

## What is Being Done?

The program multiplies:

$$AC := AC \times N$$

- **AC** initially contains the multiplicand (say  $Y$ ).
- **N** is stored in memory location **mult**.
- Multiplication is performed via repeated addition:  $AC + AC + \dots$  (**N** times).

### Program Logic:

- Adds **AC** to a product register **N** times.
- A loop:
  - Adds current **AC** to **prod**.
  - Decreases **N** (in **mult**) by 1.
  - Tests if **N** == 0 using a BZ (Branch if Zero).
  - If not, loops back.

## Key Memory Locations

Label	Purpose
one	Constant value 1
mult	Multiplier $N$
ac	Initial AC value (i.e., $Y$ )
prod	Final product / result

## Assembly Program Breakdown

Line	Label	Instruction (Assembly)	Explanation
0	one	00...01	Memory location holding constant 1
1	mult	N	Multiplier $N$ stored here
2	ac	00...00	Backup of AC (value $Y$ )
3	prod	00...00	To store the product (starts at 0)
4		ST ac	Store initial AC ( $Y$ ) into ac
5	Loop	LD mult	Load $N$ into AC
6		BZ exit	Exit loop if $N == 0$
7		LD one	Load 1
8		MOV DR, AC	Move 1 into DR
9		LD mult	Reload $N$
10		SUB	$N := N - 1$
11		ST mult	Store updated $N$
12		LD ac	Load multiplicand ( $Y$ )
13		MOV DR, AC	Move $Y$ into DR
14		LD prod	Load current product
15		ADD	Add $Y$ to product
16		ST prod	Store updated product
17		BRA Loop	Jump back to loop start
18	exit		End of program

## Limitations of Accumulator-Based CPUs

### 1. Only one accumulator:

- Intermediate results must be moved to and from memory repeatedly.

### 2. Few CPU registers:

- Values like 1,  $N$ ,  $Y$ , and product are fetched repeatedly from memory.

### 3. Slower execution:

- More memory accesses = slower performance.
- If values could stay in dedicated CPU registers, the program would:
  - Run faster
  - Be shorter (fewer instructions)

## Summary

- This program multiplies two numbers via repeated addition using an accumulator and data register.
- It demonstrates the single-address, register-scarce design of accumulator-based CPUs.
- The architecture leads to inefficient memory use, limiting performance in data-heavy computations.