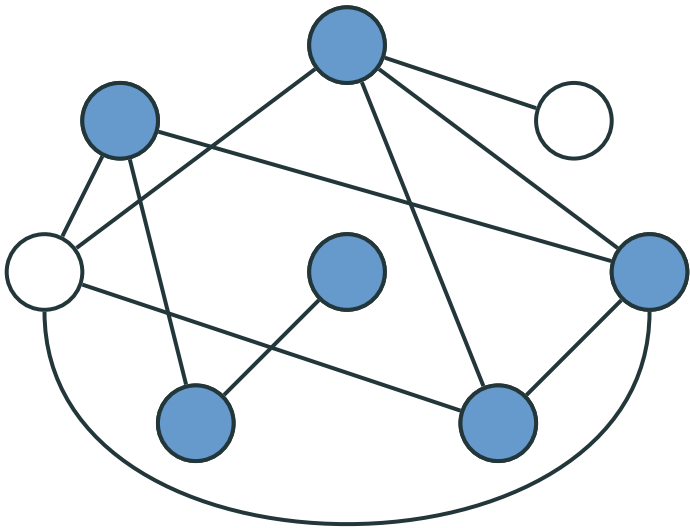
A Vertex
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Not a Mini-
mal VC

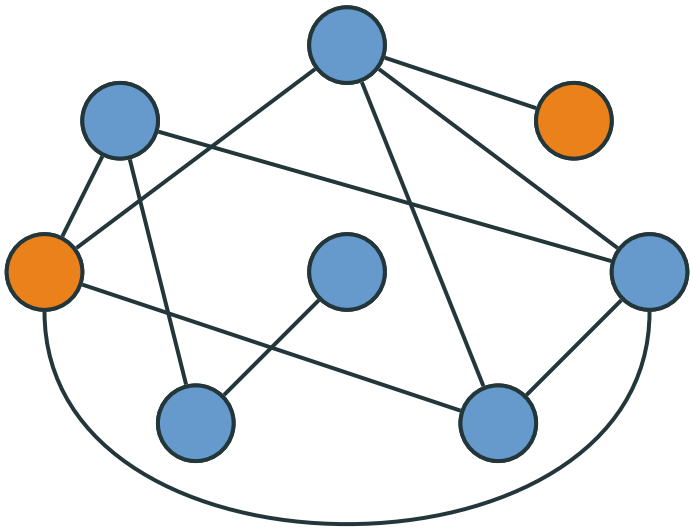


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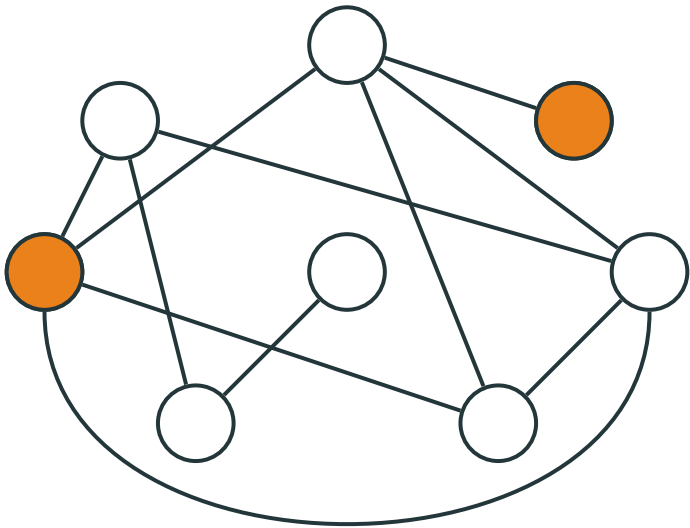


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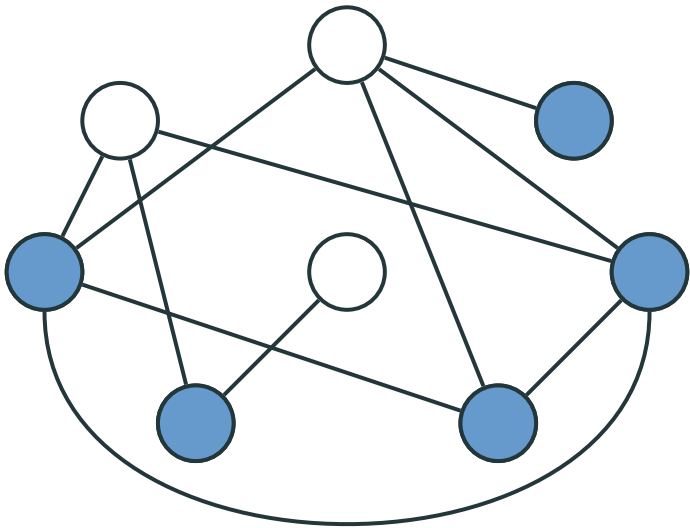
Vertex Covers: Examples

An Independent Set



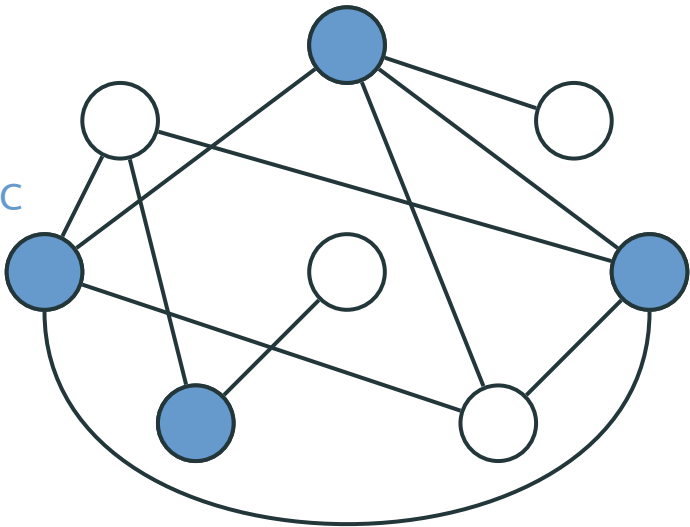
Vertex Covers: Examples

A Minimal VC



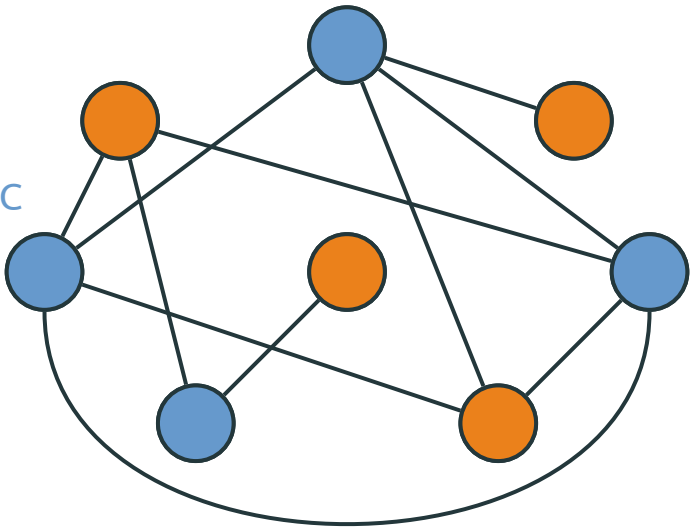
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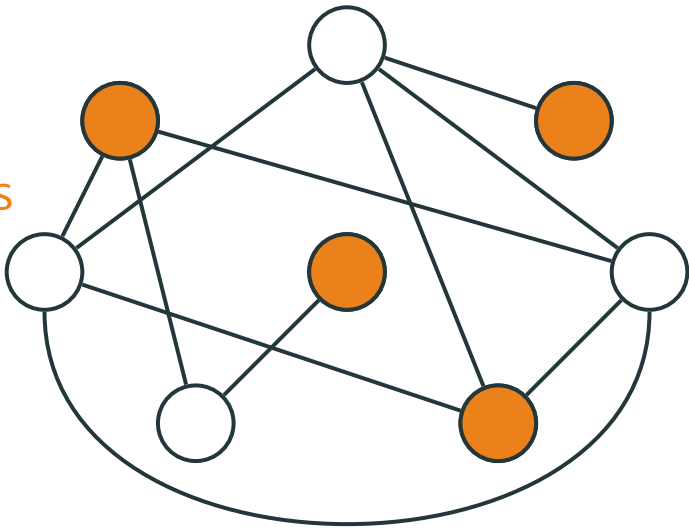
Vertex Covers: Examples

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Vertex Covers: Examples

A Maximum IS



Vertex Covers and IS's

Fact

*A set of vertices is a **Vertex Cover** if and only if its complement is an **Independent Set***

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Corollary

For every graph G on n vertices:

$$\beta(G) + \alpha(G) = n .$$

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Indeed, if an edge (u, v) is not covered by a matching, then it can be added to it (contradicts its maximality).

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Theorem (König, 1931)

*In a bipartite graph, the number of edges in a **Maximum Matching** equals the number of vertices in a **Minimum Vertex Cover**.*

König's Theorem. Proof

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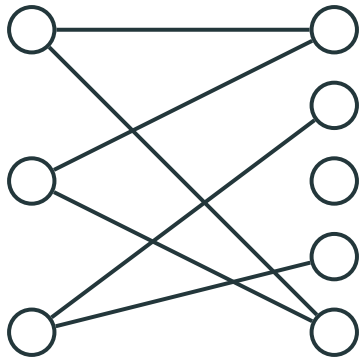
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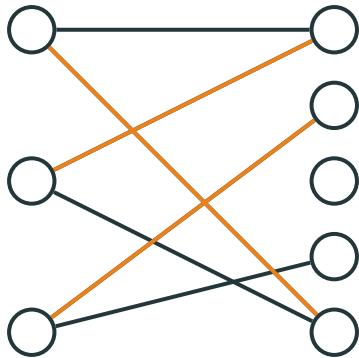
Proof:

- A VC must contain at least 1 vertex from each edge of a Matching

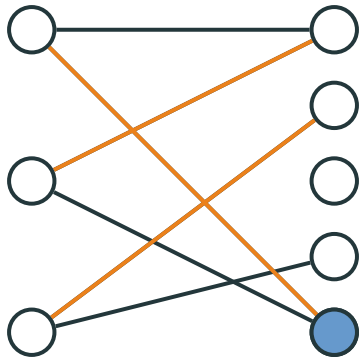
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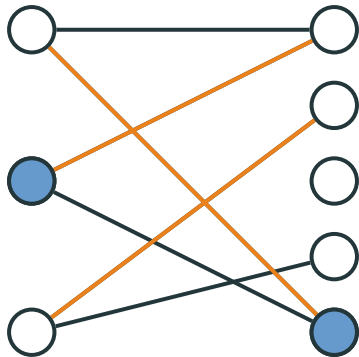
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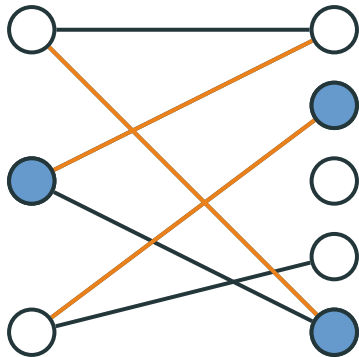
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- We'll show that $|\text{Min VC}| \leq |\text{Max Matching}|$

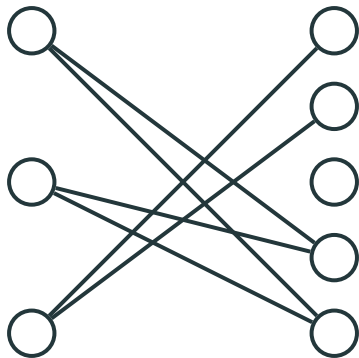
König's Theorem from Hall's Theorem

- A Min VC C . $L_C = L \cap C$, $R_C = R \cap C$

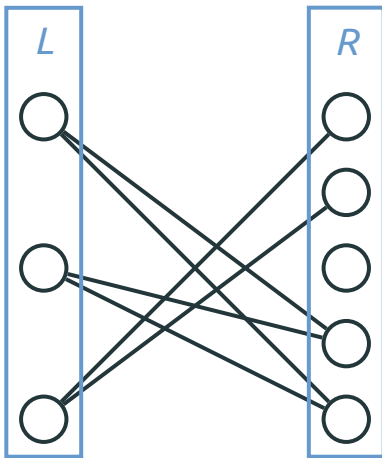
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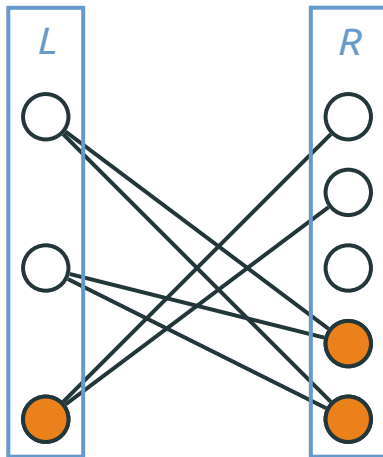


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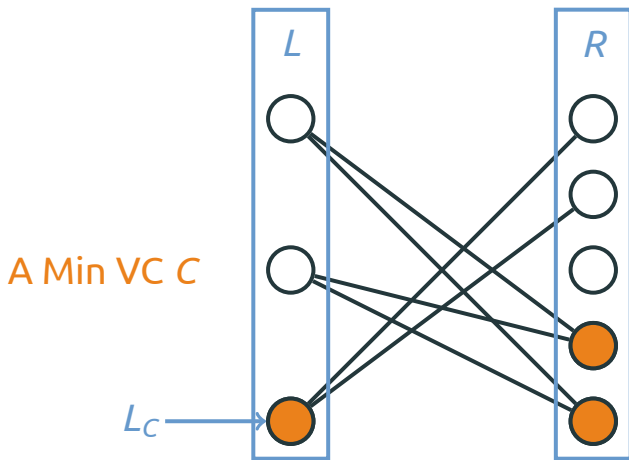


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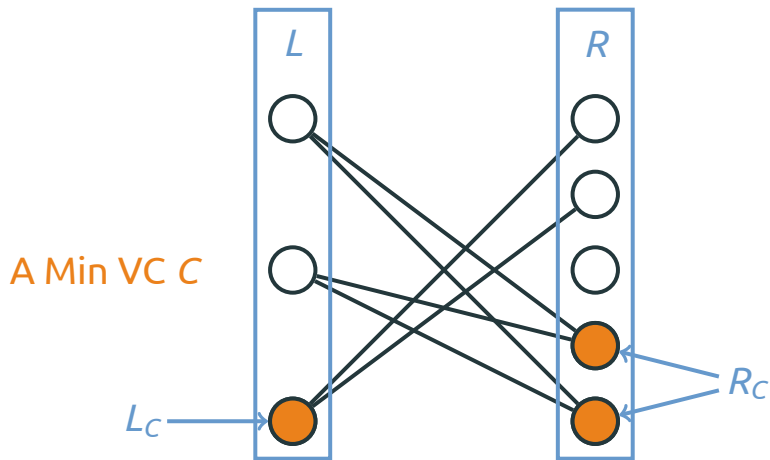
A Min VC C



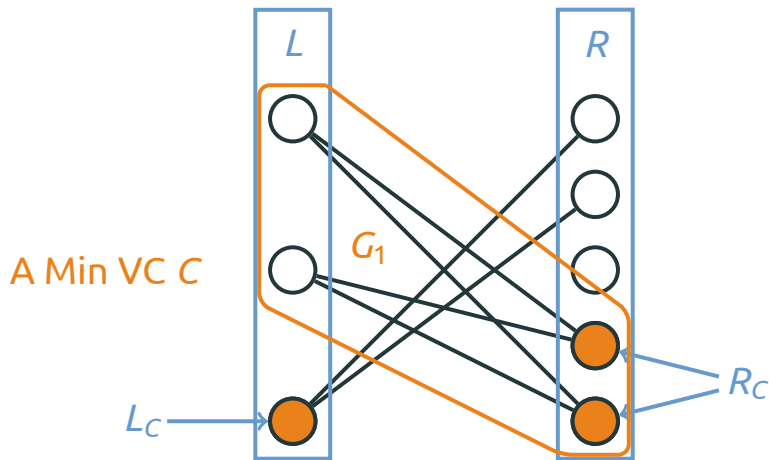
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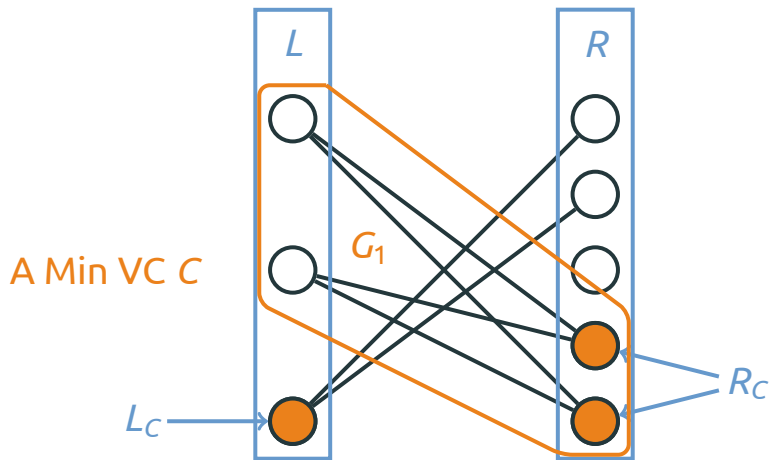
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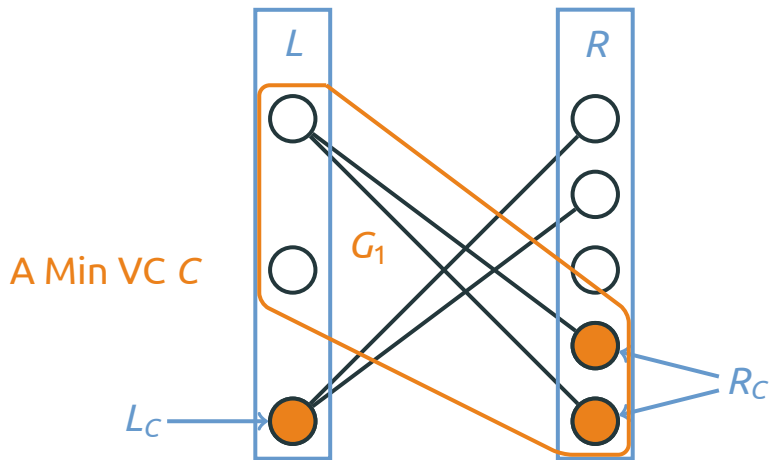
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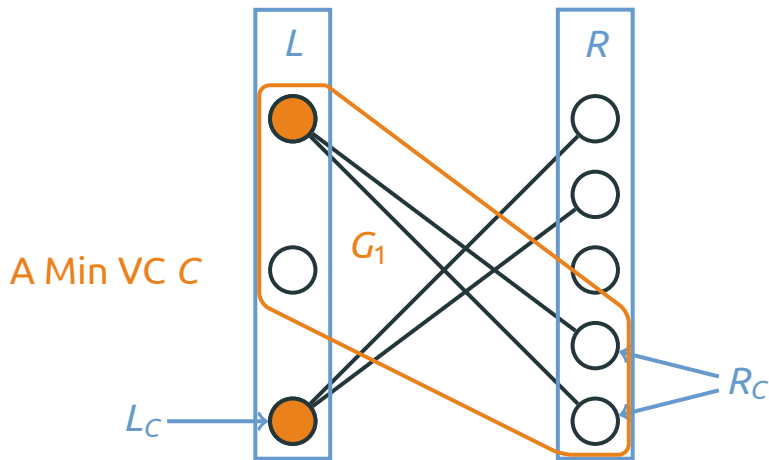
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- Similarly, G_2 – the subgraph on $L_C \cup (R \setminus R_C)$ has a matching of size $|L_C|$
- Total matching size $|L_C| + |R_C| = |C|$