

COMP26120 Algorithms and Data Structures
Topic 2: Data Structures

Data Structures vs Data Types

Dr. Thomas Carroll
thomas.carroll@manchester.ac.uk
All information on Blackboard

Learning Outcomes

- Recall what an Abstract Data Type (ADT) is
- Understand what a Data Structure is
- Explain the difference between ADT and Data Structure
- Understand how common ADTs can be implemented with Linked Lists and Arrays

Abstract Data Type

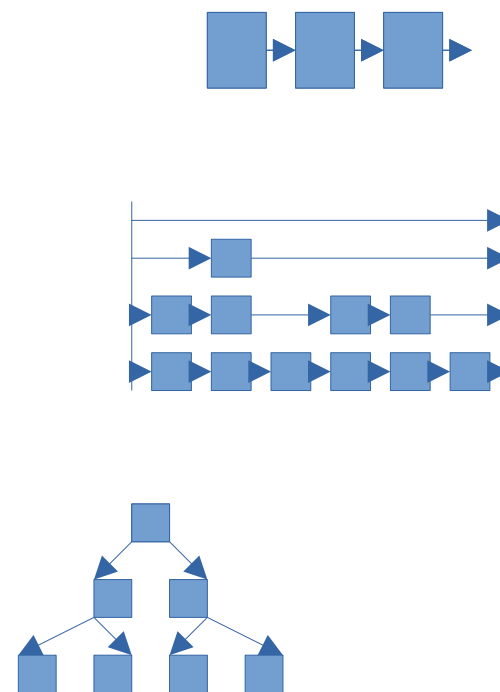
- A theoretical concept which, for some given *data type* defines:
 - The **possible values** it can take
 - The **operations** that can be performed
 - The **behaviour** of these operations
- A **list** might have the following operations:
 - add Adds an element to the end of the list
 - head Returns the first element of the list
 - tail Returns the sublist *after* the head

Data Structure

- A **data structure** is a way to **store and organise** data to facilitate access and modifications
- Data structures **implement** ADTs
- Eg:
 - Linked Lists and Dynamic Arrays can implement the List ADT



Implemented By



Now...

- See how...
 - Queue
 - Stack
 - Set
 - Dictionary
- Can be implemented by...
 - Dynamic Array
 - Linked List

But First...

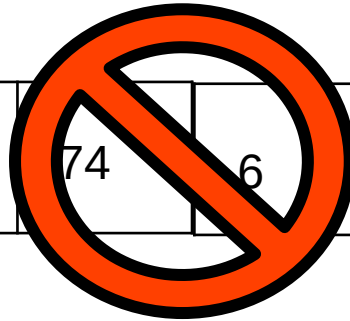
- What is a Linked List?
- What is a Dynamic Array?

Dynamic Array

Let's insert the following k numbers to the back of the array:
5,89,41,57,74,6,9

```
for i = 0 to k-1  
  A[i] = next number
```

5	89	41	57	74	6	9			
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Dynamic Array

init:

```
allocate N cells as A
len = N;
size = 0;
```

get(i):

```
if i > size then
    raise Error
else
    return A[i]
```

add(v):

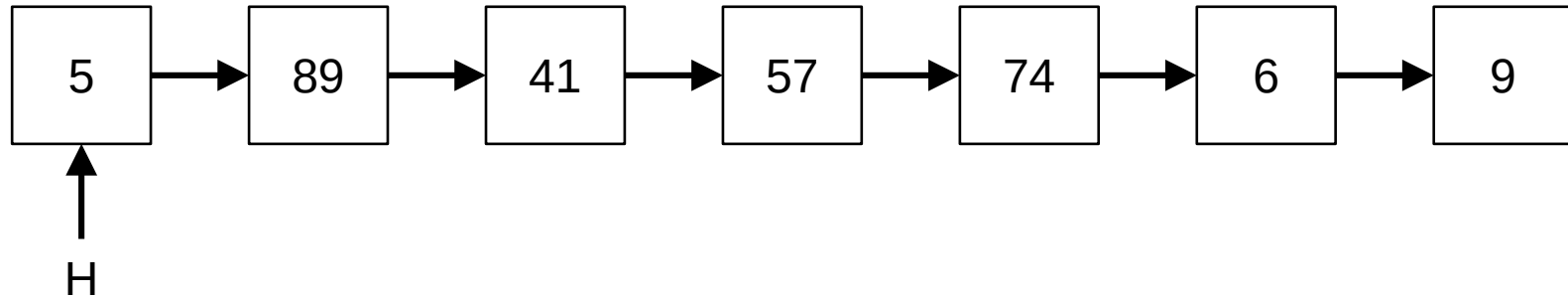
```
if size = len-1 then
    allocate 2N cells as B
    for i = 0 to N: B[i] = A[i]
    A=B
A[size] = v
size = size +1
```

set(i,v):

```
if i > size then
    raise Error
else
    A[i] = v
```

Linked List

Let's insert the following k numbers to the linked list:
5,89,41,57,74,6,9



Linked List

```
init:  
    head = new node;
```

```
addFront(v):  
    n = new node; n.value = v;  
    n.next = head; head = n;
```

```
addBack(v):  
    n = new node; n.value = v;  
    n.next = NULL;  
    if head == NULL: head = n;  
    else:  
        curr = head;  
        while curr.next != NULL:  
            curr = curr.next  
  
        curr.next = n;
```

Dynamic Arrays vs Linked Lists

- Memory footprint:
 - Dynamic arrays have $O(n)$ wasted space (plus the expensive copy!)
 - Linked Lists have space overhead for pointer
- Access:
 - Dynamic arrays have $O(1)$ access
 - Linked Lists have $O(n)$ access

Queue

- First In, First Out (FIFO)

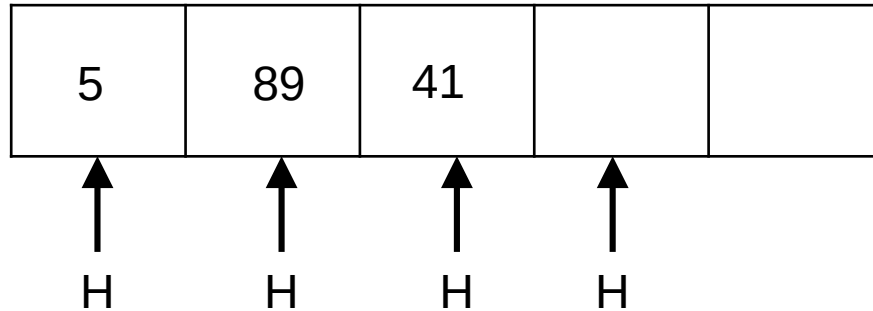


Stack

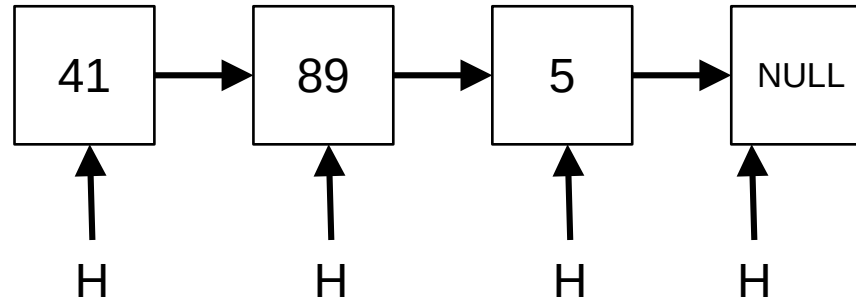
- First In, Last Out (FILO)



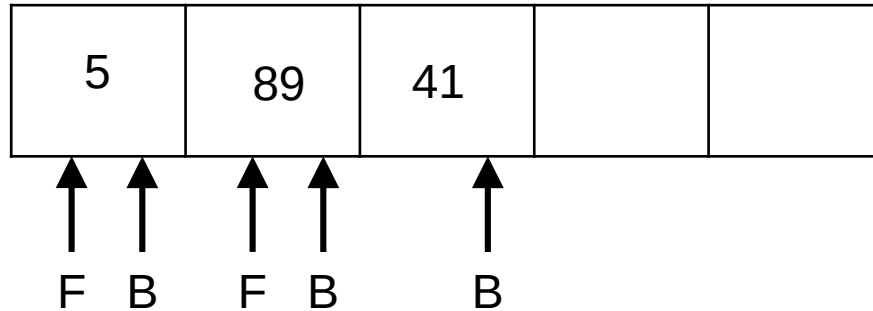
Stack with (Dynamic) Array



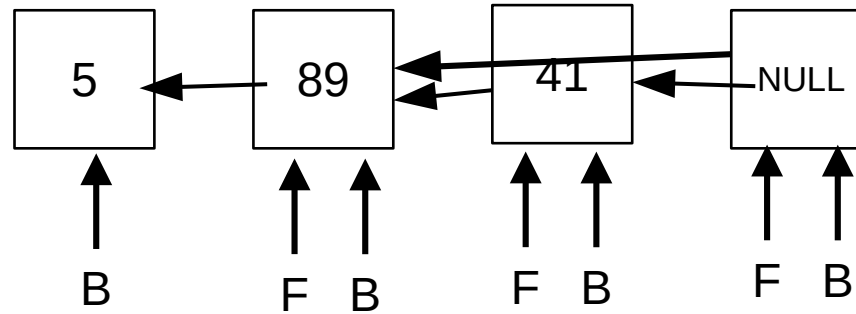
Stack with Linked List



Queue with (Dynamic) Array



Queue with Linked List



Sets and Dictionaries

Set – *an unordered collection*

`add(o)` : Adds `o` to the set

`remove(o)` : Removes `o` from the set

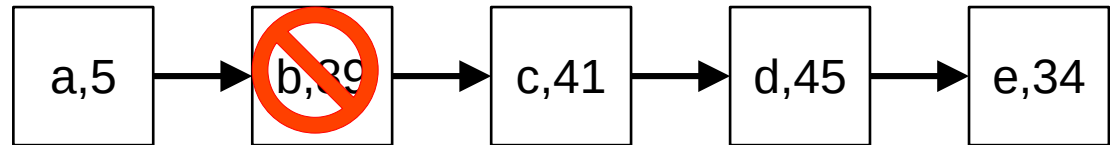
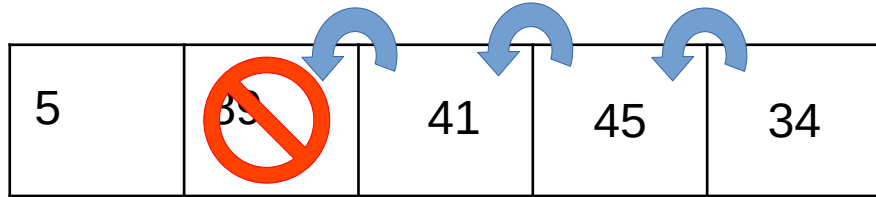
`find(o)` : Returns `TRUE` if `o` is in the set

Dictionary – generalisation of arrays to non-int indices

`insert(k,v)`: Inserts value `v` at `k`

`remove(k)`: removes entry at `k`

`find(k)`: returns value of that at `k`



`remove(89) / remove(b)`

Priority Queues and Disjoint Sets

Priority Queue – queue that is ordered by “priority”

`insert(k,e)`: Insert element `e` with key `k`

`RemoveMin()`: Remove and return element with smallest key (maybe not unique)

Disjoint Set – Collection of disjoint sets that can be merged

`add(x)` : Create new set containing `x`

`find(x)` : find the set containing `x`

`union(x,y)`: merge sets of `x` and `y`