





$\therefore$  LR(0) parsing table

Rules:

$S \rightarrow (L) \text{ --- ①}$

$S \rightarrow a \text{ --- ②}$

$L \rightarrow L, S \text{ --- ③}$

$L \rightarrow S \text{ --- ④}$

	Action					Go to	
	(	)	a	,	\$	S	L
$I_0$	$S_2$		$S_3$	Accept			
$I_1$							
$I_2$						$I_5$	$I_4$
$I_3$	$R_2$	$R_2$	$R_2$	$R_2$	$R_2$		
$I_4$		$S_6$		$S_7$			
$I_5$	$R_4$	$R_4$	$R_4$	$R_4$	$R_4$		
$I_6$	$R_2$	$R_1$	$R_3$	$R_1$	$R_1$	$I_8$	
$I_7$	$S_2$		$S_3$			$I_8$	
$I_8$	$R_3$	$R_3$	$R_3$	$R_3$	$R_3$		



Answer to the question no - 02 (LR0)

Sol<sup>m</sup>:

Given that

$$E \rightarrow E + T / T \quad \text{let } E' \rightarrow E$$

$$T \rightarrow i$$

$\therefore$  Augmented Grammar

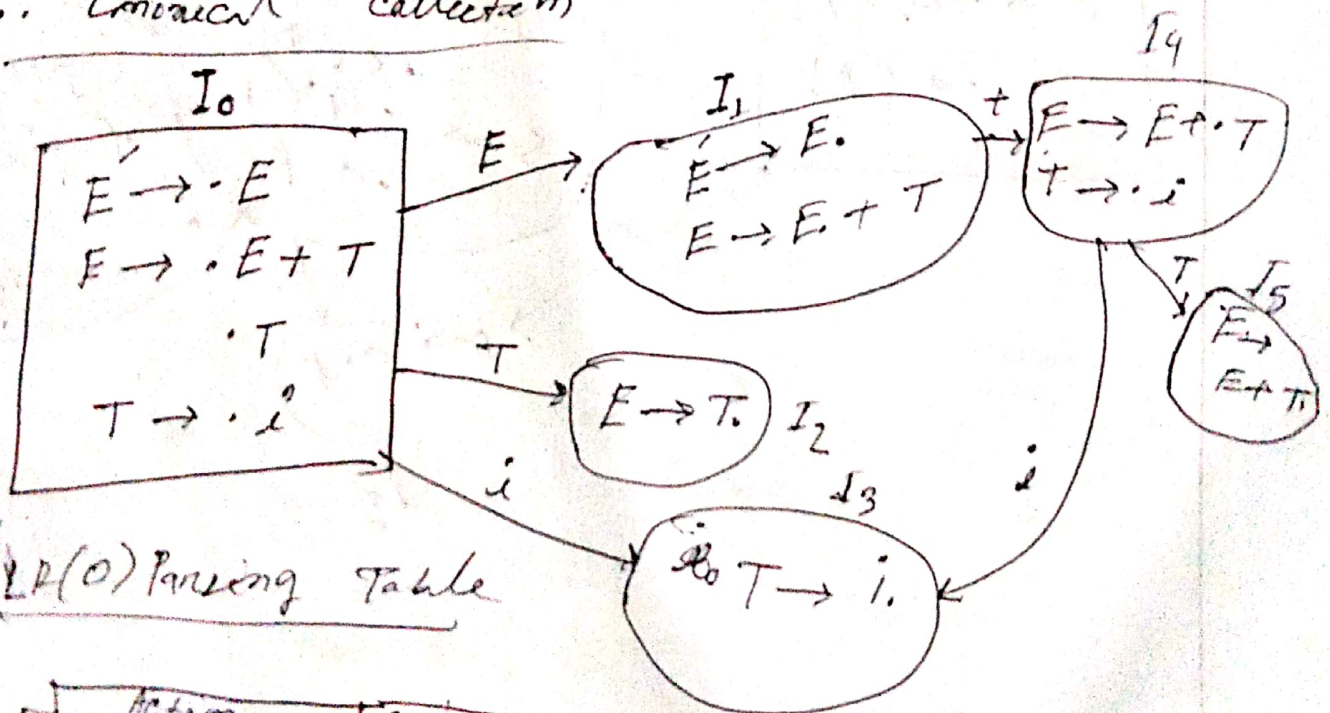
$$E' \rightarrow \cdot E$$

$$E \rightarrow \cdot E + T$$

$$\quad \quad \quad \cdot T$$

$$T \rightarrow \cdot i$$

$\therefore$  Canonical collection



LR(0) Parsing Table

	Action			Goto	
	+	i	\$	E	T
$I_0$	$S_0$	$S_0$	Accept	$I_1$	$I_2$
$I_1$	$S_1$				
$I_2$	$R_2$	$R_2$	$R_2$		
$I_4$					$I_5$
$I_5$	$R_1$	$R_1$	$R_1$		

Parser

$$E \rightarrow E + T \quad \text{--- (1)}$$

$$E \rightarrow T \quad \text{--- (2)}$$

$$T \rightarrow i \quad \text{--- (3)}$$



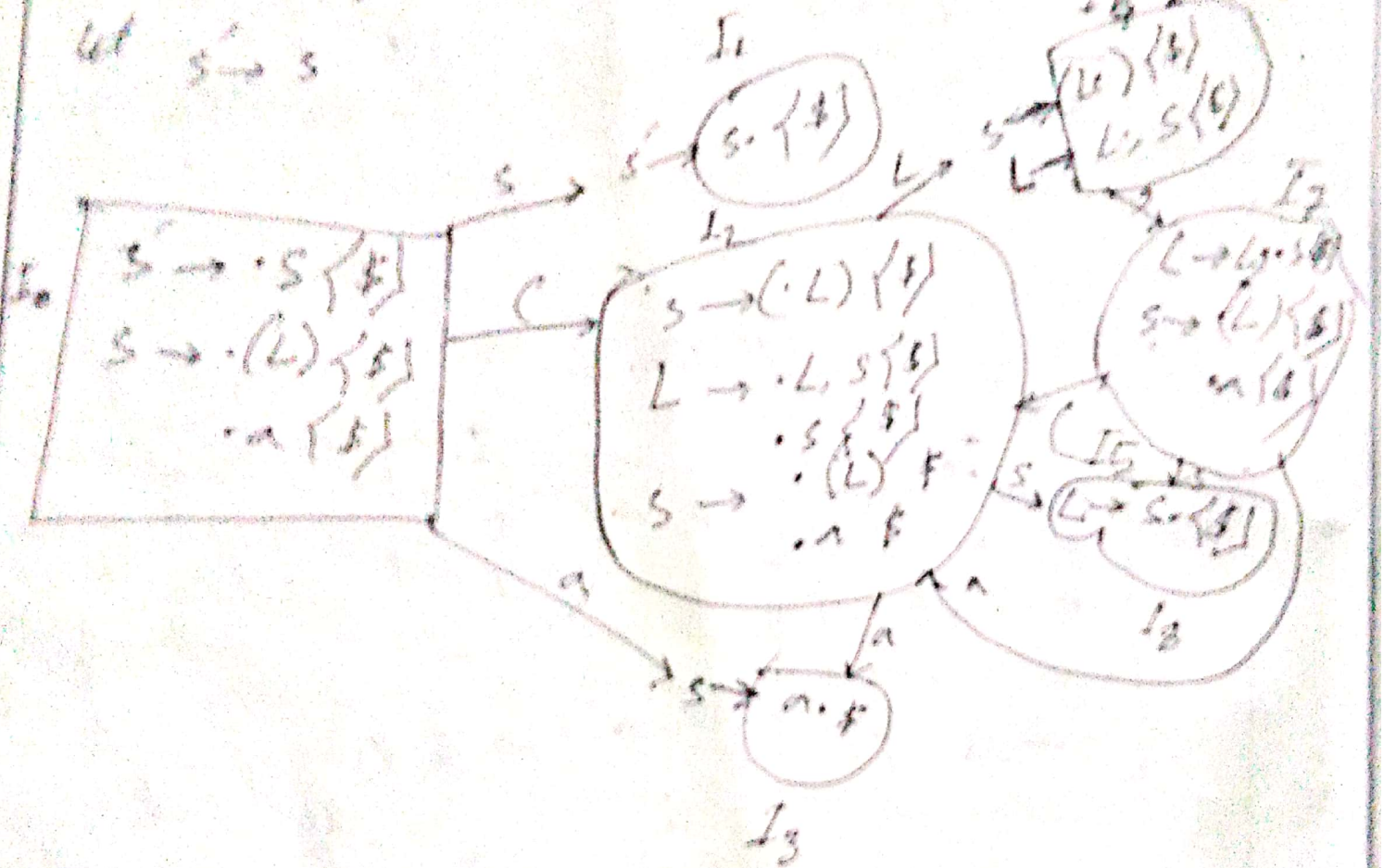
Answer to the question no-21 (LP1)

Given that,

$S \rightarrow (L) / a$

$L \rightarrow L S / S$

Let  $S' \rightarrow S$





## CLP(1) Parsing Table:

	Action					Go to	
	(	)	a	?	\$		
I <sub>0</sub>	S <sub>2</sub>		S <sub>3</sub>				
I <sub>1</sub>							
I <sub>2</sub>	S <sub>2</sub>		S <sub>3</sub>			I <sub>3</sub>	I <sub>4</sub>
I <sub>3</sub>					R <sub>2</sub>		
I <sub>4</sub>		S <sub>6</sub>		S <sub>7</sub>			
I <sub>5</sub>					R <sub>4</sub>		
I <sub>6</sub>					R <sub>1</sub>		
I <sub>7</sub>			S <sub>3</sub>			I <sub>8</sub>	
I <sub>8</sub>					R <sub>3</sub>		

## LALR

	Action					Go to	
	(	)	a	?	\$	S	L
I <sub>0</sub>	S <sub>2</sub>		S <sub>3</sub>				
I <sub>1</sub>							
I <sub>2</sub>	S <sub>2</sub>		S <sub>3</sub>			I <sub>4</sub>	I <sub>5</sub>
I <sub>3</sub>							
I <sub>4</sub>		S <sub>6</sub>		S <sub>7</sub>			
I <sub>5</sub>					R <sub>4</sub>		
I <sub>6</sub>					R <sub>1</sub>		
I <sub>7</sub>			S <sub>3</sub>			I <sub>8</sub>	
I <sub>8</sub>					R <sub>3</sub>		



Answer to the question no-002 (LP1)

Soln:

Given that

$$E \rightarrow E + T / T$$

$$T \rightarrow i$$

$$\text{let } E' \rightarrow E$$

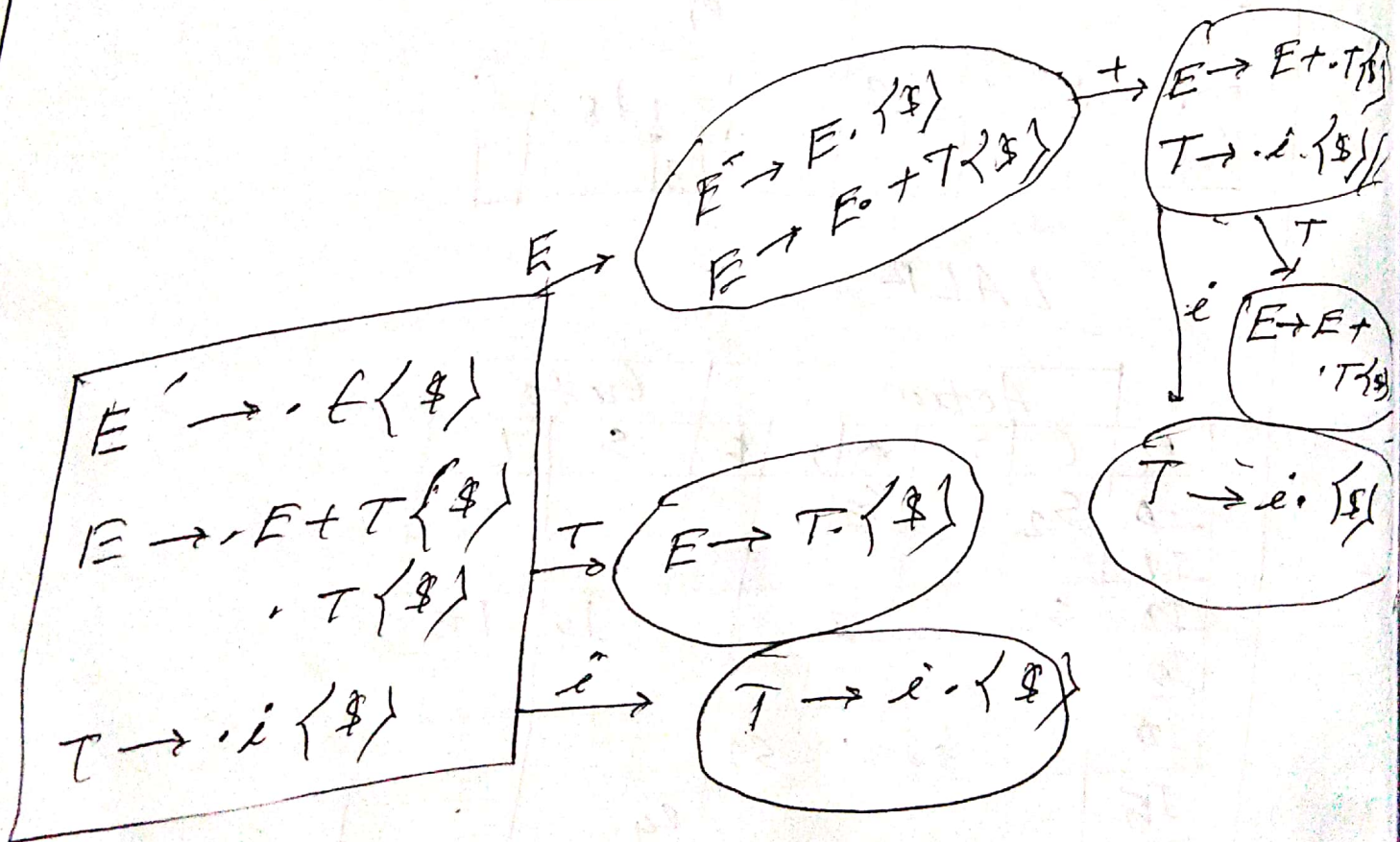
$\therefore$  Augmented Grammar

$$E' \rightarrow \cdot E \{ \$ \}$$

$$E \rightarrow \cdot ( + T \{ \$ \}$$

$$\cdot T \{ \$ \}$$

$$T \rightarrow \cdot i \{ \$ \}$$



CLR(1)

Rules

$$E \rightarrow E + T \rightarrow (1)$$

$$E \rightarrow T \rightarrow (2)$$

$$T \rightarrow i \rightarrow (3)$$

	Action			Goto	
	+	i	\$	E	T
I <sub>0</sub>		s <sub>3</sub>		I <sub>1</sub>	I <sub>2</sub>
I <sub>1</sub>	s <sub>4</sub>				
I <sub>2</sub>			r <sub>2</sub>		
I <sub>3</sub>			r <sub>3</sub>		
I <sub>4</sub>		s <sub>3</sub>			
I <sub>5</sub>			r <sub>1</sub>		

LALR

	Action			Goto	
	+	i	\$	E	T
I <sub>0</sub>		s <sub>3</sub>		I <sub>1</sub>	I <sub>2</sub>
I <sub>1</sub>	s <sub>4</sub>				
I <sub>2</sub>			r <sub>2</sub>		
I <sub>3</sub>			r <sub>3</sub>		
I <sub>4</sub>		s <sub>3</sub>			
I <sub>5</sub>			r <sub>1</sub>		