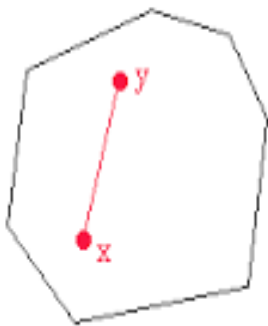


## Convex Hull (due 30 Oct 2020)

A convex hull is the smallest convex polygon that will enclose a set of points. In a convex polygon a line joining any two points in the polygon will lie completely within the polygon. One way to visualize a convex hull is as follows: imagine there are nails sticking out over the distribution of points. Place a rubber band such that it encircles all the nails. The figure described by the rubber band is a convex hull.

### Convex / Non-convex Polygons



A convex polygon



A non-convex polygon

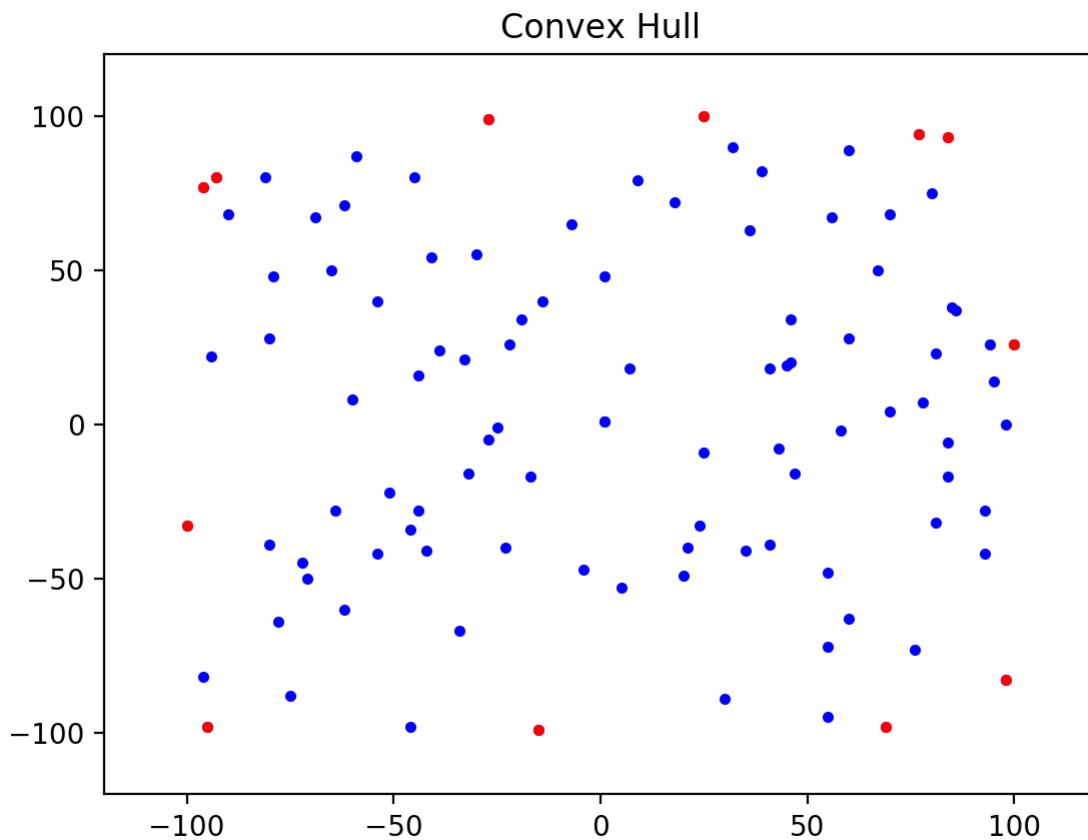
**Input:** You will read your data from standard input. The format of the input will be as given in [hull.in](#). The first line will be a single integer  $n$  denoting the number of points, where  $n$  will be in the range 3 to 100 inclusive. It will be followed by  $n$  lines of data. Each line will have the  $x$  and  $y$  coordinates of a point. The coordinates of the points will be integers in the range  $(-200, 200)$ .

**Output:** For the given input you will print the vertices of the convex hull starting at the extreme left point and going clockwise around the convex hull. You will print your output to standard out as given in the following format [hull.out](#).

```
Convex Hull
(-100, -33)
(-96, 77)
(-93, 80)
(-27, 99)
(25, 100)
(77, 94)
(84, 93)
(100, 26)
(98, -83)
(69, -98)
(-15, -99)
(-95, -98)
```

Area of Convex Hull = 37218.0

### Convex Hull for Input Data



There are many algorithms for solving the convex hull problem. You will have to implement the *Graham's scan* algorithm that is of order  $O(n \log n)$ . The algorithm is given to you in pseudo code.

Input: A set of point objects in the x-y plane.

Output: A list of point objects that define the vertices of the convex hull in clockwise order.

1. Sort the points by x-coordinates resulting in a sorted sequence  $p_1 \dots p_n$ .
2. Create an empty list *upper\_hull* that will store the vertices in the upper hull.
3. Append the first two points  $p_1$  and  $p_2$  in order into the *upper\_hull*.
4. For  $i$  going from 3 to  $n$
5.   Append  $p_i$  to *upper\_hull*.
6.   While *upper\_hull* contains three or more points and the last three points in *upper\_hull* do not make a right turn do
7.     Delete the middle of the last three points from *upper\_hull*
8. Create an empty list *lower\_hull* that will store the vertices in the lower hull.
9. Append the last two points  $p_n$  and  $p_{n-1}$  in order into *lower\_hull* with  $p_n$  as the first point.

10. For i going from n - 2 downto 1
11. Append p\_i to lower\_hull
12. While lower\_hull contains three or more points and the last three points in the lower\_hull do not make a right turn do
13. Delete the middle of the last three points from lower\_hull
14. Remove the first and last points for lower\_hull to avoid duplication with points in the upper hull.
15. Append lower\_hull to upper\_hull and call it the convex\_hull
16. Return the convex\_hull.

Two points p (px, py) and q (qx, qy) define a straight line. If add a third point r (rx, ry) how do you know whether r is to the right or left of the line defined by p and q. There is a simple solution to it. Evaluate the following determinant:

```

1  px  py
1  qx  qy
1  rx  ry

```

If the determinant is zero then the three points are in a straight line. The sign of the determinant will decide if the point is to the right or left of the line. Find that for yourself.

The computation of the area of a polygon given its vertices involves computing another determinant. If the vertices of a polygon are given by [p1, p2, ..., pn] where n is greater than or equal to 3, then compute the determinant given by the coordinates of the vertices:

```

x1  y1
x2  y2
..  ..
xn  yn

```

$$\text{det} = (x_1 * y_2 + x_2 * y_3 + \dots + x_n * y_1 - y_1 * x_2 - y_2 * x_3 - \dots - y_n * x_1)$$

$$\text{area} = (1/2) * \text{abs}(\text{det})$$

Here is the template of the file [Hull.py](#) that you will be submitting. You may **NOT** change the names of the functions but you may add as many helper functions as needed. You will follow the [standard coding conventions](#) in Python. Your file will have a header as follows:

```

# File: Hull.py

# Description:

# Student Name:

# Student UT EID:

# Partner Name:

# Partner UT EID:

# Course Name: CS 313E

# Unique Number:

# Date Created:

```

# Date Last Modified:

**You can always add more functions than those listed.** Mac users will run their code on the command line as follows:

```
python3 Hull.py < hull.in
```

Windows users will run their code on the command line as follows:

```
python Hull.py < hull.in
```

For this assignment you may work with a partner. Both of you must read the paper on [Pair Programming](#) and abide by the ground rules as stated in that paper. If you are working with a partner then only one of you will submit the code. Make sure that in the header in Canvas that you have your name and UT EID and your partner's name and UT EID. If you are working alone then you will just have your name and your UT EID.

Use the [Canvas](#) system to submit your **Hull.py** file. We should receive your work by 11 PM on Friday, 30 Oct 2020. There will be substantial penalties if you do not adhere to the guidelines. Remember Python is case sensitive. The name of your file must match exactly what we have specified.

- Your Python program should have the proper header.
- Your code must run before submission.
- You should be submitting your file through the web based *Canvas* program. We will not accept files e-mailed to us.

## References

- [Convex Hulls](#)
- [Convex Hulls Algorithms](#)