

DM552 exercises

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1. Under what conditions on xs and ys does the following equation hold?

$$[x \mid x \leftarrow xs, y \leftarrow ys] \equiv [x \mid y \leftarrow ys, x \leftarrow xs]$$

2. Define appropriate versions of the library functions

```
repeat    :: a → [a]
repeat x  = xs where xs = x : xs
take      :: Int → [a] → [a]
take n _ | n ≤ 0 = []
take _ []      = []
take n (x : xs) = x : take (n - 1) xs
replicate :: Int → a → [a]
replicate n      = take n ∘ repeat
```

for the following type of binary trees:

```
data Tree a = Nil | Node a (Tree a) (Tree a)
```

3. You are given the following data types:

```
data Direction = L | R deriving Show
data Tree a = Nil | Node a (Tree a) (Tree a) deriving Show
```

The direction data type will be used to define a path in the tree. Define the following functions:

```
elementAt :: [Direction] → Tree a → Maybe a
```

which - for a given path and tree - returns the value at the node reached, if the path is followed (it returns *Nothing* in case the path is invalid, or specifies a *Nil*-node).

$$\text{modifyAt} :: (a \rightarrow a) \rightarrow [\text{Direction}] \rightarrow \text{Tree } a \rightarrow \text{Tree } a$$

which - for a given path and tree - applies a function to the value at the node. Invalid paths will just leave the tree unchanged.

4. Define a function

$$\text{toTree} :: [a] \rightarrow \text{Tree } a$$

which converts a list to a tree, filling each level in the tree from left to right.

[1, 2, 3, 4, 5] would be converted to

$$\text{Node } 1 (\text{Node } 2 (\text{Node } 4 \text{ Nil Nil}) (\text{Node } 5 \text{ Nil Nil})) (\text{Node } 3 \text{ Nil Nil})$$

HINT: First define a function which splits a list into levels: [1, 2, 3, 4, 5, 6, 7, 8, 9, 10] becomes [[1], [2, 3], [4, 5, 6, 7], [8, 9, 10]]. Then define left- and right descends on this level list representation of the tree. Using these functions, one can make the conversion to the *Tree*-datatype.

5. You are given types to represent expressions in a small interpreted language.

```

data Expression =
  Var String    -- Variable
  | Val Int     -- Integer literal
  | Op Expression Bop Expression -- Operation
  deriving (Show, Eq)
  -- Binary (2-input) operators
data Bop =
  Plus
  | Minus
  | Times
  | Divide
  | Gt
  | Ge
  | Lt
  | Le

```

| *Eq*
deriving (*Show*, *Eq*)

We store the currently defined variables in a function (see the last slides from lecture 3)

type *State* = *String* → *Maybe Int*

Implement an evaluator for expressions:

evalE :: *State* → *Expression* → *Either String Int*

The evaluator either returns an error message or the evaluated value - for boolean outputs, the convention is that 0 is false and 1 is true.