Python Introduction Tutorial

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Code Explanation

Program 1

```
from stack import ListStack
values = ListStack()

for i in range(16):
   if i % 3 == 0:
   values.push(i)
   elif i % 4 == 0:
   values.pop()

print(values)
```

Class Definition (ListStack):

- Purpose: This class defines a simple stack data structure.
- Methods:
 - push: Adds an item to the top of the stack.
 - pop: Removes and returns the top item from the stack. If the stack is empty, it returns None.

Main Code Execution:

- A ListStack object, values, is created.
- A for loop iterates from 0 to 15 (inclusive).
- During each iteration:
 - If i is a multiple of 3 (i % 3 == 0), i is added to the stack.
 - If i is a multiple of 4 (i % 4 == 0) and not a multiple of 3, the top item of the stack is removed.
- The state of the stack is printed at the start of each loop iteration and after the loop concludes.

Loop Iterations and Stack State

- Start of loop (initial state): []
- i = 0: [0] (0 is pushed)
- i = 1: [0]
- i = 2: [0]
- i = 3: [0, 3] (3 is pushed)
- i = 4: [0] (3 is popped)
- i = 5: [0]
- i = 6: [0, 6] (6 is pushed)
- i = 7: [0, 6]
- i = 8: [0] (6 is popped)
- i = 9: [0, 9] (9 is pushed)
- i = 10: [0, 9]
- i = 11: [0, 9]
- i = 12: [0, 9, 12] (12 is pushed)
- i = 13: [0, 9, 12]
- i = 14: [0, 9, 12]
- i = 15: [0, 9, 12] (Since i is a multiple of 3, 15 is pushed but it's not reflected in the loop's output)

Final State of the Stack

Final state of the stack after the loop: [0, 9, 12, 15]

Program 2 Explanation and Output

Program 2

The square Function:

This function takes an integer **n** as input and performs operations on a stack (ListStack). It aims to manipulate the stack in a specific pattern and return the modified stack.

• Initialization: A ListStack named SquareStack is created, and a variable i is set to 0.

• Building the Stack:

- The outer loop iterates through odd numbers from 1 up to n, with a step of 2.
- For each odd number j:
 - * The first inner loop pushes incrementing values of i onto the stack j times, incrementing i each time.
 - * The second inner loop pops elements from the stack j-1 times.
- This pattern creates a sequence of square numbers.

• Final Adjustment:

- A while loop ensures that no number greater than **n** remains in the stack.

Examples

- For j = 1: The stack has 0 pushed (1 time), no pop occurs. (i becomes 1).
- For j = 3: The stack has 1, 2, 3 pushed, then 2, 3 popped. Only 1 remains. (i becomes 4).
- For **j** = 5: The stack has 4, 5, 6, 7, 8 pushed, then 5, 6, 7, 8 popped. Only 4 remains. (**i** becomes 9).
- For j = 7: The stack has 9, 10, 11, 12, 13, 14, 15 pushed, then 10, 11, 12, 13, 14, 15 popped. Only 9 remains. (i becomes 16).
- For j = 9: The stack has 16, 17, 18, 19, 20, 21, 22, 23, 24 pushed, then 17, 18, 19, 20, 21, 22, 23, 24 popped. Only 16 remains. (i becomes 25).
- For j = 11: The stack has 25, 26, 27, 28, 29, 30, 31, 32, 33, 34, 35 pushed, then 26, 27, 28, 29, 30, 31, 32, 33, 34, 35 popped. Only 25 remains. (i becomes 36).

Output:

When the function is called with n=100, the output is a stack containing the elements:

This output reveals that the function generates and stores square numbers up to and including n, provided n itself is a square number.

Translation of Infix to Postfix

i) Infix to Postfix:

1.
$$V \times W \times (X+Y) - Z$$

Derivation:

Step 1:
$$X+Y \to XY+$$

Step 2: $VW \times XY+ \to VW \times XY+*$
Step 3: $VW \times XY+*-Z \to VW \times XY+*Z-$
Postfix: $VW * XY + *Z-$

2. $X - (Y + W \times (Z/V))$

Derivation:

Step 1:
$$Z/V \to ZV/$$

Step 2: $W \times ZV/ \to WZV/*$
Step 3: $Y + WZV/* \to YWZV/* + +$
Step 4: $X - YWZV/* + \to XYWZV/* + -$
Postfix: $XYWZV/* + -$

3. $A - B \times (C + D/E)$

Derivation:

Step 1:
$$D/E \to DE/$$

Step 2: $C + DE/ \to CDE/+$
Step 3: $B \times CDE/+ \to BCDE/+*$

Step 4: $A - BCDE/*+ \rightarrow ABCDE/+*-$

Postfix: ABCDE/ + *-

4. $A/(B \times C) - (D+E)$

Derivation:

Step 1: $B \times C \rightarrow BC *$

Step 2: $A/BC* \rightarrow ABC*/$

Step 3: $D + E \rightarrow DE +$

Step 4: $ABC/*-DE+ \rightarrow ABC/*DE+-$

Postfix: ABC * /DE + -

Translation of Postfix to Infix

ii) Postfix to Infix:

Expression (a): ABC - D * +

- 1. Postfix: ABC D * +
- 2. Steps:
 - Start from left, AB (operands), and C (operand), followed by an operator -. In infix, this becomes (B-C).
 - Now, we have D(B-C). The next operator is *, which applies to D and (B-C), making it D*(B-C).
 - Finally, the last operator + applies to the whole expression and another operand D, resulting in A + (B C) * D.
 - Rearranging for clarity: (A + ((B C) * D)).
- 3. **Infix:** (A + ((B C) * D))

Expression (b): XYZ + AB - *-

- 1. Postfix: XYZ + AB *-
- 2. Steps:
 - Start with XY, then Z, followed by +, translating to (Y + Z).
 - Next, we encounter AB, and -, translating to (A B).

- The * operator applies to (Y + Z) and (A B), giving (Y + Z) * (A B).
- Finally, the operator is applied between X and the entire (Y + Z) * (A B) expression, resulting in X ((Y + Z) * (A B)).
- 3. **Infix:** X ((Y + Z) * (A B))

Expression (c): AB + CD - /E +

- 1. Postfix: AB + CD /E +
- 2. Steps:
 - Starting with AB+, this becomes (A+B).
 - Moving on, we have CD-, translating to (C-D).
 - The / operator then applies to (A+B) and (C-D), resulting in ((A+B)/(C-D)).
 - Finally, the + operator combines this result with E, yielding (((A+B)/(C-D)) + E).
- 3. **Infix:** (((A+B)/(C-D))+E)

Expression (d): AB + C - DE * +

- 1. Postfix: AB + C DE * +
- 2. Steps:
 - Start with AB+, converting to (A+B).
 - Next, we see C, and then a operator, which applies to (A+B) and C, giving ((A+B)-C).
 - Following this, DE* translates to (D*E).
 - The final + operator combines ((A+B)-C) and (D*E), resulting in (((A+B)-C)+(D*E)).
- 3. **Infix:** (((A+B)-C)+(D*E))

Code Execution and Output

Initial Configuration

- Initial Stack S: [3, 6, 78, 9] (Top $\to 9$)
- Initial Queue Q:

Execution Steps

1. First Call to CheckElement:

- Pops 9 from S: S = [3, 6, 78], Q = [9]
- Since $9 \neq 78$, calls CheckElement again.

2. Second Call to CheckElement:

- Pops 78 from S: S = [3, 6], Q = [9, 78]
- Since 78 = 78, prints "78 is an element of S."

3. Restoring the Stack:

- Move all elements from Q back to S.
- S = [3, 6, 9, 78], Q = [], but the last two numbers order in S is reversed.
- ullet Move elements from S to Q and back to S to restore original order.
- Final S = [3, 6, 78, 9] (original order).

Key Points

- \bullet Stack S holds the elements, queue Q is used for temporary storage to maintain order.
- Function uses recursion to check each element in S.
- \bullet Double transfer between S and Q restores original order of elements in S.

Output

before: [3, 6, 78, 9] 78 is an element of S. after: [3, 6, 78, 9]

1 Introduction

This document outlines the design and implementation of a discrete event simulation for a single-server queuing system, known as the M/M/1 queue. The simulation aims to analyze customer waiting times and queue length in a system where customer arrivals and service times follow an exponential distribution.

M/M/1 Queue

The M/M/1 queue is a fundamental model in queueing theory, which is the study of waiting lines. It is defined by the following characteristics:

- 1. M (First M): Stands for "Markovian" arrival process. Arrivals occur according to a Poisson process, with times between arrivals being exponentially distributed.
- 2. M (Second M): Indicates a "Markovian" service process. The service times are also exponentially distributed.
- 3. 1: Represents the number of servers in the system, which is one in this case.

Key Characteristics

- Single Server: Only one server is available for serving customers.
- Exponential Inter-Arrival and Service Times: Both arrivals and services follow an exponential distribution.
- **First-Come**, **First-Served**: The service discipline is usually assumed to be first-come, first-served.
- Infinite Queue Capacity: The queue can accommodate an unlimited number of customers.
- Steady-State Analysis: The system is often analyzed in a steady state, where its properties do not change over time.

Performance Metrics

- Average Number of Customers in the System (L): Includes both customers being served and those waiting in the queue.
- Average Number of Customers in the Queue (Lq): The average number of customers waiting for service.
- Average Time a Customer Spends in the System (W): Total time spent by a customer in the system.
- Average Time a Customer Spends in Queue (Wq): Average waiting time for a customer before their service starts.
- Server Utilization Factor (ρ): The fraction of time the server is busy.

2 Simulation Objectives

- 1. Calculate the average waiting time for each customer.
- 2. Determine the queue length at the end of the simulation period.

3 Simulation Parameters

- Total simulation duration in discrete time units.
- Customer arrival rate (λ) , following an exponential distribution.
- Service rate (μ) , also following an exponential distribution.

4 Implementation Details

4.1 Classes and Methods

- Customer Class: Private data fields for customer ID and arrival time.
- Server Class: Methods for starting and stopping service.
- **Simulation Class:** Runs the simulation, handles events like customer arrivals, service initiation, and service completion.

4.2 Simulation Process

- 1. Initialize with $\lambda = 9$, $\mu = 10$, and total time = 480 minutes.
- 2. Run the simulation in a loop, processing events at each time step.
- 3. Generate customers, place them in a queue, and have the server process them.
- 4. Calculate the total waiting time and queue length at the end.

5 Results and Analysis

- Run the simulation multiple times (1000 times) to average the results.
- Compare the empirical results with theoretical values:
 - Average queue length $L = \frac{\rho}{1-\rho}$ where $\rho = \frac{\lambda}{\mu}$.
 - Average waiting time $T = \frac{\rho}{\mu \lambda}$.

6 Conclusion

This simulation provides insights into the behavior of queues under various arrival and service rates, which is crucial for numerous real-world applications.