



# LOADING BALANCE PROBLEM

Greedy loading balance strategy

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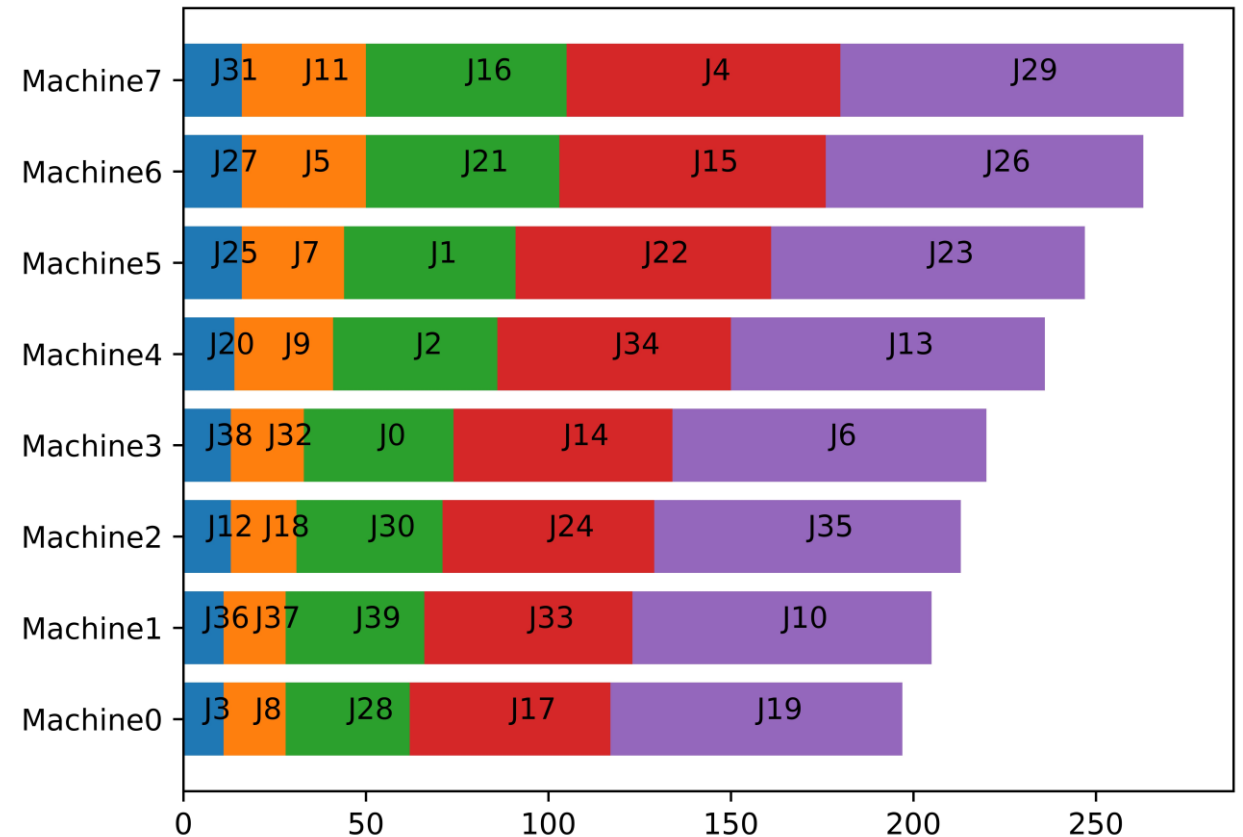


**SUSTech**

Southern University  
of Science and  
Technology

# LOADING BALANCE PROBLEM

This problem's target is assign jobs to computing machines to minimizes the total time consumption.



# GREEDY LOADING BALANCE STRATEGY

This strategy will assign a job to the machine with the smallest load in an arbitrary order of jobs.

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## Algorithm 2 Loading Balance Problem

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**Require:**  $m$  : number of machines,  $job\_times$  : The time of each jobs

**Ensure:**

```
1: function greedy_workload_balance( $m, job\_time$ )
2:    $M \leftarrow []$ 
3:   for  $i=1$  to  $m$  do
4:      $L_i \leftarrow 0$ 
5:      $M(i) \leftarrow []$ 
6:   end for
7:   for  $j=1$  to  $n$  do
8:      $i \leftarrow \operatorname{argmin}_k L_k$ 
9:      $M(i) \leftarrow M(i).append(j)$ 
10:     $L_i \leftarrow L_i + t_j$ 
11:   end for
12:   return  $M$ 
13: end function
```

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# EFFECT OF ALGORITHM

If the theoretical optimal makespan of the job queue is  $T^*$ , then the greedy algorithm's makespan  $T$  will not be worse than  $2T^*$ .

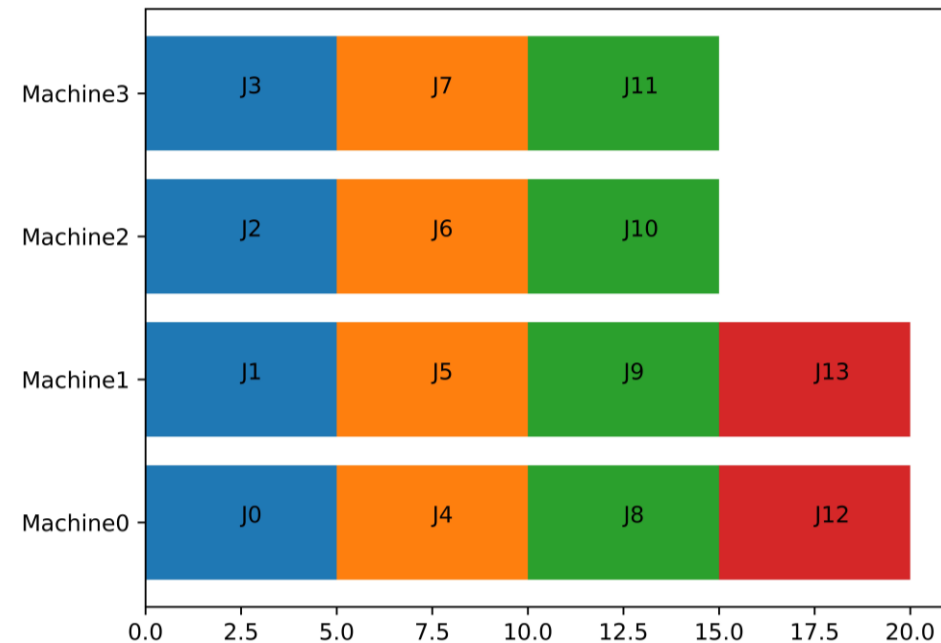
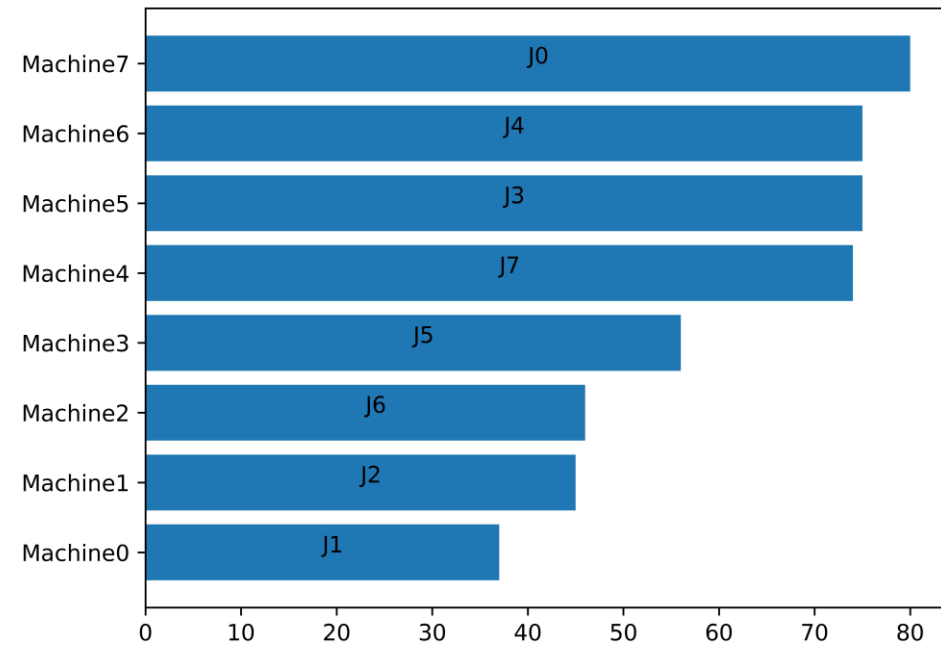
$$T^* \leq T \leq 2T^*$$

# THE FIRST EXAMPLE

You want an example that the makespan  $T$  equal to the  $T^*$ ? It's so easy!

What if the number of the jobs equal to the jobs of machine?

What if that each jobs need the time?



# THE SECOND AND THIRD EXAMPLE

These two questions can be answered by one case.

Let's suppose that there are  $m$  machines, and the minimum unit of time is 1.

What if we have  $m$  jobs with time  $m$ , and  $m$  jobs with time 1?



Let's consider two scenarios

1. Let the jobs queue be:  $[m, m, m, \dots, m, 1, 1, \dots, 1, 1]$
2. Let the jobs queue be:  $[m, 1, m, 1, m, 1, \dots, m, 1, m, 1]$



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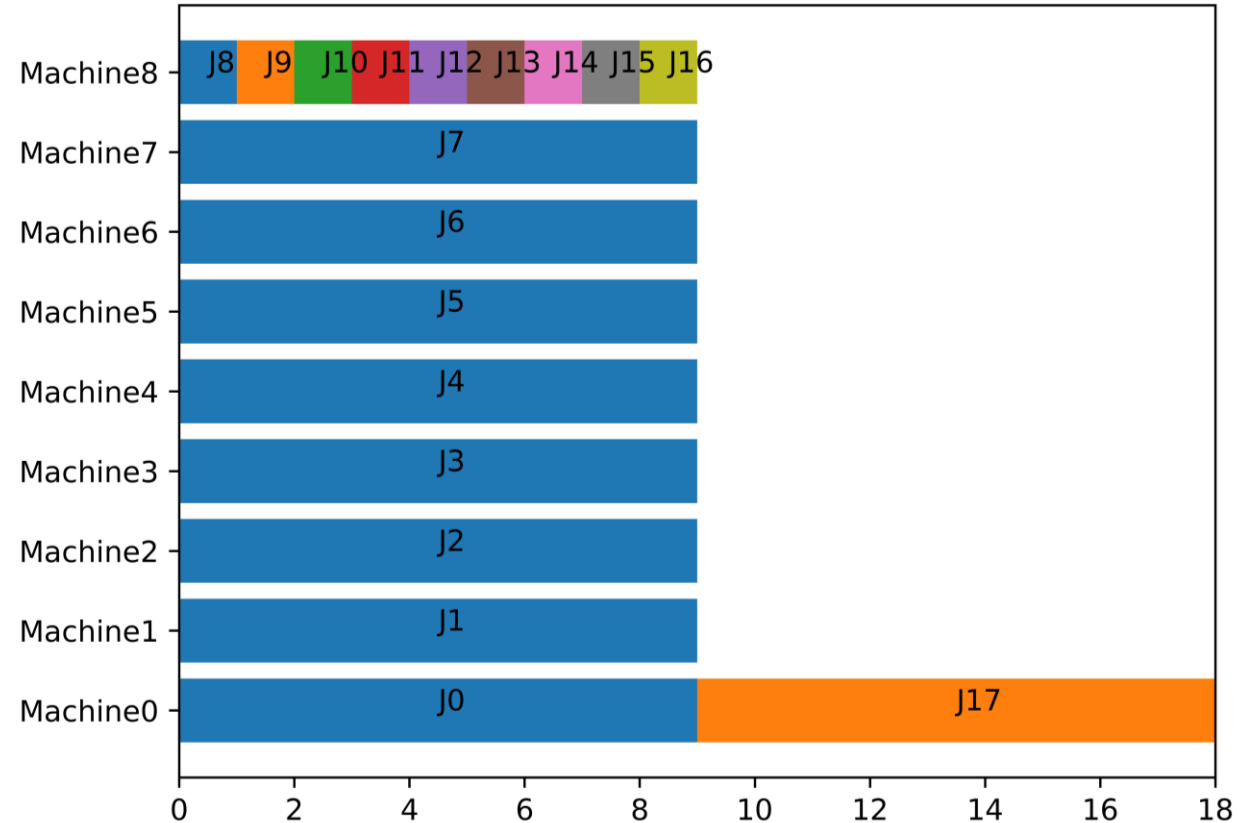


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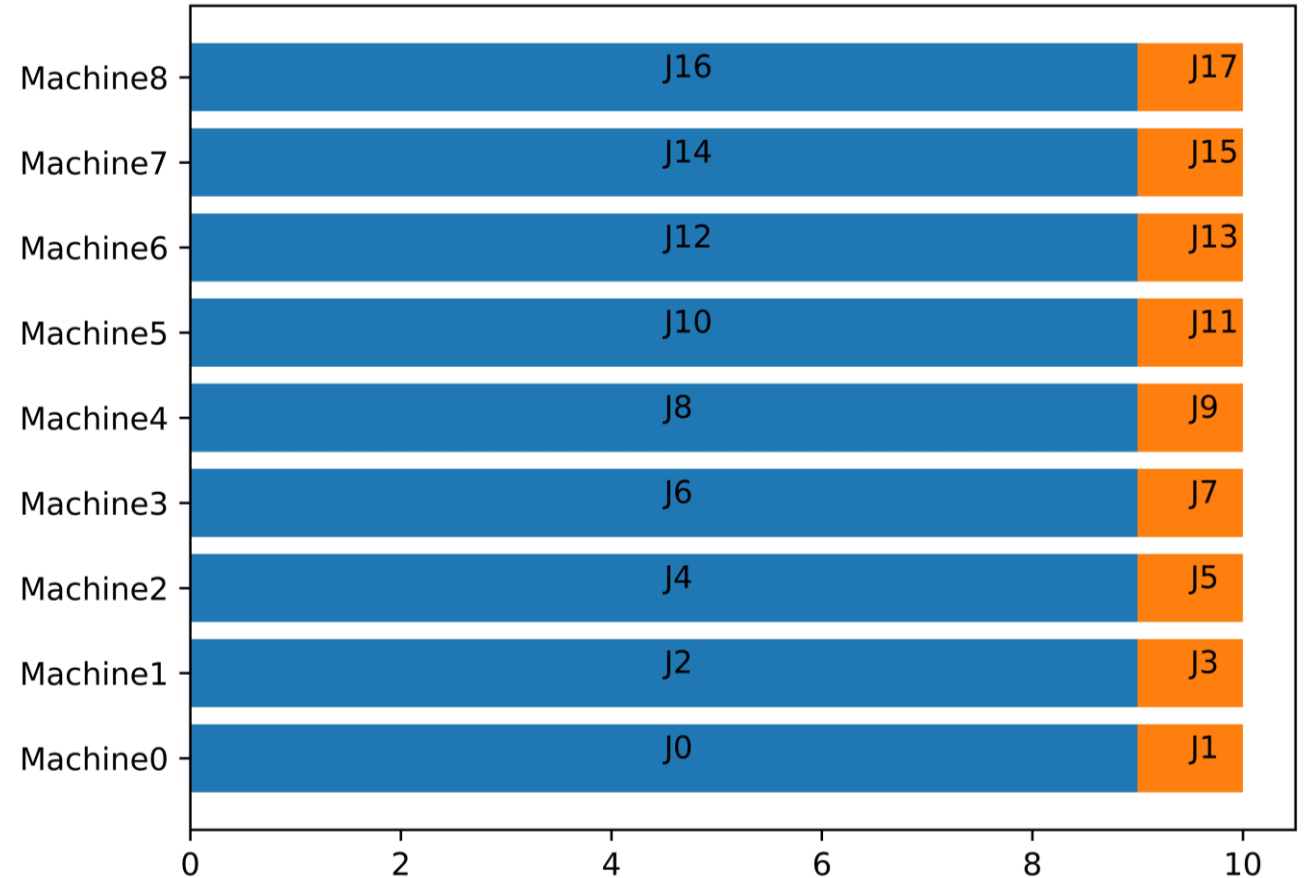


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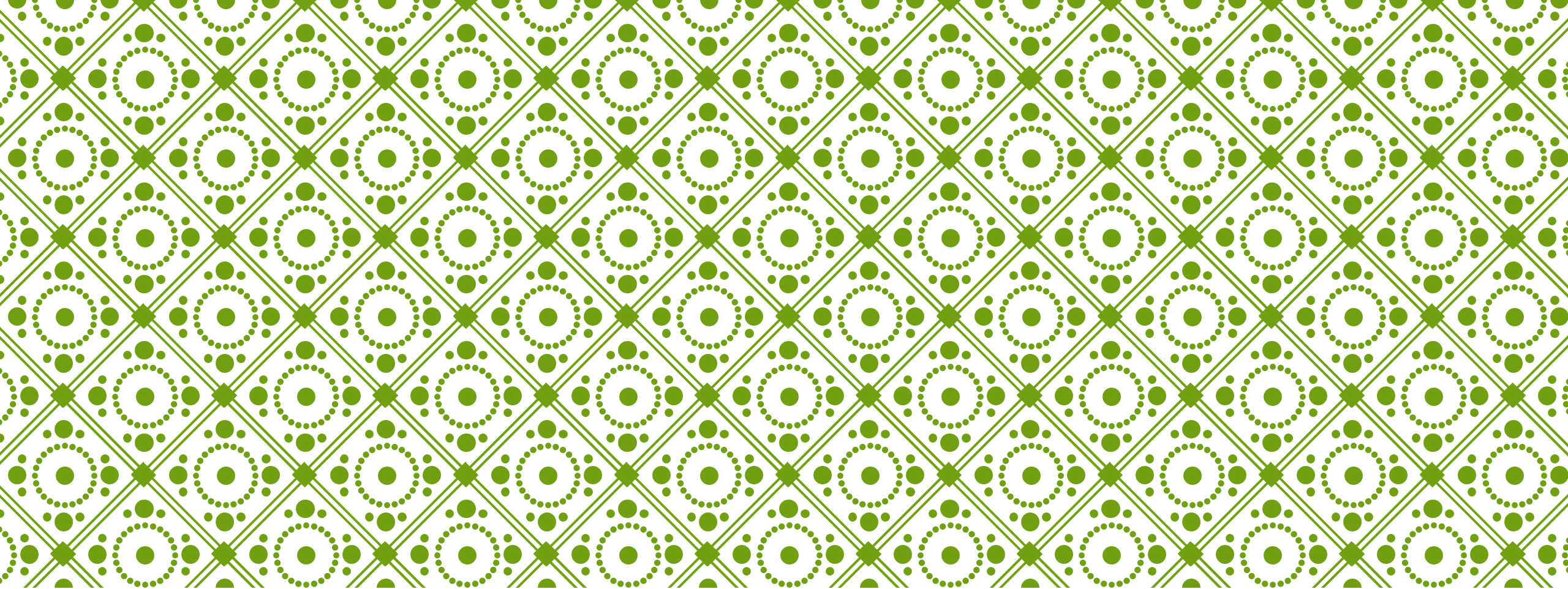
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For the two examples above:

1. When the job queue is  $[m, 1, m, 1, m, 1, \dots, m, 1, m, 1]$ , the result of the greedy algorithm is optimal,  $T=T^*=m+1$
2. When the job queue is  $[m, m, m, \dots, m, 1, 1, \dots, 1, 1, m]$ , it will get the worst solution of the greedy algorithm,  $T=2m$

So, for these jobs, the result can equal to  $T^*$ , and also can equal to  $2T^*-2$ , which is dependent on the order of the jobs.





# THANK YOU !

Q&A