

# Vehicle routing for the communication of time-dependent information

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Presentation Roadef

- The Vehicle Routing Problem (VRP) [1] is one of the most extensively studied problems in operations research due to its methodological interest and practical relevance in many fields such as transportation, logistic, telecommunications, and production.
- This paper presents a vehicle handles the physical information gathering at each node and with technological advances comes the need to incorporate wireless delivery, such that the vehicle must choose whether to go to a node to collect or receive information in the form wireless.

# Description of the Problem

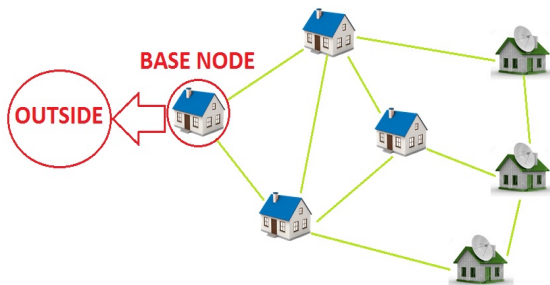


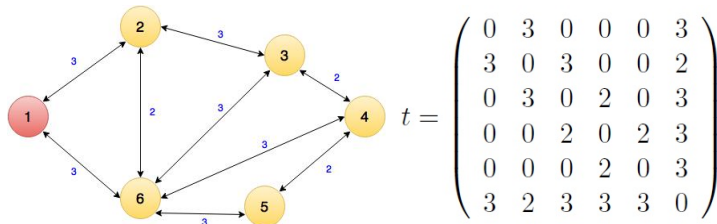
Figure: stations

# Description of the Problem

- (v) The problem is to define routes for the vehicle and also to decide how to collect the information in a way that the amount of information collected to a finite time  $T$  is maximized. on the other hand it is important to guarantee that from time to time the information collected will be sent to the outside.

# Preliminary Model

- (i) Consider a directed graph  $D = (V, A)$ . We have the set  $V$  consisting of  $n$  nodes representing  $n$  stations and a certain amount of  $m$  paths connecting stations represented by the set of arcs  $A$ .
- (ii) Each road  $(i, j) \in A$  has a certain weight  $t_{ij}$  representing the time it takes to travel from  $i$  to  $j$ .

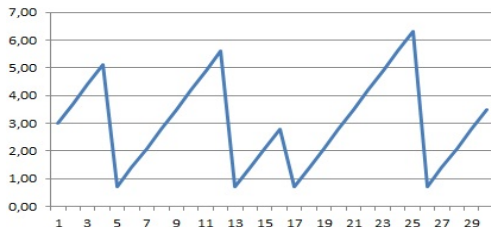


# Preliminary Model

We suppose that the amount of information in node  $j$  in time  $k$ ,  $q_{jk}$ , is proportional to the elapsed time from the last extraction.

$$q_{jk} = (k - t_{last})r_j \quad (r_j \text{ is the constant of proportional})$$

$$q_{j1} = C_j \quad (C_j \text{ is the amount of initial information})$$



**Figure:**  $q_{jk}$  with  $r_j = 0.4$ ,  $C_j = 3$ , collecting or sending wireless in  $k = 4, 12, 16, 25$

# Wireless

The time for transmission depends:

- on the amount of information transmitted.
- the square of the distance between nodes
- other factors for example: the equipment installed in the nodes, obstacles between stations, etc

## Wireless delivery time information

$$\tau_{ijk} = \alpha_{ij} d_{ij}^2 q_{jk} \quad (1)$$

also we can make (M) multiple sending to the same node and the total information time corresponds to the highest.

# Preliminary Model

For each  $i, j \in V$  We define the following decision variables

$$x_{ijk} = \begin{cases} 1 & \text{If the vehicle goes from } i \text{ to } j \text{ in time } k, (i, j) \in A \\ 0 & \text{c.c.} \end{cases}$$

$$w_{jikl} = \begin{cases} 1 & \text{if the information will be sent from node } j \\ & \text{to node } i \text{ at time } k, \text{ and it will take a while } l \\ 0 & \text{c.c.} \end{cases}$$

$$y_{ik} = \begin{cases} 1 & \text{the vehicle is ready to go from the node } i \text{ to a neighboring} \\ & \text{in time } k, \text{ after receiving wireless information} \\ 0 & \text{c.c.} \end{cases}$$



$$\text{minimize } \left\{ \sum_{k=1}^T \max_{j \in V} \{q_{jk}\} \right\}$$

The objective function is to sum every moment the maximum amount information in the nodes

**Subject to:**

$$\sum_{(1,s) \in A} x_{1st_1s} = 1 \quad (2)$$

$$\sum_{(s,1) \in A} x_{s1T} = 1 \quad (3)$$

(1) the vehicle starts from node 1 to a neighboring node.

(2) the vehicle ends at node 1

$$q_{j1} = C_j, j \in V \quad (4)$$

$$q_{jk+1} = q_{jk} \left(1 - \sum_{i \in V(j)} x_{ijk} - \sum_{p=1}^r \sum_{i \in V/\{j\}} w_{jikp}\right) + r_j, j \in V, \quad (5)$$

the constraint (3) amount of initial information in the node  $j$  and (4) means that  $q_{jk+1} = r_j$  if the vehicle has gone through  $j$  or if sent information from  $j$  (wireless) at time  $k$ , in otherwise

$$q_{jk+1} = q_{jk} + r_j$$

$$x_{ijk} \leq \sum_{(j,p) \in A}^{k+t_{jp} \leq T} x_{jp(k+t_{jp})} + \sum_{p=1}^S \sum_{u \in V}^{u \neq j} w_{uj(k+1)p}, \quad \forall (i, j) \in A, \quad (6)$$

$$\sum_{l=1}^S \sum_{j \in V}^{j \neq i} w_{jikl} \leq M \quad \forall i \in V, \quad \forall k \quad (7)$$

$$1 \geq \sum_{(i,j) \in A} x_{ijk} + \left(\frac{1}{M}\right) \sum_{p=1}^S \sum_{(j,i) \in V \times V}^{j \neq i} w_{jikp}, \quad \text{for each } k, \quad (8)$$

$$\alpha_{ji}d_{ij}^2q_{jk}w_{jikl} - l \leq 1, \quad (j, i) \in V \times V (j \neq i) \forall k, \forall l \quad (9)$$

$$y_{ik} \leq \sum_{(i,j) \in A}^{k+t_{ij} \leq T} x_{ij(k+t_{ij})}, \quad \forall i \in V, \forall k; \quad (10)$$

$$w_{jikl} \leq \sum_{(p,i) \in A} x_{pi(k-1)}, \quad \forall (j, i), \forall k \quad (11)$$

$$(1/M) \sum_{l=1}^S \sum_{j \neq i}^{j \in V} w_{jikl} \leq \sum_{s=1}^T y_{i(k+s-1)}; \quad (12)$$

$$x_{ijk}, y_{ik}, w_{jikl} \in \{0, 1\}, \quad (13)$$

# Numerical Results

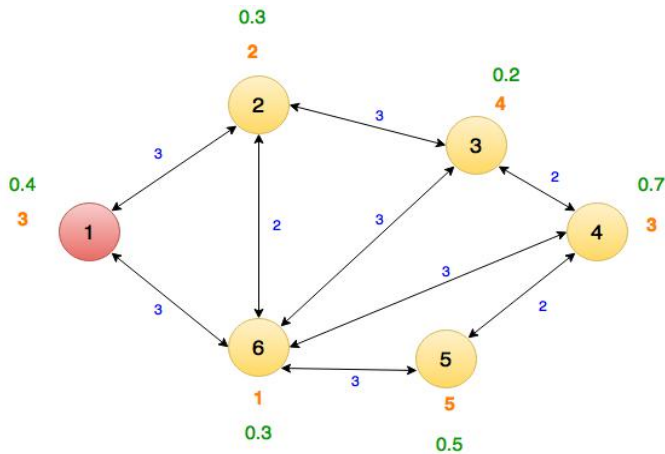
$$V = \{1, 2, 3, 4, 5, 6\}, \quad C = \{3, 2, 4, 3, 1, 5\}$$

$$r = \{0.4, 0.3, 0.2, 0.7, 0.5, 0.3\} \quad S = 3, M = 2, T = 25$$

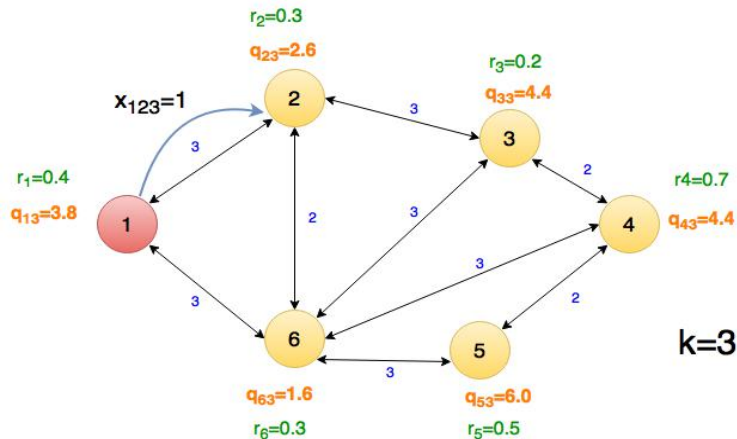
$$t = \begin{pmatrix} 0 & 3 & 0 & 0 & 0 & 3 \\ 3 & 0 & 3 & 0 & 0 & 2 \\ 0 & 3 & 0 & 2 & 0 & 3 \\ 0 & 0 & 2 & 0 & 2 & 3 \\ 0 & 0 & 0 & 2 & 0 & 3 \\ 3 & 2 & 3 & 3 & 3 & 0 \end{pmatrix} \quad d = \begin{pmatrix} 100 & 300 & 200 & 400 & 400 & 400 \\ 300 & 0 & 3 & 5 & 6 & 7 \\ 200 & 3 & 0 & 2 & 3 & 3 \\ 400 & 5 & 2 & 0 & 2 & 3 \\ 400 & 6 & 3 & 2 & 0 & 3 \\ 400 & 7 & 3 & 3 & 3 & 0 \end{pmatrix}$$

$$\alpha = \begin{pmatrix} 0 & 0.01 & 0.01 & 0.01 & 0.01 & 0.01 \\ 0.01 & 0 & 0.01 & 0.01 & 0.01 & 0.01 \\ 0.01 & 0.01 & 0 & 0.01 & 0.01 & 0.01 \\ 0.1 & 0.1 & 0.1 & 0 & 0.1 & 0.1 \\ 0.01 & 0.01 & 0.01 & 0.01 & 0 & 0.01 \\ 0.01 & 0.01 & 0.01 & 0.01 & 0.01 & 0 \end{pmatrix}$$

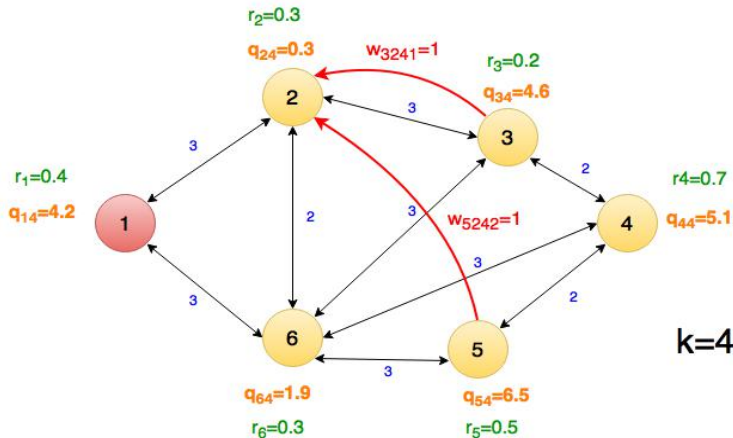
# Numerical Results



# Experiment

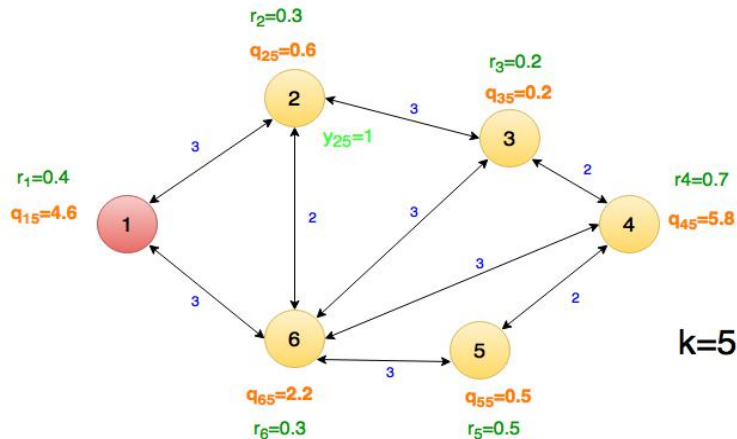


# Experiment

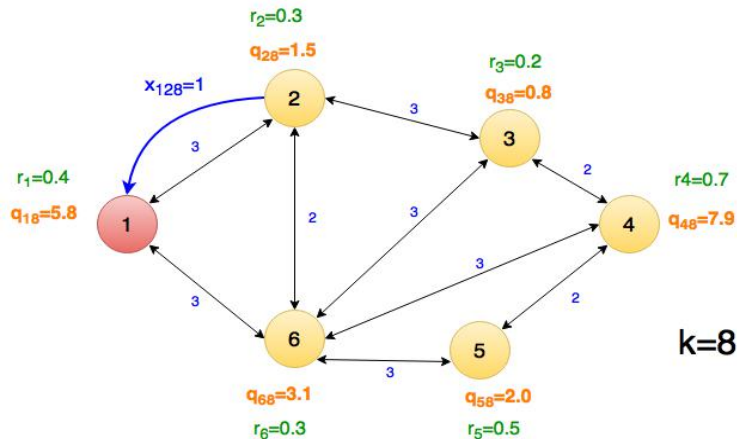




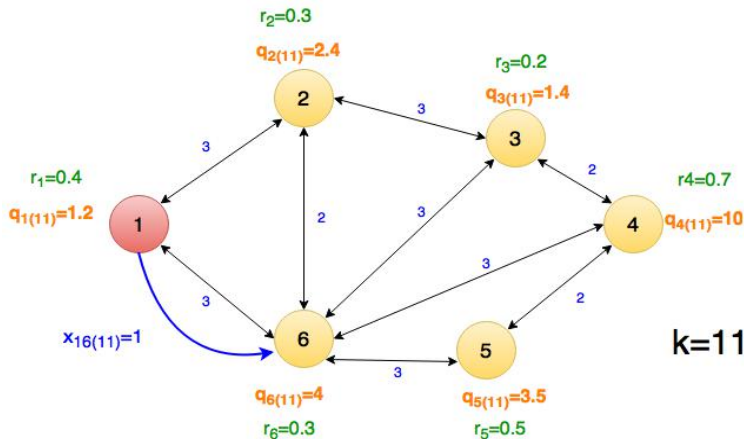
# Experiment



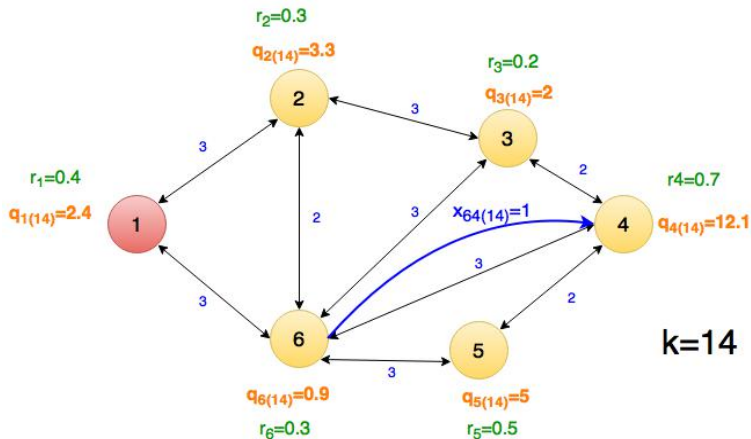
# Experiment



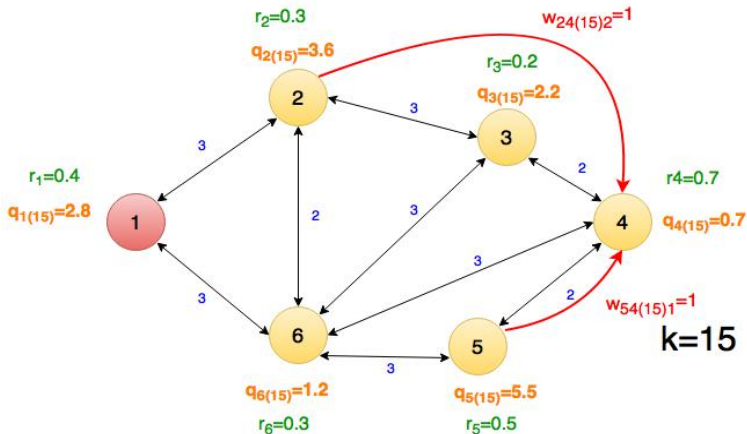
# Experiment



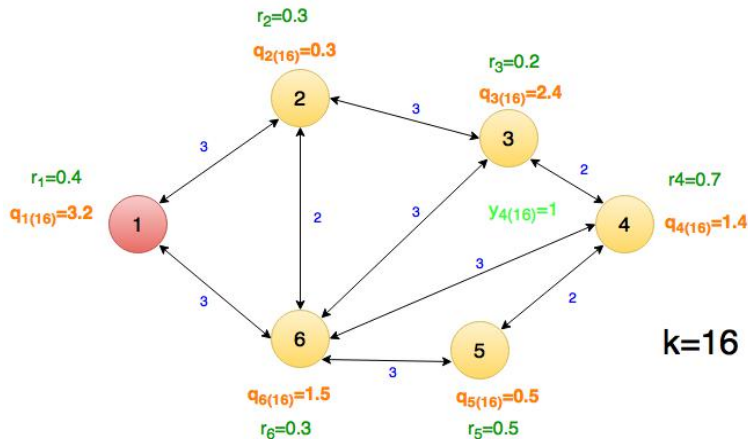
# Experiment



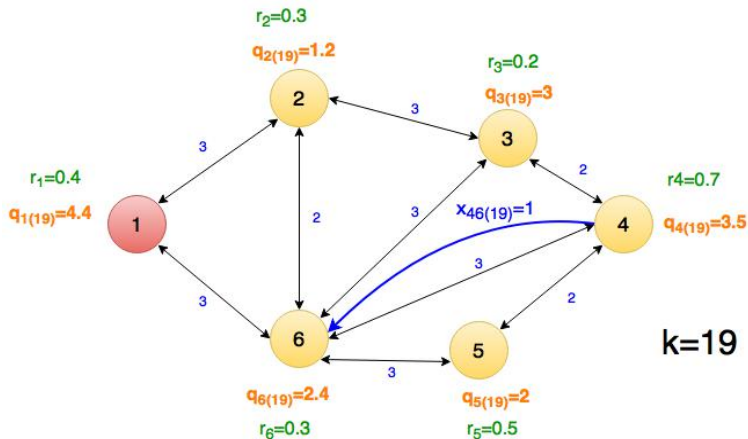
# Experiment



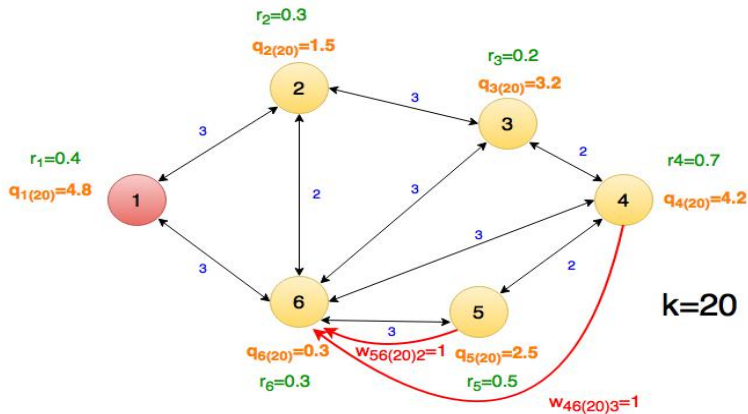
# Experiment



# Experiment

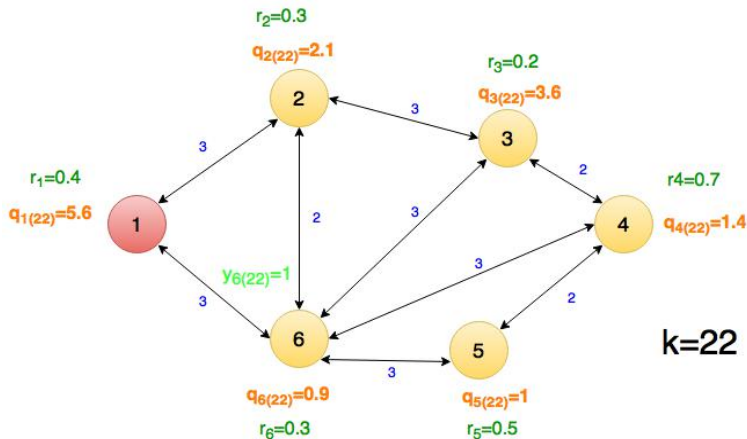


# Experiment

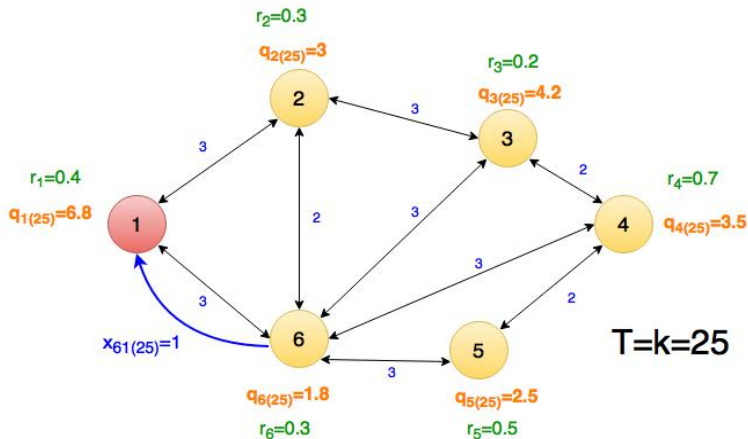




# Experiment



# Experiment



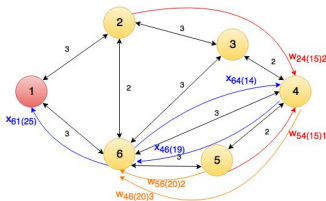
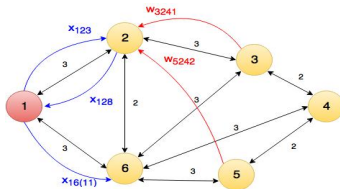
# Comparison of objective functions

$$\text{minimize} \left\{ \sum_{k=1}^T \max_{j \in V} \{q_{jk}\} \right\}$$

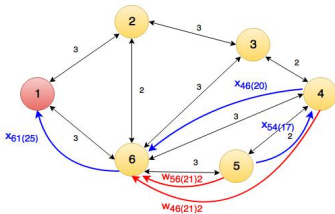
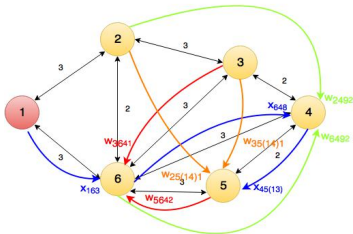
$$\text{maximize} \left\{ \sum_{k=1}^T \sum_{(i,j) \in A} x_{ijk} q_{jk} + \sum_{k=1}^T \sum_{l=1}^S \sum_{(j,i) \in V \times V} w_{jikl} q_{jk} \right\}$$

$$\text{minimize} \left\{ \max_{j,k} \{q_{jk}\} \right\}$$

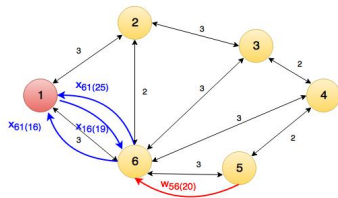
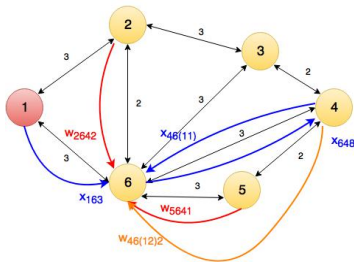
$$o.f1 : minimize \left\{ \sum_{k=1}^T \max_{j \in V} \{q_{jk}\} \right\}$$



$$o.f2 : maximize \left\{ \sum_{k=1}^T \sum_{(i,j) \in A} x_{ijk} q_{jk} + \sum_{k=1}^T \sum_{l=1}^S \sum_{(j,i) \in V \times V} w_{jikl} q_{jk} \right\}$$

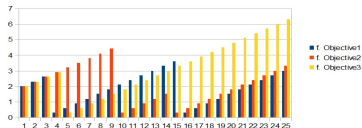


$$o.f3 : minimize \left\{ \max_{j,k} \{q_{jk}\} \right\}$$

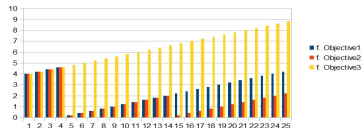


# Amount of information in the nodes

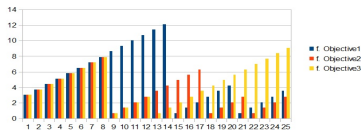
Node2



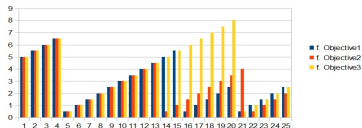
Node3



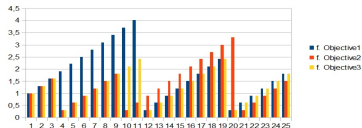
Node4



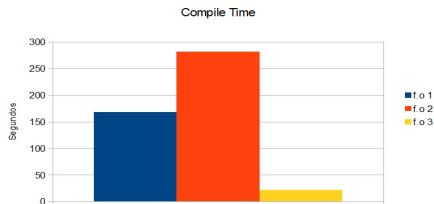
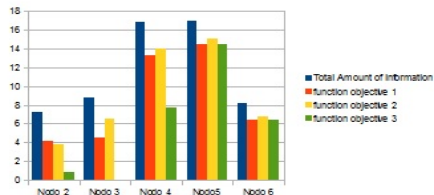
Node5



Node6



# Amount of information in the nodes





# Improvements

- Including constraints that allow the wireless transmission from nodes in a given neighborhood and not from primary nodes in the network.
- Compare the solutions obtained with other objective function.
- Eliminate some non necessary nodes in  $V'$ .

Thank you for  
your attention