非递归下降分析器

• 消除间接左递归方法

例:

$$S o Qc|c$$

 $Q o Rb|b$
 $R o Sa|a$

过程:

- (1)代入化简,将间接左递归变成直接左递归
- 将R的产生式代入Q的产生式,得到

Q o Sab|ab|b

将Q的产生式代入S的产生式,得到

S o Sabc|abc|bc|c

- (2)消除直接左递归,得到的结果与Q和R无关,可删去Q和R的产生式
- S o abcS'|bcS'|cS'

 $S' o abcS' | \varepsilon$

• 简化文法,仅包含while循环的文法分析

注: 忽略空格

- $(0)program \rightarrow block$
- $(1)block \rightarrow \{stmts\}$
- $(2)stmts
 ightarrow stmt \ stmts$
- (3)stmts
 ightarrow arepsilon
- $(4)stmt \rightarrow id = expr;$
- $(5)stmt \rightarrow while(bool)stmt$
- $(6)stmt \rightarrow block$
- $(7)bool \rightarrow expr\ bool'$
- $(8)bool' \rightarrow < expr$
- $(9)bool' \rightarrow <= expr$
- $(10)bool' \rightarrow > expr$
- $(11)bool' \rightarrow >= expr$
- $(12)bool' \rightarrow \varepsilon$
- $(13)expr
 ightarrow term\ expr''$
- $(14)expr'' \rightarrow expr' \ expr''$
- $(15)expr'' \rightarrow \varepsilon$
- $(16)expr' \rightarrow +term$
- $(17)expr' \rightarrow -term$
- $(18)term \rightarrow factor \ term''$
- $(19)term'' \rightarrow term' term''$
- $(20)term'' \rightarrow \varepsilon$
- $(21)term' \rightarrow *factor$
- $(22)term' \rightarrow /factor$
- $(23) factor \rightarrow (expr)$
- $(24) factor \rightarrow id$
- (25) factor
 ightarrow num

• 手工构造分析表

			id	- 1		while	,		4	-		>=				,	num
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program	$program \rightarrow block$																
block	$block \rightarrow \{stmts\}$																
stmts	$stmts \rightarrow stmt\ stmts$	$stmts \rightarrow \varepsilon$	$stmts \rightarrow stmt \ stmts$			$stmts \rightarrow stmt \ stmts$											
stmt	$stmt \rightarrow block$		$stmt \rightarrow id = expr;$			while(boot)stmt											
bool			$bool \rightarrow expr bool'$				$bool \rightarrow expr bool'$										$bool \rightarrow expr bool'$
bool*								$bool' \rightarrow \varepsilon$	$boot' \to < expr$	$boot' \rightarrow <= expr$	$bool' \rightarrow> expr$	$bool' \rightarrow>= expr$					
expr			$expr \rightarrow term \ expr^*$				$capr \to term \; capr^{\nu}$										$expr \rightarrow term expr^o$
expr'													$expr' \rightarrow +term$	$expr' \rightarrow -term$			
expr-								$expr^n \to \varepsilon$	$expv^{\alpha} \rightarrow \varepsilon$	$expv^{\alpha} \rightarrow \varepsilon$	$expv^A \rightarrow \varepsilon$	$expv^n \rightarrow \varepsilon$	$expr^n \rightarrow expr' expr''$	$expr'' \rightarrow expr' expr''$			
term			$term \to factor term^a$				$term \to factor\ term^{\kappa}$										$term \rightarrow factor term^o$
term'															$term' \rightarrow *factor$	$term' \rightarrow /factor$	
term"								$term^v \to \varepsilon$	$term^v \to \varepsilon$	$term^{v} \rightarrow \varepsilon$	$term^v \rightarrow \varepsilon$	$term^{o} \rightarrow \varepsilon$			$term'' \to term' term''$	$term'' \to term' term''$	
factor			$factor \rightarrow id$		т		$factor \rightarrow (expr)$										$factor \rightarrow num$

• 求First集、Follow集

```
First(program) = \{ \{ \} \}
First(block) = \{ \{ \} \}
First(stmts) = \{id, while, \{, \varepsilon\}\}\
First(stmt) = \{id, while, \{\}\}
First(bool) = \{(,id,num)\}
First(bool') = \{<, <=, >, >=, \varepsilon\}
First(expr) = \{(, id, num)\}
First(expr') = \{+, -\}
First(expr'') = \{+, -, \varepsilon\}
First(term) = \{(, id, num)\}
First(term') = \{*,/\}
First(term'') = \{*,/,\varepsilon\}
First(factor) = \{(, id, num)\}
Follow(program) = \{\#\}
Follow(block) = \{ \#, id, while, (, \} \}
Follow(stmts) = \{ \} \}
Follow(stmt) = \{id, while, (, \} \}
Follow(bool) = \{\}
Follow(bool') = \{\}
Follow(expr) = \{), <, <=, >, >= \}
Follow(expr') = \{+, -, \}, <, <=, >, >= \}
Follow(expr'') = \{\}, <, <=, >, >= \}
Follow(term) = \{+, -, \}, <, <=, >, >= \}
Follow(term') = \{*,/,+,-,),<,<=,>,>=\}
Follow(term'') = \{+, -, \}, <, <=, >, >= \}
Follow(factor) = \{*,/,+,-,),<,<=,>,>=\}
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