

#### **TREES**

- What is Tree?
- Binary Trees
- Types of Binary Trees
- Binary Tree Traversal
- Properties of Binary Trees



## Objectives:

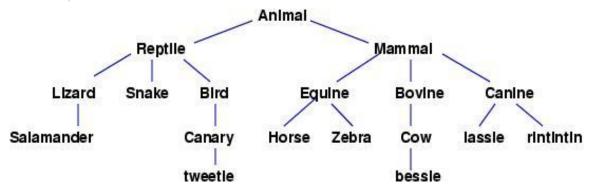
- Identify the terms in determining a tree,
- Differentiate binary trees from expressions trees
- Differentiate the operations that can be performed in a binary tree
- Enumerate the types of binary tree traversal



#### **TREES**

- root is on the top and the leaves are at a bottom
- display data items in a hierarchical structure
- organize information in database systems
- also work as ordered array and a linked list

#### Example:

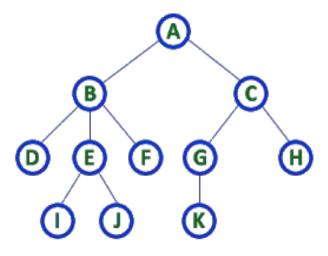




#### **TREES**

a tree is composed of node connected by edges or lines

#### Example:



#### TREE with 11 nodes and 10 edges

- In any tree with 'N' nodes there will be maximum of 'N-1' edges
- In a tree every individual element is called as 'NODE'



#### **TREES:** Basic Terminologies

• Node. what trees composed of connected by edges or lines that contain a value or represent a data structure.

 Root Node. the node at the topmost of the tree and where the tree commonly begins

- Edges.Lines.Paths. lines connecting to nodes describing their relationship
- Leaf Nodes. nodes that are at the bottommost level of the tree and so, they have no children
- Internal Nodes/Child Nodes. nodes that are inside a tree or the nodes below a given node which also known as sibling nodes.



#### **TREES**: Basic Terminologies

• **Subtree.** It is viewed as a complete tree itself, which consists of children and its children's children and so on

 Visiting. A node is said to be visited when a program control arrives at the node to perform operations such as checking the value of one of its data fields, or simply to display it.

 Traversing. To visit all the nodes in a specific order is called traverse or more commonly known as walking the tree, and the action is walk.

Levels. These are generations that a tree has started at the root.

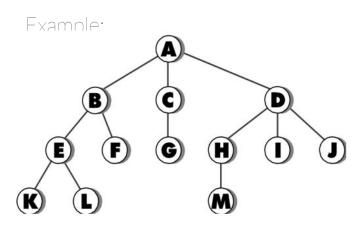


#### **TREES:** Basic Terminologies

- Degree. This refers to the number of nodes in a subtree.
- Depth of the tree. This pertains to the highest level of the tree.
- **Keys.** These represent the value inside a node represented in a circle.



#### **TREES**



- Depth of Tree.
- Degree.

Degree of A is 3 Degree of B is 2 Degree of C is 1 Degree of K is 0 The maximum degree is 3

Root Node.

H, I, J, K, L, M

K, L, M, F,

Level 1 - A

G, I, J

B, E, F

E, K, L

H, M

D, H, I, J

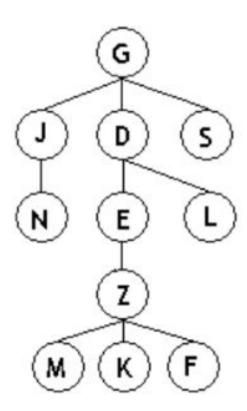
C, G

- Internal Nodes.
- Leaves or Leaf Nodes.
- Levels.

Subtree.



#### **Exercise Problem**



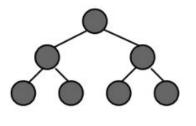
- Root Node
- Depth of the tree
- Nodes at level 4
- Degree of D
- Children of Z
- Total Nodes in the Tree
- Parent of F
- Leaf Nodes
- Child nodes
- Maximum degree of the tree

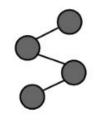


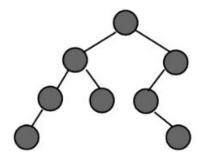
#### **BINARY TREE**

- if every node in a tree can have at most two children
  - left subtree
  - right subtree

can contain only left or right subtree or no nodes at all, in which case it's a least







- two applications of binary trees:
  - binary expression tree
  - binary search tree



#### **Binary Search Tree**

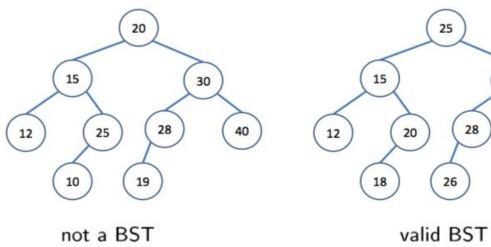
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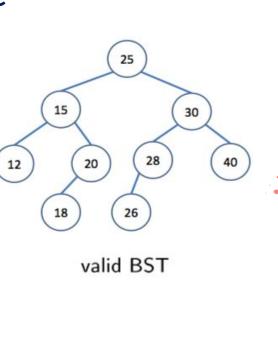
 Hierarchical file structure in a computer system is one of the common trees

Hierarchical file structure is not a binary tree

Binary search trees are trees such that any node has a key which is notified more than the key in its right child node and no less than the key in its left child node

#### **Binary Search Tree**



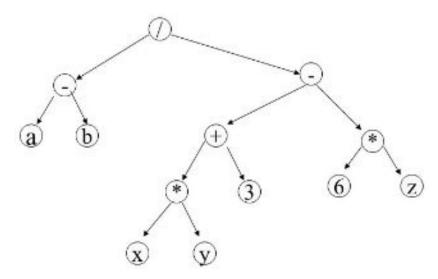




• is a binary tree used in representing an arithmetic expression

• specifically combinations of operators, operands, and order of evaluation

Example: (a-b) / ((x\*y+3)-(6\*z))





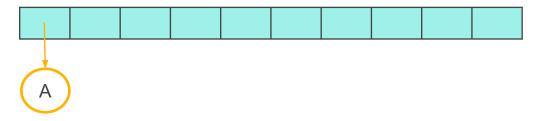
#### Conducting an Expression Tree:

- Read the symbols one at a time
- If the symbol is an operand:
  - Create a one-node tree
  - Push a pointer to the node onto the stack
- If the symbol is an operator:
- Pop two pointers from stack. These pointers represent address of the root nodes of two trees, T1 and T2.
- Form a new tree whose root is the operator and whose and right children are T1 and T2, respectively.
  - Push the pointer to the new tree onto the stack.
- Repeat steps 1 thru 3 until the last symbol has been read and processed.

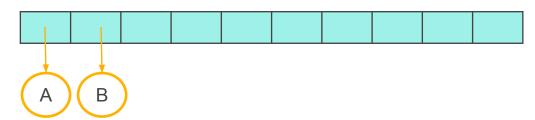


Example Expression: AB+CDE+\*/

- 1. Read the first symbol: "A"
- 2. Create a one-node tree for "A" and push pointer to "A" onto a stack.



- 3. Read next symbol: "B"
- 4. Create a one node tree for "B" and push a pointer to "B" into the stack.

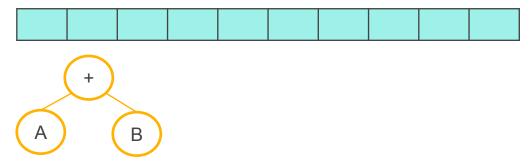


Example Expression: AB+CDE+\*/

5. Read next symbol: "+"

6. Pop two pointers from the stack: Create a tree whose root is "+". Make "A" the left

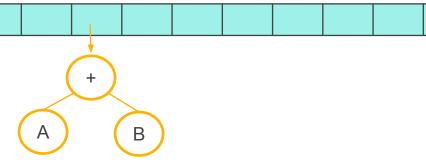
subtree and "B" the right subtree.



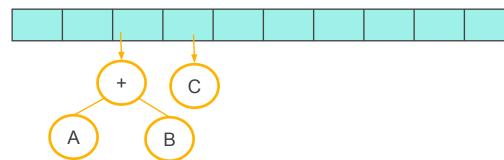


Example Expression: AB+CDE+\*/

7. Push the pointer to the new tree into the stack.



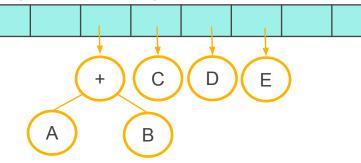
- 8. Read next symbol: "C"
- 9. Create a one node tree for "C" and push a pointer to "C" into the stack.





Example Expression: AB+CDE+\*/

10. Repeat the same process for the next two symbols: "D" and "E".

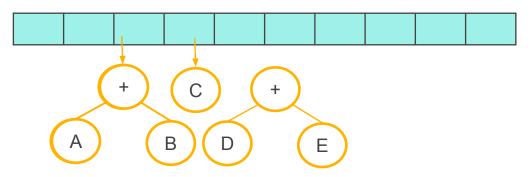


11. Read the next symbol: "+".

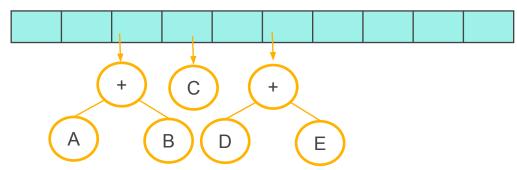


Example Expression: AB+CDE+\*/

12. Pop two pointers from the stack: Create a tree whose root is "+". Make "D" the left subtree and "E" the right subtree.



13. Push the pointer to the new tree into the stack.

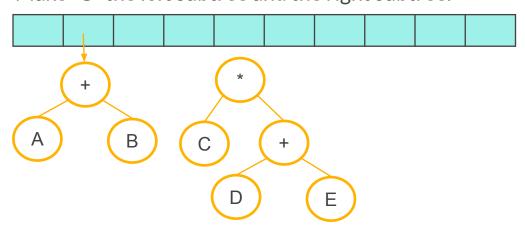




Example Expression: AB+CDE+\*/

12. Read the next symbol: "\*"

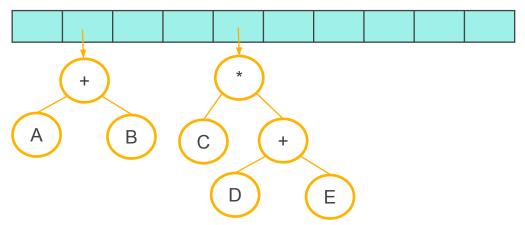
13. Pop two pointers from the stack: Create a tree whose root is "\*". Make "C" the left subtree and the right subtree.





Example Expression: AB+CDE+\*/

16. Push the pointer to the new tree into the stack.



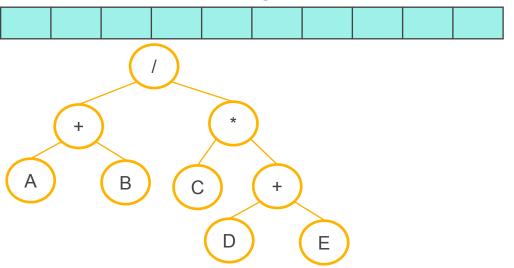
17. Read the next symbol: "/"



Example Expression: AB+CDE+\*/

18. Pop two pointers from the stack: Create a tree whose root is "/".

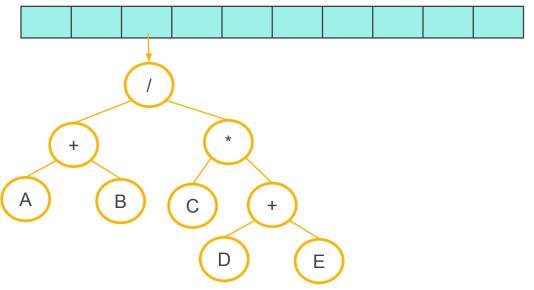
Make the left subtree and the right subtree..





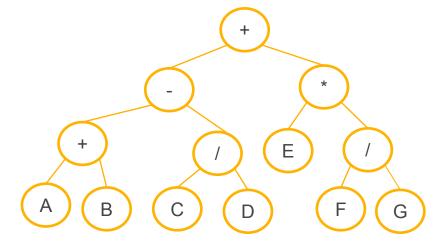
Example Expression: AB+CDE+\*/

19. Push the pointer to the new tree into the stack.



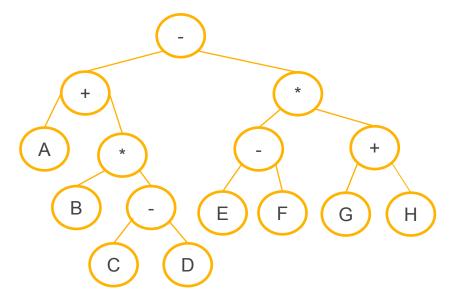


((A + B) - (C/D)) + (E\*(F/G))





(A+(B\*(C-D)))-(E-F\*(G+H))





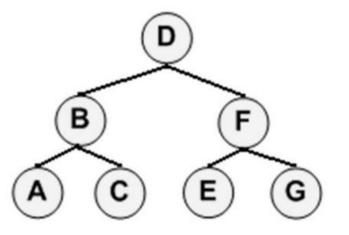
# Binary Tree Traversal

- to visit all the nodes of the tree in order
- not commonly used as finding, inserting, deleting nodes
- useful in some circumstances and simpler than deleting of nodes
  - Inorder
  - Preorder
  - Postorder



Binary Tree Traversal: Inorder Traversal

- If the node has a left subtree:
  - Traverse the left subtree in preorder (recursive call). Once completed, proceed to step 2.
  - Otherwise, proceed to step 2.
- Read the root node
- If the node has a right subtree:
  - Traverse the right subtree in preorder (recursive call).

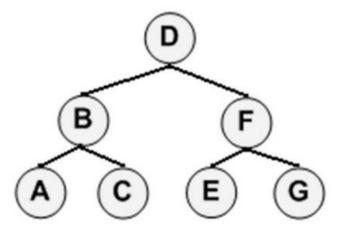


Answer: ABCDEFG



Binary Tree Traversal: Preorder Traversal

- Read the root node.
- If the node has a left subtree:
  - Traverse the left subtree in preorder (recursive call). Once completed, proceed to step 3.
  - Otherwise, proceed to step 3. Read the root node
- If the node has a right subtree:
  - Traverse the right subtree in preorder (recursive call).

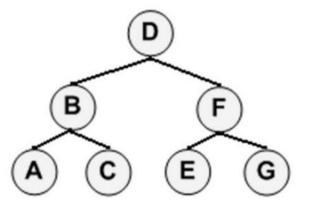


Answer: DBACFEG



Binary Tree Traversal: Postorder Traversal

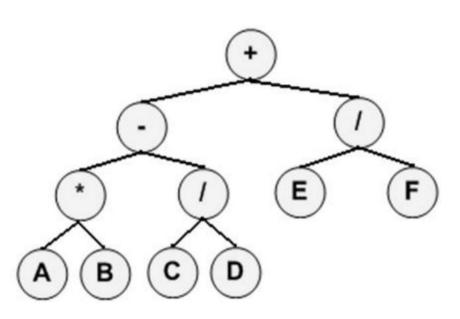
- If the node has a left subtree:
  - Traverse the left subtree in preorder (recursive call). Once completed, proceed to step 2.
  - Otherwise, proceed to step 2.
- If the node has a right subtree:
  - Traverse the right subtree in preorder (recursive call). Once completed, proceed to step 3.
  - Otherwise, proceed to step 3.
- Read the root node



Answer: A C B E G F D



Binary Tree Traversal: Preorder Traversal



Inorder: A \* B - C/D + E/F

Preorder: +-\*AB/CD/EF

Postorder: AB\*CD/-EF/+







