

Linked List

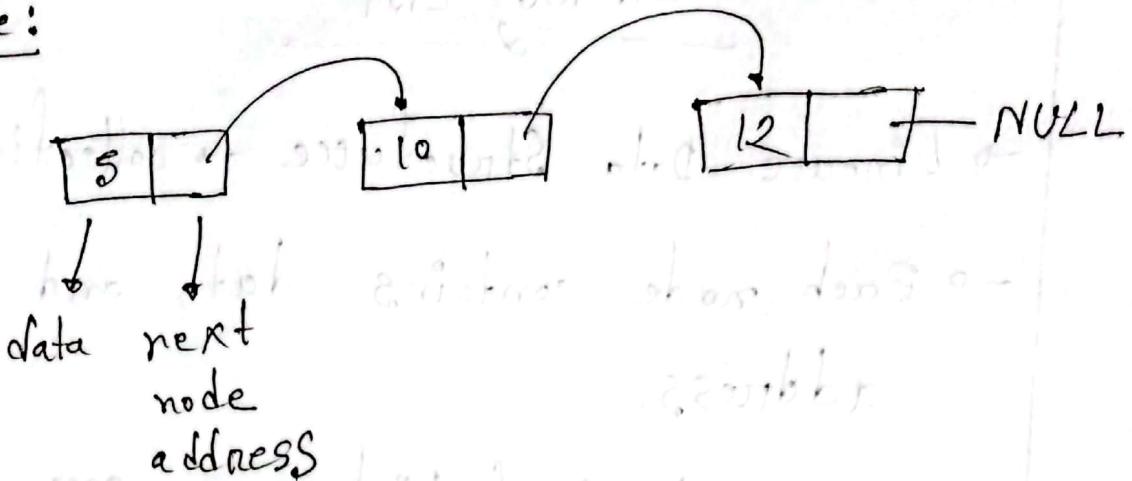
- Linear Data Structure → collection of nodes
- Each node contains data and next node address.
- Using Linked List we can insert or remove any elements of any position in the linked list.

Linked List

Array

1. Data Structure	Non-contiguous	contiguous
2. Memory Allocation	individual elements	whole array
3. Insertion/Deletion	Efficient	Inefficient
4. Access	Sequential	Random

Example:



class Node { encapsulate a Node }

public:

int data;

Node* next;

// constructor

Node (int data){

this → data = data;

this → next = NULL;

}

Access:

newnode → data

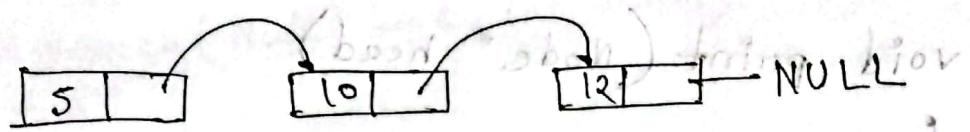
newnode → next

Types of Linked List

- Singly Linked List
- Doubly Linked List
- Circular Linked List
- Circular Doubly Linked List

Singly Linked List

Ex:



Operations on Singly Linked List:

- Traversal
- Searching
- Length
- Insertion:
 - at the beginning (head)
 - at the end (tail)
 - at specific position.

- Deletion:
 - from the beginning (Head)
 - from the end (tail)
 - at specific node

(beginning)

Traversal:

```
void print(Node* head);
```

```
{
```

```
// start from the beginning
```

```
Node* temp = head;
```

```
// traverse until null pointer
```

```
while (temp != NULL)
```

```
{
```

```
cout << temp->data << " "
```

```
temp = temp->next;
```

```
}
```

```
cout << endl;
```

Time Complexity : $O(N)$

Space Complexity : $O(1)$

call function:

```
print(head);
```

Searching in Singly Linked List:

```
bool search(Node* head, int target)
{
    Node* temp = head;
    while(temp != NULL)
    {
        if(temp->data == target)
        {
            return true;
        }
        temp = temp->next;
    }
    return false;
}
```

call function:

```
search(head, 5);
```

Time Complexity : O(N)
Space Complexity : O(1)

Finding Length in singly Linked List:

```
int Length(Node* head)
{
    int cnt = 0;
    Node* temp = head;
    while (temp != NULL)
    {
        cnt++;
        temp = temp->next;
    }
    return cnt;
}
```

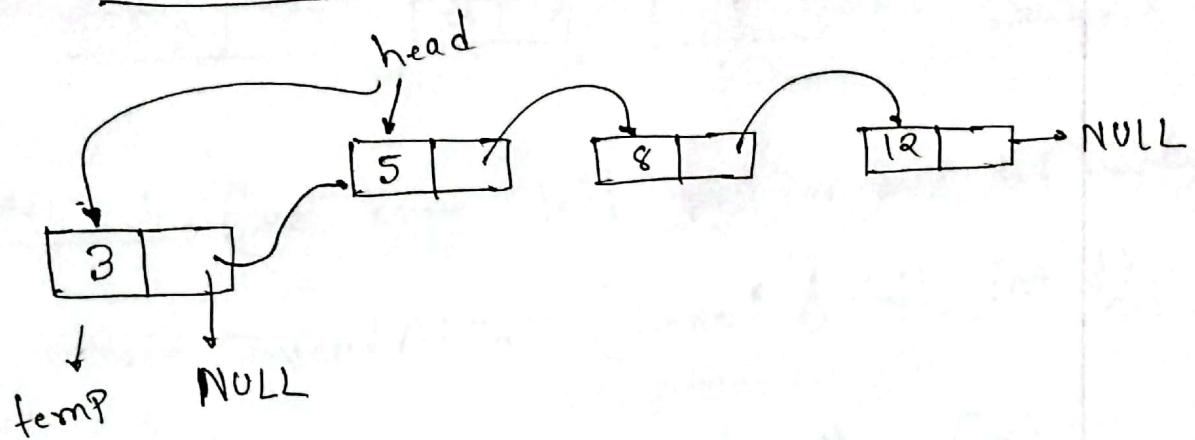
//call functions

```
Length(head);
```

Time Complexity: $O(N)$
Space Complexity: $O(1)$

Inserstion in Singly Linked List:

1. Insert at the beginning:



```
void InsertAtHead (Node* &head, int d)
```

```
{
```

```
    Node* temp = new Node(d); → create  
    new node
```

```
    temp->next = head;
```

```
    head = temp;
```

```
}
```

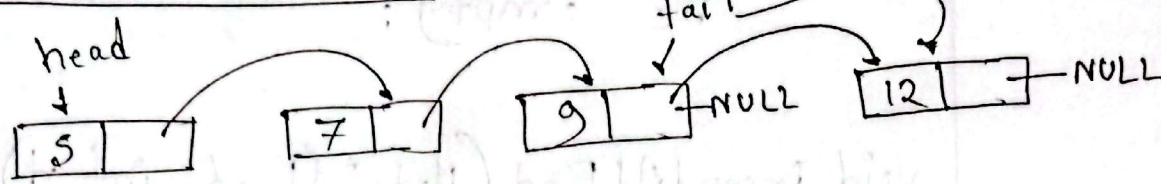
Time Complexity: $O(1)$

Space Complexity: $O(1)$

call function:

```
InsertAtHead (head, 3);
```

2. Insert at the End :



Case - 1: (we know tail position and Not empty)

void InsertAtTail (Node* &tail, int d)

```
{  
    Node* temp = new Node(d);
```

```
    tail->next = temp;
```

```
    tail = temp;  
}
```

call function:

InsertAtTail (tail, d);

Case-2: we don't know tail and it can
be empty:

void InsertAtEnd (Node* &head, int d)

{
 Node *temp = new Node(d);

 if (empty list)

 if (head == NULL)

 {
 temp->next = NULL;
 head = temp;
 return;
 }

T.C = O(N)
S.C = O(1)

 Node *curr = head;

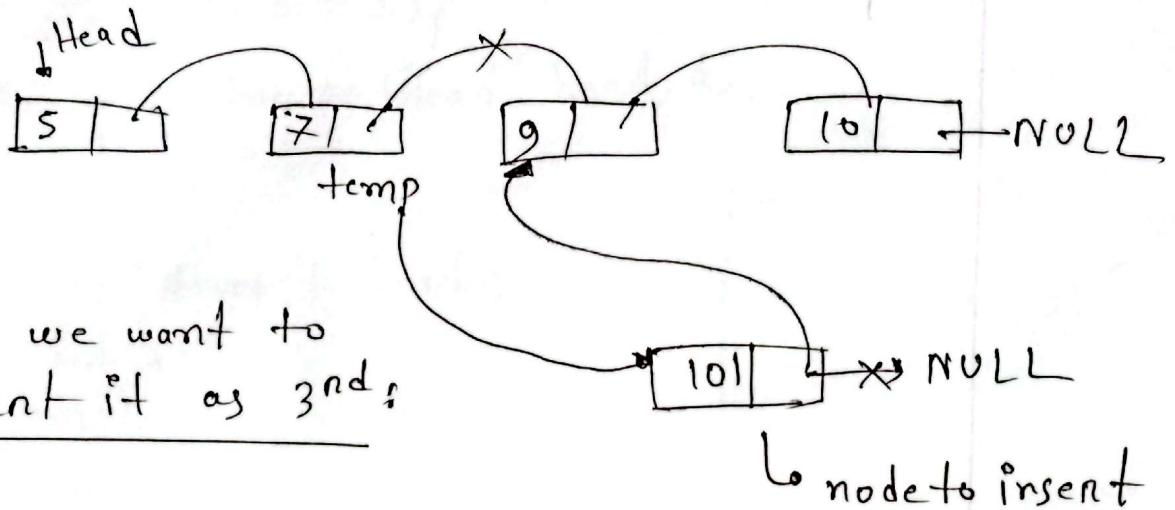
 while (curr->next != NULL)

 {
 curr = curr->next;
 }

 curr->next = temp;

}

Insert at a Specific Position of the singly
Linked List:



Suppose we want to
insert it as 3rd:

$n \rightarrow$ position

traverse $(n-1)$

```

void InsertAtPosition(Node* &head, Node* &tail, int pos, int d)
{
    if (pos == 1) {
        InsertAtHead(head, d);
        return;
    }

    Node* temp = head;
    int cnt = 1;
    while (cnt < pos - 1) {
        temp = temp->next;
        cnt++;
    }

    if (temp->next == NULL) {
        InsertAtTail(tail, d);
        return;
    }

    Node* nodeToInsert = new Node(d);
    nodeToInsert->next = temp->next;
    temp->next = nodeToInsert;
}

```

T.C = O(N)
 S.C = O(1)

call function: InsertAtPosition (head, tail, 5, 7);

Deletion in Singly Linked List

④ First Node :



Node^{*} temp = head;

head = head → next;

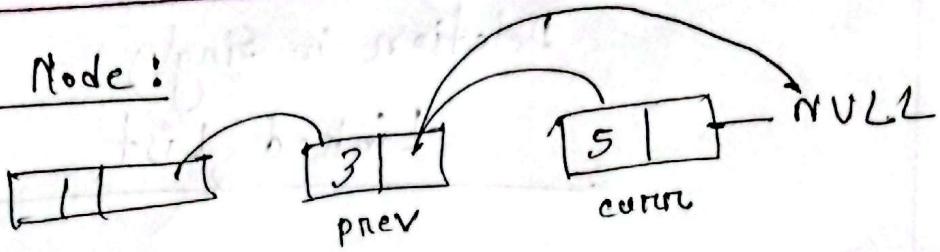
temp → next = NULL;

delete temp;

Time Complexity → O(1)

Space Complexity → O(1)

Last Node:



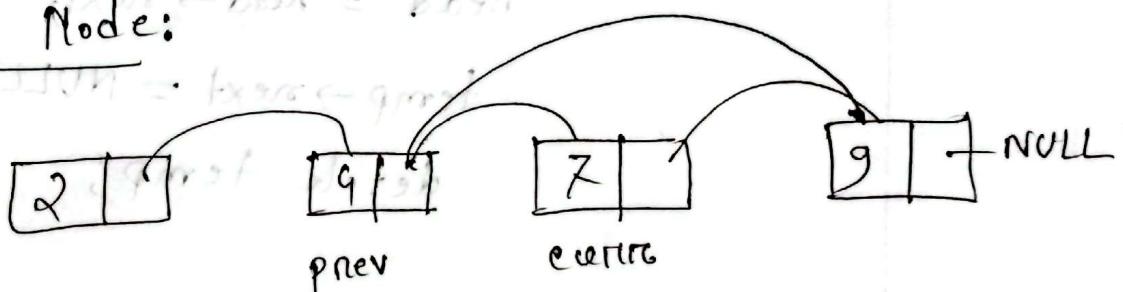
~~(1)~~ $\text{pPrev} \rightarrow \text{next} \rightarrow \text{curr} \rightarrow \text{next}$

$\text{curr} \rightarrow \text{next} \rightarrow \text{NULL}$

~~delete curr;~~

T.C $\rightarrow O(N)$
S.C $\rightarrow O(1)$

Middle Node:



(1) $\text{pPrev} \rightarrow \text{next} \rightarrow \text{curr} \rightarrow \text{next}$

$\text{curr} \rightarrow \text{next} \rightarrow \text{NULL}$

~~delete curr;~~

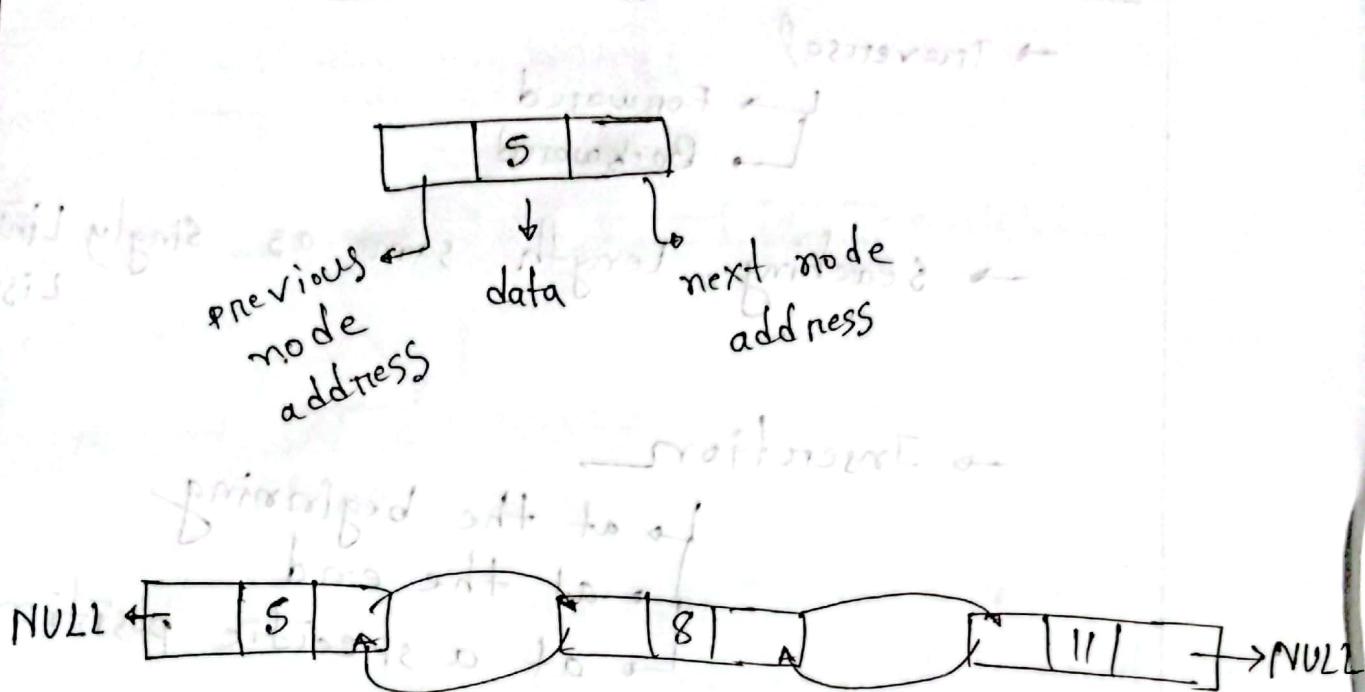
T.C $\rightarrow O(N)$
S.C $\rightarrow O(1)$

```

void DeleteNode (Node* &head, int pos)
{
    if (pos == 1) {
        Node* temp = head;
        head = head->next;
        temp->next = NULL;
        delete temp;
    }
    else {
        Node* curr = head;
        Node* prev = NULL;
        int cnt = 1;
        while (cnt < pos) {
            prev = curr;
            curr = curr->next;
            cnt++;
        }
        prev->next = curr->next;
        curr->next = NULL;
        delete curr;
    }
}

```

Doubly Linked List



```
class Node {  
public:  
    Node* pprev;  
    int data;  
    Node* next;  
    Node (int d){  
        this->data = d;  
        this->pprev = NULL;  
        this->next = NULL;  
    }  
}
```

Operations:

→ Traversal

- └─ Forward
- └─ Backward

→ Searching, Length same as singly Linked List

→ Insertion

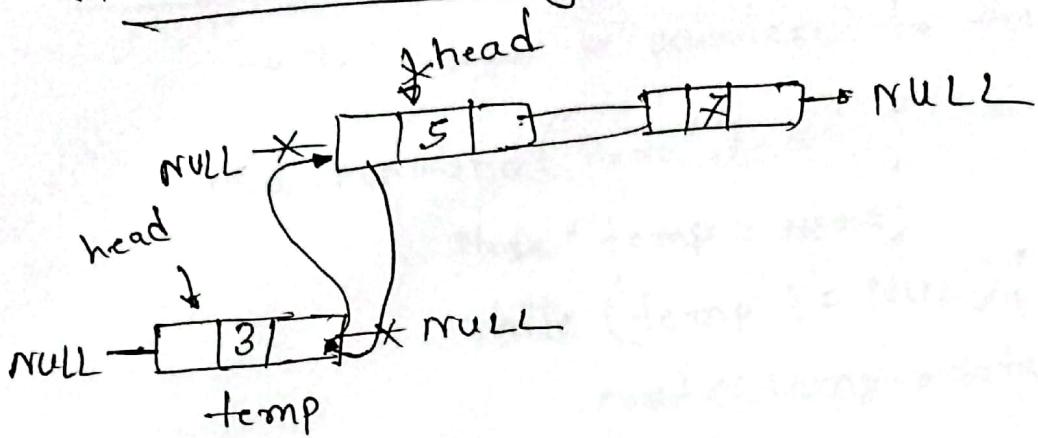
- └─ at the beginning
- └─ at the end
- └─ at a specific position

→ Deletion

- └─ from the beginning
- └─ from the end
- └─ at a specific position

④ Insertion in Doubly Linked List:

④ At the beginning:



→ $\text{Node}^* \text{temp} = \text{new Node}(d)$
→ $\text{temp} \rightarrow \text{next} = \text{head}$
→ $\text{head} \rightarrow \text{prev} = \text{temp}$
→ $\text{head} = \text{temp}$

④ If head is NULL, then,

```
if (head != NULL) {  
    head → prev = temp  
}  
}
```

□ Traversal in Doubly Linked List:

□ Forward Traversal:

→ Initialize a pointer to the head.

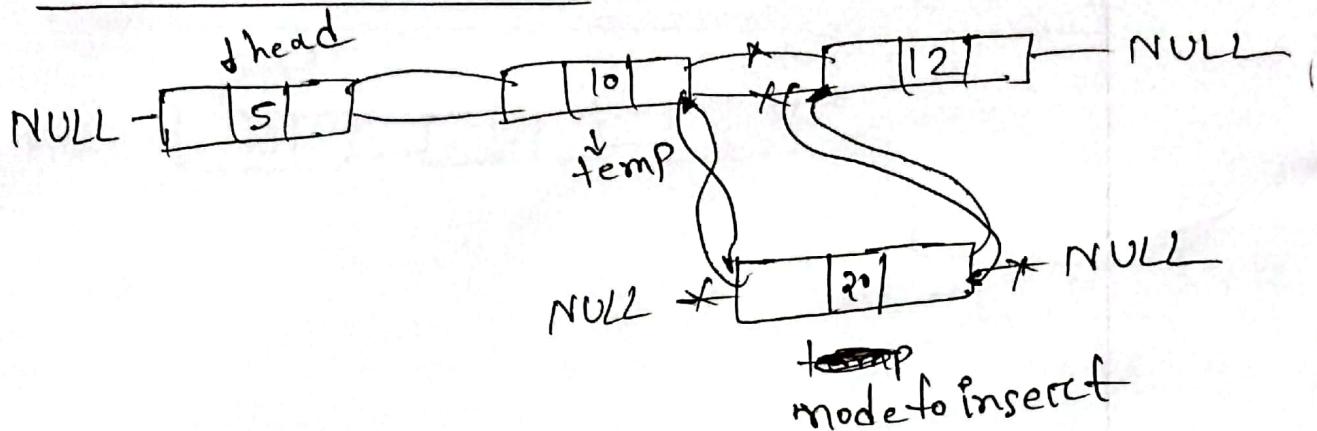
```
void Forward(Node *head){  
    Node *temp = head;  
    while (temp != NULL){  
        cout << temp->data << " ";  
        temp = temp->next;  
    }  
    cout << endl;  
}
```

□ Backward Traversal:

→ Initialize a pointer to the tail

```
void Backward(Node *tail){  
    Node *temp = tail;  
    while (temp != NULL){  
        cout << temp->data << " ";  
        temp = temp->prev;  
    }  
    cout << endl;  
}
```

④ Insert at middle:



$\rightarrow \text{Node}^* \text{node to insert} = \text{new Node}(d)$

$\text{node to insert} \rightarrow \text{next} = \text{temp} \rightarrow \text{next}$

$\text{node to insert} \rightarrow \text{prev} = \text{node to insert}$

$\text{temp} \rightarrow \text{next} \rightarrow \text{prev} = \text{node to insert}$

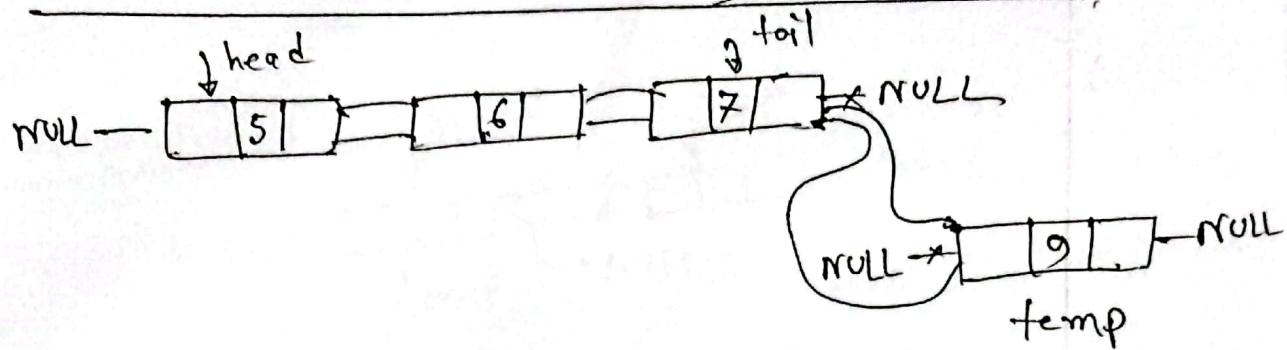
$\text{temp} \rightarrow \text{next} = \text{node to insert}$

$\text{node to insert} \rightarrow \text{prev} = \text{temp}$

④ If $\text{pos} = 1$, Insert At head

$\text{pos} = \text{Last}$ Insert At tail

④ Insert At the End of Doubly Linked List



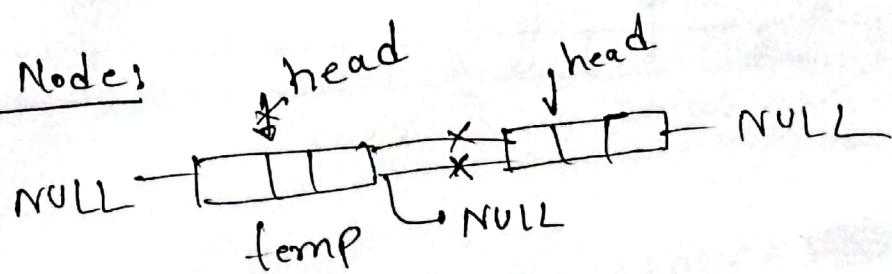
→ `Node *temp = new Node(d)`
→ `tail → next = temp`
→ `temp → prev = tail`
→ `tail = temp`

⑤ For empty list,

```
if (tail == NULL) {  
    tail = temp;  
    return;  
}
```

Delete

First Nodes



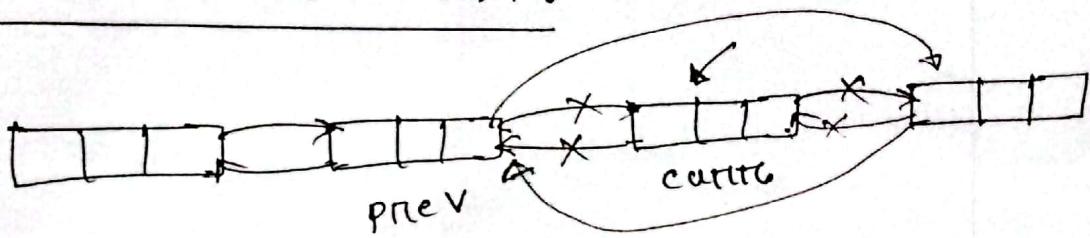
$\rightarrow \text{temp} \rightarrow \text{next} \rightarrow \text{pnew} = \text{NULL}$

$\rightarrow \text{head} \Rightarrow \text{temp} \rightarrow \text{next}$

$\hookrightarrow \text{temp} \rightarrow \text{next} = \text{NULL}$

delete temp;

Delete the Middle or Last:



$\text{curr} \rightarrow \text{next} \rightarrow \text{prev} = \cancel{\text{curr}} \text{prev}$

$\text{curr} \rightarrow \text{prev} = \text{NULL}$

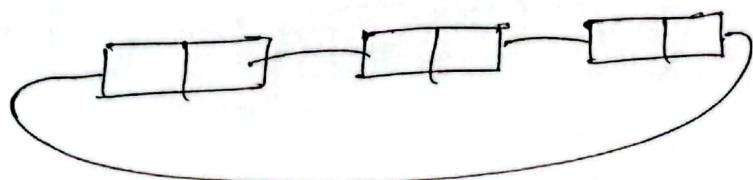
$\text{prev} \rightarrow \text{next} = \text{curr} \rightarrow \text{next}$

$\text{curr} \rightarrow \text{next} = \text{NULL}$

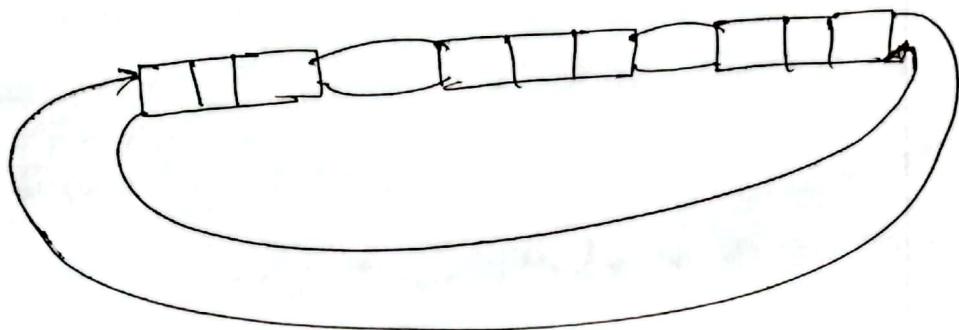
`delete curr;`

Circular Linked List

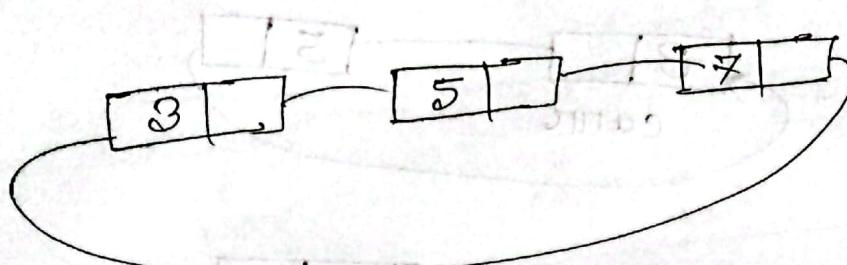
Singly :



Doubly :



Circular Singly Linked List



Insert :

(D) shall move = arrst * shall o-

⇒ Empty list :

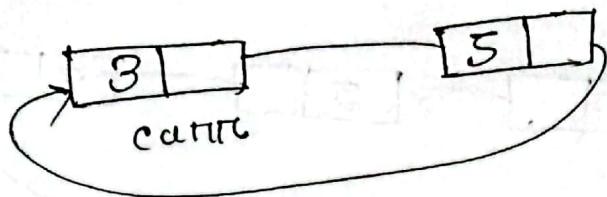
arrst = arrst * new Node (d)
newNode = new Node (d)

arrst = arrst * tail = newNode)

tail → next = newNode



Non-empty:



i/p → element

here = 3

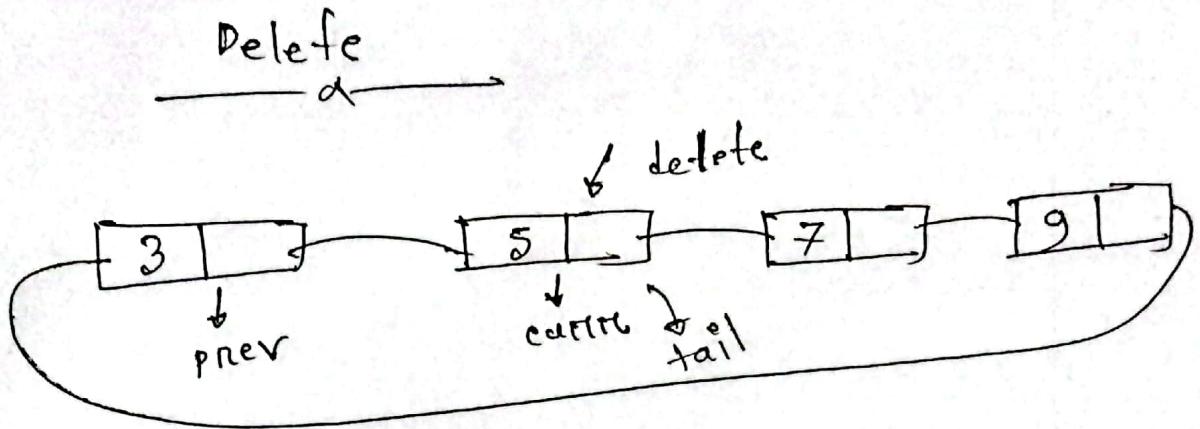
temp

→ Node * temp = new Node(4);

(2) ~~old~~ curr → ~~old~~ temp → next = curr → next

~~old~~ curr → next = temp





i/p element = 5, pnext → next = current → next

1 Node pnext = current
 tail = NULL

2 Node tail = current
 tail = pnext

current → next = NULL;

delete current