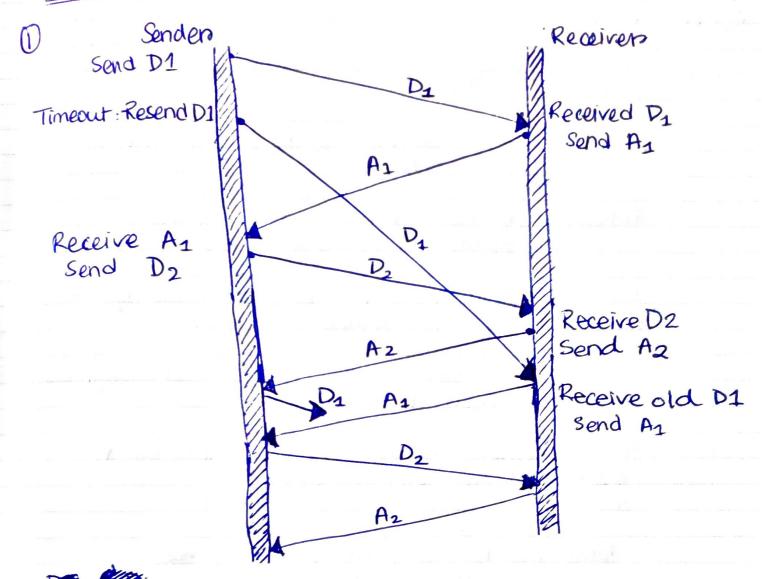
Assignment2-Paret 1 Mehezabin Ahamed 8524484



In the diagram, D1 & D2 are data with sequence no. 1 & 2. and same for A1 & A2 acknowledgements.

2. a Bandwidth = 16bps, segment size = 576 octets.

1 octet = 8 bits. 16b = 109 bits 60 ms = 0.06s

Window Size = $10\cancel{9}0.06 \approx 13,000$ segments. 576×8

Following timeout, window size halves. window size ~ 6500 segments

trom this point, there will be exponential growth before the congestion threshold is passed.

window size increases to $2^{13} = 8192$ segments, where RTT = 13.

Now, it follows additive increases until we veach the full window size.

This takes 13000 - 8192 = 4808 RTT.

Total RTT = 4808 + 13 = 4821

Total time to reach full window = 4821 x 0.06
= 289.26 seconds

(b) Window size = $\frac{10^9 \times 000}{16 \times 10^3 \times 8} \approx 460$ segments.

Threshold when halved after timeout = 230 seg $2^8 = 256$ segments, where RTT=8

For additive increase, 460-256=204 RTT Total RTT= 212 RTT

Total time to reach full window size,

=> 212 x0.06.
= 12.72 seconds.

$$\mu = \frac{1.5 \times 10^6}{4000}$$
 bits/s = 375 packets/sec

Utilization,
$$f = \frac{2}{n} = \frac{300}{375} = 0.8$$

mean
$$\not\approx$$
 of packets in queue,
 $\Rightarrow \frac{f}{1-f} = \frac{0.8}{1-0.8} = 4$

Probability that given has
(1-
$$f$$
)(f^2)
$$= (1-0.8)(0.8^2)$$

$$= 0.128$$

(2) packets

$$(1-9)(9^3)$$

 $=(1-0.8)(0.8^3)$
 $=0.102$

The input rate per (second) =
$$\frac{50 \text{ gal}}{60 \text{ sec}} = \frac{576}{60 \text{ sec}} = \frac{576}{60 \text{ sec}} = \frac{576}{60 \text{ sec}} = \frac{215}{60} =$$