非递归下降分析器

• 消除间接左递归方法

例:

$$S \rightarrow Qc|c$$

$$Q \rightarrow Rb|b$$

$$R \rightarrow Sa|a$$

过程:

(1)代入化简,将间接左递归变成直接左递归

将R的产生式代入Q的产生式,得到

Q o Sab|ab|b

将Q的产生式代入S的产生式,得到

S o Sabc|abc|bc|c

(2)消除直接左递归,得到的结果与Q和R无关,可删去Q和R的产生式

• 简化文法, 仅包含while循环的文法分析

注: 以空格分隔各符号(终结符与非终结符)

- $(0)program \rightarrow block$
- $(1)block
 ightarrow \{\ stmts\ \}$
- $(2)stmts
 ightarrow stmt \ stmts$
- (3)stmts
 ightarrow arepsilon
- $(4)stmt \rightarrow id = expr;$
- $(5)stmt \rightarrow while \ (\ bool\)\ stmt$
- $(6)stmt \rightarrow block$
- (7)bool
 ightarrow expr bool'
- $(8)bool' \rightarrow < expr$
- $(9)bool' \rightarrow <= expr$
- $(10)bool' \rightarrow > expr$
- $(11)bool' \rightarrow >= expr$
- $(12)bool' \rightarrow \varepsilon$
- $(13)expr
 ightarrow term\ expr''$
- $(14)expr'' \rightarrow expr' \ expr''$
- (15)expr'' o arepsilon
- $(16)expr' \rightarrow + term$
- $(17)expr' \rightarrow -term$
- $(18)term \rightarrow factor \ term''$
- $(19) term'' \rightarrow term' \ term''$
- $(20)term'' \rightarrow \varepsilon$
- $(21)term' \rightarrow *factor$
- $(22)term' \rightarrow / factor$
- $(23) factor \rightarrow (\ expr\)$
- $(24) factor \rightarrow id$
- (25) factor
 ightarrow num

• 手工构造分析表

	-{	3	id	-		white	()	4		-	>			•	- /	num
program	$program \rightarrow block$																
block	$block \rightarrow \{strats\}$																
stmts	$stmts \rightarrow stmt\ stmts$	$stmts \rightarrow \varepsilon$	$stmts \rightarrow stmt \ stmts$			$stmts \rightarrow stmt\ stmts$											
stmt	$stmt \rightarrow block$		$stmt \rightarrow id = expr;$			while(bool)stmt											
bool			$bool \to expr bool'$				$bool \to expr bool^c$										bool → expr bool'
bool'								$bool' \to \varepsilon$	$toot' \rightarrow < expr$	$bool' \rightarrow <= expr$	$bool' \rightarrow> expr$	$bool' \rightarrow>= expr$					
expr			$expr \rightarrow term \ expr^*$				$capr \to term \; capr^*$										$expr \rightarrow term expr^o$
expr'													$expr' \rightarrow +term$	$expr' \rightarrow -term$			
expr"					$expv^o \rightarrow \varepsilon$			$expr^{\alpha} \rightarrow \varepsilon$	$expr^A \rightarrow \epsilon$	$expv^A \rightarrow \varepsilon$	$expr^4 \rightarrow \epsilon$	$expv^A \rightarrow \varepsilon$	$expe^{\mu} \rightarrow expr' expe^{\mu}$	$expe^{\mu} \rightarrow expr' expe^{\mu}$			
term			$term \rightarrow factor\ term^s$				$term \to factor\ term^{\vee}$										$term \rightarrow factor term^{\alpha}$
term'															$term' \rightarrow *factor$	$term' \rightarrow /factor$	
term"					$term^v \to \varepsilon$			$term^v \to \varepsilon$	$term^v \rightarrow \varepsilon$	$term^v \rightarrow \varepsilon$	$term^{\nu} \rightarrow \varepsilon$	$term^v \rightarrow \varepsilon$	$term^o \rightarrow e$	$term^o \rightarrow e$	$term^{\nu} \rightarrow term' \ term^{\nu}$	$term^{\sigma} \rightarrow term^{\prime} \; term^{\sigma}$	
factor			$factor \rightarrow id$				$factor \rightarrow (expr)$										$factor \rightarrow num$

• 求First集、Follow集

```
First(program) = \{ \{ \} \}
First(block) = \{ \{ \} \}
First(stmts) = \{id, while, \{, \varepsilon\}\}
First(stmt) = \{id, while, \{\}\}
First(bool) = \{(,id,num)\}
First(bool') = \{<, <=, >, >=, \varepsilon\}
First(expr) = \{(, id, num)\}
First(expr') = \{+, -\}
First(expr'') = \{+, -, \varepsilon\}
First(term) = \{(, id, num)\}
First(term') = \{*,/\}
First(term'') = \{*,/,\varepsilon\}
First(factor) = \{(, id, num)\}
Follow(program) = \{\#\}
Follow(block) = \{ \#, id, while, (, \} \}
Follow(stmts) = \{ \} \}
Follow(stmt) = \{id, while, (, \} \}
Follow(bool) = \{\}
Follow(bool') = \{\}
Follow(expr) = \{), <, <=, >, >=, ; \}
Follow(expr') = \{+, -, \}, <, <=, >, >=, ; \}
Follow(expr'') = \{\}, <, <=, >, >=, ; \}
Follow(term) = \{+, -, \}, <, <=, >, >=, ; \}
Follow(term') = \{*,/,+,-,\}, <, <=,>,>=,;\}
Follow(term'') = \{+, -, \}, <, <=, >, >=, ; \}
Follow(factor) = \{*,/,+,-,),<,<=,>,>=,;\}
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