Compiling from F_i^+ to JavaScript

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Syntax of F_i^+

Types
$$A,B,C ::= \top \mid \bot \mid \mathbb{B} \mid X \mid A \to B \mid \forall X*A. \ B \mid \{\ell:A\} \mid A \& B$$

Type indices
$$T ::= \mathbb{B} \mid \overrightarrow{T} \mid T^{\forall} \mid \{\ell : T\} \mid T_1 \& T_2$$

Expressions
$$e := \{\} \mid b \mid x \mid \mathbf{fix} \ x : A. \ e \mid \lambda x : A. \ e : B \mid e_1 \ e_2 \mid \Lambda X * A. \ e : B \mid e \ A \mid \{\ell = e\} \mid e.\ell\}$$

$$| e_1,, e_2 | e : A$$

Values
$$v := \{\} \mid b \mid \lambda x : A.\ e : B \mid \Lambda X * A.\ e : B \mid \{\ell = v\} \mid v_1 ,, v_2 \}$$

$$|A| = T$$
 (Type translation)

$$|\mathbb{B}| = \mathbb{B} \qquad |X| = X \qquad |\forall X*A. \ B| = |B|^\forall \qquad |A \rightarrow B| = \overrightarrow{|B|} \qquad |\{\ell:A\}| = \{\ell:|A|\} \qquad |A \ \& \ B| = |A| \ \& \ |B|$$

$$A^{\circ}$$
 (Ordinary types)

O-Top O-Bot O-Base O-Var
$$B^{\circ}$$
 O-Arrow B° O-All B° A° A°

$$\Gamma \vdash e \Leftrightarrow A \ \leadsto J \mid z^{\pm}$$

(Type-directed compilation)

$$J\text{-}G\mathrm{EN}$$

$$\Gamma \vdash e \Leftrightarrow A \leadsto J \mid z^-$$

$$\Gamma \vdash e \Leftrightarrow A \rightsquigarrow S_0 \mid z^+$$

$$\Gamma \vdash \{\} \Rightarrow \top \rightsquigarrow \varnothing \mid z^-$$

J-TopAbs

$$\exists B \lceil$$

$$\frac{\Gamma \vdash e \Leftrightarrow A \leadsto J \mid z^{-}}{\Gamma \vdash e \Leftrightarrow A \leadsto S_{0} \mid z^{+}} \qquad \frac{\text{J-TopAbs}}{\Gamma \vdash \{\} \Rightarrow \top \leadsto \varnothing \mid z^{-}} \qquad \frac{\text{J-TopAbs}}{\Gamma \vdash \lambda x : A. \ e : B \Rightarrow A \to B \leadsto \varnothing \mid z^{-}}$$

$$\exists B$$

$$\Gamma \vdash \Lambda X * A. \ e : B \Rightarrow \forall X * A. \ B \leadsto \varnothing \mid z^-$$

$$\Gamma \vdash e \Rightarrow A \qquad]A[$$

$$\frac{\text{J-TopTAbs}}{\Gamma \vdash \Lambda X * A.\ e : B \Rightarrow \forall X * A.\ B \leadsto \varnothing \mid z^-} \qquad \frac{\text{J-TopRcd}}{\Gamma \vdash \{\ell = e\} \Rightarrow \{\ell : A\} \leadsto \varnothing \mid z^-}$$

$$T =$$

$$\Gamma \vdash b \Rightarrow \mathbb{B} \rightsquigarrow S_1 \mid z^-$$

$$x:A\in\Gamma$$

$$\begin{array}{c|c} \text{J-Base} & \text{J-Var} \\ \hline T = |\mathbb{B}| & x: A \in \Gamma \\ \hline \Gamma \vdash b \Rightarrow \mathbb{B} \rightsquigarrow S_1 \mid z^- & \hline \hline \Gamma \vdash x \Rightarrow A \rightsquigarrow S_2 \mid z^- \\ \hline \end{array} \qquad \begin{array}{c|c} \text{J-VarGen} \\ \hline x: A \in \Gamma \\ \hline \hline \Gamma \vdash x \Rightarrow A \rightsquigarrow S_2 \mid z^- \\ \hline \end{array}$$

$$r \cdot A \in$$

$$\Gamma \vdash x \Rightarrow A \rightsquigarrow \varnothing \mid x^+$$

$$\Gamma$$
, this : $A \vdash e \Leftarrow A \rightsquigarrow J \mid this$

$$\Gamma \vdash \mathbf{fix} \ this : A. \ e \Rightarrow A \bowtie S_3 \mid z^-$$

$$T = \overrightarrow{|B|}$$

$$\Gamma, x : A \vdash e \Leftarrow B \rightsquigarrow J \mid y^+$$

$$\frac{\Gamma, this: A \vdash e \Leftarrow A \leadsto J \mid this^{-}}{\Gamma \vdash \mathbf{fix} \ this: A. \ e \Rightarrow A \leadsto S_{3} \mid z^{-}} \frac{\Gamma, x: A \vdash e \Leftarrow B \leadsto J \mid y^{+}}{\Gamma \vdash \lambda x: A. \ e: B \Rightarrow A \rightarrow B \leadsto S_{4} \mid z^{-}}$$

J-App

$$\Gamma \vdash e_1 \Rightarrow A \implies J_1 \mid x^+$$
$$A \rhd B \to C$$

$$\Gamma \vdash e_2 \Leftarrow B \longrightarrow C$$

$$\Gamma \vdash e_2 \Leftarrow B \longrightarrow J_2 \mid y^+ \mid$$

$$x: A \bullet y \leadsto J_3 \mid z$$

$$\begin{array}{c|c}
\Gamma \vdash e_2 \Leftarrow B & \leadsto J_2 \mid y \\
\hline
x : A \bullet y & \leadsto J_3 \mid z \\
\hline
\Gamma \vdash e_1 e_2 \Rightarrow C & \leadsto J_1; J_2; J_3 \mid z^-
\end{array}$$

$$T = |B|^{\vee}$$

$$\Gamma, X * A \vdash e \Leftarrow B \rightsquigarrow J \mid y^+$$

$$T = |B|^{\forall}$$

$$\Gamma, X * A \vdash e \Leftarrow B \implies J \mid y^{+}$$

$$\Gamma \vdash \Lambda X * A. \ e : B \Rightarrow \forall X * A. \ B \implies S_{5} \mid z^{-}$$

J-TAPP

$$\Gamma \vdash e \Rightarrow B \implies J_1 \mid y^+$$

$$B \rhd \forall X * C_1. C_2$$

$$\Gamma \vdash A * C_1$$

$$y: B \bullet A \leadsto J_2 \mid z$$

$$\Gamma \vdash e A \Rightarrow C_2[X \mapsto A] \rightsquigarrow J_1; J_2 \mid z^-$$

J-Rcd

$$T = \{\ell : |A|\}$$

$$\Gamma \vdash e \Rightarrow A \rightsquigarrow J \mid y^+$$

$$T = \{\ell : |A|\}$$

$$\Gamma \vdash e \Rightarrow A \rightsquigarrow J \mid y^{+}$$

$$\Gamma \vdash \{\ell = e\} \Rightarrow \{\ell : A\} \rightsquigarrow S_{6} \mid z^{-}$$

J-Proj

$$\Gamma \vdash e \Rightarrow A \rightsquigarrow J_1 \mid y^+$$

$$A \rhd \{\ell : B\}$$

$$y : A \bullet \ell \rightsquigarrow J_2 \mid z$$

$$\Gamma \vdash e.\ell \Rightarrow B \rightsquigarrow J_1; J_2 \mid z^-$$

$$\jmath:Aullet\ell|\leadsto J_2\mid z$$

$$\Gamma \vdash e.\ell \Rightarrow B \rightsquigarrow J_1; J_2 \mid z^-$$

J-Merge

$$\Gamma \vdash e_1 \Rightarrow A \rightsquigarrow J_1 \mid z^{-1}$$

$$\Gamma \vdash e_2 \Rightarrow A \rightsquigarrow J_1 \mid z^{-1}$$

$$\begin{array}{c|c}
\Gamma \vdash e_2 \Rightarrow B & \rightsquigarrow J_2 \mid z^- \\
& \Gamma \vdash A * B & & \Gamma \vdash e \Leftarrow A & \rightsquigarrow J \mid z^- \\
\hline
\Gamma \vdash e_1 ,, e_2 \Rightarrow A \& B & \rightsquigarrow J_1; J_2 \mid z^- & & \hline
\end{array}$$

$$\begin{array}{c|c}
J-ANNO \\
\Gamma \vdash e \Leftarrow A & \rightsquigarrow J \mid z^- \\
\hline
\Gamma \vdash e : A \Rightarrow A & \leadsto J \mid z^-
\end{array}$$

$$\Gamma \vdash e \Leftarrow A \rightsquigarrow J \mid z$$

$$\Gamma \vdash e : A \Rightarrow A \mid \leadsto J \mid z^-$$

$$\Gamma \vdash e \Leftarrow A \leadsto J \mid z^+$$

$$\Gamma \vdash \circ \cdot \land \rightarrow \land$$

$$\Gamma \vdash e \Leftarrow B$$

$$\Gamma \vdash e \Rightarrow A \rightsquigarrow J \mid z^{-}$$

$$\Gamma \vdash e \Leftarrow A \rightsquigarrow J \mid z$$

J-SubEqGen

$$\frac{\Gamma \vdash e \Rightarrow A}{\Gamma \vdash e \Leftarrow A} \rightsquigarrow J \mid z^{+}$$

$x: A \bullet arg \leadsto J \mid z$

(Distributive application)

$$\frac{\text{A-TOP}}{x : A \bullet arg} \xrightarrow{\sim} \varnothing \mid z$$

$$\frac{A - \text{Arrow}}{T = |\overrightarrow{B}|}$$

$$x : A \to B \bullet y \leadsto S_7 \mid z$$

A-RCD
$$T = \{\ell : |A|\}$$

$$x : A \bullet arg \rightsquigarrow J_1 \mid z$$

$$x : B \bullet arg \rightsquigarrow J_2 \mid z$$

$$x : \{\ell : A\} \bullet \ell \rightsquigarrow S_9 \mid z$$

$$x : A \& B \bullet arg \rightsquigarrow J_1; J_2 \mid z$$

 $x:A<:y:B \leadsto J$

S-Top

(Coercive subtyping)

S-TOP
$$\frac{B^{\circ}}{x : A <: y : B} \longrightarrow \varnothing$$
S-BASE
$$\frac{T = |\mathbb{B}|}{x : \mathbb{B} <: y : \mathbb{B}} \longrightarrow S_{10}$$
S-ALL
$$T_{1} = |A_{2}|^{\forall} \qquad T_{2} = |B_{2}|^{\forall}$$

$$B_{2}^{\circ} \qquad B_{1} <: A_{1}$$

$$x_{2} : A_{2} <: y_{2} : B_{2} \longrightarrow J_{2}$$

$$\overline{x : \forall X * A_{1}. A_{2} <: y : \forall X * B_{1}. B_{2} \longrightarrow S_{12}}$$

S-Base

$$T_{1} = |A_{2}|^{\forall} \quad T_{2} = |B_{2}|^{\forall} \quad T_{2} = |B_{2}|^{\forall}$$

S-Arrow

 $T_1 = \overrightarrow{|A_2|}$ $T_2 = \overrightarrow{|B_2|}$ B_2° $B_1 <: A_1$

 $x_2:A_2 <: y_2:B_2 \leadsto J_2$ $\overline{x: A_1 \to A_2 <: y: B_1 \to B_2} \leadsto S_{11}$

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x:A \ \rhd \ z:C \ \vartriangleleft \ y:B \ \leadsto \ J
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return y;

};

(Coercive merging)

};

};