



### Top-Down Parsing

Lecture 5

Exercise

#### Question?

Choose the next parse state given the grammar, parse table, and current state below. The initial string is:

if true then { true } else { if false then { false } } \$

	if	then	else	{	}	true	false	\$
E	if Bthen { E } E'				3	В	В	3
E'			else { E}		3			3
В						true	false	

-	Stack	Input	F_
Current	E'\$	else { if false then { false } } \$	_
0	\$	\$	E'-
$\bigcirc$	else { E } \$	else { if false then { false } } \$	В-
$\bigcirc$	E}\$	iffalse then { false } } \$	
$\bigcirc$	else { if Bthen { E} E'}\$	else { if false then { false } } \$	

E $\rightarrow$  if B then { E } E' | B |  $\epsilon$ E' $\rightarrow$  else { E } |  $\epsilon$ B $\rightarrow$  true | false

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#### Answer!

Choose the next parse state given the grammar, parse table, and current state below. The initial string is:

if true then { true } else { if false then { false } } \$

	if	then	else	{	}	true	false	\$
E	if Bthen { E } E'				3	В	В	3
E'			else { E}		3			3
В						true	false	

		Stack	Input	F
Cu	rrent	E'\$	else { if false then { false } } \$	_
	$\bigcirc$	\$	\$	E'
	0	else { E } \$	else { if false then { false } } \$	В-
	0	E}\$	iffalse then { false } } \$	
	$\bigcirc$	else { if Bthen { E} E'}\$	else { if false then { false } } \$	

E→ if B then  $\{E\}E'$   $|B|\epsilon$   $E' \rightarrow \text{else } \{E\} | \epsilon$   $B \rightarrow \text{true } | \text{false}$ 

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#### Question?

$S \rightarrow iE + SS' \mid a$	First(S) =	Follow(S) =
$S' \rightarrow e S \mid e$	First(S') =	Follow(S') =
$E \rightarrow b$	First(E) =	Follow(E) =

	α	Ь	e	i	†	\$
5						
5'						
E						

#### Answer!

$$S 
ightharpoonup i E + S S' \mid a$$
 First(S) = { i, a } Follow(S) = { e, \$ } S' 
ightharpoonup e S \mid e First(S') = { e,  $\epsilon$  } Follow(S') = { e, \$ } Follow(E) = { t }

	а	b	e	i	†	\$
5	a			i E † S S'		
5'			<b>e S</b> , ε			3
Е		Ь				

#### Question?

$$E \rightarrow T E'$$
  
 $E' \rightarrow + T E' \mid \in T \rightarrow F T'$   
 $T' \rightarrow * F T' \mid \in F \rightarrow (E) \mid id$ 

Fol	low	(F')	=
1 01	IUWI		_

	id	+	*	(	)	\$
E						
E'						
T						
Τ'						
F						

#### Answer!

$$E \rightarrow T E'$$
  
 $E' \rightarrow + T E' \mid \in T \rightarrow F T'$   
 $T' \rightarrow * F T' \mid \in F \rightarrow (E) \mid id$ 

```
First(E,T,F) = { id, ( }
First(E') = { +, ∈ }
First(T') = { *, ∈ }
Follow(E) = { $ , ) } Follow(E') = { $ , ) }
Follow(T) = { + , $, ) } Follow(T') = { + , $, ) }
Follow(F) = { *, + , $, ) }
```

	id	+	*	(	)	\$
Е	ΤE			ΤE		
E'		+ T E'			€	€
Т	FΤ			FΤ		
T'		€	* F T'		€	€
F	id			(E)		

#### Question?

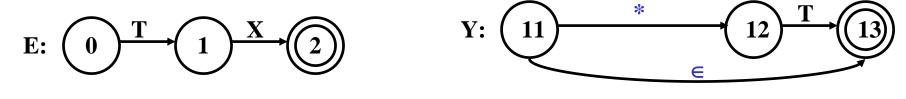
Consider the grammar

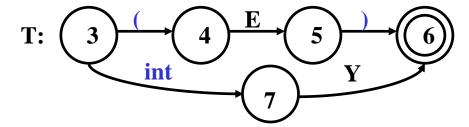
$$E \rightarrow TX$$
  $X \rightarrow + E \mid \varepsilon$   
 $T \rightarrow (E) \mid int Y$   $Y \rightarrow * T \mid \varepsilon$ 

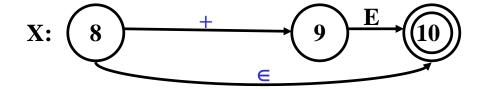
- · Convert the given grammar to a transition diagram
- · Simplify the Diagram (if it is possible)
- Write a step-by-step parsing of input 'int \* int'
- Draw the parse tree of the input

#### Answer!

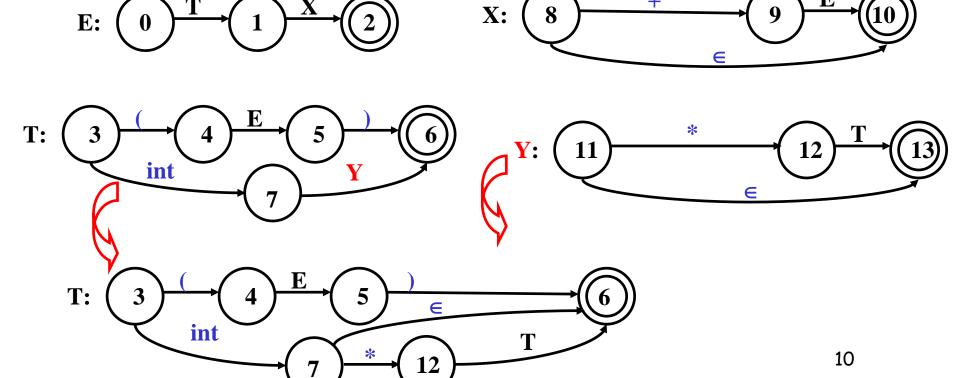
$$E \to TX T \to (E) | int Y$$
 
$$X \to +E | \in Y \to *T | \in$$



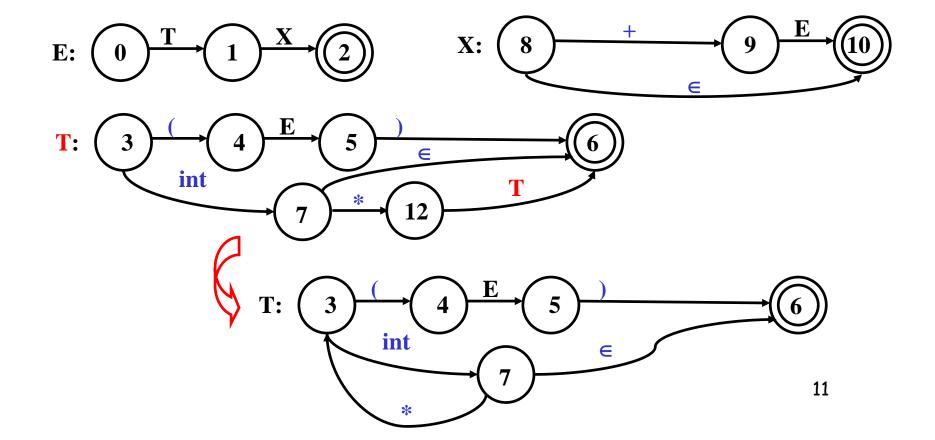




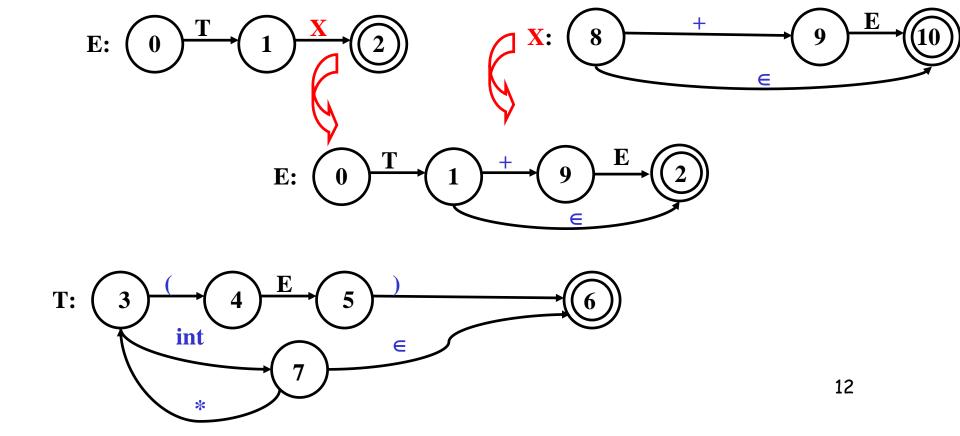
$$E \to TX \qquad X \to +E \mid \in T \to (E) \mid \text{int } Y \qquad Y \to *T \mid \in$$



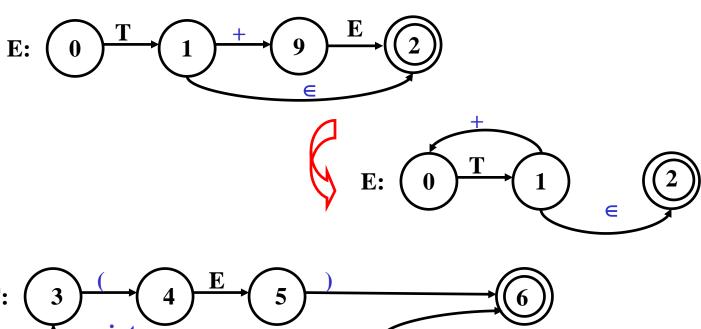
$$E \to TX T \to (E) | int Y$$
 
$$X \to +E | \in Y \to *T | \in$$

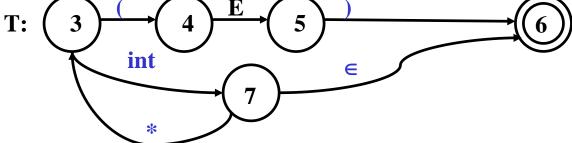


$$E \to TX T \to (E) | int Y$$
 
$$X \to +E | \in Y \to *T | \in$$



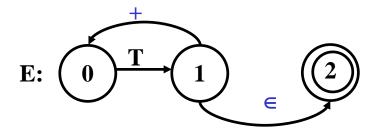
$$E \to TX \qquad X \to +E \mid \in T \to (E) \mid \text{int } Y \qquad Y \to *T \mid \in$$

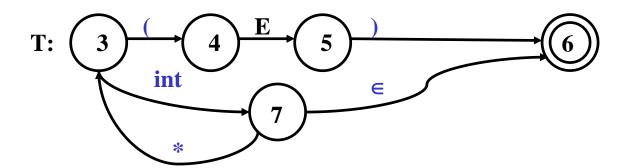


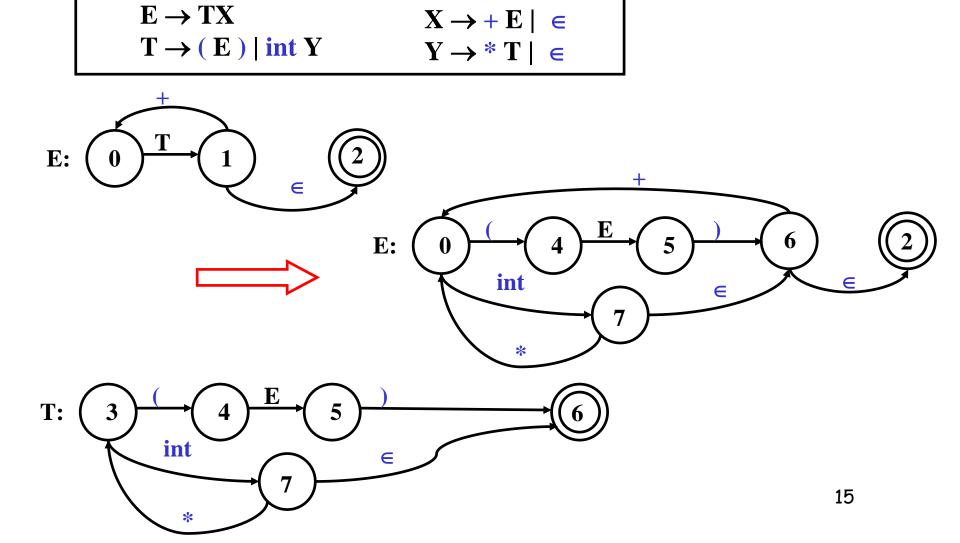


$$E \rightarrow T + E \mid T$$

$$T \rightarrow (E) \mid int * T \mid int$$







$$E \rightarrow (E) | (E) + E | int * E | int + E | int$$

