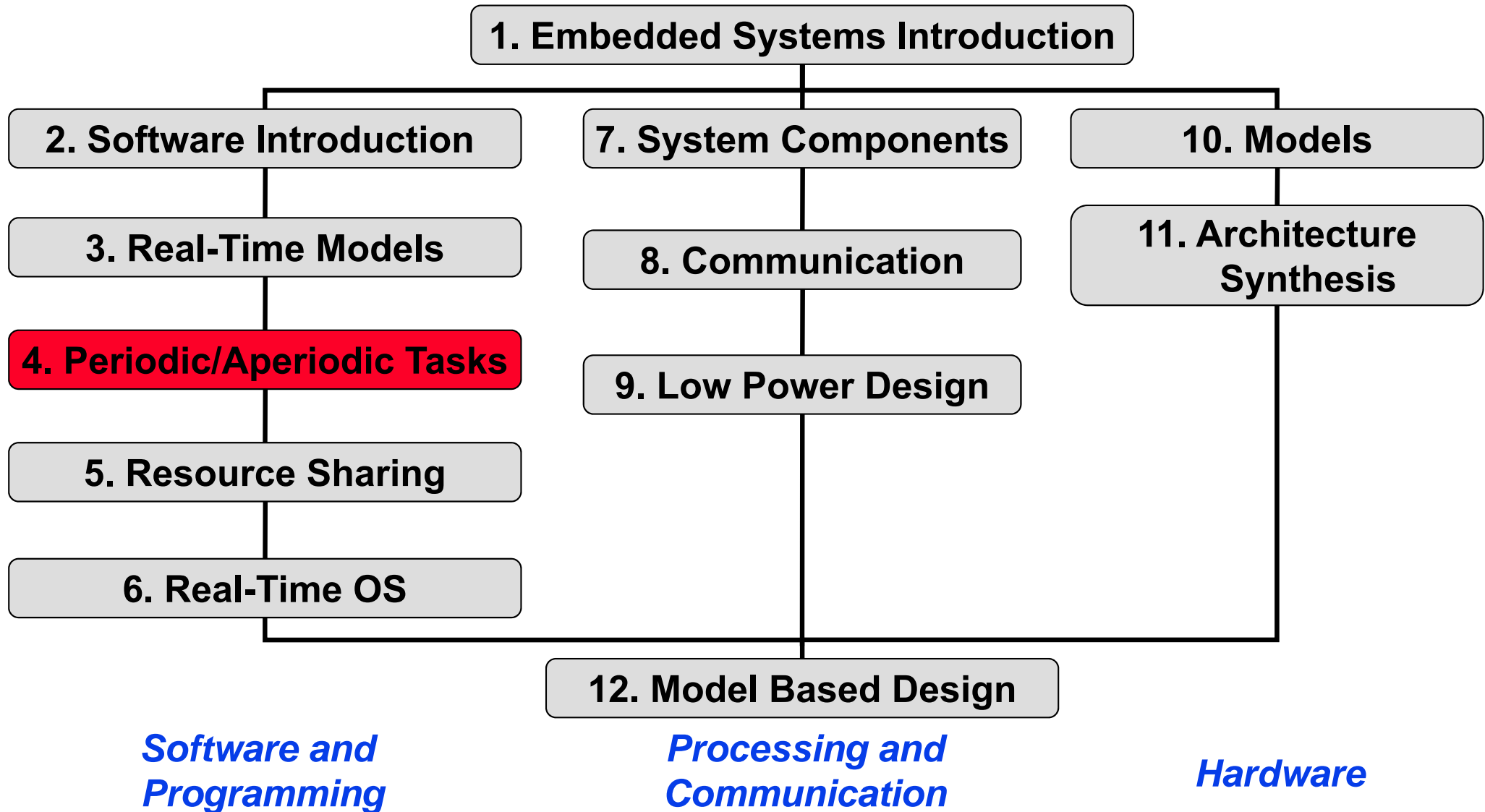


# Embedded Systems

## 4a. Example Network Processor

Lothar Thiele

# Contents of Course



# Software-Based NP

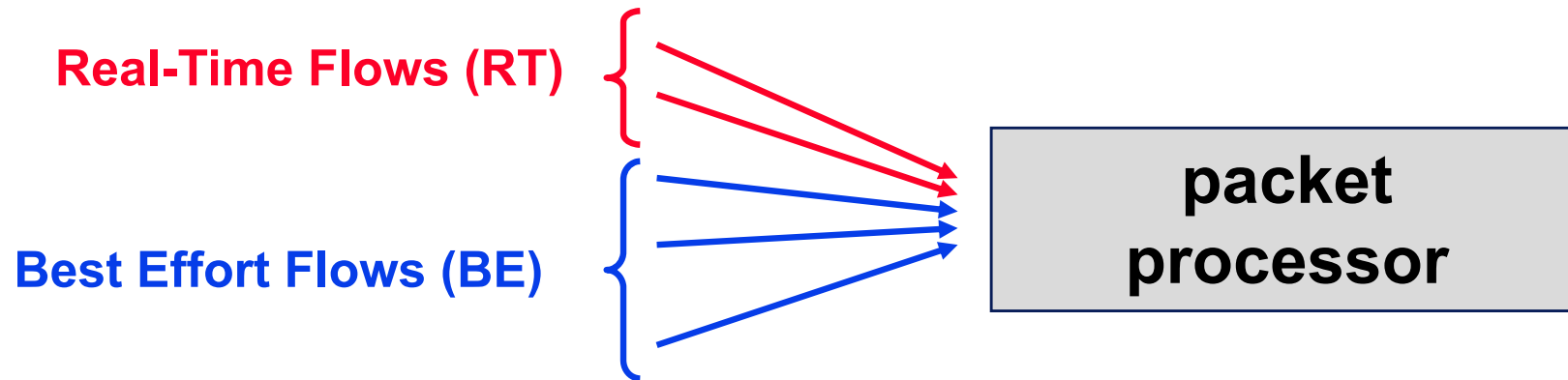
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## Network Processor: Programmable Processor Optimized to Perform Packet Processing

- ▶ How to Schedule the CPU cycles meaningfully?
  - Differentiating the level of service given to different flows
  - Each flow being processed by a different processing function

# Our Model – Simple NP

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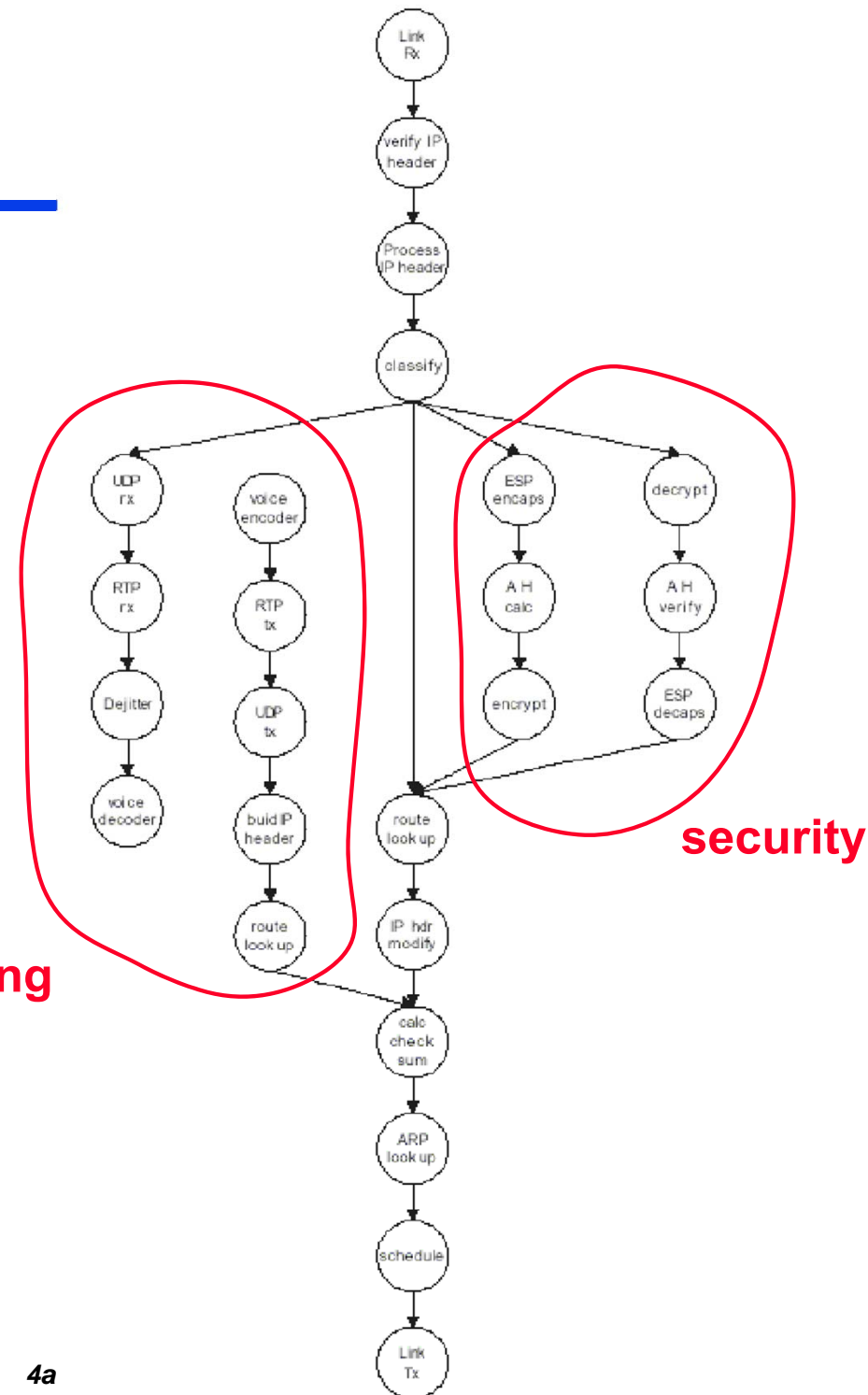
- ▶ Real-time flows have deadlines which must be met
- ▶ Best effort flows may have several QoS classes and should be served to achieve maximum throughput

# Task Model

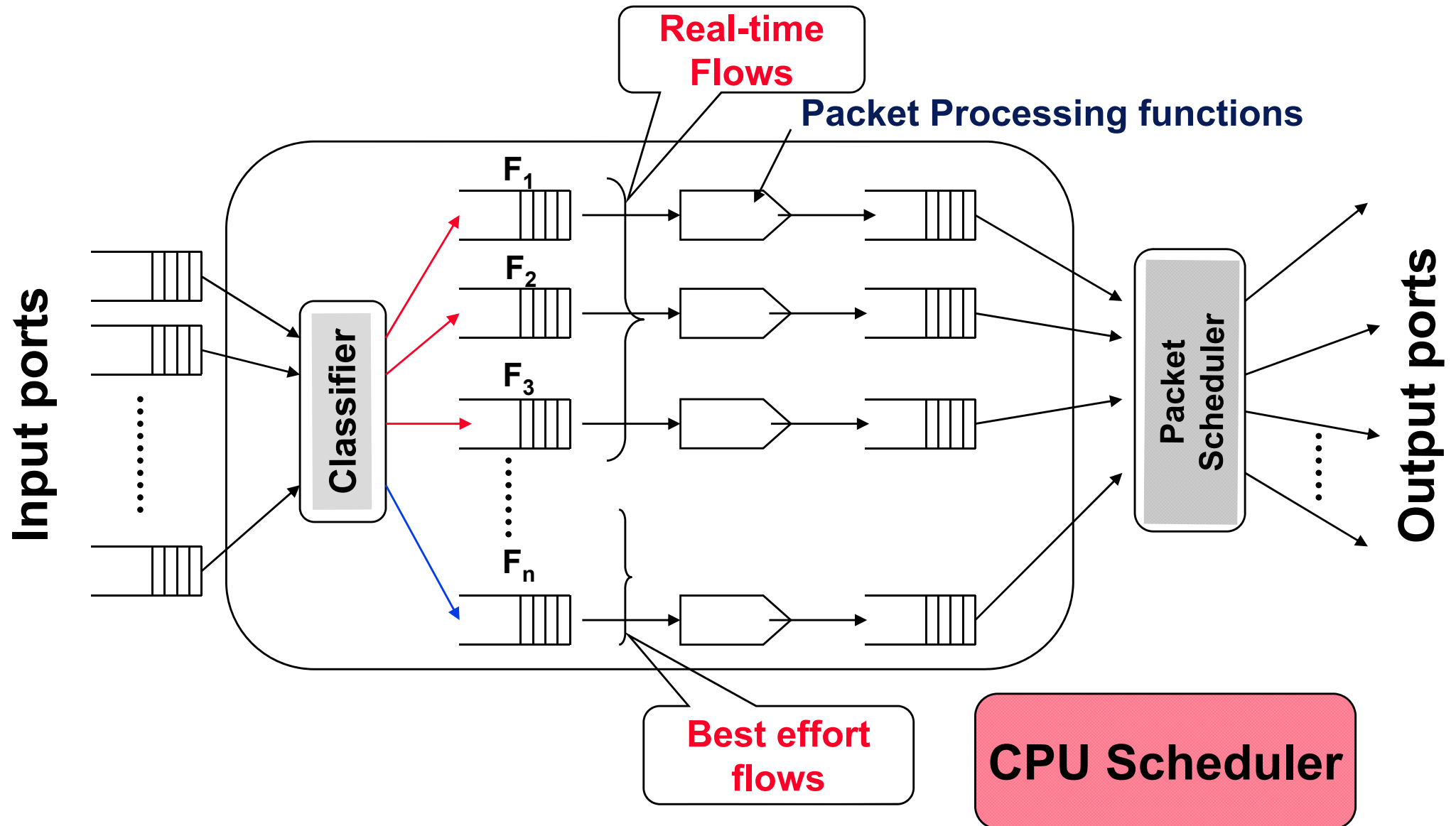
- ▶ Packet processing functions may be represented by directed acyclic graphs
- ▶ End-to-end deadlines for RT packets

voice processing

security



# Architecture

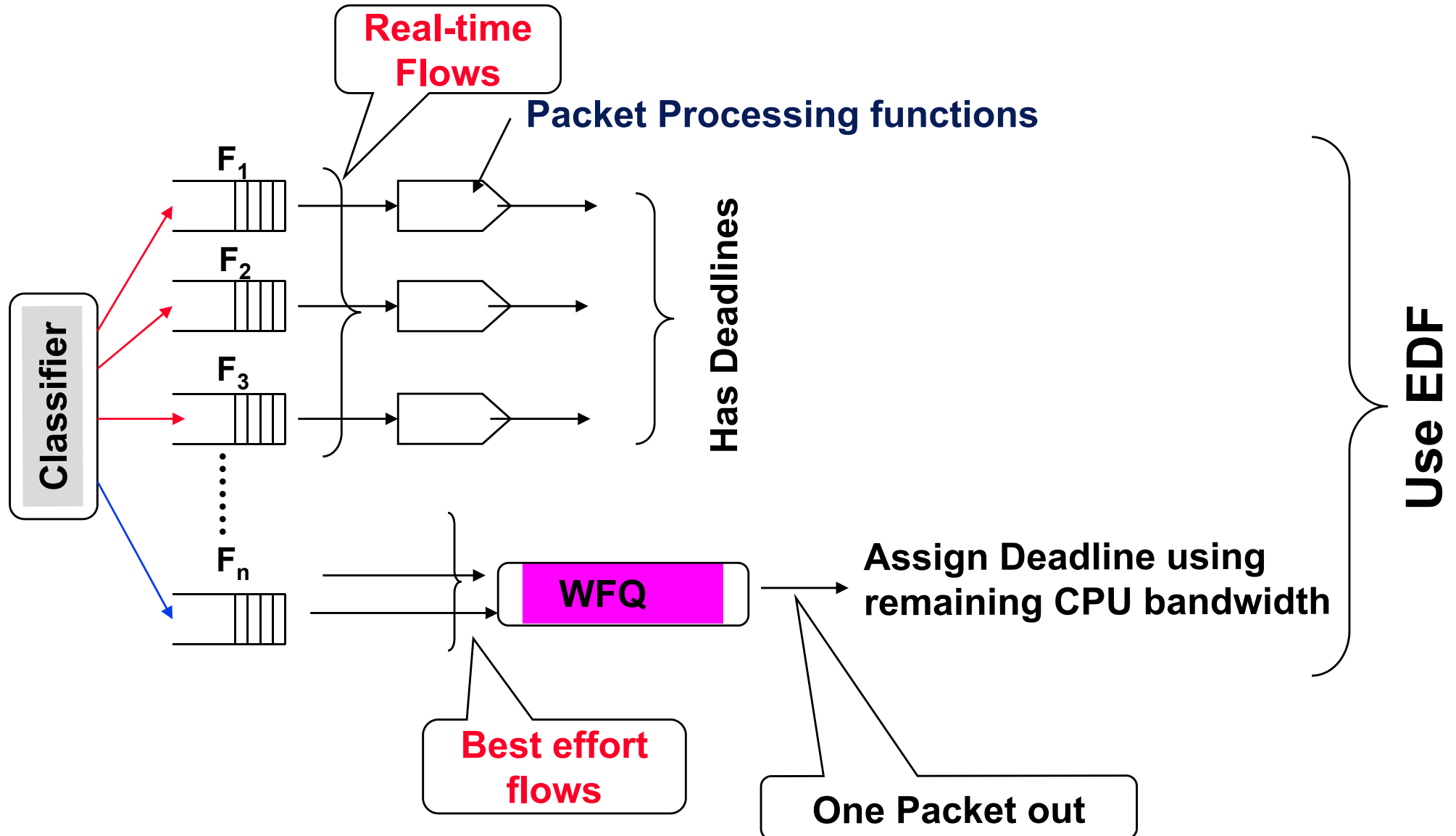


# CPU Scheduling

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- ▶ First Schedule RT, then BE (background scheduling)
  - Overly pessimistic
- ▶ Use **EDF Total Bandwidth Server**
  - EDF for Real-Time tasks
  - Use the remaining bandwidth to server Best Effort Traffic
  - WFQ (weighted fair queuing) to determine which best effort flow to serve; not discussed here ...

# CPU Scheduling





# CPU Scheduling

- ▶ As discussed, the *basis is the TBS*:

$$d_k = \max\{r_k, d_{k-1}\} + c_k / U_s$$

computation demand of best effort packet

deadline of best effort packet

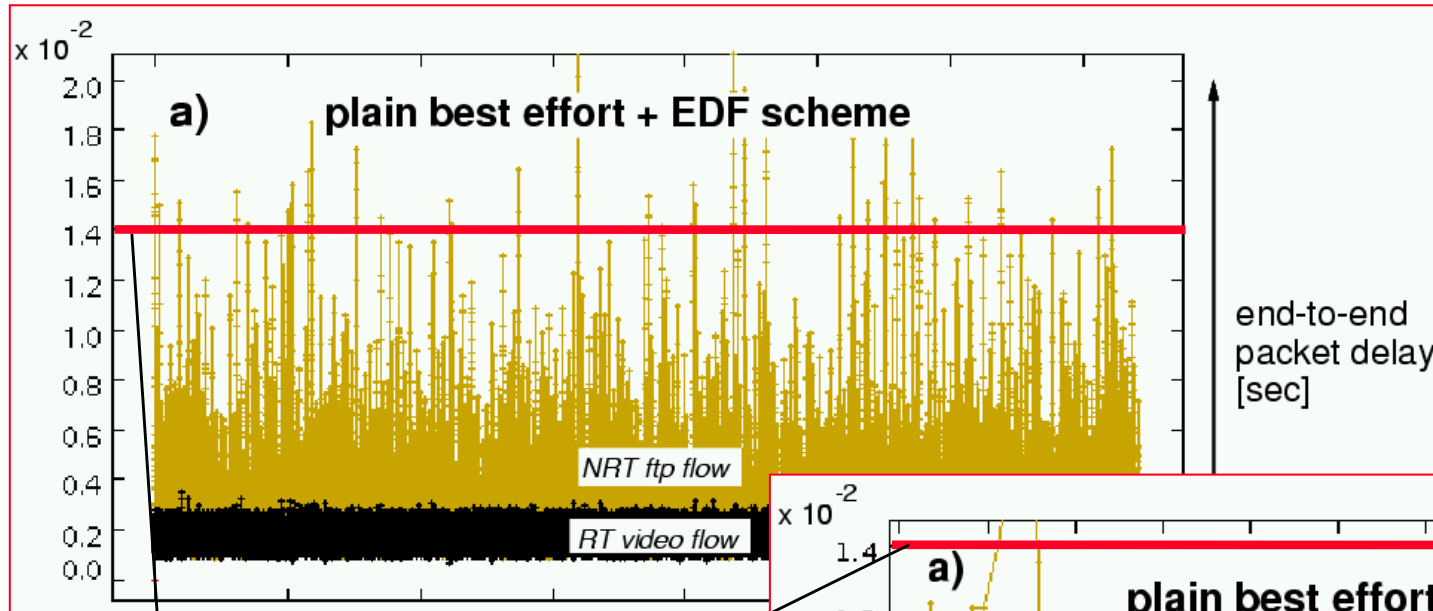
arrival of best effort packet

utilization by real-time flows

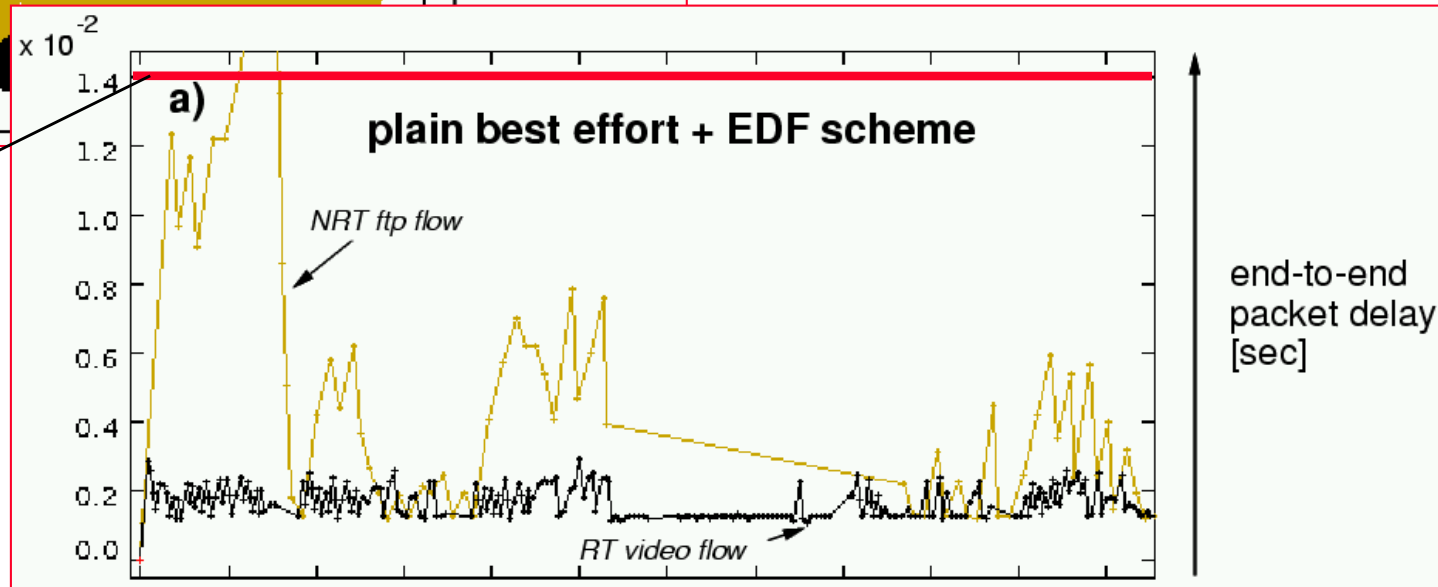
- ▶ *But*: utilization depends on time (packet streams) !
  - Just taking upper bound is too pessimistic
  - Solution with time dependent utilization is (much) more complex – BUT IT HELPS ...

# CPU Scheduling

## ► Before



deadline  
RT flows



# CPU Scheduling

## ► After

