```
Minclude <string.h>
Fdefine MAXPAROLA 30
#define MAXRIGA 80
   int treq[MAXPAROLA]; /* vettore di contatoni
delle frequenze delle lunghezze delle perole
   char riga[MAXRIGA] ;
lint i, inizio, lunghezza
```

# **Synchronization**

# **Introduction to Synchronization**

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## Introduction

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#### Introduction

Critical Section (CS) or Critical Region (CR)

These are sections of code or variables that are accessed by multiple threads. Proper synchronization ensures that only one thread accesses these sections at a time to prevent inconsistencies.

A section of code, common to multiple threads, in which each thread can read and write shared objects

Critical Sections (CS): These are the parts of your program that access shared resources and must be executed by only one thread at a time.

- Access to CS is subject to race conditions
  - The result depends on the execution order of the processes instructions

Race Conditions: This happens when the outcome of a program depends on the sequence or timing of uncontrollable events such as the scheduling of threads.

# **Example**

FIFO, Queue, Circular Buffer

```
void enqueue (int val) {
  if (n>SIZE) return;
  queue[tail] = val;
  tail=(tail+1)%SIZE;
  n++;
  return;
}

register = n
  register = register + 1
  n = register
```

```
Code Explanation:
```

The enqueue function adds an element to the queue.

The dequeue function removes an element from the queue.

```
int dequeue (int *val) {
  if (n<=0) return;
  *val=queue[head];
  head=(head+1)%SIZE;
  n--;
  return;
}

register = n
  register = register - 1
  n = register</pre>
```

Even if enqueue and dequeue operate on the different ends of the queue, the variable n is shared

#### Shared Variable (n):

Even though enqueue and dequeue operate on different ends of the queue, they both modify the shared variable n, which tracks the number of elements in the queue.

A race condition can occur if both functions try to modify n simultaneously, leading to incorrect updates.

Race condition:
Increments and decrements can
be lost

If enqueue and dequeue are called simultaneously:

enqueue reads n, increments it, and stores it back.

dequeue reads n, decrements it, and stores it back.

Without proper synchronization, both may read the same value of n and update it incorrectly.

#### **Critical sections**

Preventing Race Conditions: Use access protocols to enforce mutual exclusion (ensuring only one thread accesses the critical section at a time).

Code that runs after leaving the critical section to signal that the critical section is now free for other threads to use.

# Release Section: Race conditions could be prevented with an access protocol that enforces mutual exclusion for each CS

Reservation Section:

Code that runs before entering the critical section to check if the critical section is available. If another thread is using it, the current thread waits.

- > Before a CS, there should be a reservation section
  - The reservation code must block (lock out) the P (or T) if another P (or T) is using its CS
- > After the CS, there should be a release section
  - The release possibly unlocks another P (or T) which was waiting in the "reservation" code of its CS

Non-critical

# **Access protocol**

```
T_i
                                                                        while
           while (TRUE)
                                                                                     (TRUE)
                                                   Structure of Critical Section Protection:
                                                   Reservation Code: Check and wait if
                                                   necessary until the critical section is
                reservation code
                                                                            reservation code
                                                   Critical Section Code: The actual code
                Critical Section
                                                                            Critical Section
                                                   that accesses shared resources.
                                                   Release Code: Signal that the critical
                                                   section is now free.
                release code
                                                                            release code
Sections: These are parts
or your code that do not access shared esources non critical section
                                                                            non critical section
and do not require
synchronization
```

- Every Critical Section is protected by an
  - > Enter code (reservation, or prologue)
  - > Exit code (release, or epilogue)
- Non-critical sections should not be protected
- OSs provide appropriate primitives

Locks and Unlocks:

In the reservation code, threads try to acquire a lock (using mutex.lock()). If the lock is already taken, they wait,

In the release code, threads release the lock (using mutex.unlock()), allowing other waiting threads to proceed.