

# Homework 2

January 27, 2024

Please submit your HW on Canvas; include a PDF printout of any code and results, clearly labeled, e.g. from a Jupyter notebook. For coding problems, we recommend using Julia, but you can use other languages if you wish. It is due Friday February 2nd by 11:59pm EST.

## Problem 1

Continue reading the draft course notes (linked from <https://github.com/mitmath/matrixcalc/>). Find another place that you found confusing, in a different chapter from your answer in homework 1, and write a paragraph explaining the source of your confusion and (ideally) suggesting a possible improvement.

(Any other corrections/comments are welcome, too.)

## Problem 2

The [course notebook on finite differences](#) includes, without derivation, a mysterious four-line Julia function called `stencil` that can compute finite-difference rules for an arbitrary number of points. In particular, if you want to compute the  $m$ -th derivative of a smooth (analytic) scalar function  $f(x)$  at  $x_0$ , it returns the weights  $w_k$  of an  $n$ -point ( $n > m$ ) finite-difference rule from evaluating  $f$  at points  $x_k$  for  $k = 1 \dots n$ :

$$f^{(m)}(x_0) \approx \sum_{k=1}^n w_k f(x_k)$$

by solving the system of equations  $Aw = e_{m+1}$ , where  $e_j \in \mathbb{R}^n$  is the Cartesian unit vector in the  $j$ -th direction and  $A$  is an  $n \times n$  matrix with entries  $A_{ij} = \frac{(x_j - x_0)^{i-1}}{(i-1)!}$ . Here, you will analyze and derive this technique.

1. Let  $x_0 = 0$ . According to the notes, you can then compute  $f^{(m)}(y) \approx \frac{1}{h^m} \sum_{k=1}^n w_k f(y + hx_k)$  for an arbitrary point  $y$  and an arbitrary step-size scaling factor  $h$  (which can be made smaller and smaller to reduce truncation errors). Derive this formula (via the chain rule).
2. Evaluate the `stencil` function (or its equivalent in another language if you want to re-implement it) for  $x_0 = 0$  and  $x = [0, 1]$  with  $m = 1$ . Check that the resulting  $w$  corresponds to the familiar forward-difference approximation from class. (You could alternatively solve  $Aw = e_{m+1}$  analytically here, since it is  $2 \times 2$ .) In Julia, you can pass `0//1` for  $x_0$  and it will return exact rational weights.
3. Now evaluate it for  $x_0 = 0$  (`0//1` in Julia for exact results) and  $x = [0, 1, 2, 3]$  with  $m = 1$ , i.e. using  $n = 4$  equally spaced points  $\geq x_0$ . Use the resulting weights, in the formula scaled by  $h$  as above, to approximate the derivative  $f'(1)$  for  $f(x) = \sin(x)$ , and plot the relative error (compared to the exact derivative) as a function of  $h$  on a log-log scale, similar to the course notebook. What power law in  $h$  does the truncation error (approximately) seem to follow? That is, what is the “order of accuracy”?
4. Derive the stencil equation  $Aw = e_{m+1}$  above: write out the first  $n - 1$  terms of the Taylor series for  $f(x_0 + \delta x)$ , and try to find a linear combination of this series evaluated at  $\delta x = x_k - x_0$  for  $k = 1 \dots n$  in such a way that you obtain  $f^{(m)}(x_0)$ .

## Problem 3

Consider the following system  $g(x, p) = 0$  of two nonlinear equations in two variables  $x \in \mathbb{R}^2$ , parameterized by three parameters  $p \in \mathbb{R}^3$ :

$$g(x, p) = \begin{pmatrix} p_1 x_1^2 - x_2 \\ x_1 x_2 - p_2 x_2 + p_3 \end{pmatrix} = \begin{pmatrix} 0 \\ 0 \end{pmatrix}.$$

For  $p = [1, 2, 1]$  this has an exact solution  $x = [1, 1]$ .

1. What are the Jacobian matrices  $\frac{\partial g}{\partial x}$  and  $\frac{\partial g}{\partial p}$ ? (That is, as defined in class, the linear operators such that  $dg = \frac{\partial g}{\partial x}dx + \frac{\partial g}{\partial p}dp$  for any change  $dx$  and  $dp$ , to first order.)
2. In Julia (or Python etc.), implement Newton's method to solve  $g(x, p) = 0$  from a given starting guess  $x$ . Using  $p = [1, 2, 1]$ , start your Newton iteration at  $x = [1, 2]$  and show that it converges rapidly to  $x = [1, 1]$  (it should converge to machine precision in  $< 10$  steps).
3. Now, consider the function  $f(p) = \|x(p)\|$ , where the “implicit function”<sup>1</sup>  $x(p)$  is a solution of  $g(x, p) = 0$ . Given a solution  $x(p)$  for some  $p$ , explain how to compute  $\nabla f$  (see the adjoint-method notes from lecture 5). Implement this algorithm in Julia (etc.), and validate it against a finite-difference approximation for  $p = [1, 2, 1]$ ,  $x(p) = [1, 1]$ , and a random small  $\delta p$  (solving for  $x(p + \delta p)$  by Newton's method starting from  $x(p)$ ).

## Problem 4

1. Suppose that  $f(A)$  is a function that maps (real)  $m \times m$  matrices to  $m \times m$  matrices, and its derivative is the linear operator  $f'(A)[dA]$ . For the Frobenius inner product  $\langle X, Y \rangle = \text{trace}(X^T Y)$ , it turns out that we typically have

$$\langle X, f'(A)[Y] \rangle = \langle f'(A^T)[X], Y \rangle,$$

which conceptually corresponds to “transposing” the linear operator  $f'(A)^T = f'(A^T)$ . Your job is to show this.

- (a) Show this for  $f(A) = A^n$  for any  $n \geq 0$ . (Hint: From the product rule, it is easy to see that  $f'(A)[dA] = \sum_{k=0}^{n-1} A^k dA A^{n-1-k}$ ; we've already seen this explicitly for several  $n$ . Combine this with the cyclic rule for the trace.)

It immediately follows that this identity also works for any  $f(A)$  described by a Taylor series in  $A$  (any “analytic”  $f$ ), such as  $e^A$ .

- (b) Show this for  $f(A) = A^{-1}$ .

(You can then compose the above cases to show that it works for any  $f(A) = p(A)q(A)^{-1}$  for any polynomials  $p$  and  $q$ , i.e. for any rational function of  $A$ . You need not do this, however.)

2. Consider the function  $f(A) = \det(\exp(A))$ .
  - (a) Write  $f'(A)[dA]$  in terms of  $\exp'(A)[dA]$ . (You learned how to compute  $\exp'$  in pset 1.)
  - (b) Using the identity from the previous part, write  $\nabla f$  in a way that can be evaluated efficiently (“reverse mode”) using only one or two evaluations of  $\exp$  (and/or  $\exp'$ ) and  $\det$ , independent of the size of  $A$ .
3. Check your answer from the previous part in Julia (or Python etc.): choose a random  $5 \times 5$   $A = \text{randn}(5, 5)$  and a random small  $dA = \text{randn}(5, 5) * 1e-8$ , compute  $df = f(A + dA) - f(A)$  and  $\nabla f$  (at  $A$ ), and verify that  $df \approx \langle \nabla f, dA \rangle$ . Compute  $\exp'(A)[dA]$  using the same technique as in pset 1.

<sup>1</sup>This  $x(p)$  can be defined uniquely in some neighborhood of a root like the one above, thanks to the implicit-function theorem.