/**

* Title: Homework 1

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* Assignment : 1

* Description: Question 1, 2 and 3

*/

Question 1:

(a) $f(n) = 20 n^4 + 20 n^2 + 5$ is $O(n^5)$ by the Big-O definition, if $f(n) \le cn^5$ for some $n \ge n_0$. If $20 n^4 + 20 n^2 + 5 \le cn^5$ then $\frac{20}{n} + \frac{20}{n^3} + \frac{5}{n^5} \le c$. Therefore, this condition holds for $n \ge n_0 = 1$ and $c \ge 45 = (20 + 20 + 5)$.

(b) Sort the array [18, 4, 47, 24, 15, 24, 17, 11, 31, 23] in ascending with selection and bubble sort.

Selection sort:

Note: "|" separates the sorted and unsorted part.

Initial array: [18, 4, 47, 24, 15, 24, 17, 11, 31, 23]

After 1st swap: [18, 4, 23, 24, 15, 24, 17, 11, 31 | 47]

After 2nd swap: [18, 4, 23, 24, 15, 24, 17, 11 | **31**, 47]

After 3rd swap: [18, 4, 23, **11**, 15, 24, 17 | **24**, 31, 47]

After 4th swap: [18, 4, 23, 11, 15, 17 | 24, 24, 31, 47]

After 5th swap: [18, 4, **17**, 11, 15 | **23**, 24, 24, 31, 47]

After 6th swap: [15, 4, 17, 11 | 18, 23, 24, 24, 31, 47]

After 7th swap: [15, 4, **11**| **17**, 18, 23, 24, 24, 31, 47]

After 8th swap: [11, 4| 15, 17, 18, 23, 24, 24, 31, 47]

After 9th swap: [4| 11, 15, 17, 18, 23, 24, 24, 31, 47]

Sorted array: [4, 11, 15, 17, 18, 23, 24, 24, 31, 47]

Bubble sort:

Note: "|" separates the sorted and unsorted part. Initial array: [18, 4, 47, 24, 15, 24, 17, 11, 31, 23]

Pass 1

Initial array: [18, 4, 47, 24, 15, 24, 17, 11, 31, 23]]

[4, 18, 47, 24, 15, 24, 17, 11, 31, 23]

[4, 18, **24, 47**, 15, 24, 17, 11, 31, 23]

[4, 18, 24, **15, 47**, 24, 17, 11, 31, 23]

[4, 18, 24, 15, **24, 47**, 17, 11, 31, 23]

[4, 18, 24, 15, 24, **17, 47**, 11, 31, 23]

[4, 18, 24, 15, 24, 17, **11, 47**, 31, 23]

[4, 18, 24, 15, 24, 17, 11, **31, 47**, 23]

[4, 18, 24, 15, 24, 17, 11, 31, 23 | 47]

Pass 2

[4, 18, 24, 15, 24, 17, 11, 31, 23 | 47]

[4, **18, 24**, 15, 24, 17, 11, 31, 23 | 47]

[4, 18, **24, 15**, 24, 17, 11, 31, 23 | 47]

[4, 18, 15, **24, 24**, 17, 11, 31, 23 | 47]

[4, 18, 15, 24, **24, 17**, 11, 31, 23 | 47]

[4, 18, 15, 24, 17, **24, 11**, 31, 23 | 47]

[4, 18, 15, 24, 17, 11, **24, 31**, 23 | 47]

[4, 18, 15, 24, 17, 11, 24, **31, 23** | 47]

[4, 18, 15, 24, 17, 11, 24, 23 | **31, 47**]

Pass 3

[**4, 18**, 15, 24, 17, 11, 24, 23 | 31, 47]

[4, **18, 15**, 24, 17, 11, 24, 23 | 31, 47]

[4, 15, **18, 24**, 17, 11, 24, 23 | 31, 47]

[4, 15, 18, **24, 17**, 11, 24, 23 | 31, 47]

[4, 15, 18, 17, **24, 11**, 24, 23 | 31, 47]

[4, 15, 18, 17, 11, **24, 24**, 23 | 31, 47]

[4, 15, 18, 17, 11, 24, **24, 23** | 31, 47]

[4, 15, 18, 17, 11, 24, 23 | **24, 31, 47**]

Pass 4

[4, 15, 18, 17, 11, 24, 23 | 24, 31, 47]

[4, **15, 18**, 17, 11, 24, 23 | 24, 31, 47]

[4, 15, **18, 17**, 11, 24, 23 | 24, 31, 47]

[4, 15, 17, **18, 11**, 24, 23 | 24, 31, 47]

[4, 15, 17, 11, **18, 24**, 23 | 24, 31, 47]

[4, 15, 17, 11, 18, **24, 23** | 24, 31, 47]

[4, 15, 17, 11, 18, 23 | **24, 24, 31, 47**]

Pass 5

[4, 15, 17, 11, 18, 23 | 24, 24, 31, 47]

[4, **15, 17**, 11, 18, 23 | 24, 24, 31, 47]

[4, 15, **17, 11**, 18, 23 | 24, 24, 31, 47]

[4, 15, 11, **17, 18**, 23 | 24, 24, 31, 47]

[4, 15, 11, 17, **18, 23** | 24, 24, 31, 47]

[4, 15, 11, 17, 18 | **23, 24, 24, 31, 47**]

Pass 6

[4, 15, 11, 17, 18 | 23, 24, 24, 31, 47]

[4, **15**, **11**, 17, 18 | 23, 24, 24, 31, 47]

[4, 11, **15, 17**, 18 | 23, 24, 24, 31, 47]

[4, 11, 15, **17, 18** | 23, 24, 24, 31, 47]

[4, 11, 15, 17 | **18, 23, 24, 24, 31, 47**]

Sorted array: [4, 11, 15, 17, 18, 23, 24, 24, 31, 47]

Question 2:

C:\Users\meli						\Debug\l	-IW1.exe							
Insertion Sort) 2	Comp 3	Count: 5	59 Move 6	Count: 7	89 8	9	9	11	11	14	15	16	17	18
lerge Sort) 2	Comp 3	Count: 5	46 Move 6	Count: 7	128 8	9	9	11	11	14	15	16	17	18
uick Sort			47 Move								4.5	4.5	47	40
2		5	6	7	8	9	9	11	11	14	15	16	17	18
	R	andom A												
Part c - Time	analysis			ort										
rray Size			Elapsed			Count		Move(
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.5000		216			5616			56192						
0000		385			9976			99802						
25000		602				155895541 155945539								
art c - Time	analysis													
rray Size			Elapsed			Count		Move(12361						
.000 .0000		2 3			55199 1205			26723						
5000		5			1894			41723						
2000		5			2610			57446						
5000		8			3341			73446						
art c - Time	analysis													
rray Size			Elapsed			Count		Move(
000		0			6911			10791						
.0000 .5000		1			1487			24103						
19000		2 4			2606! 3591:			41832 64509						
5000		3			4619			77601						
Part c - Time : Array Size	analysis		ertion S Elapsed	OI.F	Comp	Count		Move	`ount					
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.0000		0			0			19998	3					
5000		0			ø			29998						
0000		0			0			39998						
5000		1			0			49998	3					
Part c - Time	analycic	of Mon	go Sont											
art C - Time : rray Size	analysis		Elapsed		Comp	Count		Move	Count					
000		1	Liapsca		32004			12361						
.0000		4			6900			26723						
5000		3			1063			41723						
0000		4			1480			57446						
5000		4			1884	76		73446	54					
art c - Time	analysis													
rray Size			Elapsed			Count		Move(
0000		42			1249			14997						
.0000		140 263			4999			29997						
15000 10000		263 482				92500 90000		44997 59997						
5000		742				87500		74997						
		/ 12			3127			, 4557						
			ecution											

Performance of Sorting Algorithms (Random Arrays)



Figure 1: Performance of the Sorting Algorithms (Random Arrays)

Performance of Sorting Algorithms (Already Sorted Arrays)



Figure 2: Performance of the Sorting Algorithms (Already Sorted Arrays)

Algorithm	Time Complexities						
	Best	Average	Worst				
Insertion Sort	$\Omega(n)$	$\theta(n^2)$	O(n ²)				
Merge Sort	$\Omega(nlog(n))$	$\theta(nlog(n))$	O(nlog(n))				
Quick Sort	$\Omega(nlog(n))$	$\theta(nlog(n))$	$O(n^2)$				

Table 1: Time Complexities of the Sorting Algorithms

From the data it can be observed that insertion sort worked in $O(n^2)$ for random arrays but worked in o(n) for the already sorted arrays. Time elapsed is not clear for the insertion sort on already sorted arrays but it can be observed from the number of moves. So the experiment shows that for the insertion sort the best case is when array is already sorted.

Another observation is that merge sort worked in O(nlog(n)) for both random and already sorted arrays. It can said that the empirical result and the theoretical result are same.

Quick sort worked in $O(n^2)$ for already sorted array and worked in $O(n\log(n))$ for random arrays. Theoretically, it is obvious that already sorted arrays are the worst case of the quick sort when pivot is the first element. In quick sort selecting the pivot is very important, when pivot is selected near to the average it works in $O(n\log(n))$ however in our example pivot is always the first element and in the already sorted array first element is the smallest. The smallest element should not be selected as a pivot because it prohibit dividing the array two balanced sub arrays.

Question 3:

When the experiment results observed it can be seen that elapsed time of insertion sort is directly proportional to the K. This result is reasonable because when K increases, the move and comparison counts also increases. Move count and comparison count are the two things that determines the elapsed time.

For the merge and quick sort, it looks like the change of K has no significant effect on elapsed time. Both of them works in O(nlog(n)) as always. When K = 0, which means that array is already sorted, elapsed time for quick sort is very high because first element is being selected as the pivot. This cause unbalanced subarrays for the quick sort and it cause the algorithm works in $O(n^2)$.

	**********	*****Question 3*****	**************
Size: 5000			
K	Insertion Sort	Merge Sort	Quick Sort
0	0	4	30
500	4_	4	1
1000	7	2	1
1500	9	0	0
2000	18	1	0
2500	20	4	0
3000	26	0	0
3500	42	0	4
4000	45	0	4
4500	38	4	0
Size: 10000			
K	Insertion Sort	Merge Sort	Quick Sort
0	0	1	138
1000	11	7	0
2000	33	4	4
3000	83	4	0
4000	92	0	4
5000	58	0	4
6000	76	2	2
7000	129	4	0
8000	171	4	0
9000	197	8	4
Size: 15000			0:15:
K	Insertion Sort	Merge Sort	Quick Sort
0	0	10	304
1500	56	11	4
3000	56	1	6
4500	91	4	0
6000 7500	178	2	5
9000	138	6 4	2 2
	166		3
10500 12000	186 247	4 2	3 4
13500	254	2	4
Size: 20000		2	4
312e: 20000 K	Insertion Sort	Merge Sort	Quick Sort
0	0	Herge 30Pt	496
2000	73	7	4
4000	102	6	4
6000	148	7	3
8000	179	4	4
10000	204	4	4
12000	236	4	4
14000	314	4	6
16000	355	4	4
10000		7	4