

HW 1

Question 1)

a) Insertion Sort

Initial array: [9, 8, 2, 5, 7, 4] => key = 8
[9, 9, 2, 5, 7, 4] => Shifting until find smaller than 8 or beginning of array
[8, 9, 2, 5, 7, 4] => Insert 8; key = 2
[8, 8, 9, 5, 7, 4] => Shifting until find smaller than 2 or beginning of array
[2, 8, 9, 5, 7, 4] => Insert 2; key = 5
[2, 8, 8, 9, 7, 4] => Shifting until find smaller than 5 or beginning of array
[2, 5, 8, 9, 7, 4] => Insert 5; key = 7
[2, 5, 8, 8, 9, 4] => Shifting until find smaller than 7 or beginning of array
[2, 5, 7, 8, 9, 4] => Insert 7; key = 4
[2, 5, 5, 7, 8, 9] => Shifting until find smaller than 4 or beginning of array
[2, 4, 5, 7, 8, 9] => Insert 4; the end

b) Selection Sort

Initial array: [9, 8, 2, 5, 7, 4] => Find largest in unsorted part, which is 9
[4, 8, 2, 5, 7, 9] => Swap 9 with the end of the unsorted part, which is 4
[4, 8, 2, 5, 7, 9] => Find largest in unsorted part, which is 8
[4, 7, 2, 5, 8, 9] => Swap 8 with the end of the unsorted part, which is 7
[4, 7, 2, 5, 8, 9] => Find largest in unsorted part, which is 7
[4, 5, 2, 7, 8, 9] => Swap 7 with the end of the unsorted part, which is 5
[4, 5, 2, 7, 8, 9] => Find largest in unsorted part, which is 5
[4, 2, 5, 7, 8, 9] => Swap 5 with the end of the unsorted part, which is 2
[4, 2, 5, 7, 8, 9] => Find largest in unsorted part, which is 4
[2, 4, 5, 7, 8, 9] => Swap 4 with the end of the unsorted part, which is 2, the
end

c) Bubble Sort

Initial array: [9, 8, 2, 5, 7, 4] => start from index 0
[8, 9, 2, 5, 7, 4] => swap 9 with 8

[8, 2, 9, 5, 7, 4] => swap 9 with 2
 [8, 2, 5, 9, 7, 4] => swap 9 with 5
 [8, 2, 5, 7, 9, 4] => swap 9 with 7
 [8, 2, 5, 7, 4, 9] => swap 9 with 4
 [8, 2, 5, 7, 4, 9] => start from index 0
 [2, 8, 5, 7, 4, 9] => swap 8 with 2
 [2, 5, 8, 7, 4, 9] => swap 8 with 5
 [2, 5, 7, 8, 4, 9] => swap 8 with 7
 [2, 5, 7, 4, 8, 9] => swap 8 with 4
 [2, 5, 7, 4, 8, 9] => start from index 0
 [2, 5, 4, 7, 8, 9] => swap 7 with 4
 [2, 5, 4, 7, 8, 9] => start from index 0
 [2, 4, 5, 7, 8, 9] => swap 5 with 4
 [2, 4, 5, 7, 8, 9] => sorted is true, the end

d) Merge Sort

Initial array: [9, 8, 2, 5, 7, 4] => divide
 [9, 8, 2] - [5, 7, 4] => divide
 [9, 8] - [2] -- [5, 7] - [4] => divide
 [9] - [8] -- [2] ----- [5] - [7] -- [4] => start to merge
 [8, 9] - [2] -- [5, 7] - [4] => merge
 [2, 8, 9] - [4, 5, 7] => merge
 [2, 4, 5, 7, 8, 9] => copy the temp array back into original array; the end

Function Calls:

mergesort (arr, 0, 5)
 mergesort (arr, 0, 2)
 mergesort (arr, 0, 1)
 mergesort (arr, 0, 0)
 mergesort (arr, 1, 1)
 merge (arr, 0, 0, 1)
 mergesort (arr, 2, 2)

merge (arr, 0, 1, 2)
mergesort (arr, 3, 5)
mergesort (arr, 3, 4)
mergesort (arr, 3, 3)
mergesort (arr, 4, 4)
merge (arr, 3, 3, 4)
mergesort (arr, 5, 5)
merge (arr, 3, 4, 5)
merge (arr, 0, 2, 5)

e) Quick Sort

Initial array: [9, 8, 2, 5, 7, 4] => pivot = 4, put 4 into first place
[4, 8, 2, 5, 7, 9] => first unknown is 8, which is greater than 4, S2 starts with 8
[4, 8, 2, 5, 7, 9] => 2 is smaller so swap it to increase S1
[4, 2, 8, 5, 7, 9] => 5, 7 and 9 are bigger, no swap
[2, 4, 8, 5, 7, 9] => put 4 between S1 and S2
Sort S1, which has only 2 so it is already sorted
Sort S2
[8, 5, 7, 9] => pivot = 9, put 9 into first place
[9, 5, 7, 8] => first unknown is 5, which is smaller than 9, S1 starts with 5
[9, 5, 7, 8] => 7 and 8 are smaller, no swap
[8, 5, 7, 9] => put 9 between S1 and S2 (S2 has no elements)
Sort S1
[8, 5, 7] => pivot = 7, put 7 into first place
[7, 5, 8] => first unknown is 5, which is smaller than 7, S1 starts with 5
[7, 5, 8] => 8 is greater, no swap
[5, 7, 8] => put 7 between S1 and S2
Sort S1, which has only 5 so it is already sorted
Sort S2, which has only 8 so it is already sorted
[5, 7, 8, 9] => combining
[2, 4, 5, 7, 8, 9] => combining

Function Calls:

quicksort (array, 0, 5)

partition (array, 0, 5, ?)

quicksort (array, 0, 0)

quicksort (array, 2, 5)

partition (array, 2, 5, 1)

quicksort (array, 2, 4)

partition (array, 2, 4, 5)

quicksort (array, 2, 2)

quicksort (array, 5, 5)

quicksort (array, 6, 5)

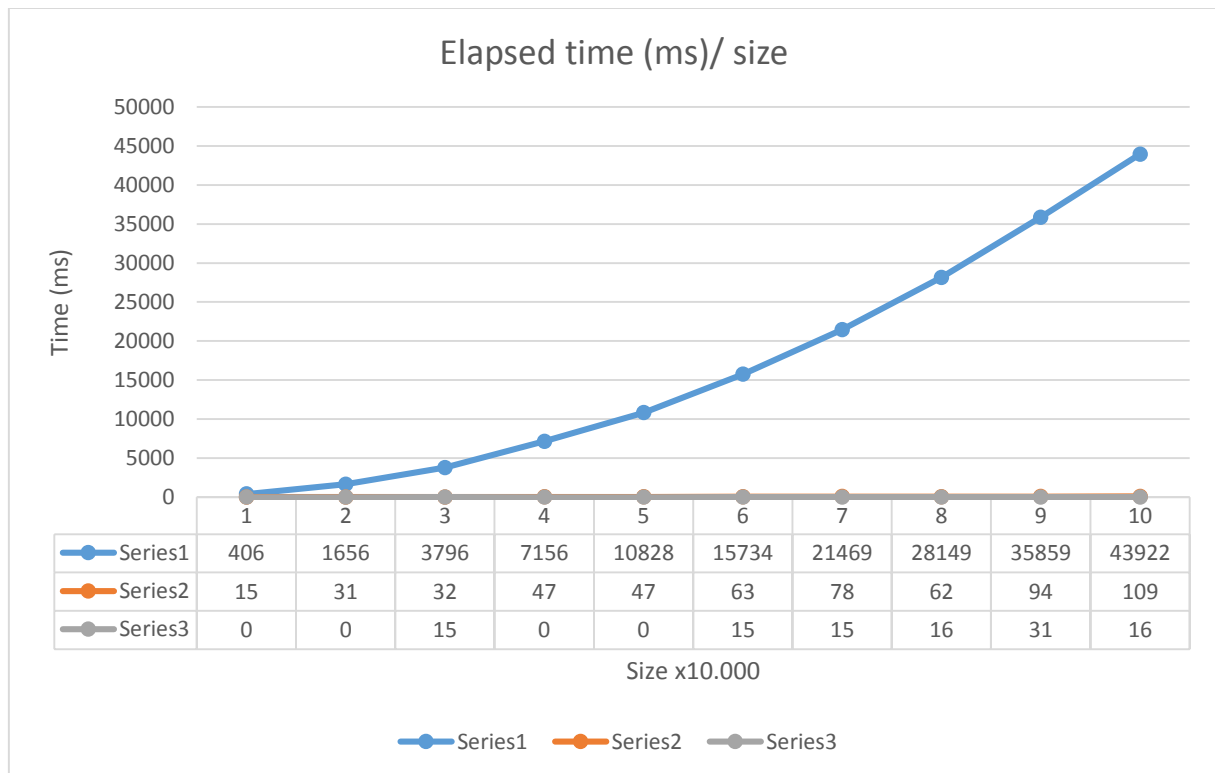
Question 2)

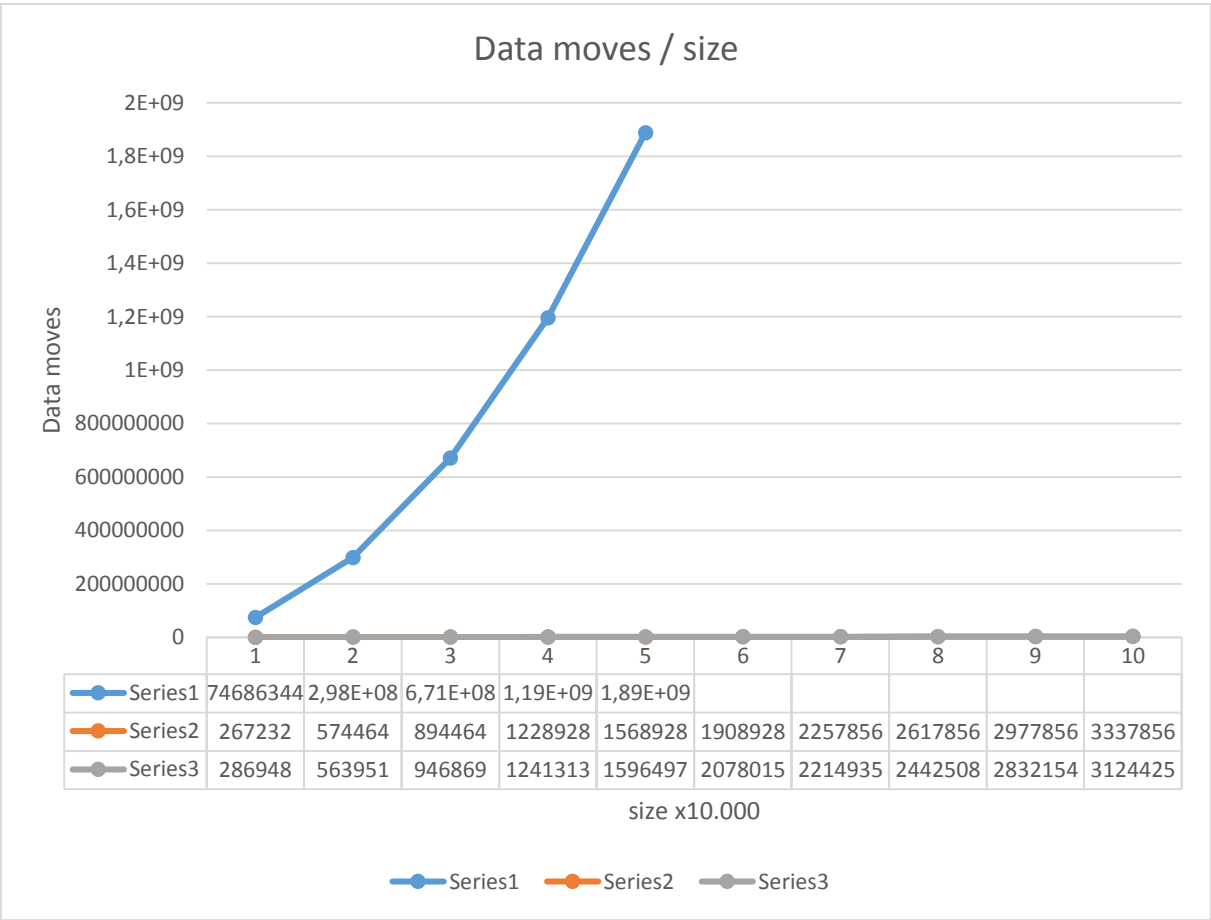
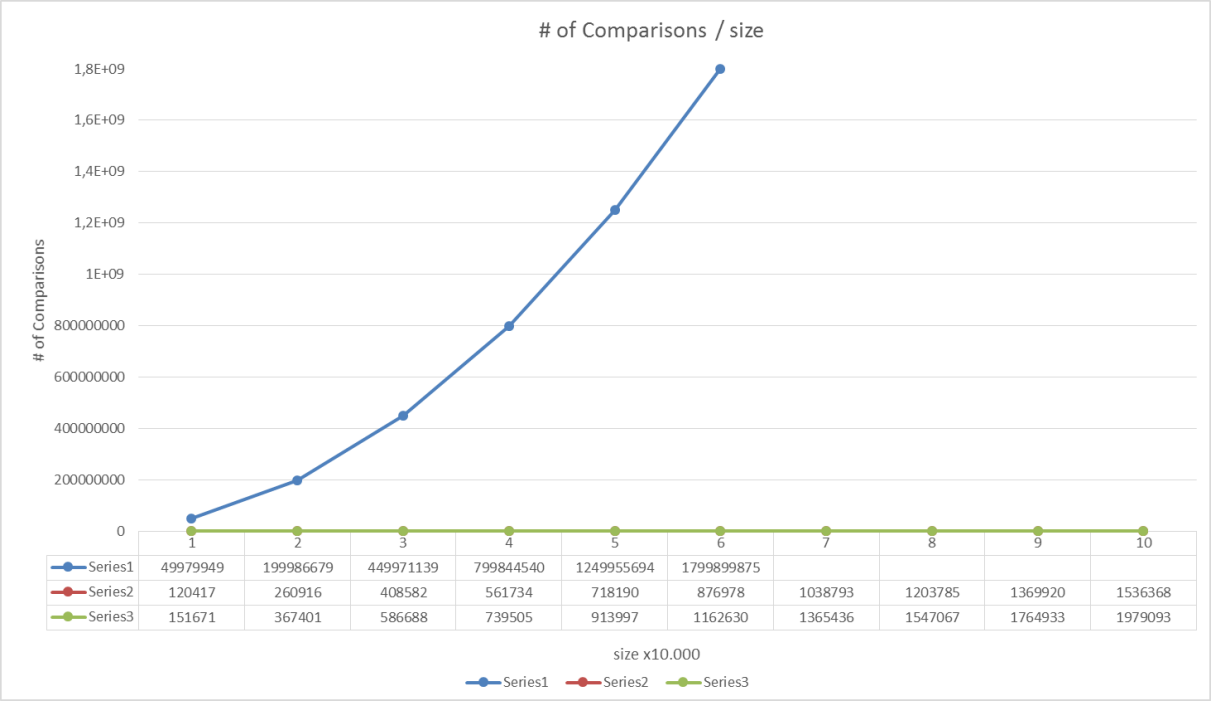
PART 1

Series 1: Bubble sort

Series 2: Merge Sort

Series 3: Quick sort



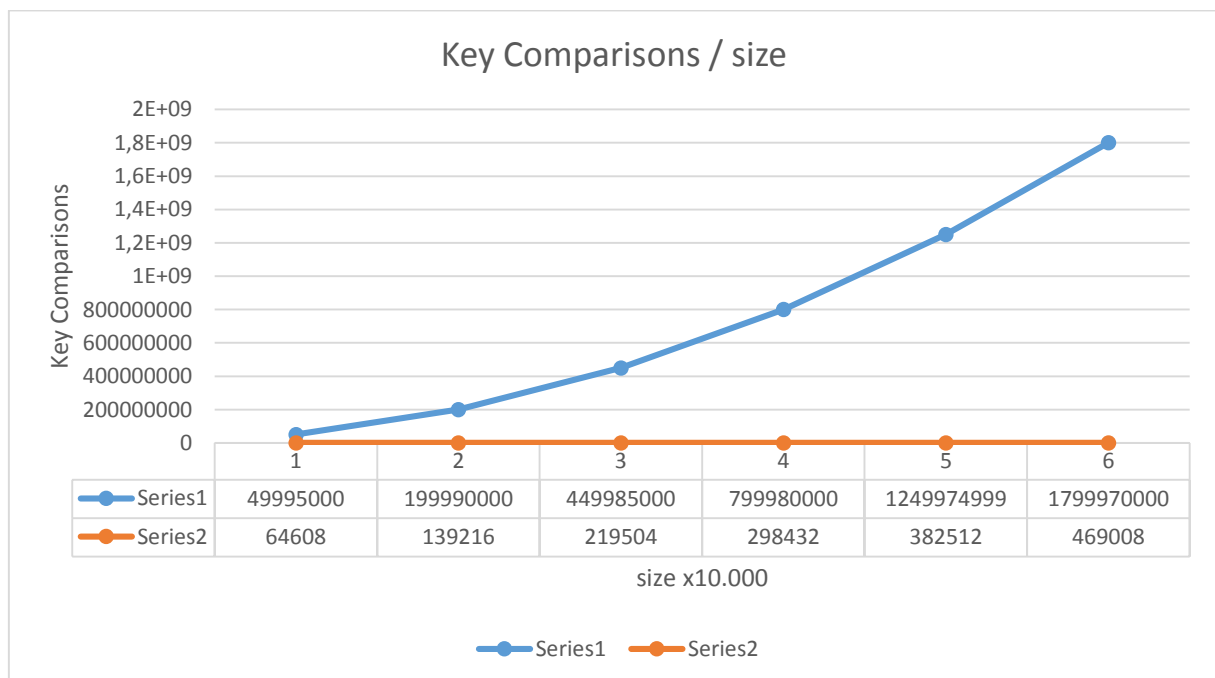
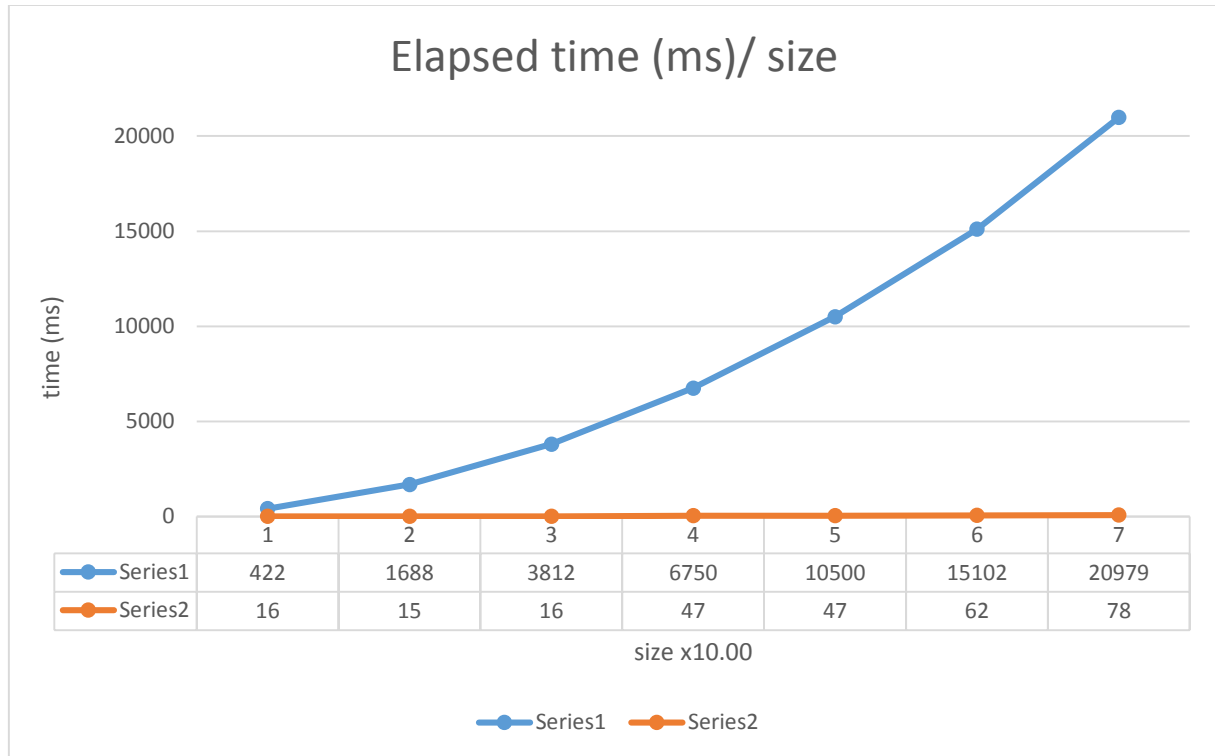


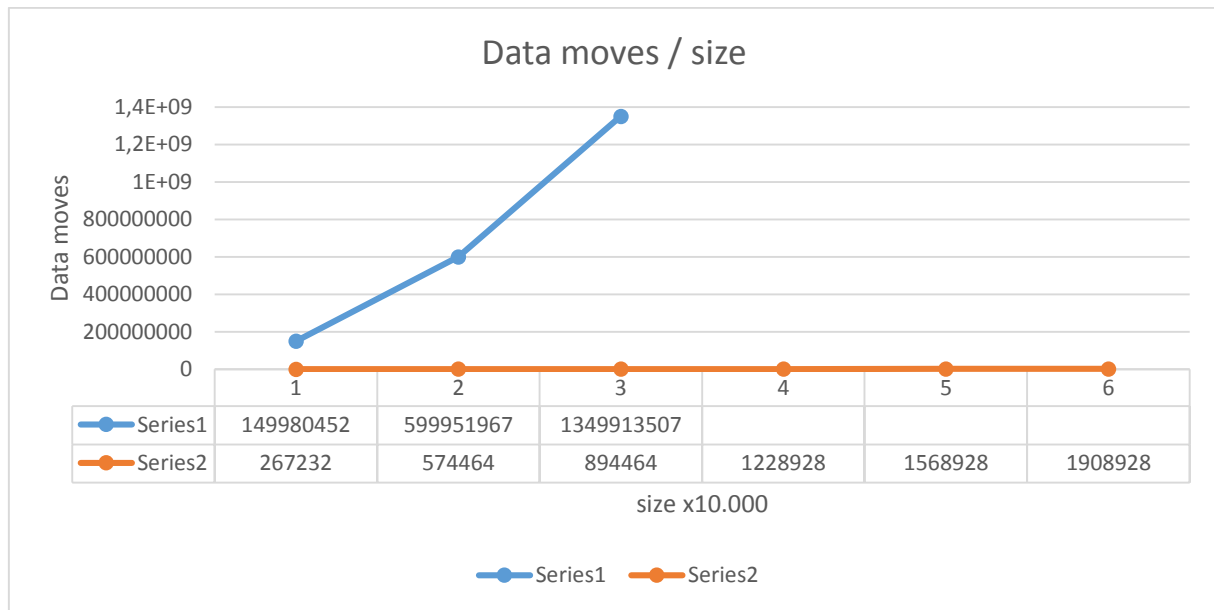
PART 2

Series 1: Bubble Sort

Series 2: Merge Sort

Note: For this part quick sort is not included in the graph since it fails after size 10000.





PART 3

Note: For this part quick sort is not included in the graph since it fails after size 10000.

Elapsed time for bubble sort is 0, number of comparisons are size - 1, and number of data moves are 0. It is best case for bubble sort, which is $O(n)$.

Number of data moves for merge sort is the same as part 1 and 2. Therefore there is no need to plot the graphs.

Question 3)

If algorithms use randomly generated numbers (average cases), they are closer to the theoretic results and quick sort seems to be best algorithm in those 3. For bubble sort worst (descending part) and average (random part) cases are $O(n^2)$ and best (ascending part) case is $O(n)$. For merge sort all situations are the same as we expect in theoretical, which is $O(n \log n)$. For quick sort, best and average (random) cases are $O(n \log n)$, however, worst (ascending and descending) is $O(n^2)$.

Part 2 and part 3, data moves for merge sort is the same for same sizes. (Actually, number of data moves for merge sort is always same for the same sizes in part 1, 2 and 3.) However, in these parts quick sort fails after size 10000. Ascending or descending orders are worst for quick sort, however for merge sort all of them is $O(n \log n)$. Even bubble sort is better than quick sort for these parts.

Part 3, elapsed time for bubble sort is 0, comparisons are size -1, and data moves are 0. It is best case for bubble sort, which is $O(n)$.

Question 4)

$$1) n! = 1.2.3.4 \dots n = [1.n][2(n-1)][3(n-2)] \dots \left[\frac{n}{2}.\frac{n}{2}\right] < \left(\frac{n^2}{4}\right)^{\frac{n}{2}} = \left(\frac{n}{2}\right)^n$$

$$\Rightarrow \log n! < \log \left(\frac{n}{2}\right)^n = n \cdot \log \left(\frac{n}{2}\right) = O(n \log n)$$

$$\therefore \log n! = O(n \log n)$$

2) It actually demonstrates the merge sort algorithm.

$$\begin{aligned} T(n) &= 2T\left(\frac{n}{2}\right) + O(n) = 2\left(2T\left(\frac{n}{4}\right) + O\left(\frac{n}{2}\right)\right) + O(n) = 4T\left(\frac{n}{4}\right) + 2O\left(\frac{n}{2}\right) + O(n) \\ &= 4T\left(\frac{n}{4}\right) + 2O(n) = 4\left(2T\left(\frac{n}{8}\right) + O\left(\frac{n}{4}\right)\right) + 2O(n) \\ &= 8T\left(\frac{n}{8}\right) + 4O\left(\frac{n}{4}\right) + 2O(n) = 8T\left(\frac{n}{8}\right) + 3O(n) = 2^k T\left(\frac{n}{2^k}\right) + kO(n) \end{aligned}$$

$$2^k = n \Rightarrow k = \log n$$

$$T(n) = nT(1) + \log n O(n) = O(n) + O(n) \log n = O(n \log n)$$

3) $f(n) = O(T(n)) \Rightarrow f(n) \leq c.T(n)$ when $n > n_0$

If we take $c = 8$ and $n_0 = 2$;

$$f(n) = 4n^5 + 3n^2 + 1 < 8.T(n) = 8.n^5 \text{ where } n > 2$$

The condition is satisfied.

$$\therefore f(n) = 4n^5 + 3n^2 + 1 = O(n^5)$$