

Cliques and Independent Sets

Alexander Golovnev

Outline

Graph Cliques

Cliques and Independent Sets

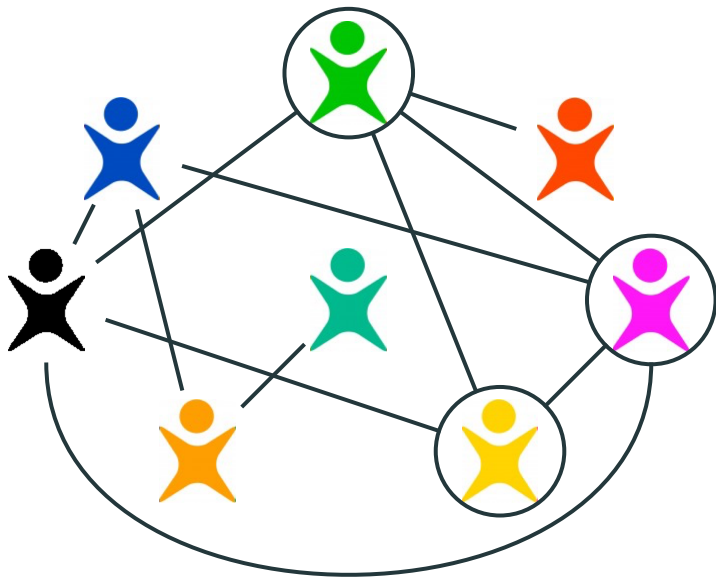
Connections to Coloring

Mantel's Theorem

Friendship Graph



Friendship Graph



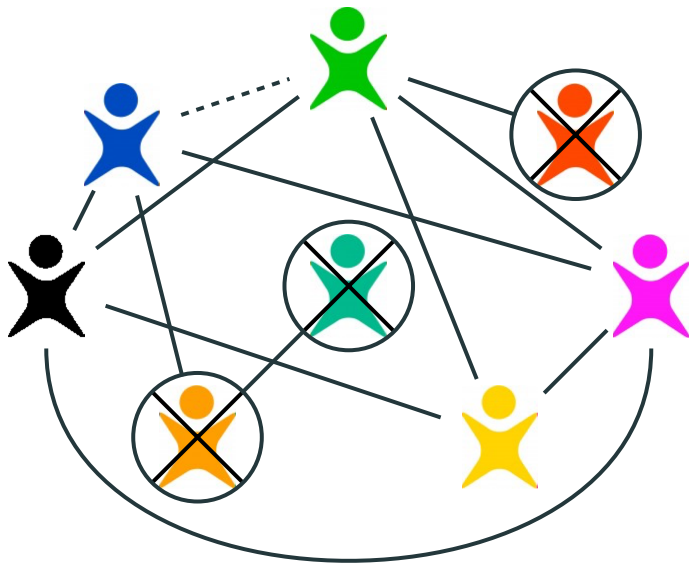
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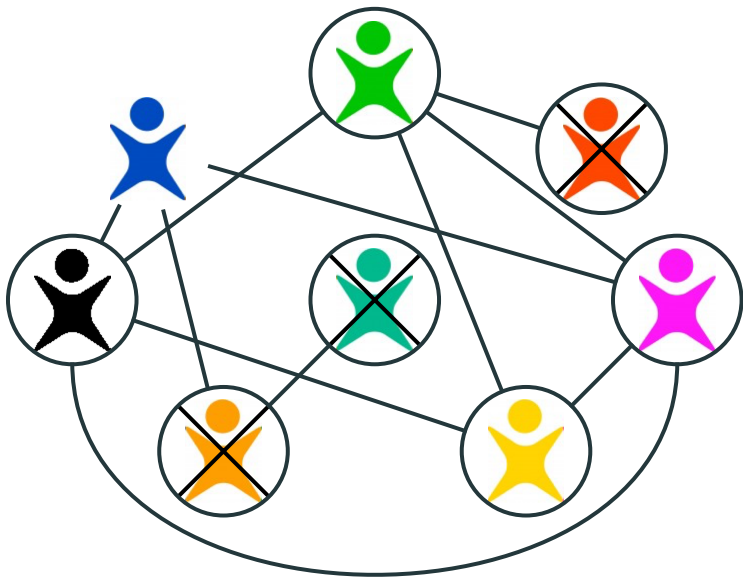
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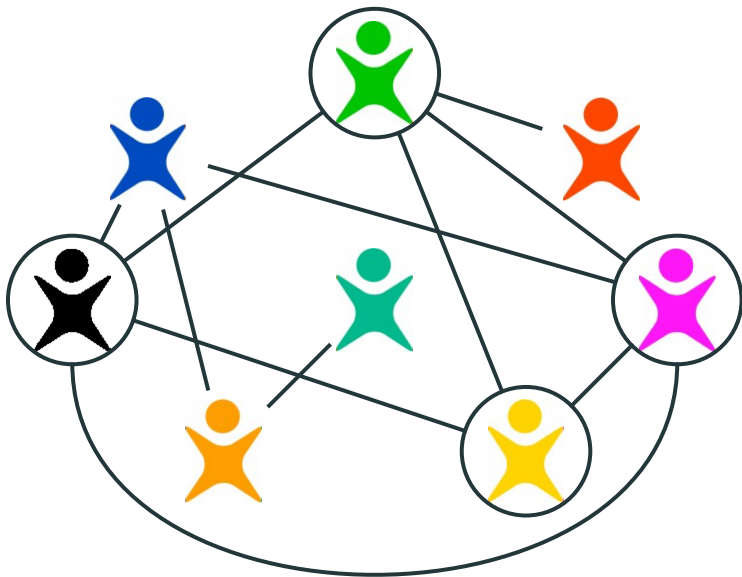
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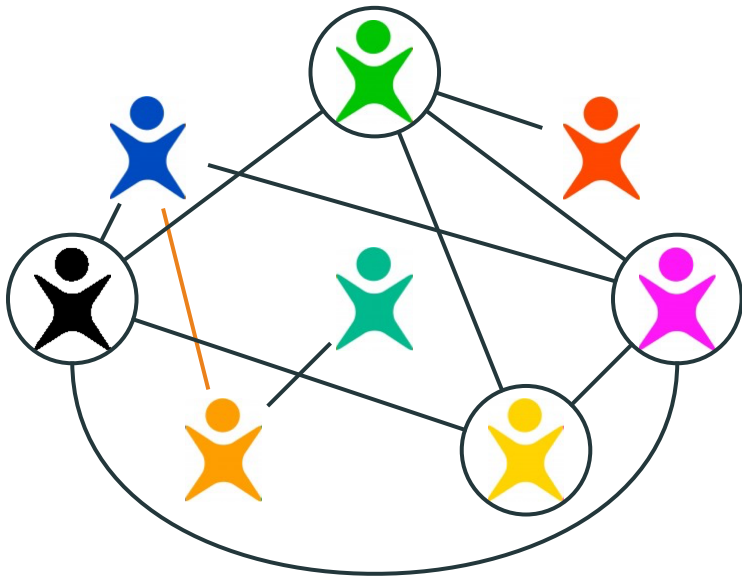
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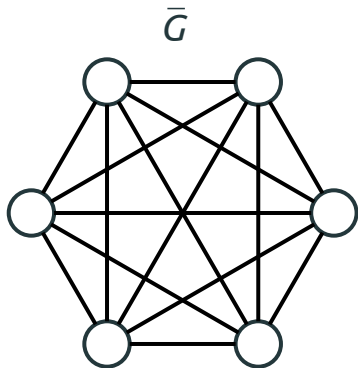
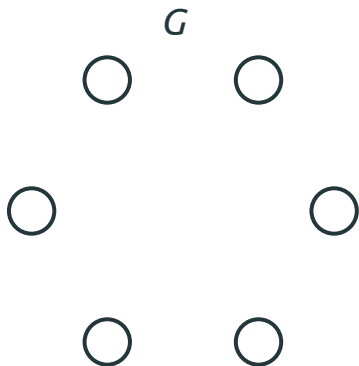
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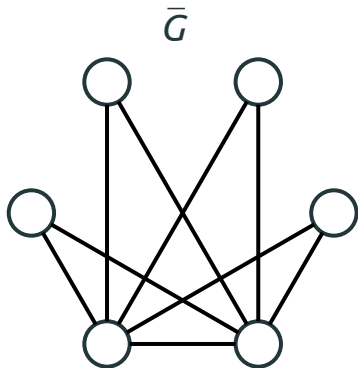
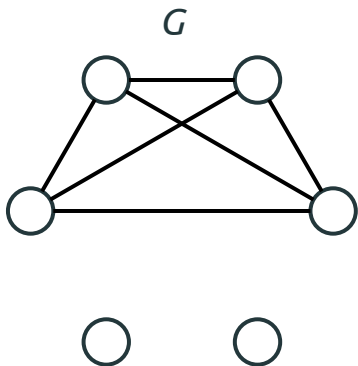
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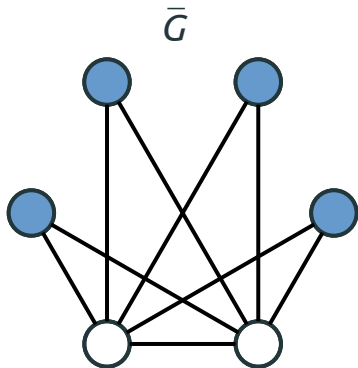
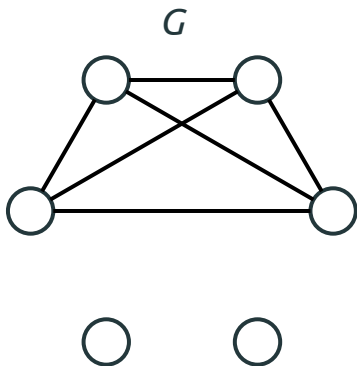
Complement Graph



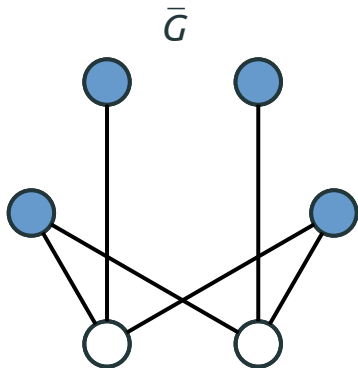
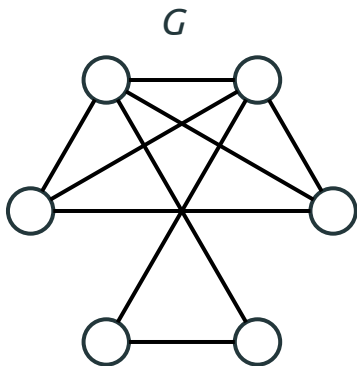
Complement Graph



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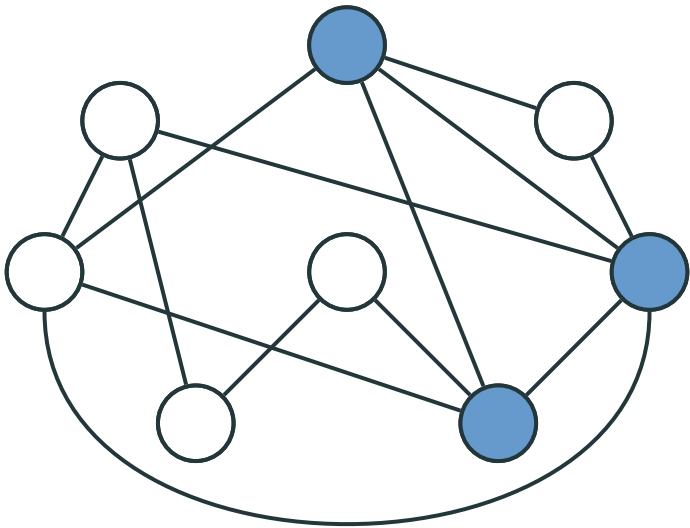
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- The **Clique Number** $\omega(G)$ of a graph G is the number of vertices in its maximum clique.

A graph with 8 nodes and 12 edges. The nodes are arranged in a roughly circular pattern. A cycle of 7 nodes is highlighted, with the 8th node connected to two nodes within this cycle. The edges are as follows: (1,2), (2,3), (3,4), (4,5), (5,6), (6,7), (7,1) forming the cycle; and (1,8), (2,8), (3,8), (4,8), (5,8), (6,8), (7,8) connecting the 8th node to the other 7 nodes.

Cliques: Examples

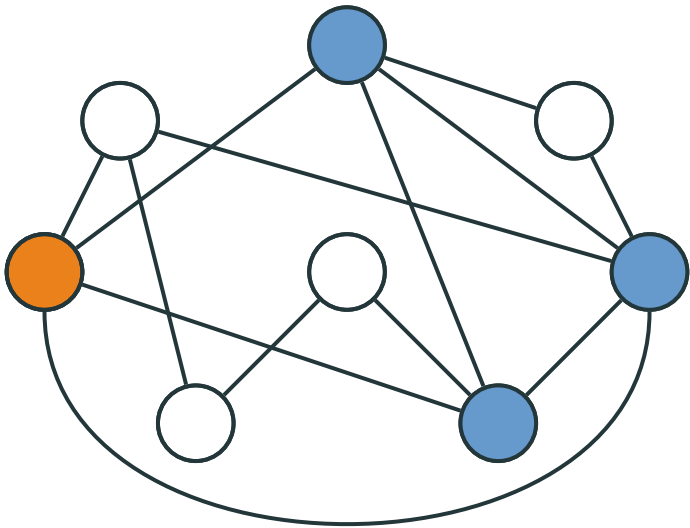
A Clique



Cliques: Examples

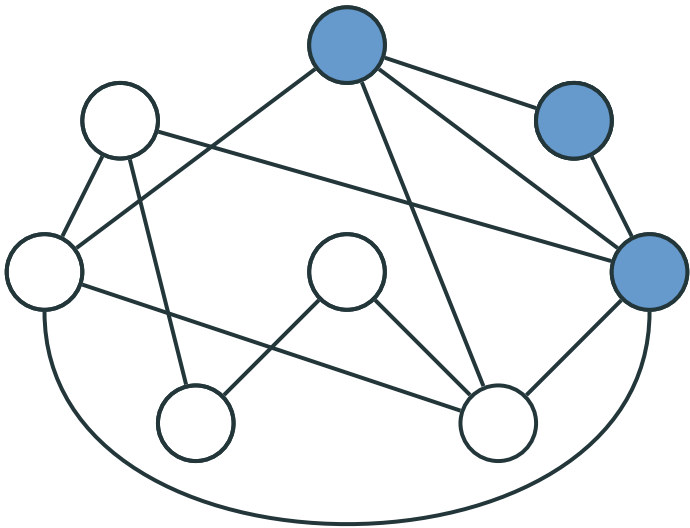
A Clique

Not a **Maxi-
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Cliques: Examples

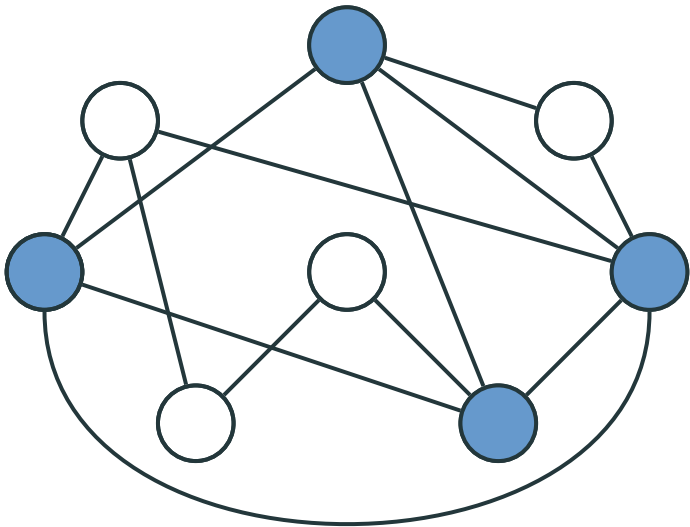
A Maximal
Clique: cannot be extended to a larger clique



Cliques: Examples

A Maximum
Clique:

there are
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than 4 ver-
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Independent Sets

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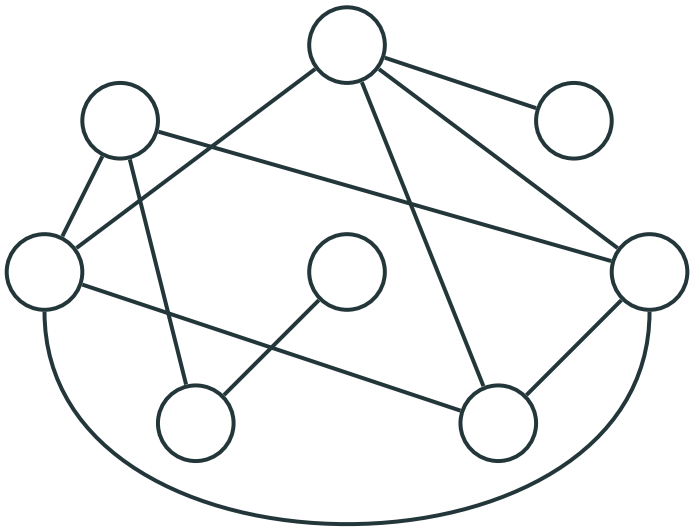
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Independent Sets

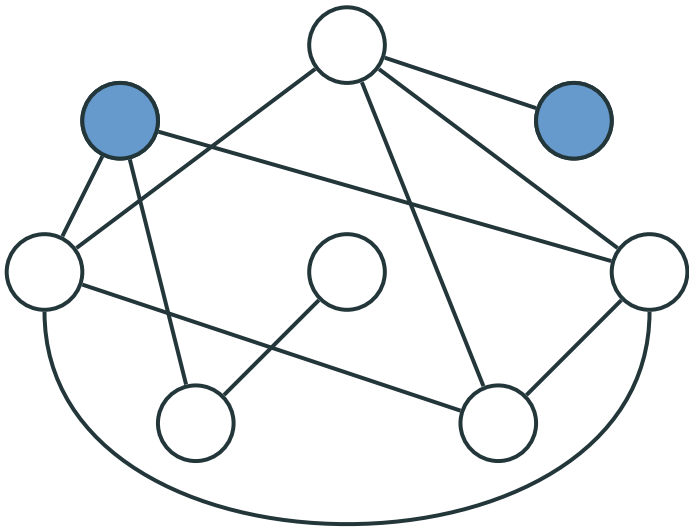
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- The **Independence Number $\alpha(G)$** of a graph G is the number of vertices in its maximum IS.

Independent Sets: Examples



Independent Sets: Examples

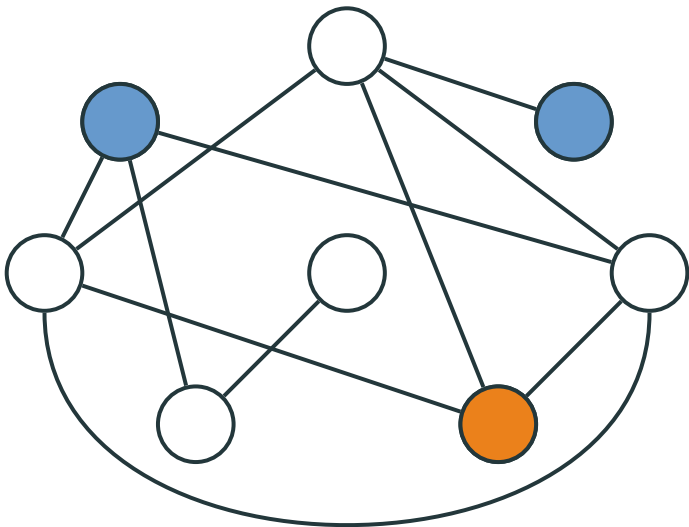
An IS



Independent Sets: Examples

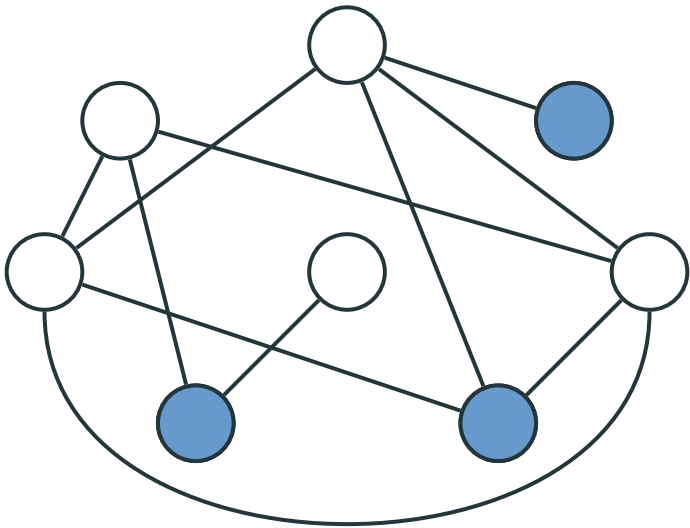
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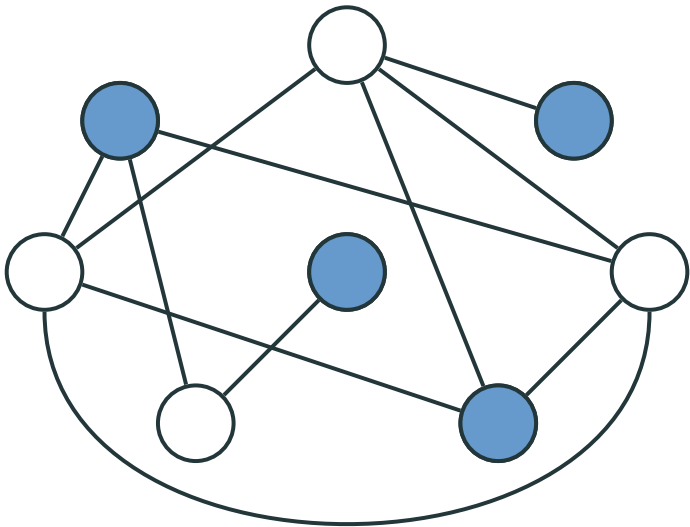
Independent Sets: Examples

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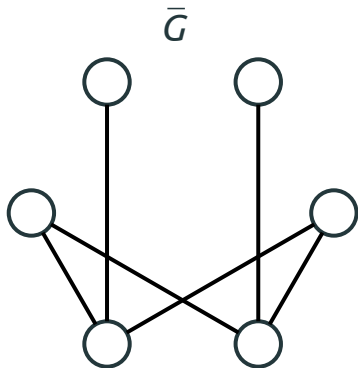
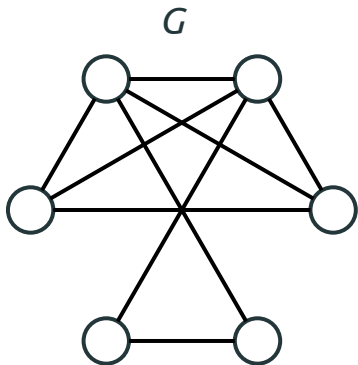


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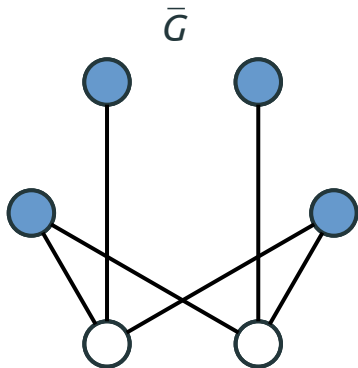
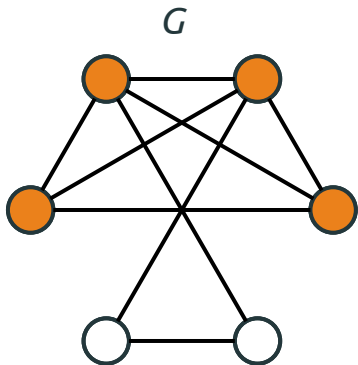
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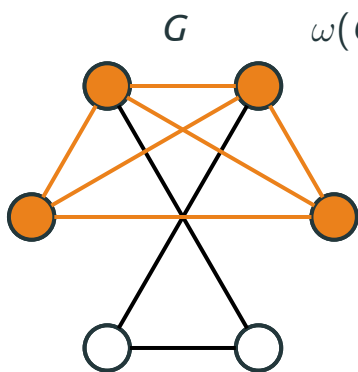
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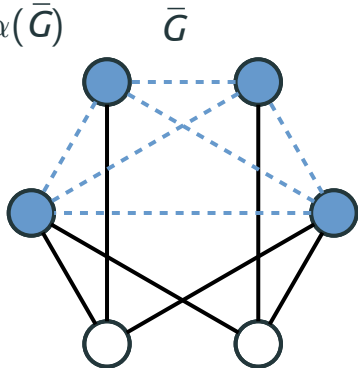
Cliques and IS's



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$$\omega(G) = \alpha(\bar{G})$$



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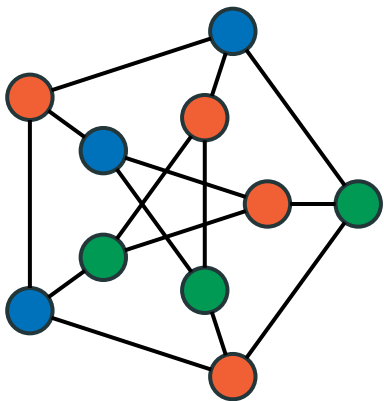
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Mantel's Theorem

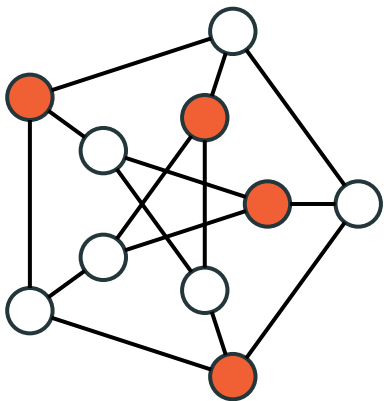
Coloring and IS's

Each color forms an independent set



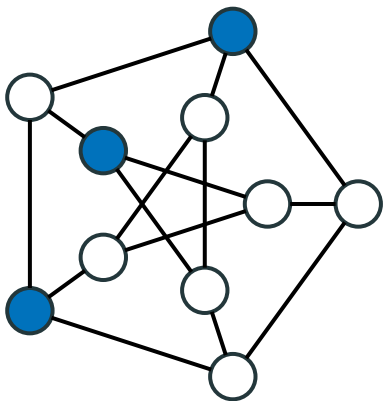
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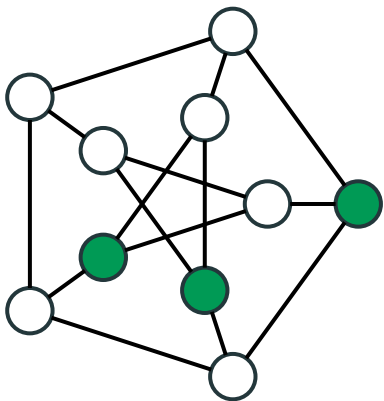
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Chromatic and Independence Numbers

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For every graph G on n vertices it holds that:

$$\chi(G) \cdot \alpha(G) \geq n .$$

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- Each IS is of size $\leq \alpha(G)$
- Therefore, $n \leq \chi(G) \cdot \alpha(G)$

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- **Base cases.** $n = 1, 2$: trivial
- **Induction hypothesis.** Holds for all graphs of size $\leq n - 2$
- **Induction step** will prove the statement for all graphs of size $\leq n$. Step of size 2, this is why we did base cases for $n = 1, 2$

Mantel's Theorem. Proof

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Tight Example: $K_{n/r, \dots, n/r}$

