

MERCURY DEVELOPER'S MEETING: ANL ACTIVITIES

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ENABLING DATA SERVICES (MOCHI)

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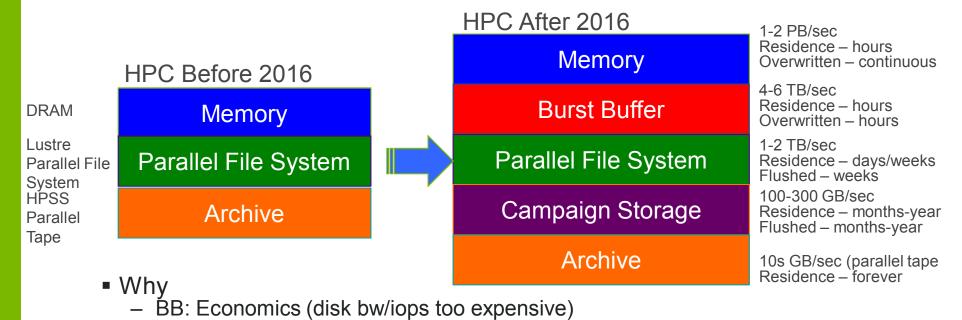
Carnegie Mellon University

The HDF Group

Los Alamos National Laboratory



ARCHITECTURAL TRENDS: MORE STORAGE/MEMORY LAYERS...



PFS: Maturity and BB capacity too small
Campaign: Economics (tape bw too expensive)
Archive: Maturity and we really do need a "forever"

Slide from Gary Grider (LANL).



WHAT COMES NEXT?

- Assumptions
 - New layers in storage hierarchy, lower latencies
 - Storage resources will be highly contended for
 - No "holy grail" emerges that solves everyone's problems
- Alternative model to the "PFS for data management"
 - Multiple services employed for different classes of data
 - Specialization for scalability/efficiency/productivity
 - In some cases, co-design with applications



Specialized data services are already here!		Provisioning	Comm.	Local Storage	Fault Mgmt. and Group Membership	Security
	ADLB Data store and pub/sub.	MPI ranks	MPI	RAM	N/A	N/A
	DataSpaces Data store and pub/sub.	Indep. job	Dart	RAM (SSD)	Under devel.	N/A
	DataWarp Burst Buffer mgmt.	Admin./ sched.	DVS/ Inet	XFS, SSD	Ext. monitor	Kernel, Inet
	FTI Checkpoint/restart mgmt.	MPI ranks	MPI	RAM, SSD	N/A	N/A
	Kelpie Dist. in-mem. key/val store	MPI ranks	Nessie	RAM (Object)	N/A	Obfusc. IDs
	SPINDLE Exec. and library mgmt.	Launch MON	TCP	RAMdisk	N/A	Shared secret

OUR GOAL

Enable composition of data services for DOE science and systems

- Application-driven
 - Identify and match to science needs
 - Traditional data roles (e.g., checkpoint, data migration)
 - New roles (e.g., equation of state/opacity databases)
- Composition
 - Develop/adapt building blocks
 - Communication
 - Concurrency
 - Local Storage
 - Resilience
 - Authentication/Authorization
- Enable rapid development of specialized services
- Don't built new services from scratch every time



SERVICE COMPONENTS AND MERCURY **INTEGRATION**

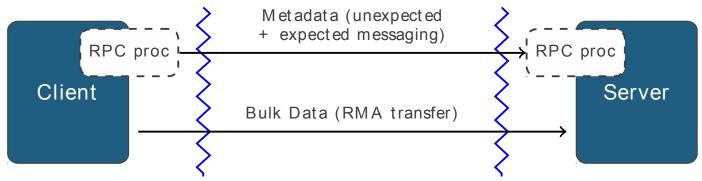


COMMUNICATION: MERCURY

https://mercury-hpc.github.io/

The Mercury RPC system is the core building block of all of the services we are building in the Mochi project.

- Portable across systems and network technologies
- Builds on lessons learned from IOFSL, Nessie, Inet, and others
- Efficient bulk data movement to complement control messages







PROGRAMMING MODELS FOR MERCURY SERVICES

Concurrency, low latency, and ease of use

- The Mercury API uses an event-driven callback model that performs well, but we need more than that to build flexible services rapidly
 - Programmability of callback continuation (i.e. avoid stack-ripping)
 - Assignment of work to CPU cores
 - Scheduling for extremely high concurrency
- We've elected to combine Mercury with user-level threading
 - Similar to co-routines, green threads, fibres, etc.
 - Hide callbacks from service developer
 - Combine latency benefits of Mercury's event model with programmability of conventional thread model

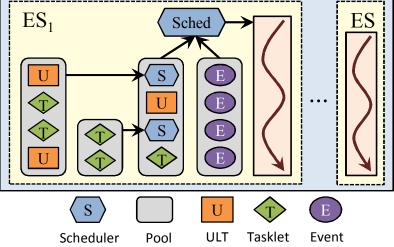


CONCURRENCY: ARGOBOTS

Argobots is a lightweight threading/tasking framework

- Features relevant to I/O services:
 - Flexible mapping of work to hardware resources (which ULTs can run on which cores)
 - Ability to delegate service work with fine granularity across those resources (lightweight tasklets and ULTs)
 - Modular, user-definable scheduling
- We developed asynchronous bindings to:
 - Mercury
 - LevelDB
 - POSIX I/O
- Working with Argobots team to identify needed functionality (e.g., idling) for our use cases

Argobots Execution Model





MARGO

Binding Mercury to Argobots

- Margo is a support library that provides Argobots-aware bindings to Mercury
 - Mercury operations are normally expressed with a post/callback model
 - Margo instead provides blocking functions that will suspend and resume the calling ULT automatically
- Internal progress loop drives Mercury progress and callbacks
 - Can run on dedicated core/thread or share resources of caller
- Incoming RPC handlers are executed as new user level threads
- "_timed()" versions of each function can be used to specify timeouts for cancellation/completion



WHAT DOES THE CODE LOOK LIKE?

A few Margo examples

- Macros are used to define wrappers for RPC handlers that will be launched as user-level-threads
 https://xgitlab.cels.anl.gov/sds/margo/blob/master/examples/my-rpc.c#L111
- Services code (both clients and servers) issue Mercury calls as blocking operations
 https://xgitlab.cels.anl.gov/sds/margo/blob/master/examples/my-rpc.c#L58
- Explicit concurrency is achieved by creating and joining user-level threads https://xgitlab.cels.anl.gov/sds/margo/blob/master/examples/client.c#L137
- Boilerplate: after starting Mercury library, set up Argobots and tell Margo what resources to use for driving progress and running RPC handlers https://xgitlab.cels.anl.gov/sds/margo/blob/master/examples/server.c#L53

HIGH-LEVEL TUNING OBSERVATIONS

Observations from Argobots/Margo/Mercury/CCI stack

- CPU usage
 - Trade-off between latency and how frequently you are willing to busy-spin on a CPU core
 - CCI level
 - Reduce CPU usage by using CCI poll feature
 -DNA_CCI_USE_POLL:BOOL=ON), recently upstreamed in Mercury
 - Argobots level
 - Reduce CPU usage with abt-snoozer scheduler (optional add-on for use with Margo and other support libraries that block on I/O)
- Bandwidth
 - So far so good for large transfers



HIGH-LEVEL TUNING OBSERVATIONS (CONT)

Observations from Argobots/Margo/Mercury/CCI stack

- Latency
 - Pay close attention to callback scheduling and context switching points in progress loop (Margo in our case)
 - Disable checksumming in Mercury (-DMERCURY_USE_CHECKSUMS:BOOL=OFF)
 - Tcmalloc reduces memory management latency
- Ongoing
 - Isolating sources of variability
 - Investigating memory allocation and reuse
 - Testing more scenarios

We need to standardize on how to measure and tune more of these things...





MERCURY DEVELOPER'S MEETING: ANL EXAMPLES AND BENCHMARKING

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DEPLOYMENT SCENARIOS

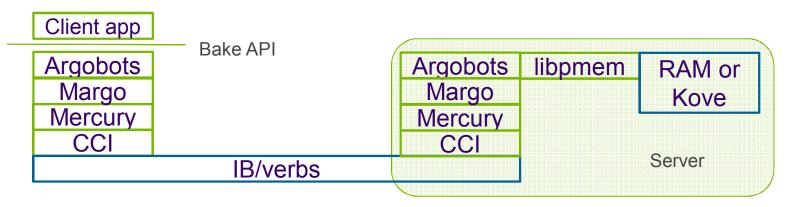
When, where, and how to run composable data services

- Use cases will depend on both the system and the application, so we need to be flexible
- Examples:
 - Co-located alongside (or even within) application processes
 - Set-aside nodes in a single job
 - Co-scheduled across jobs
 - Persistent (or semi-persistent) services
- Storage may reside on compute nodes, burst buffer nodes, or off-system
- Duration may be transient, on-demand, or long-running
- Some of these scenarios will rely on resource management integration



BUILDING BLOCKS FOR STORAGE SERVICES

 In-progress low-level storage service, "BAKE", to serve as a basis for future data services

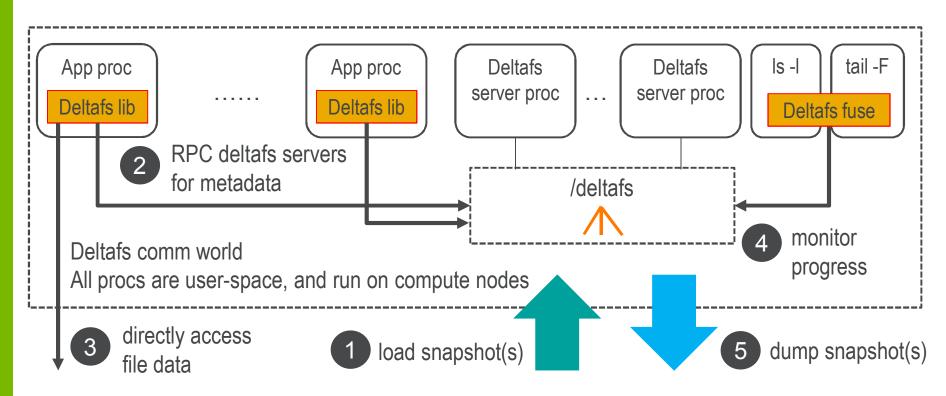


- Looks like a lot of components
- ... but data copies and context switches are rare
- RDMA transfer directly from client app address space to storage device that has been mapped into server address space
- Argobots/Margo/Mercury handle concurrency without any OS threads



TRANSIENT FILE SYSTEM VIEWS: DELTAFS

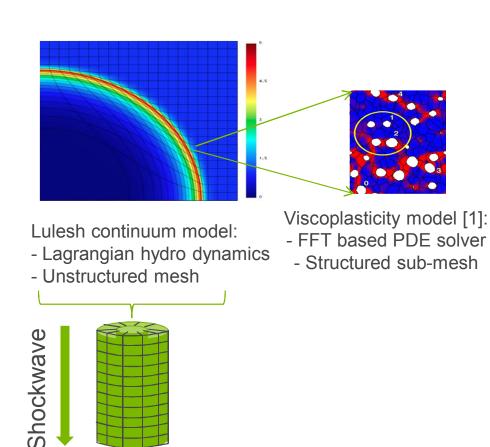
Supporting legacy POSIX I/O in a scalable way.



Credit: Qing Liu, CMU



CO-DESIGNING COUPLED MODELS



- Future applications are exploring the use of multi-scale modeling
- As an example: Loosely coupling continuum scale models with more realistic constitutive/response properties
 - e.g., Lulesh from ExMatEx
- Fine scale model results can be cached and new values interpolated from similar prior model calculations

R. Lebensohn et al, Modeling void growth in polycrystalline materials, Acta Materialia, http://dx.doi.org/10.1016/j.actamat.2013.08.004.



CO-DESIGNING A FINE SCALE MODEL DATABASE

Goals

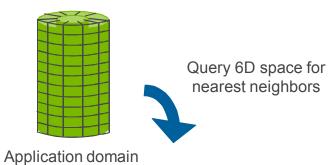
- Minimize fine scale model executions
- Minimize query/response time
- Load balance DB distribution

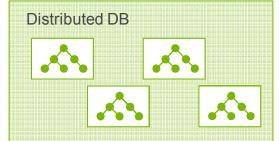
Approach

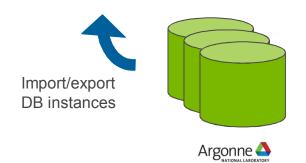
- Distributed approx. nearest-neighbor query
- Partitioned spatial data structures across nodes
- Import/export to persistent store

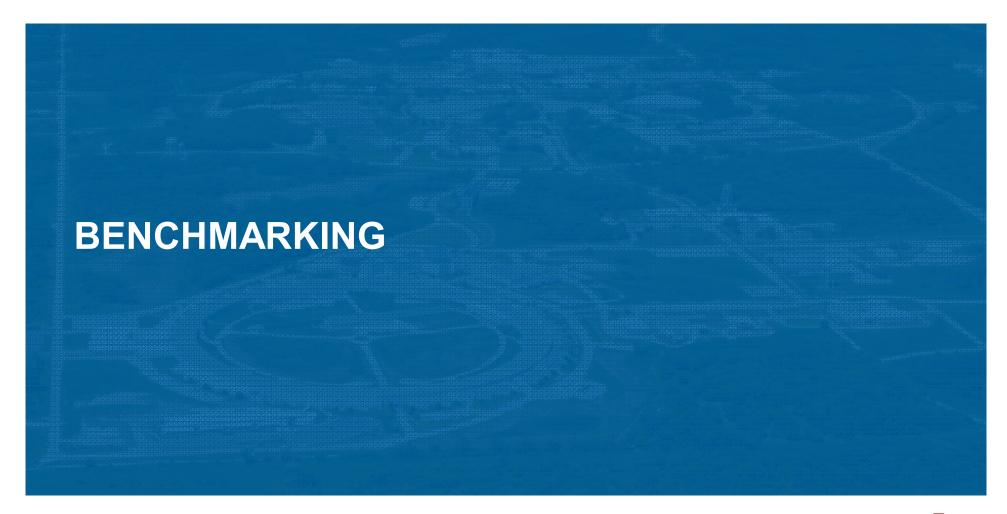
Status

- Mercury-based, centralized in-memory DB service
- Investigating distributed, incremental nearest-neighbor indexing







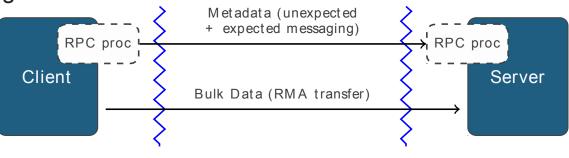




MICROBENCHMARKS

Initial code at https://github.com/mercury-hpc/mercury-benchmarks

- RPC round-trip latency
- Bulk transfer rate
- Concurrent RPCs/bulks
- Overheads (not yet directly measured)
 - RPC data (un)marshalling
 - HG/NA work queue management
 - NA → plugin translation



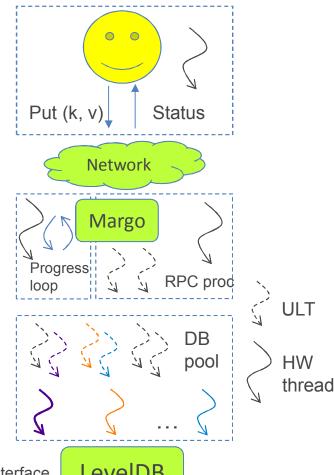
Network Abstraction Layer



BENCHMARKING COMMON DESIGN PATTERNS

Capturing Mercury's surrounding ecosystem

- Example: LevelDB wrapper
- Mixed RPC+bulk workflow
 - k/v pair represented as a bulk handle
 - Server receives "put" RPC, RDMA reads using k/v handle
- Multithreading scenarios
 - HG progress, DB op setup / dispatch
- Interaction with other layers
 - Argobots ULT creation, context switching, thread handoff



NOTE: LevelDB doesn't have an async interface, otherwise could unify the RPC proc and DB pools









TOPICS FOR AFTERNOON DISCUSSION

- Multi-user environments (secure messaging)
- Group membership (dynamic, fault tolerant service participation)
- Transport roadmaps (libfabric)
 - currently using BMI from PVFS, CCI as prototyping vehicles
- Standardizing set of benchmarks?
- Leveraging LNET experience



RESOURCES

- Mercury: https://github.com/mercury-hpc/mercury
- Mercury benchmarks: https://github.com/mercury-hpc/mercury-benchmarks
- Argobots: https://collab.cels.anl.gov/display/ARGOBOTS/Argobots+Home
- Margo: https://xgitlab.cels.anl.gov/sds/margo
- abt-snoozer: https://xgitlab.cels.anl.gov/sds/abt-snoozer





