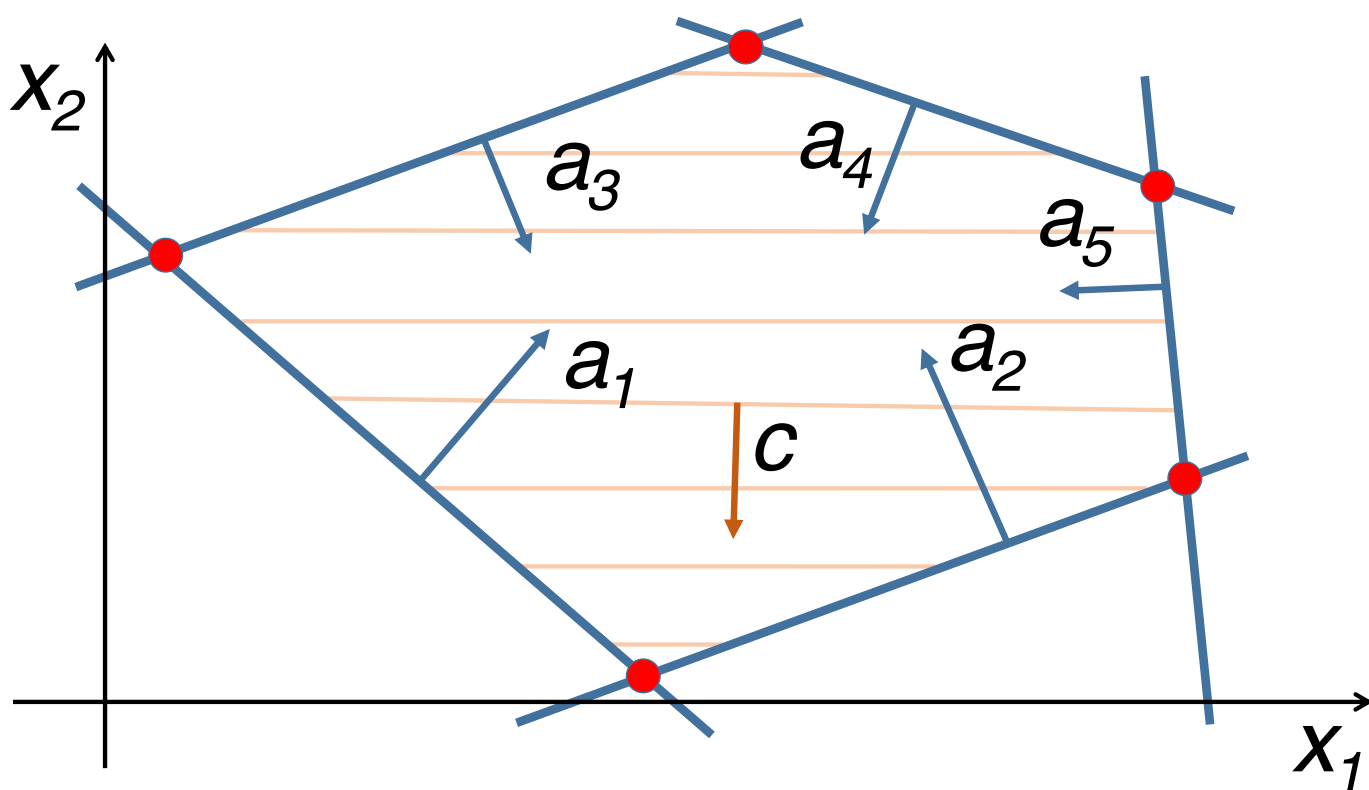


Introduction to Linear Programming

What is LP

Generally speaking, all problems with linear objective and linear equalities/inequalities constraints could be considered as Linear Programming. However, there are some widely accepted formulations.

$$\begin{array}{ll} \min_{x \in \mathbb{R}^n} & c^\top x \\ \text{s.t.} & Ax \leq b \end{array} \quad (\text{LP.Basic})$$



for some vectors $c \in \mathbb{R}^n$, $b \in \mathbb{R}^m$ and matrix $A \in \mathbb{R}^{m \times n}$. Where the inequalities are interpreted component-wise.

Standard form

This form seems to be the most intuitive and geometric in terms of visualization. Let us have vectors $c \in \mathbb{R}^n$, $b \in \mathbb{R}^m$ and matrix $A \in \mathbb{R}^{m \times n}$.

$$\begin{array}{ll} \min_{x \in \mathbb{R}^n} & c^\top x \\ \text{s.t.} & Ax = b \\ & x_i \geq 0, \quad i = 1, \dots, n \end{array} \quad (\text{LP.Standard})$$

Canonical form

$$\begin{aligned} \min_{x \in \mathbb{R}^n} c^\top x & \quad (\text{LP. Canonical}) \\ \text{s.t. } Ax & \leq b \\ x_i & \geq 0, \quad i = 1, \dots, n \end{aligned}$$

Real world problems

Diet problem

Imagine, that you have to construct a diet plan from some set of products: 🍌 🍰 🍗 🥚 🐟 . Each of the products has its own vector of nutrients. Thus, all the food information could be processed through the matrix W . Let also assume, that we have the vector of requirements for each of nutrients $r \in \mathbb{R}^n$. We need to find the cheapest configuration of the diet, which meets all the requirements:

$$\begin{aligned} \min_{x \in \mathbb{R}^p} c^\top x \\ \text{s.t. } Wx & \geq r \\ x_i & \geq 0, \quad i = 1, \dots, p \end{aligned}$$



$$W \in \mathbb{R}^{n \times p},$$

Proteins
Carbs
Fats
Calories
Vitamin D

Requirements

$$r \in \mathbb{R}^n$$

$c \in \mathbb{R}^p$ - cost per 100 g

$$\begin{aligned} \min_{x \in \mathbb{R}^p} c^\top x \\ Wx & \geq r \end{aligned}$$

How to retrieve LP

Basic transformations

Inequality to equality by increasing the dimension of the problem by m .

$$Ax \leq b \leftrightarrow \begin{cases} Ax + z = b \\ z \geq 0 \end{cases}$$

unsigned variables to nonnegative variables.

$$x \leftrightarrow \begin{cases} x = x_+ - x_- \\ x_+ \geq 0 \\ x_- \geq 0 \end{cases}$$

Chebyshev approximation problem

$$\min_{x \in \mathbb{R}^n} \|Ax - b\|_\infty \leftrightarrow \min_{x \in \mathbb{R}^n} \max_i |a_i^\top x - b_i|$$

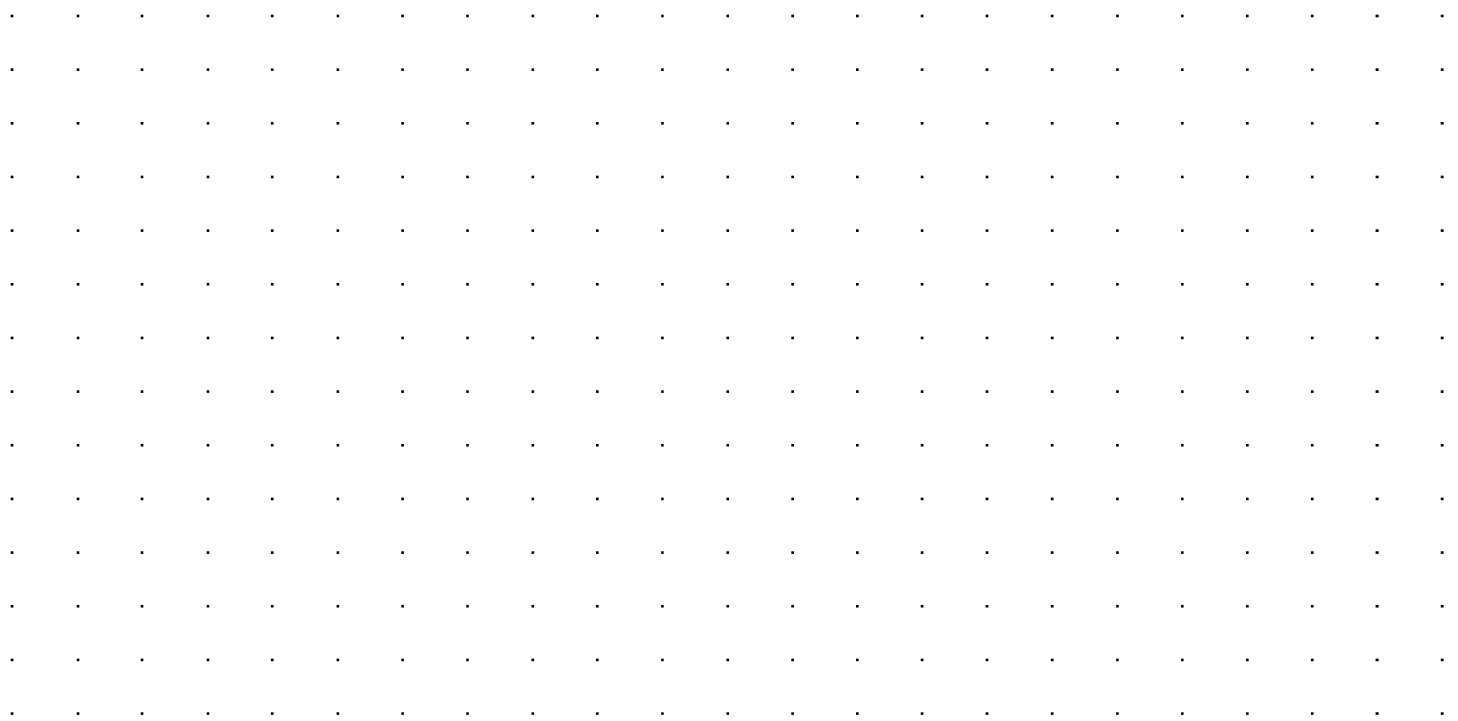
$$\begin{aligned} & \min_{t \in \mathbb{R}, x \in \mathbb{R}^n} t \\ & \text{s.t. } a_i^\top x - b_i \leq t, \quad i = 1, \dots, n \\ & \quad -a_i^\top x + b_i \leq t, \quad i = 1, \dots, n \end{aligned}$$

l_1 approximation problem

$$\min_{x \in \mathbb{R}^n} \|Ax - b\|_1 \leftrightarrow \min_{x \in \mathbb{R}^n} \sum_{i=1}^n |a_i^\top x - b_i|$$

$$\begin{aligned} & \min_{t \in \mathbb{R}^n, x \in \mathbb{R}^n} \mathbf{1}^\top t \\ & \text{s.t. } a_i^\top x - b_i \leq t_i, \quad i = 1, \dots, n \\ & \quad -a_i^\top x + b_i \leq t_i, \quad i = 1, \dots, n \end{aligned}$$

Idea of simplex algorithm

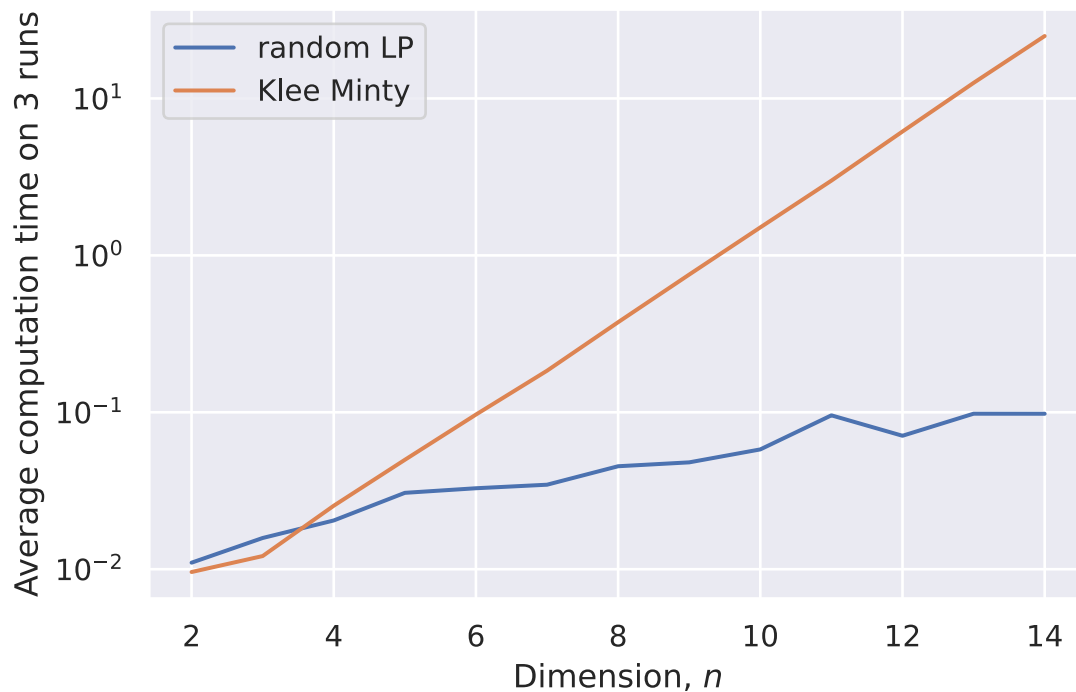


Convergence

Klee Minty example

In the following problem simplex algorithm needs to check $2^n - 1$ vertexes with $x_0 = 0$.

$$\begin{aligned} & \max_{x \in \mathbb{R}^n} 2^{n-1}x_1 + 2^{n-2}x_2 + \dots + 2x_{n-1} + x_n \\ \text{s.t. } & x_1 \leq 5 \\ & 4x_1 + x_2 \leq 25 \\ & 8x_1 + 4x_2 + x_3 \leq 125 \\ & \dots \\ & 2^n x_1 + 2^{n-1}x_2 + 2^{n-2}x_3 + \dots + x_n \leq 5^n \quad x \geq 0 \end{aligned}$$



Summary

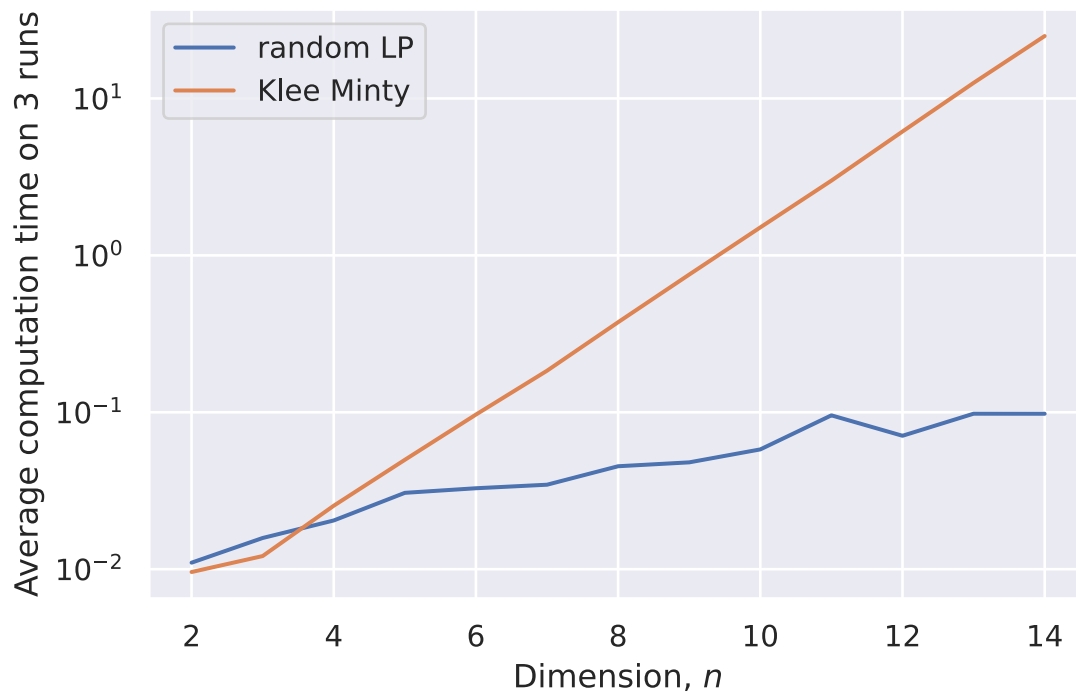
- A wide variety of applications could be formulated as the linear programming.
- Simplex algorithm is simple, but could work exponentially long.
- Khachiyan's ellipsoid method is the first to be proved running at polynomial complexity for LPs. However, it is usually slower than simplex in real problems.
- Interior point methods are the last word in this area. However, good implementations of simplex-based methods and interior point methods are similar for routine applications of linear programming.

Code

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Materials

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