Seeded Algorithm

Jiaxin Hu

March 28, 2022

This note combines previous results of error control and the non-iterative clean up. Theorem 1 is the theoretical guarantee for the complete seeded matching Algorithm 1.

Theorem 1 (Guarantee for Algorithm 1)**.** *Let $\rho = \sqrt{1-\sigma^2}$ and $s_0 = C(\log n^{1/4} + 1)^{1/(m-1)}$. Suppose $\sigma \leq c/s_0^{1/3}$ for sufficiently small constant c. Choose thresholds $\xi \geq c_1\sqrt{s_0}$ with universal positive constant c_1 and $\zeta \leq \sqrt{\sigma/n^{m-1}}$. Algorithm 1 recover the true permutation π^* with probability tends to 1.*

Remark 1 (Compared with previous results.)**.** Compared with the result in note 0306_22_proof, we relax the conditions from $\sigma < c/\log^{1/3(m-1)} n$ to $\sigma \leq c/\log^{1/3(m-1)} n^{1/4}$. The exponent 1/4 over n comes from the choice of $r_0 = \mathcal{O}(n^{3/4})$ in Theorem 2. The range of r_0 is determined by Theorem 3, however, there may exist some problem in Theorem 3. See the Fixme below.

Proof of Theorem 1. Based on Theorem 2 and Theorem 3, the output $\hat{\pi}$ of Sub-Algorithm 1 and Sub-Algorithm 2 fully recovers the true permutation if the number of seeds s satisfying $s^{m-1} \gtrsim \log n^{1/4} + 1$ and we take $r_0 = \mathcal{O}(n^{3/4})$.

[FIXME (Jiaxin): The Theorem 3 indicates we can choose $r_0 = \mathcal{O}(n - \sqrt{n}\log^{1/2(m-1)} n)$. However, if we take $r_0 \asymp n - n^{1/2+\epsilon}$, for some small $\epsilon \in (0, 1/2)$, we will have $s_0^{m-1} \gtrsim \log\left(\frac{1}{n^{1/2-\epsilon}+1} + 1\right) \asymp \log n^{\epsilon - 1/2}$. Then by Theorem 1, when n goes larger, we need fewer seeds and thus has a looser upper bound for σ, which is counter-intuitive and contradicts to the result in Ding et al. (2021). Therefore, I guess there may be some problems I did not recognize in Theorem 3.]

Hence, we only need to show the set

$$\mathcal{S} = \{(i,k) \in [n]^2 \colon a_i, b_k \geq \xi, d_p(\mu_i, \nu_k) \leq \zeta\},$$

where

$$a_i = \frac{1}{\sqrt{n^{m-1}}} \sum_{\omega \in [n]^{m-1}} \mathcal{A}_{i,\omega}, \quad b_k = \frac{1}{\sqrt{n^{m-1}}} \sum_{\omega \in [n]^{m-1}} \mathcal{B}_{k,\omega},$$

with proper thresholds ξ and ζ has enough true pairs and no fake pairs with high probability.

Note that for fake pair $(i,k) \in [n]^2$, i.e, $i \neq \pi^*(k)$, we have

$$\mathbb{P}(a_i \geq \xi, b_k \geq \xi) = \mathbb{P}(a_i \geq \xi)\mathbb{P}(b_k \geq \xi) = Q^2(\xi),$$

Algorithm 1 Gaussian tensor matching with seed improvement

Input: Gaussian tensors $\mathcal{A}, \mathcal{B} \in \mathbb{R}^{n^{\otimes m}}$, threshold ξ, ζ.

Calculate the distance statistics $d_p(\mu_i, \nu_k)$ for each pair of $(i,k) \in [n]^2$.

Obtain the high-degree set $\mathcal{S} = \{(i,k) \in [n]^2 \colon a_i, b_k \geq \xi, d_p(\mu_i, \nu_k) \leq \zeta\}$, where $a_i = \frac{1}{\sqrt{n^{m-1}}} \sum_{\omega \in [n]^{m-1}} \mathcal{A}_{i,\omega}, \quad b_k = \frac{1}{\sqrt{n^{m-1}}} \sum_{\omega \in [n]^{m-1}} \mathcal{B}_{k,\omega}$.

if there exists a permutation π_0 such that $\mathcal{S} = \{(i, \pi_0(i)) : i \in [n]\}$ **then**

    Run Sub-Algorithm 1 with seed π_0 and obtain output π_1.

    Run Sub-Algorithm 2 with π_1 and obtain output $\hat{\pi}$.

else

    Output error.

end if

Output: Estimated permutations $\hat{\pi}$ or error.

Sub-Algorithm 1: Seeded matching

Input: Gaussian tensors $\mathcal{A}, \mathcal{B} \in \mathbb{R}^{n^{\otimes m}}$, seed $\pi_0 : S \mapsto T$.

For $i \in S^c$ and $k \in T^c$, obtain the similarity matrix $H = [\![H_{ik}]\!]$ as $H_{ik} = \sum_{\omega \in S^{m-1}} \mathcal{A}_{i,\omega} \mathcal{B}_{k,\pi_0(\omega)}$.

Find the optimal bipartite permutation $\tilde{\pi}_1$ such that

$$\tilde{\pi}_1 = \underset{\pi : S^c \mapsto T^c}{\arg\max} \sum_{i \in S^c} H_{i,\pi(i)}.$$

Let π_1 denote the matching on $[n]$ such that $\pi_1|_S = \pi_0$ and $\pi_1|_{S^c} = \tilde{\pi}_1$.

Output: Estimated permutations $\hat{\pi}_1$.

Sub-Algorithm 2: Non-iterative clean-up

Input: Gaussian tensors $\mathcal{A}, \mathcal{B} \in \mathbb{R}^{n^{\otimes m}}$ and permutation $\pi_1 : [n] \mapsto [n]$.

For each pair $(i,k) \in [n]^2$, calculate $W_{ik} = \sum_{\omega \in [n]^{m-1}} \mathcal{A}_{i,\omega} \mathcal{B}_{k,\pi_1(\omega)}$.

Sort $\{W_{ik} : (i,k) \in [n]^2\}$ and let $\hat{S}$ denote the set of indices of largest n elements.

if there exists a permutation $\hat{\pi}$ such that $\hat{S} = \{(i, \hat{\pi}(i)) : i \in [n]\}$ **then**

    Output $\hat{\pi}$.

else

    Output error.

end if

Output: Estimated permutations $\hat{\pi}$ or error.

where Q is the complementary CDF of normal distribution. For true pair $(i,k) \in [n]^2$, i.e, $i = \pi^*(k)$, we have

$$\begin{aligned}
\mathbb{P}(a_i \geq \xi, b_k \geq \xi) &= \mathbb{P}(a_i \geq \xi, \sqrt{1-\sigma^2} a_i + \sigma z_i) \\
&\geq \mathbb{P}(a_i \geq \xi/\sqrt{1-\sigma^2}, z_i \geq 0) \\
&\geq \frac{1}{2} Q(\xi/\sqrt{1-\sigma^2}) \\
&\geq Q(\xi) \exp(-C\sigma^2\xi^2),
\end{aligned}$$

where C is a positive constant.

Take $\zeta \leq \sqrt{\sigma/n^{m-1}}$. To let $\mathcal{S}$ satisfy the conditions for in Theorem 2, we need

1. $\mathcal{S}$ has s true pairs with high probability (the expectation of the true pairs in $\mathcal{S}$ is larger than

s)

$$nQ(\xi)\exp(-C\sigma^2\xi^2) \geq s; \tag{1}$$

2. $\mathcal{S}$ has no fake pairs (the expectation of the fake pairs in $\mathcal{S}$ converges to 0 as $n \to \infty$)

$$n^2 Q^2(\xi) C_2 \exp\left(-\frac{1}{\sigma}\right) = o(1). \tag{2}$$

Take $\xi \geq c_1\sqrt{s}$. By inequality (1), we have $Q(\xi) \geq \frac{s}{n}\exp\left(Cc_1^2\sigma^2 s\right)$. Pluging the inequality for $Q(\xi)$ into the inequality (2), we have

$$C_2 s^2 \exp\left(2Cc_1^2\sigma^2 s - \frac{1}{\sigma}\right) = o(1),$$

which implies $\sigma \leq \frac{c}{s^{1/3}}$ with small constant c such that $2Cc_1^2 c^2 - \frac{1}{c^2} < 0$.

□

Useful Theorems and Lemmas for the proof of Theorem 1

Theorem 2 (Guarantee for Sub-Algorithm 1)**.** *Suppose the seed π_0 corresponds to s true pairs and no fake pairs. Assume $s^{m-1} \gtrsim \log n - \log(r_0+1) + 1$. The output π_1 of seeded matching Sub-Algorithm 1 has at most r_0 errors for $r_0 \in \mathbb{N} \cap [0, n-s]$.*

Proof of Theorem 2. See note 0323_22_seeded. □

Theorem 3 (Guarantee for Sub-Algorithm 2)**.** *Suppose the input permutation π_1 has at most r fake pairs such that $(n-r)^{(m-1)/2} \gtrsim n^{(m-1)/4}\log^{1/4} n + \log^{1/2} n$. Then, the output of non-iterative clean up Sub-Algorithm 2 is equal to the true permutation with a high probability; i.e., $\hat{\pi} = \pi^*$ with a high probability as $n \to \infty$.*

Proof of Theorem 3. See note 0321_22_cleanup. □

Lemma 1 (Tail bounds for the product of normal variables)**.** *Consider the correlated pairs of normal variables (X_i, Y_i) for $i \in [n]$, where $X_i, Y_i \sim N(0,1)$. Let $M = \frac{1}{n}\sum_{i\in[n]} X_i Y_i$. If $cov(X_i, Y_i) = \rho > 0$, then we have*

$$\mathbb{P}\left(|M - \rho| \geq t\right) \leq 4\exp\left(-\min\left\{\frac{1}{32\rho^2}, \frac{1}{16(1-\rho^2)}\right\} nt^2\right) \leq 4\exp\left(-\frac{nt^2}{32}\right),$$

for constant $t \in [0, \min\{2\rho, 2\sqrt{2}\sqrt{1-\rho^2}\}]$. If $cov(X_i, Y_i) = 0$, then, we have

$$\mathbb{P}\left(|M| \geq t\right) \leq 2\exp\left(-\frac{nt^2}{4}\right),$$

for constant $t \in [0, \sqrt{2}]$.

Proof of Lemma 1. See note 0306_22_proof. □

References

Ding, J., Ma, Z., Wu, Y., and Xu, J. (2021). Efficient random graph matching via degree profiles. *Probability Theory and Related Fields*, 179(1):29–115.