

Лекция 9

Методы проекции градиента.

$$\min_{x \in \mathbb{R}^n} f(x) \rightarrow \text{loc. p-угл}$$

$$\begin{aligned} \text{GD: } x^{k+1} &= x^k - h \nabla f(x^k) = \\ &= \underset{x \in \mathbb{R}^n}{\operatorname{argmin}} \left\{ f(x^k) + \langle \nabla f(x^k), x - x^k \rangle + \right. \\ &\quad \left. + \frac{1}{2h} \|x - x^k\|_2^2 \right\} \end{aligned}$$

$$\min_{x \in Q} f(x) \rightarrow \text{loc. p-угл} \rightarrow \text{loc. мн-бо}$$

$$\begin{aligned} \text{PGD: } x^{k+1} &= \pi_Q(x^k - h \nabla f(x^k)) = \\ &= \underset{x \in Q}{\operatorname{argmin}} \left\{ f(x^k) + \langle \nabla f(x^k), x - x^k \rangle + \right. \\ &\quad \left. + \frac{1}{2h} \|x - x^k\|_2^2 \right\} \end{aligned}$$



$$\min_{x \in Q} f(x) \rightarrow \text{loc. мн-бо}$$

$$|f(y) - f(x)| \leq M_p \|y - x\|_p, \quad p \in [1, 2]$$

$\| \cdot \|_p$.

$f(x) - \text{лок. п-угл}$:

1-смкно-бескнчнл. омн. р-рнлн

$$V(y, x) = f(y) - f(x) - \langle \nabla f(x), y - x \rangle$$

grubepunkt

$$J(x) = \frac{1}{2} \|x\|_2^2$$

↑
1- частные
отн. 2- полные

$$V(y, x) = \frac{1}{2} \|y - x\|_2^2$$

$$x^{k+1} = \underset{x \in Q}{\operatorname{argmin}} \left\{ f(x) + \langle \nabla f(x^k), x - x^k \rangle + \frac{1}{2} \|x - x^k\|_2^2 \right\}$$

$$\frac{1}{h} V(x, x^k)$$

$$\bar{x}^n = \frac{1}{n} \sum_{k=0}^{n-1} x^k$$

§ 2 Гауссов
метод

$$f(\bar{x}^n) - f(x_*) \leq \frac{M_p R_p}{\sqrt{n}}$$

$$\begin{aligned} p &= 2 \\ J(x) &= \frac{1}{2} \|x\|_2^2 \\ V(y, x) &= \frac{1}{2} \|y - x\|_2^2 \end{aligned} \quad \left| \begin{array}{l} M_p \rightarrow |f(y) - f(x)| \leq \\ \leq M \|y - x\|_2 \\ R_p = \|x^0 - x_*\|_2 \end{array} \right. \quad \downarrow \quad \sqrt{2 V(x_*, x^0)}$$

Пример (одномерный
на симплексе)

норма

$$Q = S_n(1) = \{x \in \mathbb{R}_+^n : \sum_{i=1}^n x_i = 1\}$$

$$A = n \begin{bmatrix} \ddots & \ddots & \ddots \\ \vdots & \vdots & \vdots \\ \ddots & \ddots & \ddots \\ \hline & n & n \end{bmatrix}$$

$$\inf_{x \in S_n(1)} \left\{ \frac{1}{2} \langle x, Ax \rangle - \langle b, x \rangle \right\}$$

$$J(x) = \sum_{i=1}^n x_i \ln x_i, \quad x^0 = (1/n, \dots, 1/n)$$

$$V(y, x) = \sum_{i=1}^n y_i \ln(y_i/x_i)$$

$$\begin{aligned} x^{k+1} &= \underset{x \in S_n(1)}{\operatorname{argmin}} \left\{ f(x^k) + \langle \nabla f(x^k), x - x^k \rangle + \right. \\ &\quad \left. + \frac{1}{h} \sum_{i=1}^n x_i \ln(x_i/x_i^k) \right\} \\ x_i^{k+1} &= \frac{x_i^k \exp(-h \partial f(x^k)/\partial x_i)}{\sum_{j=1}^n x_j^k \exp(-h \partial f(x^k)/\partial x_j)} \quad h = \frac{\varepsilon}{M_p^2} \end{aligned}$$

$\# \nabla f(x^k) \rightarrow O(n^2)$

$\# \text{Числосчмб "проекнчнх базисов"}$

$$O(n)$$

$$M_p = \max_{x \in S_n(1)} \|\nabla f(x)\|_q \quad \frac{1}{p} + \frac{1}{q} = 1$$

$$\begin{array}{l} p=1 \\ q=\infty \end{array} \quad \max_{x \in S_n(1)} \|Ax-b\|_\infty \leq \max_{i,j} |A_{ij}| + \max_i |b_i|$$

$$\begin{aligned} R_p &= \sqrt{2V(x_0, x^0)} \leq \max_{x \in S_n(1)} \sqrt{2 \sum_{i=1}^n x_i \ln(x_i/n)} \leq \\ &x_0 = (1, 0, \dots, 0) \quad \leq \sqrt{2 \ln n} \end{aligned}$$

$$M_p R_p = \underbrace{\left(\max_{i,j} |A_{ij}| + \max_i |b_i| \right)}_{O(1)} \sqrt{2 \ln n}$$

$$p=2 \quad M_2 = \max_{x \in S_n(1)} \|Ax-b\|_\infty \simeq \max_{k \in [n]} \|A\|_2 + \|b\|_2$$

$$R_2 = \max_{x \in S_n(1)} \|x^0 - x\|_2 \simeq O(1)$$

$$M_2 R_2 \sim \sqrt{n} \quad \text{vs} \quad M_1 R_1 \sim \sqrt{\ln n}$$

Гор. оптимизация

$$\min_{x \in Q} f(x)$$

$$|f(y) - f(x)| \leq$$

$$\leq M \|y - x\|$$

$$x^{k+1} = \arg \min_{x \in Q} \left\{ f(x^k) + \langle \nabla f(x^k), x - x^k \rangle + \frac{1}{h} V(x, x^k) \right\}$$

$$h = \varepsilon / M^2$$

1-е кратн.

беск.
отн. н. н.

$\nabla f(x)$ — не определен

Например, $Ax = b \rightarrow$ разг.

A — опред. разг.

$\nabla f(x, z)$ — смеш. разг.

1) $E_z \nabla f(x, z) \equiv \nabla f(x)$, где бар $x \in Q$

2) $E \left[\|\nabla f(x, z)\|_*^2 \right] \leq M^2$

$\|\cdot\|_*$ — смеш. норма $\propto \|\cdot\|$, т.е.

$$\|y\|_* = \sup_{\|x\| \leq 1} \langle y, x \rangle$$

смешанное
нормальное
случай

$$x^{k+1} = \arg \min_{x \in Q} \left\{ f(x^k) + \langle \nabla f(x^k), x - x^k \rangle + \frac{1}{h} V(x, x^k) \right\}$$

Пример $Q = \mathbb{R}^n$: $x^{k+1} = x^k - h \nabla f(x^k, z^k)$, z^k i.i.d.
SGD

$$\underbrace{E[f(\bar{x}^n)] - f(x_*)}_{\text{Несправедл. нер.}} \leq \frac{MR}{\sqrt{n}}$$

$R = \sqrt{2V(x_*, x^0)}$

Пример (оптимизация в задачах)

норма \downarrow $Q = S_n(1) = \{x \in \mathbb{R}_+^n : \sum_{i=1}^n x_i = 1\}$

 $A = n \begin{bmatrix} \ddots & \ddots & \ddots \\ \vdots & \ddots & \ddots \\ 1 & \cdots & 1 \end{bmatrix}$
 $p=1$

$$f(x) = \min_{x \in S_n(1)} \left\{ \frac{1}{2} \langle Ax, x \rangle - \langle b, x \rangle \right\}$$
 $d(x) = \sum_{i=1}^n x_i \ln x_i, \quad x^0 = (1/n, \dots, 1/n)$

$$x_i^{k+1} = \frac{x_i^k \exp(-h \delta f(x^k, \zeta)/\delta x_i)}{\sum_{j=1}^n x_j^k \exp(-h \delta f(x^k, \zeta)/\delta x_j)}$$

если
 $\nabla f(x^k, \zeta)$
 не линейн.,
 но непрм.
 время. $O(n)$

$$\nabla f(x, \zeta) = A^{(?)} - b, \quad \text{где}$$

$A^{(?)}$ — это \mathbb{R}^n -матрица, соответствующая A

$$P(\zeta = i) = x_i, \quad i = 1, n. \quad \text{если } x \in S_n(1)$$

1) $E[\nabla f(x, \zeta)] \stackrel{x}{=} \nabla f(x) \quad \text{если } Ax = b$

$$\sum_{i=1}^n (A^{(i)} - b) P(\zeta = i) = // !$$

$$= \sum_{i=1}^n A^{(i)} x_i - b \sum_{i=1}^n x_i = Ax - b$$

$$Ax = [A^{(1)}, \dots, A^{(n)}] \begin{pmatrix} x_1 \\ \vdots \\ x_n \end{pmatrix} =$$

$$= A^{(1)}x_1 + \dots + A^{(n)}x_n$$

$$2) \quad \mathbb{E}_z \|Df(x, z)\|_\infty^2 \leq M^2 \quad \text{if } \rho = 1, \frac{1}{\rho} = 1$$

$$\max_{j=1, \dots, n} \|A^{(j)} - b\|_\infty$$

$\max_{i,j} |A_{ij}| + \max_i |b_i|$
 $\Theta(L)$

$$M = \Theta(L)$$

$$R = \Theta(\sqrt{\ln n})$$

$$\mathbb{E} |f(\bar{x}^n) - f(x_*)| \leq \frac{MR}{\sqrt{n}} = \Theta\left(\frac{\sqrt{\ln n}}{\sqrt{n}}\right)$$

$$\# Df(x, z) \rightarrow \Theta(n)$$

$$\mathbb{E} |f(\bar{x}^n) - f(x_*)| \leq \epsilon$$

$$N \sim \frac{\ln n}{\epsilon^2}$$

$$T = N \cdot \underbrace{\left[\# Df(x, z) + \text{overlap} \right]}_{\Theta(n)} =$$

$$= \Theta\left(\frac{n \ln n}{\epsilon^2}\right).$$



Метод усечного градиента

(Фармер-Булбаг)

§2.6
Вариант
наг.
Булбаг
оргинал.

$$x^{k+1} = \underset{x \in Q}{\operatorname{argmin}} \left\{ f(x^k) + \langle \nabla f(x^k), x - x^k \rangle + \frac{1}{h} \|x - x^k\|_2^2 \right\}$$

$$\Rightarrow x^{k+1} - x^k = \gamma_k (y^k - x^k)$$

$$\begin{cases} x^{k+1} = (1 - \gamma_k) x^k + \gamma_k y^k, & \gamma_k = \frac{2}{k+1} \\ y^k = \underset{y \in Q}{\operatorname{argmin}} \left\{ f(x^k) + \langle \nabla f(x^k), y - x^k \rangle \right\} = \\ = \underset{y \in Q}{\operatorname{argmin}} \langle \nabla f(x^k), y \rangle \end{cases} \quad (*)$$

Теорема. Тогда $\| \nabla f(y) - \nabla f(x) \|_\infty \leq L \| y - x \|_1$

$$2R = \sup_{x, y \in Q} \| y - x \|_1 .$$

$$2 \left(\frac{1}{p} + \frac{1}{2} \right) R$$

$$f(x^N) - f(x_*) \leq \frac{2LR^2}{N+2}$$

где
услов.
метод
Геневала

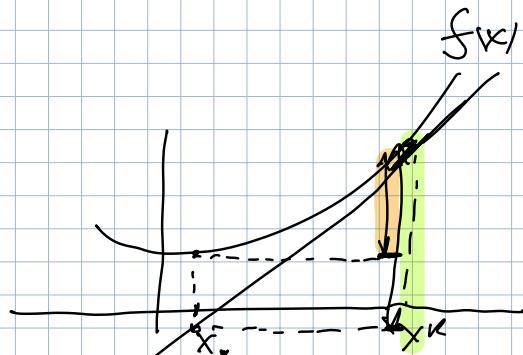
Д-6. Число 1) $\forall x, y \in Q$

$$0 \leq f(y) - f(x) - \langle \nabla f(x), y - x \rangle \leq \frac{L}{2} \| y - x \|_2^2$$

беск. ф

$$\begin{aligned} f(x^{k+1}) - f(x^k) &\leq f(x^k), x^{k+1} - x^k \rangle + \frac{1}{2} \| x^{k+1} - x^k \|_2^2 = \\ &= \gamma_k \langle \nabla f(x^k), y^k - x^k \rangle + \underbrace{\frac{L}{2} \gamma_k^2 \| y^k - x^k \|_2^2}_{\leq R^2} \leq \end{aligned}$$

$$\begin{aligned} &\leq \gamma_k \langle \nabla f(x^k), x_* - x^k \rangle + \gamma_k^2 \frac{LR^2}{2} \\ &\stackrel{(x)}{\leq} \gamma_k (f(x_*) - f(x^k)) + \gamma_k^2 \frac{LR^2}{2} \\ &\text{бесцкнм. } f \end{aligned}$$



$$\langle \nabla f(x^k), x_* - x^k \rangle \leq f(x_*) - f(x^k)$$

$$f(x^k) - f(x_*) \leq \langle \nabla f(x^k), x^k - x_* \rangle$$

$$\begin{aligned} f(x^{k+1}) - f(x^k) &\leq \gamma_k (f(x_*) - f(x^k)) + \gamma_k^2 \frac{LR^2}{2} \\ \delta_k &= \frac{f(x^k) - f(x_*)}{LR^2} \leq \frac{2}{k+2} \end{aligned}$$

$$\delta_{k+1} \leq (1 - \gamma_k) \delta_k + \frac{x^k}{2} = \left(1 - \frac{2}{k+1}\right) \delta_k + \frac{2}{(k+1)^2}$$

\nearrow
 $k=1 \quad \dots \quad 0$

$$\delta_2 \leq \frac{2}{2^2} = \frac{1}{2}$$

No изъякнм грк-тб, тмо $\delta_k \leq \frac{2}{k+2}$

Метод пакострока метод

Гарантия
метода

упр. 1.6

$$\min_{x \in S_n(1)} \frac{1}{2} \langle x, Ax \rangle - \langle b, x \rangle$$

$$Df(x) = Ax - b$$

$$(*) \quad \min_{y \in S_n(1)} \langle Df(x^k), y \rangle \Rightarrow$$

$$\begin{array}{l} \min \\ \sum y_i = 1 \\ y_i \geq 0 \end{array} \quad \sum_{i=1}^n \frac{\partial f}{\partial x_i}(x^k) y_i$$

$$i^*(k) = \arg \min_{i=1, n} \frac{\partial f}{\partial x_i}(x^k)$$

$$y^k = \begin{pmatrix} 0 \\ \vdots \\ 0 \\ \overset{i^*(k)}{\rightarrow} 1 \\ 0 \\ \vdots \\ 0 \end{pmatrix} \quad O(n)$$

$$x^{k+1} = (1-\gamma_k) x^k + \gamma_k y_k$$

$$x^k = (1-\gamma_1)(1-\gamma_2) \dots (1-\gamma_{k-1}) z^k$$

$$(1-\gamma_1) \dots (1-\gamma_k) z^{k+1} = \underbrace{(1-\gamma_1) \dots (1-\gamma_{k-1})}_{z^k} \cdot (1-\gamma_k) z^k + \gamma_k y_k$$

$$z^{k+1} = z^k + \underbrace{\frac{\gamma_k}{\prod_{i=1}^k (1-\gamma_i)} y_k}_{B_k} = z^{k+1} + \alpha_k y_k.$$

$$\beta_{k+1} = \beta_k (1 - \gamma_{k+1})$$

$$\alpha_{k+1} = \frac{\gamma_{k+1} \cdot \alpha_k}{\gamma_k (1 - \gamma_{k+1})}$$

$$\nabla f(x^k) = Ax^k - b = A\beta_k z^k - b = \beta_k A z^k - b$$

$$A z^{k+1} = A(z^k + \alpha_k y^k) = A z^k + \alpha_k A y^k = \underbrace{A z^k}_{\text{*i*, *k*}} + \underbrace{\alpha_k A y^k}_{\text{*i*, *k*}}$$

Сложность вычислений $\Theta(n)$!

$$f(x^n) - f(x_*) \leq \varepsilon \rightarrow N \approx \frac{2LR^2}{\varepsilon} = \Theta(\frac{1}{\varepsilon})$$

$$L = \max_{\|h\|_1 \leq 1} \langle h, \nabla f(x) h \rangle = \max_{i,j} |A_{ij}|$$

$$R^2 = \max_{x, y \in S_n(1)} \|y - x\|_1 \leq 2$$

Сложность вычислений $\Theta(n)$

Число итераций $\Theta(\frac{1}{\varepsilon})$

$\frac{1-\varepsilon}{\varepsilon}$
норма

$$\text{Общая трудоизрасход.} = \Theta(\frac{n}{\varepsilon}) + \Theta(n^2)$$

А если числ. огранич. нет, то как

Сложность итерации $\Theta(n^2)$

$$\nabla f(x) = Ax - b$$

число итераций

$$\Theta\left(\frac{1}{\sqrt{\epsilon}}\right) \gg$$

$$f(x^n) - f(x_*) \leq \frac{\eta R^2}{n^2}$$

$$\text{Оценка времени} = \Theta\left(\frac{n^2}{\sqrt{\epsilon}}\right)$$

Учеб. Материалы

Прим-Вып

$$\Theta\left(\frac{n^2}{\sqrt{\epsilon}}\right)$$

vs

$$\Theta\left(n^2 + \frac{n}{\epsilon}\right)$$

$$\overbrace{\Theta\left(\frac{n \ln n}{\epsilon^2}\right)}^{SGD}$$

Лекция 10

Задача с ограничениями

Системы линейных уравнений

Гаусс

$$\min_{\substack{Ax=b \\ x \in Q}} f(x) = \sum_{i=1}^n f_i(x_i)$$

$$x \in \mathbb{R}^n$$