Analysis I Solutions

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Introduction

1.1 What is analysis?

No exercises in this section.

1.2 Why do analysis?

No exercises in this section.

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Starting at the beginning: the natural numbers

2.1 The Peano axioms

No exercises in this section.

2.2 Addition

Ex. 2.2.1. Fix b and c and induct on a. For a=0 we have (0+b)+c=b+c=b+(c+0), proving the base case. Now suppose (a+b)+c=a+(b+c). Then,

$$(a++b)+c = (a+b)+++c$$

= $((a+b)+c)++$
= $(a+(b+c))++$
= $a+++(b+c)$.

Ex. 2.2.2. First we show existence by induction, starting at a=1. We have 1=0++ by definition. Now suppose there exists a b such that a=b++. Then a++=(b++)++ will do. Now for uniqueness. If a=b++ and a=b'++ then by Axiom 2.4 we have b=b'.

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- Ex. 2.2.3. (a) Since a = a + 0 we have $a \ge a$.
 - (b) Since $a \ge b$ there exists a natural number m such that a = b + m. Similarly, since $b \ge c$ there exists a natural number n such that b = c + n. Thus, a = b + m = c + (m + n) whence $a \ge c$.
 - (c) Since $a \ge b$ there exists a natural number m such that a = b + m. Similarly, since $b \ge a$ there exists a natural number n such that b = a + n. Thus, a = b + m = a + m + n whence m + n = 0 by the Cancellation law, and so m = n = 0 by Corollary 2.2.9.
 - (d) If $a \ge b$ then a = b + n for some natural number n. Thus a + c = b + c + n whence $a + c \ge b + c$. Conversely, if $a + c \ge b + c$ then a + c = b + c + n for some natural number n. Applying the Cancellation law yields a = b + n whence $a \ge b$.
 - (e) If a < b then b = a + n for some natural number n. Note that $n \neq 0$ since then we would have a = b, which by definition is not the case. Since $n \neq 0$ we can, by Lemma 2.2.10. write n = m + + for some unique natural number m. Thus b = a + m + + a + m whence $a + + \leq b$.

Conversely, if $a++ \leq b$ then b=a+++n, so b=a+n++. The point being that $n \neq 0$ by Axiom 2.3 whence $a \neq b$, or in other words, a < b.

Conversely, if b = a+d for some d > 0 then by Lemma 2.2.10 there exists a unique c such that c++=d. Thus b = a+c++=a+++c, whence $a++\leq b$ and by the above exercise a < b.

Ex. 2.2.4. We have $0 \le b$ for all b since b = b + 0. If a > b then since a + + > a we have a + + > b by transitivity. Finally, if a = b then a + + = b + + > b since for all a we have, again, a + + > a.

Ex. 2.2.5. Define Q(n) to be the property that P(m) is true for all $m_0 \leq m < n$. We shall induct on n. We have that in the base case n=0 the statement is true vacously. In fact, it is true vacously for all $n \leq m_0$. Now suppose Q(n) holds for n, or in other words, P(m) is true for all $m_0 \leq m < n$. By assumption this implies P(m+1) is true also, whence Q(n+1) is true. Thus, by the induction principle Q(n) holds for all n, whence P(m) hold for all natural numbers $m \geq m_0$.

Ex. 2.2.6. We shall induct on n. If n = 0 then P(0) is true by assumption and so P(m) is true for all $m \le 0$. Suppose now that P(n++) is true. Then by assumption P(n) is true, and so by the induction hypothesis

P(m) is true for all $m \le n$. But P(n++) was true also, so P(m) is true for all $m \le n++$.

2.3 Multiplication

Ex. 2.3.1. First we show $m \times 0 = 0$. By induction on m we have for m = 0 that $0 \times 0 = 0$ by the definition of multiplication. If $m \times 0 = 0$ then $m+++ \times 0 = (m \times 0) + 0 = 0$ by the induction hypothesis. Hence, $m \times 0 = 0$ for all m.

Next we show $m \times n++=(m \times n)+m$ by induction on m. In the base case m=0 we already have $0 \times n++=0=(m \times 0)+0$ from the paragraph above. Now suppose $m \times n++=(m \times n)+m$. Then $m++\times n++=(m+n++)+m$ by definition of multiplication, and by the induction hypothesis this equals (m+n)+m+m which in turn is equal to (m+++n)+m+++.

Finally we show multiplication is commutative. We induct on n. In the base case $m \times 0 = 0 \times m$. Suppose $m \times n = n \times m$. Then $m \times n + m = (m \times n) + m = (n \times m) + m = (n + m)$.

Ex. 2.3.2. If n=0 or m=0 then $m\times 0=0\times n=0$, respectively, by the above exercise.

Conversely, suppose m>0 and n>0. We induct on m. In the base case we have m=1 and $1\times n=n>0$. Now suppose $m\times n>0$. Then $m++\times n=(m\times n)+m>0$, by the induction principle and Proposition 2.2.8.

Ex. 2.3.3. We fix b and c and induct on a. In the base case a=0 and $(0\times b)\times c=0=0\times (b\times c)$. Now suppose abc is unambiguous. Then $((a+1)\times b)\times c=abc+bc$ by the Distributive law. But also, $(a+1)\times (b\times c)=abc+bc$ by the Distributive law.

Ex. 2.3.4. We have $(a+b)^2 = (a+b)(a+b) = a(a+b) + b(a+b) = a^2 + 2ab + b^2$.

Ex. 2.3.5. We fix q and induct on n. In the base case n=0 we have m=0 and r=0. Then n=0=0q+0 and $0 \le r < q$ since q is positive. Now suppose n=mq+r and consider n+1. Then n+1=mq+r+1. If r+1 < q we are done. Otherwise r=q and so n=(m+1)q+0.