

## 1 Introduction

## 2 Preliminaries

**Definition 2.1 (Markov chains)** A Markov Chain  $C$  is a tuple  $(Q, \delta)$  where  $Q$  is a finite set of states and  $\delta : Q \rightarrow \mathbb{D}(Q)$  is a probabilistic transition function, where  $\mathbb{D}(Q)$  is the set of all probabilistic distribution on a finite set  $Q$ . A probabilistic distribution on  $Q$ , is a function  $f : Q \rightarrow \mathbb{Q}_{\geq 0}$  such that  $\sum_{q \in Q} f(q) = 1$ .

A non-empty finite word  $\varrho = p_1 p_2 \dots p_n$  over  $Q$  is defined as a run of Markov Chain. The probability of the run is  $\prod_{0 \leq i < n} \delta(p_i, p_{i+1})$ .  $\varrho$  reaches  $q$ , if  $q = p_i$  for some  $0 \leq i \leq n$ .

The probability of eventually reaching a set of state  $T \subseteq Q$  in  $C$  starting from  $q_0$  can be denoted as  $\mathbb{P}_{q_0}^C[\Diamond T]$ .  $\Diamond T$  is equivalent to  $Q \cup T$ ,  $Q$  until  $T$ . Since not necessary every state in  $Q$  is reached before  $T$ , the probability  $\mathbb{P}_{q_0}^C[S \cup T]$  with  $S \subseteq Q$  is created. If  $q_0 \in T$ , the probability is 1. Otherwise, the probability is calculated as follows:

$$\sum \left\{ \prod_{0 \leq i < n} \delta(q_i, q_{i+1}) \mid q_0 \dots q_{n-1} \in (S \setminus T) \ \& \ q_n \in T \text{ for } n \geq 1 \right\}$$

**Lemma 1** Consider a Markov Chain  $C = (Q, \delta)$  set of states  $U, T \subseteq Q$ , and a state  $q_0 \in Q \setminus U$ . If  $\mathbb{P}_{q_0}^C[(Q \setminus U) \cup T] = 0$ , then

$$\mathbb{P}_{q_0}^C[\Diamond T] = \sum_{u \in U} \mathbb{P}_{q_0}^C[(Q \setminus U) \cup u] \mathbb{P}_{q_0}^C[\Diamond T]$$

**Definition 2.2**

**Theorem 2 ()** Let  $A$

**Proof 2.1** Here you write the proof of the theorem.

## 3 Temporal logics

## 4 Conclusion

## References