

Transaction Processing

An Introduction

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Intro

Fundamentals of Database Systems

Chapter 20

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Learning Objectives

- To understand what a Transaction is
- To understand what a Database Item is
- To understand Concurrency and Interleaving
- To understand how Database Items are managed by the DBMS
- To understand why we need concurrency control
- The lost update and dirty read problem

Transaction Processing

Any database system that requires high availability and fast response time for hundreds of concurrent users

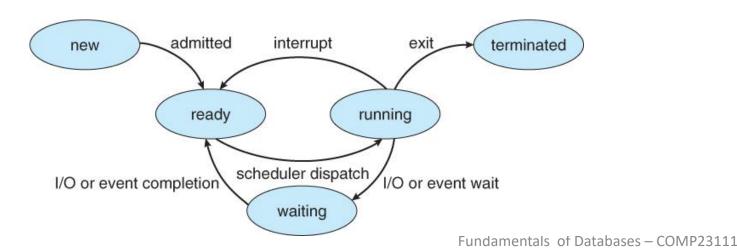
- A transaction is a mechanism for describing logical units of database processing
- Transaction processing systems are systems with large databases and hundreds of concurrent users
 - Airline Reservations
 - Banking Systems
 - Credit Card Processing
 - Online retail Purchases
 - Stock Markets
 - Supermarket Checkouts
 - and many more

Single User vs Multiuser Systems

- One criteria for classifying a database system is the number of users that can access the database concurrently
- A DBMS is single user if only one user at a time can use the system
- A DBMS is multiuser if the system is access by multiple users at the same time (concurrently)
- Concurrency is achieved using the concept of multiprogramming
 - Allows the operating system to execute multiple processes at the same time
 - But... A single core CPU can only execute at most one process at the same time, so how can concurrency be achieved?

Interleaving

- Process A runs, at some time later is switched out (or on I/O)
- Process B runs, at some point later is switched out (or on I/O)
- Process A is switched in and continues, at some point later is switched out (or on I/O)
- The process repeats for the number of processes
- So although we think these processes are running at the same time, they are actually interleaving



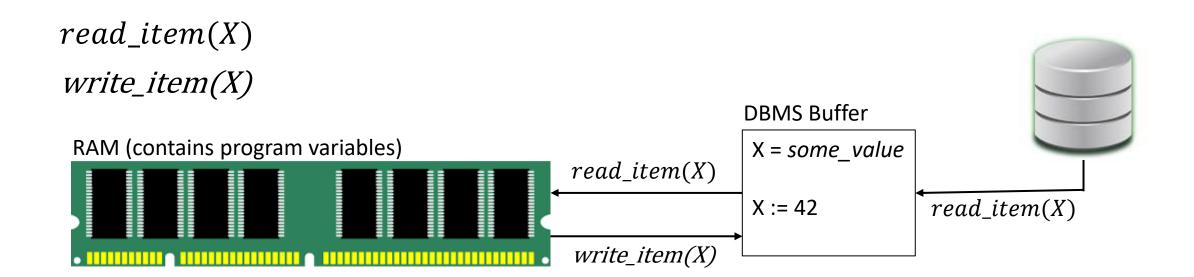
Transactions

- executing programs that forms a logical unit of database processing
- include one or more database access operations
- can be embedded within the application program or within the SQL server
- two classification of transactions

Database Items

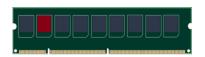
- The database model that is used to present transaction processing is relatively simple
- A database is seen as a collection of named data items
- There are various levels of granularity that could be represented

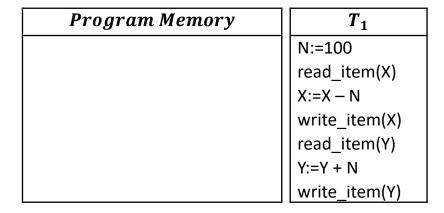
Read and Write Operations and DBMS Buffers

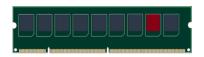


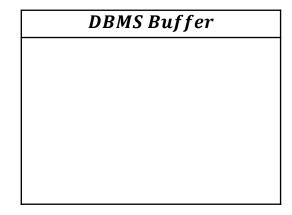
If the buffers are full and we need another buffer. The contents of a buffer will be swapped out to disk to make room. Normally using LRU.

A Simple Example with Context











Customer Accounts						
Current (main)	£100					
Savings	£1500					
Credit Card	-£1200					

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Concurrency

- The Lloyds Bank Group has 27,000,000 customers
- It is likely that several customers will access their account at the same time
- It would not be ideal to wait for one transaction to finish before another began
 - They would not have 27 million customers if we did
- So we allow concurrent access to the DBMS
- The control of this concurrent access are mainly concerned with Transactions
 - And recovery mechanisms when something goes wrong

A Simple Example with Context



Program Memory	T_1

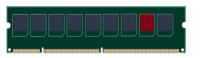
Program Memory	T_2

Program Memory	T_3

Program Memory	T_4

• • •

Program Memory	T_n
	•••



DBMS Buffer					



Supermarket Stock					
Beans	1065				
Milk	1887				
Apples 899					

Concurrency Control

 Several Problems can occur when concurrent transactions execute in an uncontrolled manner

- The Lost Update
- The Temporary Update (Dirty Read)
- The Incorrect Summary
- The Unrepeatable Read

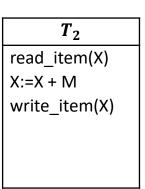
Airline Reservations Database

- A record is stored for each airline flight
 - Number of reserved seats
- When a program is written to interact with the database is uses flight number, flight date and the number of seats to be booked as parameters
 - The same program can be used to execute many different transactions each with a different
 - Flight number
 - Date
 - Number of seats to be booked

Airline Reservations Database

- T_1 shows a transfer of N reservations from one flight whose number of reserved seats is stored in X, to another flight whose number of reserved seats is stored in Y
- T_2 shows the reservation of M seats for the first flight
- These transactions are running concurrently

T_1
read_item(X)
X:=X - N
write_item(X)
read_item(Y)
Y:=Y + N
write_item(Y)



The Lost Update Problem



-	They	are	inter	leaved	to	allo)W	concu	ırrer	ncy
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- (Our values	, X	= 80,	N	= 5,	M =	= 4,	Y = 50	where
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- X and Y are the number of reservations for a given flight
- N is the number of seats to transfer to another flight
- M is the number of seats to reserve for a given flight

Question, what is the final value of X? 84

Question, should be the final value of X? 79

T_1	T_2
read_item(X)	
X:=X-N	
	read_item(X)
	X:=X + M
write_item(X)	
read_item(Y)	
	write_item(X)
Y:=Y + N	
write_item(Y)	

COMMITTED and ABORTED

- Whenever a transaction is submitted to a DBMS for execution, the system is responsible for making sure that all the operations in the transaction are completed successfully
 - COMMITTED
- If something goes wrong the database must not be updated
 - ABORTED
- A transaction is a logical unit of instructions, if the transaction fails, all changes for that transaction must be rolled back

The Temporary Update Problem

- Two transactions run at approximately the same time
- They are interleaved to allow concurrency
- Our values, X = 80, N = 5, M = 4, Y = 50
 - X is the number of reservations for a given flight
 - N is the number of seats to transfer to another flight
 - M is the number of seats to reserve for a given flight

Question, what is the final value of X? 80

Question, should be the final value of X? 79

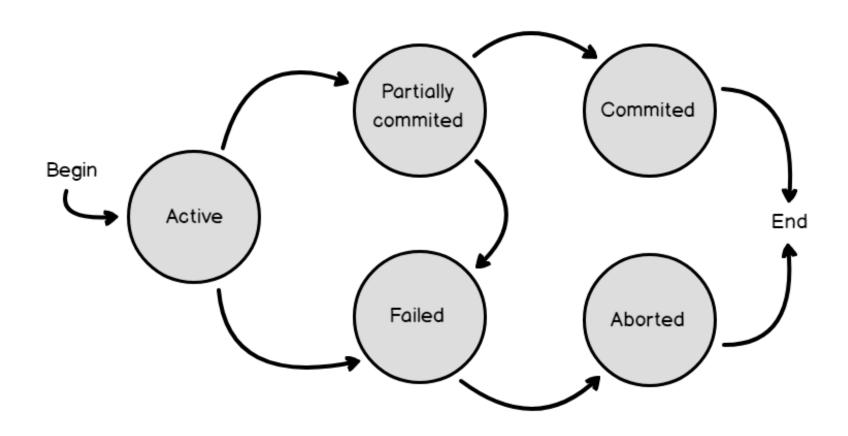
T_1	T_2
read_item(X)	
X:=X - N	
write_item(X)	
read_item(Y) Transaction fails	read_item(X) X:=X + M write_item(X)

where

What could possibly go wrong?

- A Computer Failure (System Crash)
- Transaction Error
- Logical Errors or Exception Conditions detected by the Transaction
- Disk Failure
- Physical Problems and Catastrophes

Transaction States and Additional Operations



We have learnt

- what a Transaction is
- what a Database Item is
- what meant by Concurrency and Interleaving
- an abstract understand how Database Items are managed
- why we need concurrency control
- The lost update and dirty read problem



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