# Lockstep

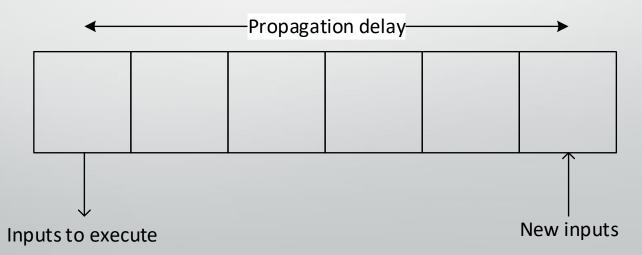
Study and implementation of a network library for distributed deterministic simulations

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- Inputs for the simulation are generated by clients
- Each client computes the new state of the simulation by applying inputs
- All the clients need a coherent view of the state
- Case study: Multiplayer RTS

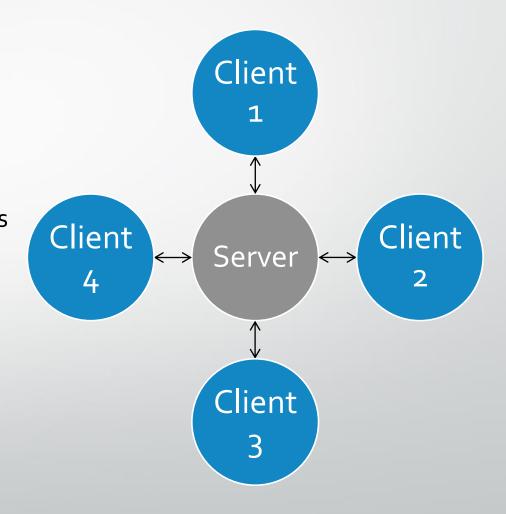
#### The FrameQueue

- Time is divided in chunks called *frames*. Each input is assigned to a frame
- All clients execute the same input during the same frame
- By scheduling new inputs for a later frame, we can introduce enough delay to propagate it to all clients and keep executing without interruptions
- The delay introduced should be dimensioned for the application requirements and network characteristics



#### Network architecture

- Client-Server architecture
  - Each clients sends its inputs to the server
  - The server forwards to each client others' inputs
- Better performance and scalability
  - Assuming better client-server connection
  - P2P would easily saturate client's upload and increase delays

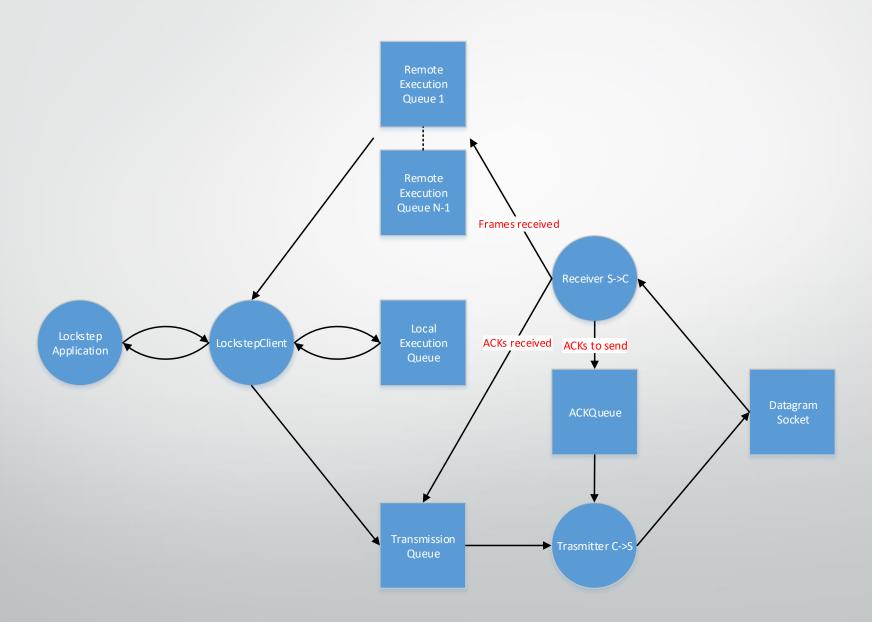


## Network protocol

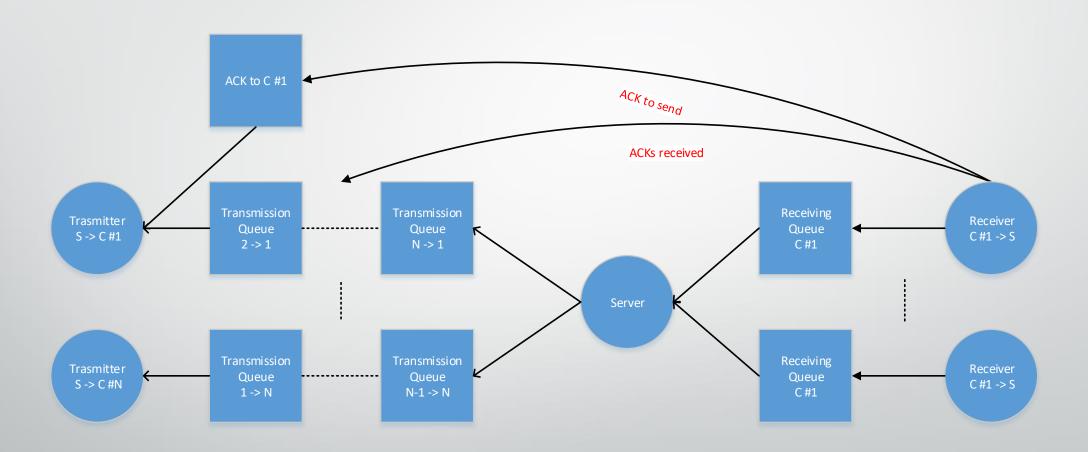
- Custom protocol based on UDP
  - Reliability
  - Minimum delay
- LockstepCommand: interface implementing Serializable
  - Implemented by the application command
  - Using Java serialization leads to high bandwidth usage
  - Minor solution by adding GZIP compression of outgoing packets

Software demonstration

#### Client Structure



#### Server structure



## Client-side Synchronization (1)

- N clients  $C_i$ , each one with N execution queues  $E_j$
- $E_j$  stores every received input generated by  $C_j$
- ullet  $C_i$  and  $E_i$  keep track of the next frame number, let's say k
- $lue{c}_i$  needs command k from every  $E_j$  to execute frame k
- If some command is missing in its queue, client stops and wait

#### Client-side Synchronization (2)

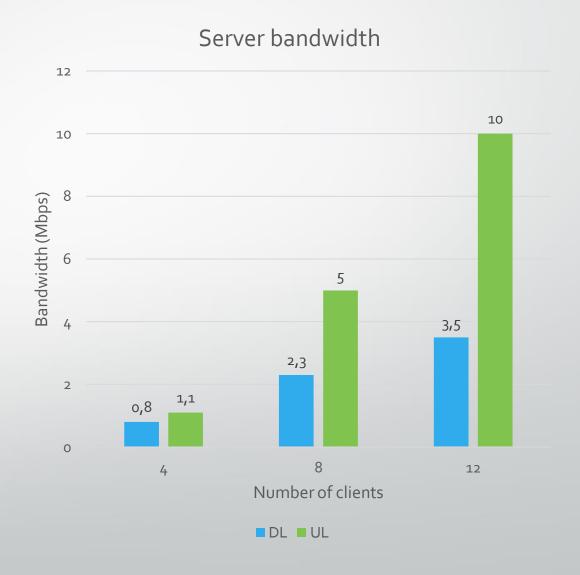
- Synchronization is achieved with a Semaphore
- $E_i$  releases a permit when it has the next command k
- ullet  $C_i$  tries to acquire N permits from semaphore
  - Failure to acquire: signal the application to suspend simulation, wait for all permits
  - Permits available: continue execution, resume simulation signal if suspended

## Server-side Synchronization

- N execution queues  $E_j$  for storing input received by  $C_j$
- Semaphore is used
- ullet  $E_i$  releases a permit as soon as it receives a command
- Server waits on semaphore
- When input is available, server forwards it to transmitters
- Clients waits for in-order command, server waits for any command

## Scalability

- The protocol has low bandwidth needs for clients
  - Receives O(N) messages
  - Sends O(1) messages
- Server instead has high bandwidth needs
  - Receives O(N) messages
  - Sends  $O(N^2)$  messages



## Future developments

- Receivers and transmitters are mostly idle
- Use idle time to improve serialization
  - Key objective: reducing bandwidth
- Optional P2P architecture
  - For high-bandwidth clients (LAN)
- Adaptive input-to-execution delay
- Security
  - Authentication
  - Integrity
  - State validation (anti-cheat)

