



THE UNIVERSITY OF QUEENSLAND  
A U S T R A L I A

OPTIMIZING PERFORMANCE  
IN GAUSSIAN PROCESSES

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## INTRODUCTION

**Origins of group theory.** The mathematical field of group theory has its origins in the early 19th century. At the time, mathematicians were investigating the solutions to polynomial equations. That is, solutions to equations of the form

$$a_n x^n + a_{n-1} x^{n-1} + \dots + a_0 = 0.$$

Full solutions to polynomial equations of low degrees (i.e.  $n \leq 4$ ) had already been formulated [Rig96]. These include the familiar *quadratic formula*, which has been known since antiquity. The formula tells us that the solutions to a general quadratic equation  $ax^2 + bx + c = 0$  are given by

$$x = \frac{-b \pm \sqrt{b^2 - 4ac}}{2a}.$$

The full solutions to any cubic ( $n = 3$ ) or quartic ( $n = 4$ ) polynomial equation were also known. These are given by the lesser-known Cardano's formula and Ferrari's method, respectively. We say that a polynomial is *solvable by radicals* if one can write all of its solutions in terms of its coefficients combined with the algebraic operations; addition, subtraction, multiplication, division, powers and radicals (i.e.  $k^{\text{th}}$  roots).

In the 1830s, the mathematician Évariste Galois provided an elegant method to prove that a general polynomial of degree  $n \geq 5$  is not solvable by radicals. Galois understood that to every polynomial one could associate a *Galois group*, a new mathematical object at the time. The Galois group was the first object in a class of mathematical objects that we call *groups* today. We say that the pair  $(G, \circ)$  is a group, where  $G$  is a set and  $\circ: G \times G \rightarrow G$  is a binary operation on  $G$ , when three conditions are satisfied:

- Associativity:  $g \circ (h \circ k) = (g \circ h) \circ k$  for all  $g, h, k \in G$ .
- Existence of an identity: there exists some  $1_G \in G$  such that  $1_G \circ g = g \circ 1_G = g$  for all  $g \in G$ .
- Existence of inverses: for every  $g \in G$ , there exists some  $g^{-1} \in G$  such that  $g \circ g^{-1} = g^{-1} \circ g = 1_G$ .

Examples of groups that are likely familiar to the reader include  $(\mathbb{Z}, +)$ , the integers under addition,  $(\mathbb{R}^+, \times)$ , the positive real numbers under multiplication, and  $(\mathbb{Z}/n\mathbb{Z}, +)$ , the integers modulo  $n$  under addition. Some geometric examples of groups are the *dihedral groups*. These groups are generated by the  $m$  symmetries associated to the regular  $m$ -sided polygon (i.e. a polygon with all interior angles and all side lengths the same). Then each dihedral group contains  $2m$  elements ( $m$  reflections and  $m$  rotations) with the group operation of composition of reflections and rotations.

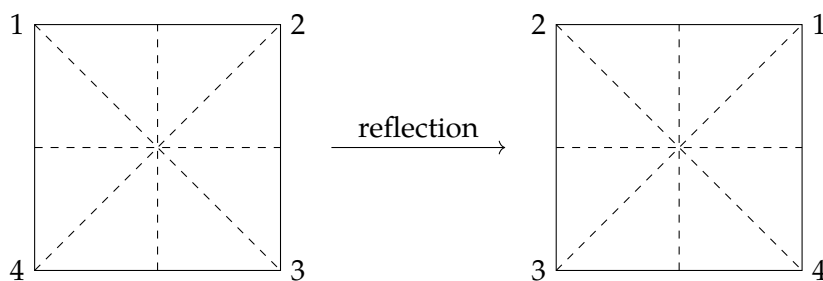


FIGURE 1. The symmetries of a square and a reflection about the vertical line of symmetry.

This gives us an intuitive understanding of groups: they encode the symmetries of mathematical objects.

**What is representation theory?** The study of groups yields insight into geometric objects. The action of the dihedral group on the  $m$ -gon serves as example of a group acting on a geometric object. More generally, we can consider the action of a group on some object. Specifically, we say that a group  $G$  *acts* on a set  $X$  if, for each  $g \in G$ , there is a map  $\cdot : G \times X \rightarrow X$  satisfying  $1_G \cdot x = x$  and  $g \cdot (h \cdot x) = (gh) \cdot x$  for all  $x \in X$ . Alternatively, one can view this as a group homomorphism  $\rho : G \rightarrow \text{Sym}(X)$ , where  $\text{Sym}(X)$  is the symmetric group associated to  $X$ , i.e. the group of permutations of elements of  $X$ .

Now we linearise the setting above by requiring that  $X = V$  is a *vector space*. Then we say that  $G$  acts *linearly* on  $V$  if there exists a group homomorphism  $\rho : G \rightarrow \text{GL}(V)$ . We call  $(V, \rho)$  a *representation* of  $G$ , and  $\rho$  is often suppressed from notation. We see that  $G$  acts on  $V$  in the sense that  $\rho(g) : V \rightarrow V$  is a linear invertible map on  $V$ . We may denote  $\rho(g)(v)$  by  $g \cdot v$  as before.

Representation theory is concerned with understanding and classifying linear actions of groups. The general situation of representation theory is as follows. If the group  $G$  acts on a vector space  $V$ , then we say that a vector subspace  $W \subseteq V$  is a *subrepresentation* of  $V$  if it is invariant under the action of  $G$ . A representation is called *irreducible* if its only proper subrepresentation is the trivial representation  $W = \{0\}$ . The primary goals of representation theory are finding all irreducible representations of  $G$ , and to decompose a given representation into its irreducible components.

We can think of irreducible representations as the building blocks of all other representations. This is a common idea in mathematics, seen in other areas. For instance, in number theory, the building blocks of integers are primes and, in group theory, the building blocks of groups are simple groups.

Writing a general representation in terms of irreducible components is not always possible. We call a representation *decomposable* if we can write it as the direct sum of irreducible representations. A lot can be said about the case where the representation of a finite group is over a field whose characteristic not dividing the order of the group. In this case, *Maschke's theorem* tells us that these representations are always decomposable [Lan02]. In particular, complex representations of a finite group are always decomposable.

**Gelfand Pairs.** Henceforth, we assume some knowledge of abstract algebra from the reader. Let  $G$  be a finite group and  $K \leq G$  a subgroup. The pair  $(G, K)$  is called a *Gelfand pair* if the induced representation  $\text{Ind}_K^G \mathbf{1}$  is multiplicity-free. Here  $\mathbf{1}$  denotes the trivial (1-dimensional) complex representation of  $K$ , and multiplicity-free means that any irreducible representation appears in the decomposition of  $\text{Ind}_K^G \mathbf{1}$  at most once (up to isomorphism).

Gelfand pairs play an important role in representation theory [Mus93], analysis [Kor80, Mor18], combinatorics [BI84], number theory [Gro91, Ter99] and probability [CSST20, Dia88]. One of our objectives is to give a detailed study of Gelfand pairs of finite groups. A main theorem of this thesis is the following:

**Theorem 1.** (*Gelfand's Trick*) Let  $G$  be a finite group and  $K$  a subgroup of  $G$ . Suppose  $\varphi : G \rightarrow G$  is an involutive anti-automorphism (i.e. a bijective anti-homomorphism) such that  $K\varphi(x)K = KxK$  for all  $x \in G$ . Then  $(G, K)$  is a Gelfand pair.

The theorem above is proved using the *Hecke algebra*. There are multiple constructions of Hecke algebras in the literature [CMHL03, CSST20].

**Types of Hecke algebras.** Another objective of this thesis is to present these a priori different Hecke algebras and resolve their apparent discrepancies. For instance, one way to define the Hecke algebra is as a convolution algebra of  $K$ -bi-invariant complex-valued functions  $f: G \rightarrow \mathbb{C}$  on a group. Another way to define the Hecke algebra is as the algebra generated by  $n - 1$  variables  $T_1, \dots, T_{n-1}$  subject to a *quadratic relation*  $T_i^2 = (q - 1)T_i + q$  and a *braid relation*

$$\underbrace{T_i T_j T_i \dots}_{m_{ij} \text{ terms}} = \underbrace{T_j T_i T_j \dots}_{m_{ij} \text{ terms}}$$

Here  $m_{ij}$  is the  $ij^{\text{th}}$  entry in the *Coxeter matrix* associated to the *Weyl group* of  $G$ . The name ‘braid relation’ is due to a method of visualising the *symmetric group*  $S_n$ . If  $n$  is a positive integer then the group  $S_n$  is the collection of bijections on the set  $\{1, 2, \dots, n\}$  to itself, with the group operation of composing functions. A natural method of visualising elements and multiplication in this group is via *braid diagrams*. For instance, if  $\sigma = (1\ 2)(3\ 5\ 4)$  and  $\pi = (1\ 2\ 4\ 6\ 5\ 3)$  are permutations in  $S_6$  (written in cycle notation), then we may visualise these elements and their product  $\pi\sigma = (1\ 4)(5\ 6)$  in the following manner:

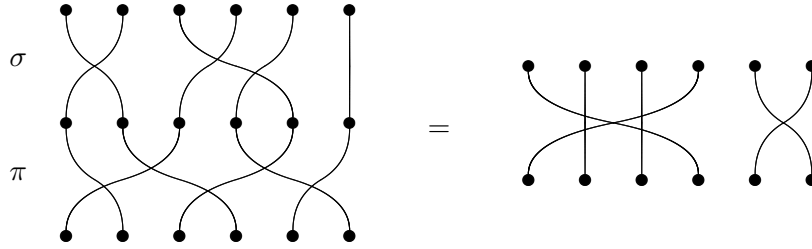


FIGURE 2. A braid diagram visualising the multiplication  $\pi\sigma = (1\ 4)(5\ 6)$ .

**Why study Hecke algebras?** The Hecke algebra arises naturally when one wishes to compute certain irreducible representations of a group [RW21, CMHL03]. Consider a finite group  $G$  and a normal subgroup  $N \triangleleft G$ . If  $G$  acts linearly on a vector space  $V$  (i.e.  $V$  is a representation of  $G$ ), then there is a natural action of  $G$  on the subrepresentation  $V^N$ , the space of vectors in  $V$  that are fixed by  $N$ . Under this action,  $N$  will clearly act trivially on  $V^N$ . This yields a representation of the quotient group  $G/N$ . After some representation theoretic arguments, one arrives at the conclusion that

$$\left\{ \begin{array}{c} \text{Irreducible representations} \\ \text{of } G \text{ with } N\text{-fixed vectors} \end{array} \right\} \xleftrightarrow{1:1} \left\{ \begin{array}{c} \text{Irreducible representations} \\ \text{of } G/N \end{array} \right\}.$$

It is a straightforward exercise that a complex representation of the group  $G/N$  is the same as a representation of the algebra  $\mathbb{C}[G/N]$ , the *group algebra* of  $G/N$ .

What happens when we do not require a normal subgroup of  $G$ ? Consider an arbitrary subgroup  $K$  of a finite group  $G$ . Now  $G/K$  is no longer necessarily a group, so  $G/K$  and  $\mathbb{C}[G/K]$  no longer necessarily make sense. We ask ourselves: what acts on  $V^K$ ? The action of  $G$  on  $V$  is not well-defined on  $V^K$  since  $K$  is no longer normal. It is not obvious how we could study irreducible  $G$ -representations with  $K$ -fixed vectors. We are able to salvage the situation with the help of the Hecke algebra.



For  $g \in G$ , define the Hecke operator  $[KgK] := \frac{1}{|K|} \sum_{x \in KgK} x \in \mathbb{C}[G]$ , which acts on  $V^K$  by

$$[KgK] \cdot v := \frac{1}{|K|} \sum_{x \in KgK} x \cdot v.$$

Now define the Hecke algebra  $\mathcal{H}(G, K)$  to be the space of functions  $f: G \rightarrow \mathbb{C}$  that are constant on  $K$ -double cosets. The indicator functions  $\chi_{KgK}$  form a basis of this space and we can uniquely associate the indicator functions  $\chi_{KgK}$  to the Hecke operators  $[KgK]$ . We see that, through the Hecke operators, we have defined an action of  $\mathcal{H}(G, K)$  on  $V^K$ . This answers our question of what acts on  $V^K$ . Through another representation-theoretic exercise, one can conclude that

$$\left\{ \begin{array}{c} \text{Irreducible representations} \\ \text{of } G \text{ with } K\text{-fixed vectors} \end{array} \right\} \xleftrightarrow{1:1} \left\{ \begin{array}{c} \text{Irreducible representations} \\ \text{of } \mathcal{H}(G, K) \end{array} \right\}.$$

An immediate example of the utility of this result is as follows. It is easy to show that if  $\mathcal{H}(G, K)$  is commutative, then all of its irreducible finite-dimensional representations are one-dimensional [EGH<sup>+</sup>11]. The commutativity of the Hecke algebra turns out to be an important property which will be investigated throughout this thesis.

**Contents of this thesis.** In Chapter 1, we begin our study of the Hecke algebra. First, we investigate the convolution algebra of all complex-valued functions on  $G$  and its ideal of  $K$ -right-invariant complex-valued functions. This is followed by results describing the relationship between the induced representation and its associated Hecke algebra. We use these results to prove Theorem 1. This allows us to write down simple proofs that  $\text{Ind}_K^G \mathbf{1}$  is multiplicity-free for certain choices of  $G$  and  $K$ . Namely,  $(G, K)$  with  $G$  commutative,  $(G, K)$  with  $[G:K] = 2$ ,  $(S_{n+m}, S_n \times S_m)$ , and  $(O_{n+1}(\mathbb{F}_q), O_n(\mathbb{F}_q))$  for  $q$  odd.

In Chapter 3, we generalise the discussion of Chapter 1 to the case of a non-trivial *character*  $\sigma: K \rightarrow \mathbb{C}^\times$ . Here our goal is to obtain a twisted analogue of Theorem 1. To this end, we describe the basis of the Hecke algebra using the idea of *relevant orbits*. We state and prove the generalisation of Theorem 1. We apply the new theorem to a particular representation, the *Gelfand–Graev representation* of  $\text{GL}_n(\mathbb{F}_q)$ , to show that it is multiplicity-free.

In Chapter ??, we investigate the Hecke algebra of Chapter 1 under the particular choice of  $G = \text{SL}_n(\mathbb{F}_q)$  and  $K = B(\mathbb{F}_q)$ , the *Borel subgroup* of  $G$ , i.e. the subgroup of upper-triangular matrices. The *Weyl group* associated to  $G$  is introduced and shown to be isomorphic to  $S_n$ . Next, we perform some elementary matrix calculations which yields the surprising result above: the Hecke algebra may be written in terms of  $n - 1$  generators subject to the quadratic relation and the braid relations associated to  $W$ . This leads to a concluding discussion of Hecke algebras generated by any finite *Coxeter group*.

In Chapter 2, we generalise the results of earlier chapters to the case where  $G$  is no longer finite, but instead a locally compact topological group. This allows for an extension of the theory we have developed to more general groups and their Hecke algebras. To do this, we discuss how one can impose a topological structure on a group and supply examples to give some intuition for these types of groups. To define Hecke algebras of these groups, we require some measure theory. In particular, the convolution product on the Hecke algebra is defined in terms of an integral with respect to the *Haar measure*. We spend some time developing the theory of Haar measures for this purpose. We conclude with a discussion of how to

recover the Hecke algebra of a finite group from this new definition. In this chapter, we shall denote the Hecke algebra by  $C_c(K \backslash G / K)$  to emphasise the non-finiteness of  $G$ .

In Chapter ??, we take a look at some specific Hecke algebras of locally compact topological groups. In particular, we restrict our attention to the general linear group over a non-archimedean local field  $k$  and its ring of integers  $\mathcal{O}$ . We look at the *Spherical Hecke algebra*, formed when one considers  $G = \mathrm{GL}_n(k)$  and  $K = K^\circ := \mathrm{GL}_n(\mathcal{O})$ , and the *Iwahori–Hecke algebra*, formed when one considers  $G = \mathrm{GL}_n(\mathcal{O})$  and  $K = I$ , the *Iwahori subgroup*. In order to investigate these algebras, we must develop an understanding of these fields. We detail their definition, classification and structure.

The contents of this thesis may be visualised with the following diagram.

$$\begin{array}{ccccc}
 \text{Ch. 3} & & \text{Ch. 1} & & \text{Ch. 2} & & \text{Ch. ??} \\
 \mathcal{H}(G, K, \sigma) & \xrightarrow{\sigma=1} & \mathcal{H}(G, K) & \xleftarrow{G \text{ finite}} & C_c(K \backslash G / K) & \xrightarrow{K=I} & C_c(I \backslash G / I) \\
 & & \downarrow \begin{array}{c} G \text{ Lie type} \\ \text{over } \mathbb{F}_q \end{array} & & \downarrow K=K^\circ & & \\
 & & \mathcal{H}_q(W, S) & & C_c(K^\circ \backslash G / K^\circ) & & \\
 & & \text{Ch. ??} & & \text{Ch. ??} & & 
 \end{array}$$

FIGURE 3. The relationship diagram of this thesis.

**Directions for future research.** We assume the reader is familiar with the contents of this thesis. The modern study of Hecke algebras is largely focused on the *Iwahori–Hecke algebra*, which is also known as the *affine Hecke algebra*. This algebra is central to the study of representations of *reductive groups* over non-archimedean local fields (e.g. groups such as  $\mathrm{GL}_n$ ,  $\mathrm{SL}_n$ ,  $\mathrm{Sp}_{2n}$  over fields such as  $\mathbb{Q}_p$  or  $\mathbb{F}_q((t))$ ).

Some topics relevant to the Iwahori–Hecke algebra include *Bernstein’s presentation*, the *Iwahori–Matsumoto presentation* and the *Satake isomorphism* [HKP09]. Properties of the Iwahori–Hecke algebra such as these presentations may be viewed as a consequence of the *universal unramified principal series module*, which we now describe.

Fix a “nice” (i.e. split and connected) reductive group  $G$  (e.g.  $\mathrm{SL}_n$ ) over a non-archimedean local field  $k$  with ring of integers  $\mathcal{O}$ . Then write  $A$  to mean a *split maximal torus* of  $G$  and write  $N$  to mean the *unipotent radical* of a Borel subgroup of  $G$  that contains  $A$ . Also recall  $I$  is the Iwahori subgroup of  $G$  given in Chapter ??.

The universal unramified principal series module  $M$  is given by  $C_c(A(\mathcal{O})N \backslash G / I)$ . It is a right module over the Iwahori–Hecke algebra under convolution. Furthermore, a basis of the Iwahori–Hecke algebra is parameterised by the *affine Weyl group*  $\widetilde{W}$ . We may write  $\widetilde{W} \cong W \ltimes \Lambda^\vee$ , where  $\Lambda^\vee$  is the *coroot lattice* of  $G$ . Then  $\mathbb{C}[\Lambda^\vee]$  is the corresponding group algebra over  $\mathbb{C}$ . Then  $M$  is also a left module over  $\mathbb{C}[\Lambda^\vee]$ .

## 1. GAUSSIAN PROCESSES

The aim of this chapter is to review some essential mathematical machinery required for us to understand the core concepts of Gaussian Processes.

## 2. KRYLOV SUBSPACE METHODS

In this section we will focus on how iterative methods, in particular a class of iterative methods called Krylov Subspace methods, may be used to solve a linear system  $Ax = b$ . While non-iterative methods exist to solve such systems virtually all of them carry an unwieldy runtime of  $\mathcal{O}(n^3)$  for a system of  $n$  parameters. Even for current computer systems, this renders many common matrix problems untractable. Consequently the focus of solving linear systems has shifted towards iterative methods. While iterative methods typically demand certain structural properties of the matrices, such as symmetry and positive definiteness, this generally is not a problem since the majority of large matrix problems that, by nature, endow these systems with the desired properties. For example, in the context of this paper the Gram matrices used to solve linear systems in Gaussian Processes possess both symmetry and positive definiteness. There are also a number of other properties of iterative methods which make them rather attractive to users. To start, iterative Krylov subspace methods are guaranteed to converge to an exact solution within a finite number of iterations and even if the method is prematurely stopped before reaching an exact solution, the approximation obtained on the final iteration will in some sense be a good enough estimate of our exact solution. Furthermore, unlike most non-iterative methods, Krylov subspace methods do not require an explicit form of the matrix  $A$  and instead only requires some routine or process for computing  $Ax$ .

**2.1. Krylov Subspaces.** We will motivate the Krylov subspaces by observing their usefulness in solving linear systems. To this end, consider the problem of solving the linear system

$$(1) \quad Ax^* = b$$

where no explicit form of  $A$  is available and instead one must draw information from  $A$  solely through a routine that can evaluate  $Av$  for any  $v$ . How could this routine be utilized in such a manner to provide with a solution to equation 1? Before answering this, consider the following theorem

**Theorem 2.** For  $A \in \mathbb{K}^{n \times n}$  if  $\|A\| = q < 1$  then  $\mathbb{1} - A$  is invertible and its inverse admits the following representation

$$(\mathbb{1} - A)^{-1} = \sum_{k=0}^{\infty} A^k.$$

[Ber96]

Consider a matrix for which  $\|A\| < 2$ , it follows that  $\|\mathbb{1} - A\| < 1$  meaning  $\mathbb{1} - (\mathbb{1} - A)$  is invertible and  $A^{-1} = (\mathbb{1} - (\mathbb{1} - A))^{-1} = \sum_{k=0}^{\infty} (\mathbb{1} - A)^k$ . Thinking back to equation 1 for any  $x_0 \in \mathbb{K}^n$  we have

$$\begin{aligned} x^* &= A^{-1}b = A^{-1}(Ax^* - Ax_0 + Ax_0) \\ &= x_0 + A^{-1}r_0 \\ &= x_0 + \sum_{k=0}^{\infty} (\mathbb{1} - A)^k \end{aligned}$$

where  $r_0 = Ax^* - Ax_0$ . A natural question that arises is that can we find a closed form solution of the above equation? To answer this question we need to enlist the help of the Cayley-Hamilton theorem.

**Theorem 3** (Cayley-Hamilton). Let  $p_n(\lambda) = \sum_{i=0}^n c_i \lambda^i$  be the characteristic polynomial of the matrix  $A \in \mathbb{K}^{n \times n}$ , then  $p_n(A) = 0$ . **THIS NEEDS A CITATION**

The Cayley-Hamilton theorem implies that

$$\begin{aligned} 0 &= c_0 + c_1 A + \dots + c_{n-1} A^{n-1} + c_n A^n \\ 0 &= A^{-1} c_0 + c_1 + \dots + c_{n-1} A^{n-2} + c_n A^{n-1} \\ A^{-1} &= \alpha_0 + c_1 + \dots + \alpha_{n-1} A^{n-2} + \alpha_n A^{n-1} \end{aligned}$$

where  $\alpha_i = -c_i/c_0$ . This demonstrates that  $A^{-1}$  can be represented as a matrix polynomial of degree  $n-1$ . This means that  $\sum_{k=0}^{\infty} (\mathbb{1} - A)^k$  indeed possess a closed form solution namely

$$x^* = x_0 + A^{-1} r_0 = \alpha_0 + c_1 + \dots + \alpha_{n-1} A^{n-2} + \alpha_n A^{n-1}.$$

This also shows that  $x^* \in \text{l.s} \{r_0, Ar_0, A^2 r_0, \dots, A^{n-1} r_0\}$ . One idea for finding a solution to equation 1 is to use our routine for evaluating  $Av$  to iteratively compute new basis elements for the space generated by  $\{r_0, Ar_0, A^2 r_0, \dots, A^{n-1} r_0\}$  and at each step carefully choosing a  $x_k$  such that  $x_k$  approaches  $x^*$ , in some form. The subspace constructed using this technique is so important that it has its own name.

**Definition 4** (Krylov Subspace). The Krylov Subspace of order  $k$  generated by the matrix  $A \in \mathbb{K}^{n \times n}$  and the vector  $v \in \mathbb{K}$  is defined as

$$\mathcal{K}_k(A, v) = \text{l.s} \{r_0, Ar_0, A^2 r_0, \dots, A^{k-1} r_0\}$$

for  $k \geq 1$  and  $\mathcal{K}_0(A, v) = \{0\}$ .

For the purposes of solving equation 1 it is of much interest to understand how  $\mathcal{K}_k(A, v)$  grows for larger and larger  $k$  since a solution for equation 1 will be present in a Krylov Subspace that cannot be grown any larger. In other words, an exact solution can be constructed once we have extracted all the information from  $A$  through multiplication of  $r_0$ . The following theorem provides information on how exactly the Krylov Subspace grows as  $k$  increases.

**Theorem 5.** There is a positive called the grade of  $v$  with respect to  $A$ , denoted  $t_{v,A}$ , where

$$\dim(\mathcal{K}_k(A, v)) = \begin{cases} k, & k \leq t \\ t, & k \geq t \end{cases}$$

Theorem 5 essentially tells us that for  $k \leq t_{v,A}$  that  $A^k v$  is linearly independent to  $A^i v$  for  $0 \leq i \leq k-1$  meaning  $\{v, Av, A^2 v, \dots, A^{k-1} v\}$  serves as a basis for  $\mathcal{K}_k(A, v)$  and that  $\mathcal{K}_{k-1}(A, v) \subsetneq \mathcal{K}_k(A, v)$ . Conversely, any new vectors formed beyond  $t_{v,A}$  will be linearly independent meaning  $\mathcal{K}_k(A, v) \subsetneq \mathcal{K}_{k+1}(A, v)$  for  $k \geq t_{v,A}$ . While  $t_{v,A}$  clearly plays a role in determining a suitable basis for which  $A^{-1}b$  lies in its importance is made abundantly clear in the following corollary.

**Corollary 6.**

$$t_{v,A} = \min \{k \mid A^{-1}v \in \mathcal{K}_k(A, v)\}$$

*Proof.* Recall from Cayley-Hamilton (theorem 3) that

$$\mathbf{A}^{-1}\mathbf{v} = \sum_{i=0}^{n-1} \alpha_i \mathbf{A}^i \mathbf{v}$$

But since  $\mathcal{K}_k(\mathbf{A}, \mathbf{v}) = \mathcal{K}_{k+1}(\mathbf{A}, \mathbf{v})$  for  $k \geq t_{\mathbf{v}, \mathbf{A}}$

$$\mathbf{A}^{-1}\mathbf{v} = \sum_{i=0}^{t-1} \beta_i \mathbf{A}^i \mathbf{v}$$

meaning  $\mathbf{A}^{-1}\mathbf{v} \in \mathcal{K}_k(\mathbf{A}, \mathbf{v})$  for  $k \geq t_{\mathbf{v}, \mathbf{A}}$ . Suppose for the sake of contradiction that this also holds for  $k = t_{\mathbf{v}, \mathbf{A}} - 1$ , that is,  $\mathbf{A}^{-1}\mathbf{v} = \sum_{i=0}^{t-2} \gamma_i \mathbf{A}^i \mathbf{v}$ . However, this gives

$$\mathbf{v} = \sum_{i=0}^{t-2} \gamma_i \mathbf{A}^{i+1} \mathbf{v} = \sum_{i=0}^{t-1} \gamma_{i-1} \mathbf{A}^i \mathbf{v}$$

implying  $\{\mathbf{v}, \mathbf{A}\mathbf{v}, \mathbf{A}^2\mathbf{v}, \dots, \mathbf{A}^{t-1}\mathbf{v}\}$  are linearly dependent which means that  $\dim(\mathcal{K}_k(\mathbf{A}, \mathbf{v})) < t$ , which provides us with our contradiction.  $\square$

This machinery allows us to make a much stronger statement on the whereabouts of  $\mathbf{x}^\star$  in relation to the Krylov Subspaces.

**Corollary 7.** *For any  $\mathbf{x}_0$ , we have*

$$\mathbf{x}^\star \in \mathbf{x}_0 + \mathcal{K}_{t_{\mathbf{r}_0, \mathbf{A}}}(\mathbf{A}, \mathbf{r}_0)$$

where  $\mathbf{r}_0 = \mathbf{b} - \mathbf{A}\mathbf{x}_0$ .

**2.2. Gram-Schmidt Process and QR factorisations.** Many areas of linear algebra involving studying the column space of matrices. The  $QR$  factorisation provides us with a powerful tool to better understand the column space of a matrix as well as serving as an important factorisation mechanism for many numerical methods. Suppose that a matrix  $\mathbf{A} = [\mathbf{a}_1, \mathbf{a}_2, \dots, \mathbf{a}_n] \in \mathbb{K}^{n \times n}$  has full rank. The idea of a  $QR$  factorisation is to find an alternative orthonormal basis for  $(\mathbf{a}_i)_{i=1}^n$ , say  $(\mathbf{q}_i)_{i=1}^n$ , and to somehow relate the original matrix  $\mathbf{A}$  to a new matrix whose columns are  $(\mathbf{q}_i)_{i=1}^n$ . Consider the following procedure that allows us to find an orthonormal basis  $(\mathbf{q}_i)_{i=1}^n$  for which  $\text{l.s}\{(\mathbf{a}_i)_{i=1}^n\} = \text{l.s}\{(\mathbf{q}_i)_{i=1}^n\}$ . First set  $\mathbf{q}_1 = \frac{\mathbf{a}_1}{\|\mathbf{a}_1\|}$ , clearly  $\text{l.s}\{\mathbf{a}_1\} = \text{l.s}\{\mathbf{q}_1\}$ . Next, construct a vector  $\mathbf{q}'_2 = \mathbf{a}_2 - r_{1,2} \cdot \mathbf{q}_1$  so that  $\mathbf{q}'_2 \perp \mathbf{q}_1$ . This means

$$\begin{aligned} 0 &= \langle \mathbf{q}_1, \mathbf{q}'_2 \rangle \\ 0 &= \langle \mathbf{q}_1, \mathbf{a}_2 - r_{1,2} \cdot \mathbf{q}_1 \rangle \\ 0 &= \langle \mathbf{q}_1, \mathbf{a}_2 \rangle - r_{1,2} \cdot \langle \mathbf{q}_1, \mathbf{q}_1 \rangle \\ r_{1,2} &= \langle \mathbf{q}_1, \mathbf{a}_2 \rangle \end{aligned}$$

Since  $\mathbf{q}'_2$  may not be a unit vector we set  $\mathbf{q}_2 = \frac{\mathbf{q}'_2}{\|\mathbf{q}'_2\|}$  where  $\text{l.s}\{(\mathbf{a}_1, \mathbf{a}_2)\} = \text{l.s}\{(\mathbf{q}_1, \mathbf{q}_2)\}$ . Continuing the vector  $\mathbf{q}'_3$  is constructed so that

$$\mathbf{q}'_3 = \mathbf{a}_3 - r_{1,3}\mathbf{q}_1 - r_{2,3}\mathbf{q}_2$$

are chosen so that  $\mathbf{q}'_3$  is orthogonal to both  $\mathbf{q}_2$  and  $\mathbf{q}_1$ . This amounts to setting  $r_{1,3} = \langle \mathbf{q}_1, \mathbf{a}_3 \rangle$  and  $r_{2,3} = \langle \mathbf{q}_2, \mathbf{a}_3 \rangle$ . Similarly,  $\mathbf{q}'_3$  is normalized so that  $\mathbf{q}_3 = \frac{\mathbf{q}'_3}{\|\mathbf{q}'_3\|}$  and  $\text{l.s}\{(\mathbf{a}_1, \mathbf{a}_2, \mathbf{a}_3)\} = \text{l.s}\{(\mathbf{q}_1, \mathbf{q}_2, \mathbf{q}_3)\}$ .

Continuing in this fashion the  $k^{th}$  vector in our orthonormal basis is computed as

$$(2) \quad \mathbf{q}_k = \frac{\mathbf{a}_k - \sum_{i=1}^{k-1} r_{i,k} \cdot \mathbf{q}_i}{r_{k,k}}$$

where  $r_{i,k} = \langle \mathbf{q}_i, \mathbf{a}_k \rangle$ ,  $r_{k,k} = \|\mathbf{a}_k - \sum_{i=1}^{k-1} r_{i,k} \cdot \mathbf{q}_i\|$  and  $\text{l.s}(\{\mathbf{a}_1, \mathbf{a}_2, \dots, \mathbf{a}_k\}) = \text{l.s}(\{\mathbf{q}_1, \mathbf{q}_2, \dots, \mathbf{q}_k\})$ . This procedure is famously known as the Gram-Schmidt process [Ber96, Tre97, Dem97] and is summarized in the following algorithm.

---

**Algorithm 1:** Classical Gram-Schmidt

---

**input :** A basis  $(\mathbf{a}_i)_{i=1}^n$ .

**output:** An orthonormal basis  $(\mathbf{q}_i)_{i=1}^n$  such that  $\text{l.s}\{(\mathbf{a}_i)_{i=1}^n\} = \text{l.s}\{(\mathbf{q}_i)_{i=1}^n\}$

**for**  $k = 1$  **to**  $n$  **do**

$\mathbf{q}'_k = \mathbf{a}_k$

**for**  $i = 1$  **to**  $k - 1$  **do**

$r_{i,k} = \langle \mathbf{q}_i, \mathbf{a}_k \rangle$

$\mathbf{q}'_k = \mathbf{q}'_k - r_{i,k} \mathbf{q}_i$

**end**

$r_{k,k} = \|\mathbf{q}'_k\|$

$\mathbf{q}_k = \mathbf{q}'_k / r_{k,k}$

**end**

**return**  $(\mathbf{q}_i)_{i=1}^n$

---

Relating the column space of  $\mathbf{A}$  to the orthonormal basis  $(\mathbf{q}_i)_{i=1}^n$  in a matrix form

$$[\mathbf{a}_1, \mathbf{a}_2, \dots, \mathbf{a}_n] = [\mathbf{q}_1, \mathbf{q}_2, \dots, \mathbf{q}_n] \begin{bmatrix} r_{1,1} & r_{1,2} & \cdots & r_{1,n} \\ & r_{2,2} & & \vdots \\ & & \ddots & \vdots \\ & & & r_{n,n} \end{bmatrix}$$

or more succinctly

$$(3) \quad \mathbf{A} = \mathbf{Q}\mathbf{R}$$

where  $\mathbf{Q} = [\mathbf{q}_1, \mathbf{q}_2, \dots, \mathbf{q}_n]$  and  $(\mathbf{R})_{i,j} = r_{i,j}$  for  $i \leq j$  and  $(\mathbf{R})_{i,j} = 0$  for  $i > j$ . This is exactly the  $QR$  factorisation for a full rank matrix. Note that  $\text{Range}(\mathbf{A}) = \text{Range}(\mathbf{Q})$ . In general, any square matrix  $\mathbf{A} \in \mathbb{K}^{m \times n}$  may be decomposed as  $\mathbf{A} = \mathbf{Q}\mathbf{R}$  where  $\mathbf{Q} \in \mathbb{K}^{m \times m}$  is an orthogonal matrix and  $\mathbf{R} \in \mathbb{K}^{m \times n}$  is an upper triangular matrix. This is known as a full  $QR$  factorisation. Since bottom  $(m - n)$  rows of this  $\mathbf{R}$  consists entirely of zeros, it is often useful to partition the full  $QR$  factorisation in the following manner to shed vacuous entries

$$\mathbf{A} = \mathbf{Q}\mathbf{R} = \mathbf{Q} \begin{bmatrix} \hat{\mathbf{R}} \\ \mathbf{0}_{(m-n) \times n} \end{bmatrix} = \begin{bmatrix} \hat{\mathbf{Q}} & \mathbf{Q}' \end{bmatrix} \begin{bmatrix} \hat{\mathbf{R}} \\ \mathbf{0}_{(m-n) \times n} \end{bmatrix} = \hat{\mathbf{Q}}\hat{\mathbf{R}}.$$

This alternate decomposition is called the reduced (or sometimes the thin)  $QR$  factorization. We shall state the following two theorems on the  $QR$  factorization are stated without proof.

**Theorem 8.** Every  $A \in \mathbb{K}^{m \times n}$ , ( $m \geq n$ ) has a full QR factorisation, hence also a reduced QR factorisation. [Tre97]

**Theorem 9.** Each  $A \in \mathbb{K}^{m \times n}$ , ( $m \geq n$ ) of full rank has a unique reduced QR factorisation  $A = \hat{Q}\hat{R}$  with  $r_{k,k} > 0$ . [Tre97]

In practice the classical Gram-Schmidt process described in algorithm 1 is rarely used as the procedure becomes numerically unstable if  $(a_i)_{i=1}^n$  are almost linearly dependent. Before looking for ways to resolve these numerical instabilities a quick recap of projectors has been devised. A square matrix  $P_G$  acting on a Hilbert space  $H$  that sends  $x \in H$  to its projection onto a subspace  $G$  is called the projector onto  $G$ . If  $(q_k)_{k=1}^m$  is an orthonormal basis in  $G$  then

$$P_G = QQ^*$$

where  $Q = [q_1, q_2, \dots, q_m, 0, \dots, 0] \in \mathbb{K}^{n \times n}$ . A special class of projectors which isolates the components of a given vector onto a one dimensional subspace spanned by a single unit vector  $q$  called a rank one orthogonal projector, denoted as  $P_q$ . Each  $k$  in the classical Gram-Schmidt process  $q'_k$  using the following orthogonal projection

$$(4) \quad q'_k = P_{A_k^\perp} a_k$$

where  $A_k = \text{l.s}\{a_i\}_{i=1}^k$  and  $P_{A_1^\perp} = \mathbb{1}$  for convenience. A modified version of the Gram-Schmidt process performs the same orthogonal projection broken up as  $k-1$  orthogonal projections of rank  $n-1$  as so

$$\begin{aligned} q'_k &= P_{A_k^\perp} a_k \\ &= (\mathbb{1} - Q_k Q_k^*) a_k \\ &= \left( \prod_{i=1}^{k-1} (\mathbb{1} - q_i q_i^*) \right) a_k \\ &= (\mathbb{1} - q_1 q_1^*) (\mathbb{1} - q_2 q_2^*) \cdots (\mathbb{1} - q_{k-1} q_{k-1}^*) a_k \\ &= P_{q_k^\perp} \cdots P_{q_1^\perp} a_k \end{aligned}$$

While its clear that  $P_{A_k^\perp} a$  and  $P_{q_k^\perp} \cdots P_{q_1^\perp} a_k$  used for computing  $q'_k$  are algebraically, they differ arithmetically as the latter expression evaluates  $q'_k$  using the follow procedure

$$\begin{aligned} q_k^{(1)} &= a_k \\ q_k^{(2)} &= P_{q_1^\perp} q_k^{(1)} \\ q_k^{(3)} &= P_{q_2^\perp} q_k^{(2)} \\ &\vdots \\ q'_k &= q_k^{(k)} = P_{q_{k-1}^\perp} q_k^{(k-1)} \end{aligned}$$

Applying projections sequentially in this manner produces smaller numerical errors. The modified Gram-Schmidt process [Tre97, Dem97] is summarized in the following algorithm.



---

**Algorithm 2:** Modified Gram-Schmidt
 

---

**input** : A basis  $\{a_i\}_{i=1}^n$ .  
**output**: An orthonormal basis  $\{q_i\}_{i=1}^n$  such that  $\text{l.s}\{a_i\}_{i=1}^n = \text{l.s}\{q_i\}_{i=1}^n$   
**for**  $k = 1$  **to**  $n$  **do**  
      $q'_k = a_k$   
**end**  
**for**  $k = 1$  **to**  $n$  **do**  
      $r_{k,k} = \|q'_k\|$   
      $q_k = q'_k / r_{k,k}$   
     **for**  $i = k + 1$  **to**  $n$  **do**  
          $r_{i,k} = \langle q_k, q'_i \rangle$   
          $q_i = q_i - r_{i,k} q_k$   
     **end**  
**end**  
**return**  $\{q_i\}_{i=1}^n$

---

**2.3. Arnoldi and Lanczos Algorithm.** As a quick reminder, we are in search of an iterative process to solve the linear system  $Ax^* = b$  where no explicit form of  $A$  is available and we may only rely on a routine that computes  $Av$  for any  $v$  to extract information on  $A$ . In section 2.1 we discovered that  $x^* \in \mathcal{K}_{t_{r_0}, A}(A, r_0)$ . With many iterative methods, computing an exact value for  $x^*$  is out the question with the view that  $t_{r_0, A}$  is impractically large. We must instead resort to approximating  $x^*$  by  $x_k$  for which  $x^k \in \mathcal{K}_k(A, r_0)$  where  $k \ll t_{r_0}$ . To find an appropriate value for  $x_k$ , a good start would be to find a basis  $\mathcal{K}_k(A, r_0)$ . Definition 4 showed us that  $\{A^{i-1}r_0\}_{i=1}^k$  serves as a basis for  $\mathcal{K}_k(A, r_0)$ . However, for numericcal reasons this is a poor choice of basis since this each consecutive term becomes closer and closer to being linearly dependent. To search for a more approporate basis, set  $n = t_{r_0, A}$  so that  $x^* \in \mathcal{K}_n(A, r_0)$ . Let  $K \in \mathbb{K}^{n \times n}$  be the invertible matrix

$$K = [r_0, Ar_0, \dots, A^{n-1}r_0].$$

Since  $K$  is invertible we can compute  $c = -K^{-1}A^n r_0$  so that

$$\begin{aligned}
 AK &= [Ar_0, A^2r_0, \dots, A^n r_0] \\
 AK &= K \cdot [e_2, e_3, \dots, e_n, -c] \triangleq KC
 \end{aligned}$$

or, in a more verbose form

$$K^{-1}AK = C = \begin{bmatrix} 0 & 0 & \cdots & 0 & -c_1 \\ 1 & 0 & \cdots & 0 & -c_2 \\ 0 & 1 & \cdots & 0 & \vdots \\ \vdots & \vdots & \cdots & \vdots & \vdots \\ 0 & 0 & \cdots & 1 & -c_n \end{bmatrix}.$$

Note here that  $C$  is upper Hessenberg. While this form is simple, it is of little practical use since the matrix  $K$  is very likely to be ill-conditioned. To remedy this we can replace  $K$  with an orthogonal matrix which

spans the same space. These are exactly the properties that the  $Q$  matrix offers in the  $QR$ -factorisation of  $K$ . With this in mind let  $K = QR$  be the full  $QR$ -factorisation of  $K$ . Then

$$\begin{aligned} AQR &= AK \\ AQ &= AKR^{-1} \\ AQ &= KCR^{-1} \\ AQ &= QRCR^{-1} \\ AQ &\triangleq QH. \end{aligned}$$

Since  $R$  and  $R^{-1}$  and both upper triangular and  $C$  is upper Hessenberg,  $H$  is also upper Hessenberg. This form provides us with a  $Q$  such that the range of  $Q$  is  $\mathcal{K}_n(A, r_0)$  and

$$(5) \quad Q^T AQ = H.$$

Again, in practice, it may be very difficult to compute this entire expression forcing us to search for approximative alternatives. Consider equation 5 for which the only first  $k$  columns of  $Q$  have been computed. Let  $Q_k = [q_1, q_2, \dots, q_k]$  and  $Q_u = [q_{k+1}, q_{k+2}, \dots, q_n]$ . Then

$$\begin{aligned} Q^T AQ &= H \\ [Q_k, Q_u]^T A [Q_k, Q_u] &= \begin{bmatrix} H_k & H_{u,k} \\ H_{k,u} & H_u \end{bmatrix} \\ \begin{bmatrix} Q_k^T AQ_k & Q_k^T AQ_u \\ Q_u^T AQ_k & Q_u^T AQ_u \end{bmatrix} &= \begin{bmatrix} H_k & H_{u,k} \\ H_{k,u} & H_u \end{bmatrix} \end{aligned}$$

where  $H_k$ ,  $H_{u,k}$ ,  $H_{k,u}$  and  $H_u$  are the relevant sub matrices. This provides us with the equality

$$(6) \quad Q_k^T AQ_k = H_k$$

noting that  $H_k$  is upper Hessenberg for the same reason that  $H$  is. We know that when  $n = t_{r_0, A}$  we can find a  $Q \in \mathbb{K}^{n \times n}$  and  $H \in \mathbb{K}^{n \times n}$  that satisfies  $AQ = QH$ . However, in general, we may not be so fortunate in finding a  $Q_k \in \mathbb{K}^{n \times k}$  and  $H_k \in \mathbb{K}^{k \times k}$  so satisfy  $AQ_k = Q_k H_k$  for any  $k < n$ . Instead we can adjust this equality by adding an error  $E_k \in \mathbb{K}^{n \times k}$  so that we do get equality. Our expression now becomes

$$(7) \quad Q_k^T AQ_k = H_k + E_k.$$

A careful choice of  $E_k$  must be made to also retain equality in equation 6, meaning  $Q_k^T E_k = 0$ . Since  $\{q_i\}_{i=1}^k$  forms an orthonormal basis for  $\mathcal{K}_n(A, r_0)$ , consider the following choice of  $E_k$ ,

$$E_k = q_{k+1} h_k^T$$

where  $h_k$  is any vector in  $\mathbb{K}^k$ . Notice that

$$Q_k^T E_k = Q_k^T (q_{k+1} h_k^T) = (Q_k^T q_{k+1}) h_k^T = 0.$$

Since this holds for any  $h_k \in \mathbb{K}^k$ , to preserve sparsity and to keep this form as simple as possible we can set  $h_k = [0, 0, \dots, h_{k+1,k}]^T$ . This means  $AQ_k$  can be written as

$$(8) \quad AQ_k = Q_k H_k + q_{k+1} h_k^T$$

where

$$\mathbf{Q}_k \mathbf{H}_k = [\mathbf{q}_1, \mathbf{q}_2, \dots, \mathbf{q}_k] \begin{bmatrix} h_{1,1} & \cdots & \cdots & \cdots & h_{1,k} \\ h_{2,1} & \cdots & \cdots & \cdots & \vdots \\ 0 & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \vdots \\ 0 & \cdots & 0 & h_{k,k-1} & h_{k,k} \\ 0 & \cdots & 0 & 0 & h_{k+1,k} \end{bmatrix}.$$

Equating the  $j^{th}$  columns of equation 8 yields

$$\mathbf{A} \mathbf{q}_j = \sum_{i=1}^{j+1} h_{i,j} \mathbf{q}_i.$$

Again since  $\{\mathbf{q}_i\}_{i=1}^n$  form an orthonormal basis, multiplying both sides by  $\mathbf{q}_m$  for  $1 \leq m \leq j$  gives

$$\mathbf{q}_m^T \mathbf{A} \mathbf{q}_j = \sum_{i=1}^{j+1} h_{i,j} \mathbf{q}_m^T \mathbf{q}_i = h_{m,j}$$

and so

$$(9) \quad h_{j+1,j} \mathbf{q}_{j+1} = \mathbf{A} \mathbf{q}_j - \sum_{i=1}^j h_{i,j} \mathbf{q}_i.$$

From equation 9 we find that  $\mathbf{q}_{j+1}$  can be computed using a recurrence involving its previous Krylov factors. Notice this bears a striking resemblance to equation 2 having a virtually an identical setup to computing an orthonormal basis using the modified Gram-Schmidt process (algorithm 2). As such, values for  $\mathbf{q}_{j+1}$  and  $h_{j+1,j}$  can be evaluated using a procedure very similar to the modified Gram-Schmidt process better known as the Arnoldi algorithm [Tre97, Dem97], presented in algorithm 3.

---

**Algorithm 3:** Arnoldi Algorithm
 

---

**input :**  $A, r_0$  and  $k$ , the number of columns of  $Q$  to compute.

**output:**  $Q_k, H_k$ .

$q_1 = r_0 / \|r_0\|$

**for**  $j = 1$  **to**  $k$  **do**

$z = Aq_j$

**for**  $i = 1$  **to**  $j$  **do**

$h_{i,j} = \langle q_i, z \rangle$

$z = z - h_{i,j}q_i$

**end**

$h_{j+1,j} = \|z\|$

**if**  $h_{j+1,j} = 0$  **then**

**return**  $Q_k, H_k$

**end**

$q_{j+1} = z / h_{j+1,j}$

**end**

**return**  $Q_k, H_k$

---

When  $A$  is symmetric then  $H = T$  becomes a tridiagonal matrix, simplifying a large amount of the Arnoldi algorithm since a considerably large number of matrix elements from  $T$  can be written as

$$T = \begin{bmatrix} \alpha_1 & \beta_1 & & & \\ \beta_1 & \ddots & \ddots & & \\ & \ddots & \ddots & \ddots & \\ & & \ddots & \ddots & \beta_{n-1} \\ & & & \beta_{n-1} & \alpha_n \end{bmatrix}.$$

As before, equating the  $j^{th}$  columns of  $AQ = QT$  yields

$$(10) \quad Aq_j = \beta_{j-1}q_{j-1} + \alpha_jq_j + \beta_jq_{j+1}.$$

Again since  $\{q_i\}_{i=1}^n$  form an orthonormal basis, multiplying both sides of equation 10 by  $q_j$  gives  $q_j^T Aq_j = \alpha_j$ . A simplified version of the Arnoldi algorithm can be devised can be used to compute  $\{q_i\}_{i=1}^n$  and  $T$  for symmetric matrices known as the Lanczos algorithm [Dem97]. The Lanczos algorithm is presented in algorithm 4.

---

**Algorithm 4:** Lanczos Algorithm
 

---

**input :**  $A, r_0$  and  $k$ , the number of columns of  $Q$  to compute.

**output:**  $Q_k, T_k$ .

$q_1 = r_0 / \|r_0\|, \beta_0 = 0, q_0 = 0$

**for**  $j = 1$  **to**  $k$  **do**

$z = Aq_j$

$\alpha_j = \langle q_j, z \rangle$

$z = z - \alpha_j q_j - \beta_{j-1} q_{j-1}$

$\beta_j = \|z\|$

**if**  $\beta_j = 0$  **then**

**return**  $Q_k, T_k$

**end**

$q_{j+1} = z / \beta_j$

**end**

**return**  $Q_k, T_k$

---

**2.4. Optimality Conditions.** So far we have shown that  $x^* \in \mathcal{K}_{t_{r_0}, A}(A, r_0)$  where  $n = t_{r_0}$  is the grade of  $r_0$  with respect to  $A$ . Moreover from section 2.3 we found ways to construct a basis for  $\mathcal{K}_{t_{r_0}, A}(A, r_0)$  allowing us to generate vectors with these affine spaces, namely the Arnoldi algorithm (algorithm 3) and Lanczos algorithm (algorithm 4) for non-symmetric and symmetric systems respectively. From now on  $\mathcal{K}_{t_{r_0}, A}(A, r_0)$  will be abbreviated to  $\mathcal{K}_{t_{r_0}, A}$  when the context is clear. The question still remains however, how should one choose an  $x_k$  that best approximates  $x^*$  satisfying equation 1? Here are a few of the most well known methods for selecting a suitable  $x_k$ .

- (1) Select an  $x_k \in x_0 + \mathcal{K}_k$  which minimizes  $\|x_k - x^*\|_2$ . While this method seems like the most intuitive and natural way to select  $x_k$ , it is unfortunately of no practical use since there is not enough information in the Krylov subspace to find an  $x_k$  which matches this specification.
- (2) Select an  $x_k \in x_0 + \mathcal{K}_k$  which minimizes  $\|r_k\|_2$  (recall is the residual of  $x_k$ , that is,  $r_k = b - Ax_k$ ). This method is possible to implement. Two well known algorithms stem from this class of methods, notably MINRES (minimum residual) and GMRES (general minimum residual) which solve linear systems for symmetric and non-symmetric  $A$  respectively.
- (3) When  $A$  is a positive definite matrix it defines a norm  $\|r\|_A = (r^T A r)^{\frac{1}{2}}$ , called the energy norm. Select an  $x_k \in x_0 + \mathcal{K}_k$  which minimizes  $\|r\|_{A^{-1}}$  which is equivalent to minimizing  $\|x_k - x\|_A$ . This technique is known as the CG (conjugate gradient) algorithm.
- (4) Select an  $x_k \in x_0 + \mathcal{K}_k$  for which  $r_k \perp \mathcal{W}_k$  where  $\mathcal{W}_k$  is some  $k$ -dimensional subspace. Two well known algorithms that belong to this family of methods are SYMMLQ (Symmetric LQ Method) and a variant of GMRES used for solving symmetric and non-symmetric methods respectively.

Interestingly, when  $A$  is symmetric positive definite and  $\mathcal{W}_k = \mathcal{K}_k$  the last two selection methods are equivalent. This is stated more precisely in theorem

**Theorem 10.** *In the context of the above selection method, if  $\mathbf{A} \succ \mathbf{0}$  and  $\mathcal{W}_k = \mathcal{K}_k$  in method (4) then it produces the same  $\mathbf{x}_k$  in method (3) [Dem97].*

In fact the very last method can be used to bring together a number of different analytical aspects and unify them in a general framework known as projection methods. Selecting an  $\mathbf{x}_k$  from our Krylov subspace allows  $k$  degrees of freedom meaning  $k$  constraints must be used to determine a unique  $\mathbf{x}_k$  for selection. As seen in method (4) already, typically orthogonality constraints are imposed on the residual  $\mathbf{r}_k$ . Specifically we would like to find a  $\mathbf{x}_k \in \mathbf{x}_0 + \mathcal{K}_k$  where  $\mathbf{r}_k \perp \mathcal{W}_k$ . This is sometimes referred to as the Petrov-Galerkin (or just Galerkin) conditions. Projection methods for which  $\mathcal{W}_k = \mathcal{K}_k$  are known as orthogonal projections while methods for which  $\mathcal{W}_k = \mathbf{A}\mathcal{K}_k$  are known as oblique projections. If we set  $\mathbf{x}_k = \mathbf{x}_0 + \mathbf{z}_k$  for some  $\mathbf{z}_k \in \mathcal{K}_k$  then the Petrov-Galerkin conditions imply  $\mathbf{r}_0 - \mathbf{A}\mathbf{z}_k \perp \mathcal{W}_k$ , or alternatively  $\langle \mathbf{r}_0 - \mathbf{A}\mathbf{z}_k, \mathbf{w} \rangle = 0$  for every  $\mathbf{w} \in \mathcal{W}_k$ . To impose these conditions it will help to have an appropriate basis for  $\mathcal{K}$  and  $\mathcal{W}$ . Suppose we have access to such a basis where  $\{\mathbf{q}_i\}_{i=1}^k$  and  $\{\mathbf{w}_i\}_{i=1}^k$  are basis elements for  $\mathcal{K}$  and  $\mathcal{W}$  respectively. Let

$$\mathbf{K}_k \triangleq [\mathbf{v}_1, \mathbf{v}_2, \dots, \mathbf{v}_k] \in \mathbb{K}^{n \times k}$$

$$\mathbf{W}_k \triangleq [\mathbf{w}_1, \mathbf{w}_2, \dots, \mathbf{w}_k] \in \mathbb{K}^{n \times k}$$

then the Petrov-Galerkin conditions can be imposed as follows

$$\mathbf{K}_k \mathbf{y}_k = \mathbf{z}_k, \quad \text{for some } \mathbf{y}_k \in \mathbb{K}^k$$

$$\mathbf{W}_k^\top (\mathbf{r}_0 - \mathbf{A}\mathbf{K}_k \mathbf{y}_k) = \mathbf{0}.$$

Moreover if  $\mathbf{W}_k^\top \mathbf{A}\mathbf{K}_k$  is invertible then  $\mathbf{x}_k$  can be expressed as

$$(11) \quad \mathbf{x}_k = \mathbf{x}_0 + \mathbf{K}_k (\mathbf{W}_k^\top \mathbf{A}\mathbf{K}_k)^{-1} \mathbf{W}_k^\top \mathbf{r}_0.$$

This justifies a general form of the projection method algorithm presented in algorithm 5.

---

**Algorithm 5:** General Projection Method

---

**output:** An approximation of  $\mathbf{x}^*$ ,  $\mathbf{x}_k$ .

**for**  $k = 1, \dots$  **until convergence do**

Select  $\mathcal{K}_k$  and  $\mathcal{W}_k$

Form  $\mathbf{K}_k$  and  $\mathbf{W}_k$

Solve  $(\mathbf{W}_k^\top \mathbf{A}\mathbf{K}_k) \mathbf{y}_k = \mathbf{W}_k^\top \mathbf{r}_0$

$\mathbf{x}_k = \mathbf{x}_0 + \mathbf{K}_k \mathbf{y}_k$

**end**

**return**  $\mathbf{x}_k$

---

**2.5. Conjugate Gradient Algorithm.** From section ?? that the Petrov-Galerkin conditions for the CG algorithm used an orthogonal projection and the matrix  $\mathbf{A}$  was assumed to be positive definite. To derive the CG algorithm we can start by using some machinery that the Lanczos algorithm provides us with. Recall, the Lanczos algorithm produces the form  $\mathbf{A}\mathbf{Q}_k = \mathbf{Q}_k \mathbf{T}_k + \mathbf{q}_{k+1} \mathbf{t}_k^\top$  where  $\mathbf{t}_k \triangleq [0, 0, \dots, 0, \beta_k]^\top \in \mathbb{K}^k$  and the columns of  $\mathbf{Q}_k$  span  $\mathcal{K}_k$ . Recall that  $\mathbf{x}_k$  can be expressed as  $\mathbf{x}_k = \mathbf{x}_0 + \mathbf{K}_k (\mathbf{W}_k^\top \mathbf{A}\mathbf{K}_k)^{-1} \mathbf{W}_k^\top \mathbf{r}_0$  (equation 11) when  $\mathbf{W}_k^\top \mathbf{A}\mathbf{K}_k$  is invertible. For the CG algorithm  $\mathcal{K} = \mathcal{W}$  and  $\mathbf{A} \succ \mathbf{0}$ . Under these

conditions we can easily show that  $\mathbf{W}_k^\top \mathbf{A} \mathbf{K}_k$  is indeed invertible. This means the approximate vector can be expressed as  $\mathbf{x}_k = \mathbf{x}_0 + \mathbf{z}_k$  where  $\mathbf{z}_k \in \mathcal{K}_k$ . In terms of the Petrov-Galerkin conditions this means that  $\mathbf{z}_k$  must satisfy  $\mathbf{r}_0 - \mathbf{A} \mathbf{z}_k \perp \mathcal{W}_k$ . Furthermore since  $\mathcal{K}_k = \text{Range}(\mathbf{Q}_k)$  where  $\mathbf{Q}_k$  has full column rank then  $\mathbf{z}_k$  can be represented as  $\mathbf{z}_k = \mathbf{Q}_k \mathbf{y}$  for a unique  $\mathbf{y} \in \mathbb{K}^k$  so that

$$(12) \quad \mathbf{x}_k = \mathbf{x}_0 + \mathbf{Q}_k \mathbf{y}.$$

Coupling this with the Petrov-Galerkin conditions means

$$(13) \quad \begin{aligned} \mathbf{Q}_k^\top (\mathbf{r}_0 - \mathbf{A} \mathbf{Q}_k \mathbf{y}) &= \mathbf{0} \\ \mathbf{Q}_k^\top \mathbf{A} \mathbf{Q}_k \mathbf{y} &= \mathbf{Q}_k^\top \mathbf{r}_0 \\ \mathbf{T}_k \mathbf{y} &= \|\mathbf{r}_0\| \mathbf{e}_1. \end{aligned}$$

In the CG algorithm  $\mathbf{x}_{k+1}$  is computed as the recurrence of the following three sets of vectors

- (1) The approximate solutions  $\mathbf{x}_k$
- (2) The residual vectors  $\mathbf{r}_k$
- (3) The conjugate gradient vectors  $\mathbf{p}_k$

The conjugate gradient vectors are given the name gradient since the attempt to find the direction of steepest descent that minimizes  $\|\mathbf{r}_k\|_{\mathbf{A}^{-1}}$ . They are also given the name conjugate since  $\langle \mathbf{p}_k, \mathbf{A} \mathbf{p}_j \rangle = 0$  for  $i \neq j$ , that is, vectors  $\mathbf{p}_i$  and  $\mathbf{p}_j$  are mutually  $\mathbf{A}$ -conjugate.

Since  $\mathbf{A}$  is symmetric positive definite then so is  $\mathbf{T}_k = \mathbf{Q}_k \mathbf{A} \mathbf{Q}_k$ . We can take the Cholesky decomposition of  $\mathbf{T}_k$  to get

$$(14) \quad \mathbf{T}_k = \mathbf{L}_k \mathbf{D}_k \mathbf{L}_k^\top$$

where  $\mathbf{L}_k$  is a unit lower bidiagonal matrix and  $\mathbf{D}_k$  is diagonal written as

$$\mathbf{L}_k = \begin{bmatrix} 1 & & & \\ l_1 & \ddots & & \\ & \ddots & \ddots & \\ & & l_{k-1} & 1 \end{bmatrix}, \quad \mathbf{D}_k = \begin{bmatrix} d_1 & & & \\ & d_2 & & \\ & & \ddots & \\ & & & d_k \end{bmatrix}.$$

Combining equations 12, 13 and 14

$$\begin{aligned} \mathbf{x}_k &= \mathbf{x}_0 + \mathbf{Q}_k \mathbf{y} \\ \mathbf{x}_k &= \mathbf{x}_0 + \|\mathbf{r}_0\| \mathbf{Q}_k \mathbf{T}_k^{-1} \mathbf{e}_1 \\ \mathbf{x}_k &= \mathbf{x}_0 + \|\mathbf{r}_0\| \mathbf{Q}_k (\mathbf{L}_k \mathbf{D}_k \mathbf{L}_k^\top)^{-1} \mathbf{e}_1 \\ \mathbf{x}_k &= \mathbf{x}_0 + \left( \mathbf{Q}_k \mathbf{L}_k^{-\top} \right) (\|\mathbf{r}_0\| \mathbf{D}_k^{-1} \mathbf{L}_k^{-1} \mathbf{e}_1) \\ \mathbf{x}_k &\triangleq \mathbf{x}_0 + \tilde{\mathbf{P}}_k \tilde{\mathbf{y}}_k \end{aligned}$$

where  $\tilde{\mathbf{P}}_k = \mathbf{Q}_k \mathbf{L}_k^{-\top}$  and  $\tilde{\mathbf{y}}_k = \|\mathbf{r}_0\| \mathbf{D}_k^{-1} \mathbf{L}_k^{-1} \mathbf{e}_1$ . The matrix  $\tilde{\mathbf{P}}_k$  can be written as  $\tilde{\mathbf{P}}_k = [\tilde{\mathbf{p}}_1, \tilde{\mathbf{p}}_2, \dots, \tilde{\mathbf{p}}_k]$ . Lemma 11 shows that the columns of  $\tilde{\mathbf{P}}_k$  are  $\mathbf{A}$ -conjugate.

**Lemma 11.** *The columns of  $\tilde{\mathbf{P}}_k$  are  $\mathbf{A}$ -conjugate, in otherwise  $\tilde{\mathbf{P}}_k^\top \mathbf{A} \tilde{\mathbf{P}}_k$  is diagonal.*

*Proof.* We compute

$$\begin{aligned}
 \tilde{P}_k^\top A \tilde{P}_k &= (Q_k L_k^{-\top})^\top A (Q_k L_k^{-\top}) \\
 &= L_k^{-1} (Q_k^\top A Q_k) L_k^{-\top} \\
 &= L_k^{-1} (T_k) L_k^{-\top} \\
 &= L_k^{-1} (L_k D_k L_k^\top) L_k^{-\top} \\
 &= D_k
 \end{aligned}$$

(equation 14)

as wanted.  $\square$

Since  $L_k$  is a lower bidiagonal, setting  $\mathbf{a} \triangleq l_{k-1} \mathbf{e}_{k-1}$ , it can be written in the form

$$L_k = \begin{bmatrix} L_{k-1} & \mathbf{0} \\ \mathbf{a}^\top & 1 \end{bmatrix}$$

meaning

$$L_k^{-1} = \begin{bmatrix} L_{k-1}^{-1} & \mathbf{0} \\ \star & 1 \end{bmatrix}.$$

With this a recurrence for the columns of  $\tilde{P}_k$  can now be derived in terms of  $\mathbf{y}_k$ . To start we can show that the first  $k-1$  entries of  $\tilde{\mathbf{y}}_k$  shares the first  $k-1$  entries with  $\tilde{\mathbf{y}}_{k-1}$  and that  $\tilde{P}_k$  and  $\tilde{P}_{k-1}$  share the same first  $k-1$  columns. To start we can compute a recurrence for  $\tilde{\mathbf{y}}_k$  as follows

$$\begin{aligned}
 \tilde{\mathbf{y}}_k &= \|\mathbf{r}_0\| D_k^{-1} L_k^{-1} \mathbf{e}_1^k \\
 &= \|\mathbf{r}_0\| \begin{bmatrix} D_{k-1}^{-1} & \mathbf{0} \\ \mathbf{0} & d_k^{-1} \end{bmatrix} \begin{bmatrix} L_{k-1}^{-1} & \mathbf{0} \\ \star & 1 \end{bmatrix} \mathbf{e}_1^k \\
 &= \|\mathbf{r}_0\| \begin{bmatrix} D_{k-1}^{-1} L_{k-1}^{-1} & \mathbf{0} \\ \star & d_k^{-1} \end{bmatrix} \begin{bmatrix} \mathbf{e}_1^k \\ 0 \end{bmatrix} \\
 &= \begin{bmatrix} \tilde{\mathbf{y}}_{k-1} \\ \eta_k \end{bmatrix}
 \end{aligned}$$

To get a recurrence for the columns of  $\tilde{P}_{k-1} = [\tilde{\mathbf{p}}_1, \tilde{\mathbf{p}}_2, \dots, \tilde{\mathbf{p}}_k]$  since  $L_{k-1}^\top$  is upper triangular then so is  $L_{k-1}^{-\top}$ , thus forming the leading  $(k-1)$ -by- $(k-1)$  submatrix of  $L_k^{-\top}$ . This means that  $\tilde{P}_{k-1}$  is identical to the leading  $k-1$  columns of

$$\tilde{P}_k = Q_k L_k^{-\top} = [Q_{k-1}, \mathbf{q}_k] \begin{bmatrix} L_{k-1}^{-1} & \mathbf{0} \\ \star & 1 \end{bmatrix} = [Q_{k-1} L_{k-1}^{-1}, \tilde{\mathbf{p}}_k] = [\tilde{P}_{k-1}, \tilde{\mathbf{p}}_k].$$

Moreover rearranging  $\tilde{P}_k = Q_k L_k^{-\top}$  we get  $\tilde{P}_k L_k^\top = Q_k$ . Equating the  $k^{th}$  column yields

$$(15) \quad \tilde{\mathbf{p}}_k = \mathbf{q}_k - l_{k-1} \tilde{\mathbf{p}}_{k-1}.$$

Finally we can use

$$(16) \quad \mathbf{x}_k = \mathbf{x}_0 + \tilde{P}_k \tilde{\mathbf{y}}_k = \mathbf{x}_0 + [\tilde{P}_{k-1}, \tilde{\mathbf{p}}_k] \begin{bmatrix} \tilde{\mathbf{y}}_{k-1} \\ \eta_k \end{bmatrix} = \mathbf{x}_0 + \tilde{P}_{k-1} \tilde{\mathbf{y}}_{k-1} + \eta_k \tilde{\mathbf{p}}_k = \mathbf{x}_{k-1} + \eta_k \tilde{\mathbf{p}}_k$$



as a recurrence for  $\mathbf{x}_k$ . A recurrence for  $\mathbf{r}_k$  is easily computed as

$$(17) \quad \mathbf{r}_k = b - \mathbf{A}\mathbf{x}_k = b - \mathbf{A}(\mathbf{x}_{k-1} + \eta_k \tilde{\mathbf{p}}_k) = (b - \mathbf{A}\mathbf{x}_{k-1}) - \eta_k \mathbf{A}\tilde{\mathbf{p}}_k = \mathbf{r}_{k-1} - \eta_k \mathbf{A}\tilde{\mathbf{p}}_k$$

Altogether we are left with recurrences for  $\mathbf{q}_k$  from Lanczos,  $\tilde{\mathbf{p}}_k$  (equation 15), the residual  $\mathbf{r}_k$  (equation 15), and for the approximate solution  $\mathbf{x}_k$  (equation 16). However, further simplification can be made for a more efficient algorithm. Recall from section 2.3 that  $\mathbf{A}\mathbf{Q}_k = \mathbf{Q}_k \mathbf{T}_k + \mathbf{q}_{k+1} \mathbf{t}_k^\top$  where  $\mathbf{t}_k = [0, 0, \dots, 0, \beta_k]^\top \in \mathbb{K}^k$  meaning

$$\mathbf{r}_k = \mathbf{r}_0 - \mathbf{A}\mathbf{Q}_k \mathbf{y}_k = \mathbf{r}_0 - \mathbf{Q}_k \mathbf{T}_k \mathbf{y}_k - \langle \mathbf{t}_k, \mathbf{y} \rangle \mathbf{q}_{k+1} = -\beta_k y_k \mathbf{q}_{k+1}.$$

This tells us that  $\mathbf{r}_k$  is parallel to  $\mathbf{q}_{k+1}$  and orthogonal to all  $\mathbf{q}_i$ ,  $1 \leq i \leq k$ . This further implies that  $\mathbf{r}_k$  is orthogonal to all  $\mathbf{r}_i$ ,  $1 \leq i \leq k-1$  since they are just  $\mathbf{q}_i$  scaled by some constant factor. So replacing  $\mathbf{r}_{k-1}$  with  $\mathbf{q}_k/\eta_k$  and defining  $\mathbf{p}_k \triangleq \tilde{\mathbf{p}}_k/\gamma_k$  gives us a new set of recurrences

$$\begin{aligned} \mathbf{x}_k &= \mathbf{x}_{k-1} + \alpha_k \mathbf{p}_k \\ \mathbf{r}_k &= \mathbf{r}_{k-1} - \alpha_k \mathbf{A}\mathbf{p}_k \\ \mathbf{p}_k &= \mathbf{r}_{k-1} + \beta_k \mathbf{p}_{k-1} \end{aligned}$$

where  $\alpha_k = \eta_k/\gamma_k$ . From theorem 11 we have shown that the columns of  $\tilde{\mathbf{P}}_k$  are  $A$ -conjugate (that is  $\langle \tilde{\mathbf{p}}_i, \mathbf{A}\tilde{\mathbf{p}}_j \rangle = 0$ ,  $i \neq j$ ) and that  $\tilde{\mathbf{P}}_k^\top \mathbf{A} \tilde{\mathbf{P}}_k = \mathbf{D}_k$ . This also means that  $\langle \mathbf{r}_i, \mathbf{r}_j \rangle = 0$ ,  $i \neq j$ . Now note that from our recurrence for  $\mathbf{p}_k = \mathbf{r}_{k-1} + \beta_k \mathbf{p}_{k-1}$  that

$$\langle \mathbf{A}\mathbf{p}_k, \mathbf{p}_k \rangle = \langle \mathbf{A}\mathbf{p}_k, \mathbf{r}_{k-1} + \beta_k \mathbf{p}_{k-1} \rangle = \langle \mathbf{A}\mathbf{p}_k, \mathbf{r}_{k-1} \rangle.$$

We can now find an expression for  $\alpha_k$  as

$$\begin{aligned} \langle \mathbf{r}_{k-1}, \mathbf{r}_k \rangle &= \langle \mathbf{r}_{k-1}, \mathbf{r}_{k-1} - \alpha_k \mathbf{A}\mathbf{p}_k \rangle \\ \langle \mathbf{r}_{k-1}, \mathbf{r}_k \rangle &= \langle \mathbf{r}_{k-1}, \mathbf{r}_{k-1} \rangle - \alpha_k \langle \mathbf{p}_k, \mathbf{A}\mathbf{p}_k \rangle \\ \alpha_k &= \frac{\langle \mathbf{r}_{k-1}, \mathbf{r}_{k-1} \rangle}{\langle \mathbf{p}_k, \mathbf{A}\mathbf{p}_k \rangle}. \end{aligned}$$

Similarly, using the recurrence for  $\mathbf{p}_k$ , an expression for  $\beta_k$  can be computed as

$$\begin{aligned} \langle \mathbf{A}\mathbf{p}_{k-1}, \mathbf{p}_k \rangle &= \langle \mathbf{A}\mathbf{p}_{k-1}, \mathbf{r}_{k-1} + \beta_k \mathbf{p}_{k-1} \rangle \\ \langle \mathbf{A}\mathbf{p}_{k-1}, \mathbf{p}_k \rangle &= \langle \mathbf{A}\mathbf{p}_{k-1}, \mathbf{r}_{k-1} \rangle + \beta_k \langle \mathbf{A}\mathbf{p}_{k-1}, \mathbf{p}_{k-1} \rangle \\ \beta_k &= -\frac{\langle \mathbf{A}\mathbf{p}_{k-1}, \mathbf{r}_{k-1} \rangle}{\langle \mathbf{A}\mathbf{p}_{k-1}, \mathbf{p}_{k-1} \rangle} \end{aligned}$$

This formula requires an additional dot product which was not present before. Fortunately, this dot product can be eliminated using our recurrence for  $\mathbf{r}_k$

$$\begin{aligned} \langle \mathbf{r}_k, \mathbf{r}_k \rangle &= \langle \mathbf{r}_k, \mathbf{r}_{k-1} - \alpha_k \mathbf{A}\mathbf{p}_k \rangle \\ \langle \mathbf{r}_k, \mathbf{r}_k \rangle &= \langle \mathbf{r}_k, \mathbf{r}_{k-1} \rangle - \alpha_k \langle \mathbf{r}_k, \mathbf{A}\mathbf{p}_k \rangle \\ \alpha_k &= -\frac{\langle \mathbf{r}_k, \mathbf{r}_k \rangle}{\langle \mathbf{r}_k, \mathbf{A}\mathbf{p}_k \rangle}. \end{aligned}$$

Equating the two expressions for  $\alpha_k$  yields

$$-\frac{\langle \mathbf{r}_k, \mathbf{r}_k \rangle}{\langle \mathbf{r}_k, \mathbf{A}\mathbf{p}_k \rangle} = \frac{\langle \mathbf{r}_{k-1}, \mathbf{r}_{k-1} \rangle}{\langle \mathbf{p}_k, \mathbf{A}\mathbf{p}_k \rangle}$$

$$-\frac{\langle \mathbf{r}_k, \mathbf{r}_k \rangle}{\langle \mathbf{r}_{k-1}, \mathbf{r}_{k-1} \rangle} = \frac{\langle \mathbf{r}_k, \mathbf{A}\mathbf{p}_k \rangle}{\langle \mathbf{p}_k, \mathbf{A}\mathbf{p}_k \rangle}.$$

This means that

$$\beta_k = \frac{\langle \mathbf{r}_{k-1}, \mathbf{r}_{k-1} \rangle}{\langle \mathbf{r}_{k-2}, \mathbf{r}_{k-2} \rangle}.$$

These recurrences are computed iteratively to form the basis of the CG algorithm, seen in Algorithm 6.

---

**Algorithm 6:** CG Algorithm

---

**input :**  $\mathbf{A} \succ \mathbf{0}$ ,  $\mathbf{b}$  and an initial guess  $\mathbf{x}_0$ .

**output:** An approximation of  $\mathbf{x}^*$ ,  $\mathbf{x}_k$ .

$\mathbf{r}_0 = \mathbf{b} - \mathbf{A}\mathbf{x}_0$ ,  $\mathbf{p}_1 = \mathbf{r}_0$

**for**  $k = 1, \dots$  **until**  $\|\mathbf{r}_{k-1}\| \leq \tau$  **do**

$$\alpha_k = \frac{\langle \mathbf{r}_{k-1}, \mathbf{r}_{k-1} \rangle}{\langle \mathbf{p}_k, \mathbf{A}\mathbf{p}_k \rangle}$$

$$\mathbf{x}_k = \mathbf{x}_{k-1} + \alpha_k \mathbf{p}_k$$

$$\mathbf{r}_k = \mathbf{r}_{k-1} - \alpha_k \mathbf{A}\mathbf{p}_k$$

$$\beta_{k+1} = \frac{\langle \mathbf{r}_k, \mathbf{r}_k \rangle}{\langle \mathbf{r}_{k-1}, \mathbf{r}_{k-1} \rangle}$$

$$\mathbf{p}_{k+1} = \mathbf{r}_k + \beta_{k+1} \mathbf{p}_k$$

**end**

**return**  $\mathbf{x}_k$

---

### 3. TWISTED HECKE ALGEBRAS OF FINITE GROUPS

We have now completed our investigation of the Hecke algebra  $\mathcal{H}(G, K)$ . The aim of this section is to generalise the results of Chapter 1 to the case of a non-trivial character  $\sigma: K \rightarrow \mathbb{C}^\times$ . Here the Hecke algebra  $\mathcal{H} = \mathcal{H}(G, K, \sigma)$  is the convolution algebra of  $(K, \sigma)$ -bi-invariant functions on  $G$ . In Section 3.1, we discuss the theory of the induced representation  $\text{Ind}_K^G \sigma$ . In Section 3.2, we revisit the Hecke algebra, identify its identity and describe its basis. Notice that the results of Section ?? and Section ?? were independent of the choice  $\sigma = 1$ , so they still apply now that we are considering a non-trivial character.

In Section 3.3, we generalise Gelfand's Trick from Section ?? to the case of a non-trivial character  $\sigma$ . Naturally, we will need to reconsider the conditions that the anti-automorphism  $\varphi: G \rightarrow G$  must satisfy. As in Section ??, we will investigate these conditions and conclude with a natural statement and proof of Gelfand's Trick in the twisted case. We conclude with Section 3.4, in which we investigate the Gelfand–Graev representation and use the results of this chapter to prove that it is multiplicity-free.

**3.1. The induced representation  $\text{Ind}_K^G \sigma$ .** Suppose that  $\sigma: K \rightarrow \mathbb{C}^\times$  is a character, i.e. a group homomorphism. Consider the space

$$W := \{f: G \rightarrow \mathbb{C} \mid f(gk) = f(g)\sigma(k), \forall g \in G, \forall k \in K\} \subseteq \text{Fun}(G).$$

As in the previous section,  $W$  is called the induced representation and denoted  $\text{Ind}_K^G \sigma$ . We state and prove a lemma analogous to Lemma ??.

**Lemma 12.**  *$W$  is a left ideal of  $(\text{Fun}(G), \star)$ .*

*Proof.* We verify that  $f \star w \in W$  whenever  $w \in W$  and  $f \in \text{Fun}(G)$ . Let  $g \in G$  and  $k \in K$ . Then

$$\begin{aligned} (f \star w)(gk) &= \sum_{xy=gk} f(x)w(y) = \sum_{x \in G} f(x)w(x^{-1}gk) = \sum_{x \in G} f(x)w(x^{-1}g)\sigma(k) \\ &= \left[ \sum_{x \in G} f(x)w(x^{-1}g) \right] \sigma(k) = \left[ \sum_{xy=g} f(x)w(y) \right] \sigma(k) = (f \star w)(g)\sigma(k). \quad \square \end{aligned}$$

**3.2. The twisted Hecke algebra of a finite group  $\mathcal{H}(G, K, \sigma)$ .** The Hecke algebra  $\mathcal{H} = \mathcal{H}(G, K, \sigma)$  is the space

$$\mathcal{H} := \{f: G \rightarrow \mathbb{C} \mid f(k_1 g k_2) = \sigma(k_1) f(g) \sigma(k_2), \forall g \in G, \forall k_1, k_2 \in K\} \subseteq \text{Fun}(G).$$

The proof of Lemma 12 can be adapted to show that  $\mathcal{H}$  is a two-sided ideal in  $(\text{Fun}(G), \star)$ . As before, the identity of  $(\text{Fun}(G), \star)$  does not lie in  $\mathcal{H}$ . Nevertheless,  $\mathcal{H}$  does have an identity of its own. It is easy to verify that the identity is  $\iota_K^\sigma$ , which we define below.

$$\iota_K^\sigma: G \rightarrow \mathbb{C}, \quad \iota_K^\sigma(g) := \begin{cases} \frac{1}{|K|} \sigma(g), & \text{if } g \in K, \\ 0, & \text{else.} \end{cases}$$

Thus,  $(\mathcal{H}, \star)$  is a unital associative algebra in its own right.

We now construct a basis for  $\mathcal{H}$ . Recall that when  $\sigma = 1$ , the basis of  $\mathcal{H}$  was described by the characteristic functions of  $K$ -double cosets. To treat the case when  $\sigma \neq 1$ , we need a lemma about group actions.

Consider the finite group  $K$  acting on a set  $X$ . For each  $x \in X$ , let  $\mathcal{O}_x := \{g \cdot x \mid g \in G\}$  be the orbit containing  $x$  and let  $K_x := \{k \in K \mid k \cdot x = x\}$  be the stabiliser subgroup of  $x$  in  $K$  (also denoted as  $\text{stab}_K(x)$ ). Consider the vector space

$$V := \{f: X \rightarrow \mathbb{C} \mid f(k \cdot x) = \sigma(k)f(x), \forall k \in K, \forall x \in X\} \subseteq \text{Fun}(X).$$

An orbit  $\mathcal{O}_x$  is called  $(K, \sigma)$ -*relevant* if there exists  $f \in V$  such that  $f|_{\mathcal{O}_x}$  is non-zero. Otherwise, we say  $\mathcal{O}_x$  is  $(K, \sigma)$ -*irrelevant*. We omit mention of  $(K, \sigma)$  if it is clear from the context.

**Lemma 13.** *An orbit  $\mathcal{O}_x$  is  $(K, \sigma)$ -relevant if and only if  $\sigma(K_x) = \{1\}$ .*

*Proof.* Assume that  $\sigma(K_x) \neq \{1\}$ . Then there exists  $k \in K_x$  such that  $\sigma(k) \neq 1$ . Now recall that for  $f \in V$  we have  $f(x) = f(k \cdot x) = \sigma(k)f(x)$ . However  $\sigma(k) \neq 1$ , so  $f(x) = 0$ . Then  $f$  must be zero on  $\mathcal{O}_x$  and  $\mathcal{O}_x$  is irrelevant.

Conversely, the fact that  $\sigma$  is trivial on  $K_x$  implies that it factors through a well-defined function  $\sigma_x: \mathcal{O}_x \simeq K/K_x \rightarrow \mathbb{C}$  given by  $\sigma_x(kK_x) := \sigma(k)$ . To see that this function is well-defined, suppose that  $k_1K_x = k_2K_x$ . Then  $k_1k_2^{-1} \in K_x$ . Since  $\sigma$  is trivial on  $K_x$ , we know  $1 = \sigma(k_1k_2^{-1}) = \sigma(k_1)\sigma(k_2)^{-1}$ . Then  $\sigma(k_1) = \sigma(k_2)$  so  $\sigma_x(k_1K_x) = \sigma_x(k_2K_x)$ . It is easy to check that  $\sigma_x \in V$ . Thus,  $\mathcal{O}_x$  is relevant.  $\square$

Now suppose  $K$  is a subgroup of  $G$  acting on  $X = G$  from the left and right by translation. Then the orbit  $\mathcal{O}_x$  is nothing but the double coset  $KxK$  and  $V$  becomes the Hecke algebra  $\mathcal{H}$ . Explicitly, we have

$$\mathcal{H} = \{f: X \rightarrow \mathbb{C} \mid f(k_1 \cdot x \cdot k_2) = \sigma(k_1)f(x)\sigma(k_2), \forall k_1, k_2 \in K, \forall x \in X\} \subseteq \text{Fun}(X).$$

We can re-write this data by considering the left action of  $K \times K^{\text{op}}$  on  $X$ , where  $K^{\text{op}}$  is the group opposite to  $K$ . Then

$$\mathcal{H} = \{f: X \rightarrow \mathbb{C} \mid f((k_1, k_2) \cdot x) = \sigma(k_1)\sigma(k_2)f(x), \forall k_1, k_2 \in K, \forall x \in X\} \subseteq \text{Fun}(X).$$

A double coset  $KxK$  is relevant if it supports a non-zero function from  $\mathcal{H}$ . Let  $X_{\text{rel}}$  be a family of relevant coset representatives. Define the family of functions  $\{\chi_x\}_{x \in X_{\text{rel}}}$  by

$$\chi_x(y) := \begin{cases} \sigma(k)\sigma(k'), & \text{if } y \in KxK \text{ with } y = kxk', \\ 0, & \text{if } y \notin KxK. \end{cases}$$

One easily checks that  $\chi_x$  is well-defined. We call  $\chi_x$  a *twisted characteristic function* associated to the relevant orbit  $KxK$ . When  $\sigma = 1$ , every orbit is relevant and  $\sigma(k)\sigma(k') = 1$ , so we obtain the original characteristic functions described in Section ?? . Define the map

$$\sigma \boxtimes \sigma: K \times K \rightarrow \mathbb{C}^\times \times \mathbb{C}^\times, \quad (\sigma \boxtimes \sigma)(k_1, k_2) := (\sigma(k_1), \sigma(k_2)).$$

As a result of Lemma 13, we see that an orbit under the left action of  $K \times K^{\text{op}}$  is relevant if and only if  $(\sigma \boxtimes \sigma)(\text{stab}_{K \times K}(x)) = \{1\}$ . As in Section ??, it is not difficult to see that the twisted characteristic functions of relevant orbits form a basis of  $\mathcal{H}(G, K, \sigma)$ .

**3.3. Twisted Gelfand's Trick.** Our goal in this section is to prove the twisted analogue of Gelfand's Trick.

**Theorem 14** (Twisted Gelfand's Trick). *Suppose that  $G$  is a finite group with  $K \leq G$  as a subgroup and character  $\sigma: K \rightarrow \mathbb{C}^\times$ . Let  $\varphi: G \rightarrow G$  be an anti-automorphism such that*

- (i)  $\varphi^2 = 1$ ,
- (ii)  $\varphi(K) = K$ ,
- (iii)  $\sigma(\varphi(k)) = \sigma(k)$  for all  $k \in K$ , and
- (iv)  $\varphi(x) = x$  for all  $x \in X_{\text{rel}}$ , a family of representatives for the  $(K, \sigma)$ -relevant  $K$ -double cosets.

Then  $\mathcal{H}(G, K, \sigma)$  is commutative.

This is a true generalisation of Theorem ???. Indeed, if we consider the trivial representation  $\sigma = \mathbf{1}$ , condition (iii) is trivially satisfied, condition (iv) corresponds to the requirement that  $K\varphi(x)K = KxK$  in Theorem ??, and condition (ii) is contained in the requirement that  $K\varphi(x)K = KxK$ .

As in Section 3.1, the proof of Theorem 14 relies on the observation that an anti-homomorphism of an algebra that acts as the identity on basis elements of the subalgebra is sufficient to conclude that the subalgebra is commutative (c.f. Lemma ?? and Corollary ??). This leaves us with a question: can we rewrite the condition  $\varphi^* \chi_x = \chi_x$ ?

Recall that  $X_{\text{rel}}$  denotes a family of representatives for the relevant double cosets. Recall the twisted characteristic functions  $\{\chi_x\}_{x \in X_{\text{rel}}}$  defined in Section 3.2 given by

$$\chi_x(y) = \begin{cases} \sigma(k)\sigma(k'), & \text{if } y \in KxK \text{ with } y = kxk', \\ 0, & \text{if } y \notin KxK. \end{cases}$$

Thus,

$$(\varphi^* \chi_x)(g) = \begin{cases} \sigma(k)\sigma(k'), & \text{if } \varphi(g) \in KxK \text{ with } \varphi(g) = kxk', \\ 0, & \text{else.} \end{cases}$$

If  $\varphi: G \rightarrow G$  is an involutive homomorphism, then  $\varphi(g) = kxk'$  is equivalent to  $g = \varphi(k')\varphi(x)\varphi(k)$ . If we further suppose that  $\varphi(x) = x$  for all  $x \in X_{\text{rel}}$  and  $\varphi(K) = K$ , then  $g = \varphi(k')\varphi(x)\varphi(k)$  is equivalent to  $g = \varphi(k')x\varphi(k)$ . Thus,

$$(\varphi^* \chi_x)(g) = \begin{cases} \sigma(\varphi(k'))\sigma(\varphi(k)), & \text{if } g \in KxK \text{ with } g = \varphi(k')x\varphi(k), \\ 0, & \text{else.} \end{cases}$$

This tells us that  $\varphi^* \chi_x$  is also supported (i.e. non-zero) on  $KxK$ . Now let's also assume that  $\sigma(\varphi(k)) = \sigma(k)$  for all  $k \in K$ . Then we can easily verify that  $\varphi^* \chi_x \in \mathcal{H}(G, K, \sigma)$ . So  $\varphi^* \chi_x$  must be a multiple of  $\chi_x$ . In fact, this multiple is 1, since

$$(\varphi^* \chi_x)(x) = \chi_x(\varphi(x)) = \chi_x(x) = 1.$$

We are now ready to prove Theorem 14.

*Proof of Theorem 14.* Suppose that  $\varphi: G \rightarrow G$  is an anti-automorphism. Also suppose that  $\varphi^2 = 1$ ,  $\varphi(K) = K$ ,  $\sigma(\varphi(k)) = \sigma(k)$  for all  $k \in K$ , and  $\varphi(x) = x$  for all  $x \in X_{\text{rel}}$ . The above discussion tells us that

$\varphi^* \chi_x = \chi_x$ . These are the basis elements of  $\mathcal{H}(G, K, \sigma)$ . We apply Corollary ?? to conclude that  $\mathcal{H}(G, K, \sigma)$  is commutative.  $\square$

**3.4. The Gelfand–Graev representation.** We construct the *Gelfand–Graev representation* of  $G = \mathrm{GL}_n(\mathbb{F}_q)$ . First, consider the *unipotent radical* of  $G$ , given by

$$U(\mathbb{F}_q) := \begin{pmatrix} 1 & \mathbb{F}_q & \cdots & \mathbb{F}_q \\ 0 & 1 & \ddots & \vdots \\ \vdots & \ddots & 1 & \mathbb{F}_q \\ 0 & \cdots & 0 & 1 \end{pmatrix}.$$

Next, fix a non-trivial additive character  $\psi: \mathbb{F}_q \rightarrow \mathbb{C}^\times$  (i.e.  $\psi(a+b) = \psi(a)\psi(b)$ ). Then define a character  $\pi: U(\mathbb{F}_q) \rightarrow \mathbb{C}^\times$  by

$$\pi(x) := \psi(x_{12} + x_{23} + \cdots + x_{n-1,n}).$$

To see that  $\pi$  is a character, observe

$$\begin{aligned} \pi(xy) &= \psi((xy)_{12} + (xy)_{23} + \cdots + (xy)_{n-1,n}) \\ &= \psi\left(\sum_{k=1}^{n-1} x_{1k}y_{k2} + \sum_{k=1}^{n-1} x_{2k}y_{k3} + \cdots + \sum_{k=1}^{n-1} x_{n-1,k}y_{kn}\right) \\ &= \psi\left(\sum_{k=1}^{n-1} x_{1k}y_{k2}\right) \psi\left(\sum_{k=1}^{n-1} x_{2k}y_{k3}\right) \cdots \psi\left(\sum_{k=1}^{n-1} x_{n-1,k}y_{kn}\right) \\ &= \psi(x_{12} + y_{12})\psi(x_{23} + y_{23}) \cdots \psi(x_{n-1,n} + y_{n-1,n}) \\ &= \psi(x_{12})\psi(y_{12})\psi(x_{23})\psi(y_{23}) \cdots \psi(x_{n-1,n})\psi(y_{n-1,n}) \\ &= \psi(x_{12})\psi(x_{23}) \cdots \psi(x_{n-1,n})\psi(y_{12})\psi(y_{23}) \cdots \psi(y_{n-1,n}) \\ &= \psi(x_{12} + \cdots + x_{n-1,n})\psi(y_{12} + \cdots + y_{n-1,n}) \\ &= \pi(x)\pi(y). \end{aligned}$$

The *Gelfand–Graev representation* of  $G$  is  $\mathrm{Ind}_U^G \pi$ . In [Bum13], Bump explains, “this Gelfand–Graev representation is important because it contains *most* irreducible representations of the group; those it contains are therefore called *generic*.” Furthermore, we have the following theorem.

**Theorem 15.** *The Gelfand–Graev representation is multiplicity-free. That is,  $(\mathrm{GL}_n(\mathbb{F}_q), U(\mathbb{F}_q), \pi)$  is a twisted Gelfand pair.*

This theorem will be proven in two parts. We begin with a lemma.

**Lemma 16.** (i) *We have the Bruhat decomposition*

$$\mathrm{GL}_n(\mathbb{F}_q) = \bigsqcup_{w \in W} BwB,$$

where  $W$  is the group of all  $n \times n$  permutation matrices and  $B$  is the subgroup of all  $n \times n$  upper-triangular matrices.

(ii) We can modify the Bruhat decomposition and write

$$\mathrm{GL}_n(\mathbb{F}_q) = \bigsqcup_{m \in M} UmU,$$

where  $M$  is the group of all  $n \times n$  monomial matrices. A monomial matrix is a matrix with exactly one non-zero element in each row and column.

Before we prove Lemma 16, we recall a simple fact about matrices. Define  $x_{ij}(t) := I_{n \times n} + tE_{ij}$ , where  $1 \leq i \leq j \leq n$  and  $E_{ij}$  is the matrix of 0's except for a 1 in the  $i^{\text{th}}$  row and  $j^{\text{th}}$  column. Notice that  $x_{ij}(t) \in B$  since  $i \leq j$ . We can achieve the usual row and column operations on a matrix  $A$  by multiplying on the left or the right by some  $x_{ij}(t)$ . The following makes this statement precise.

Right-multiplying  $A$  by  $x_{ij}(t)$  corresponds to the column operation of  $C_j \mapsto C_j + tC_i$ , where  $C_k$  is column  $k$  of  $A$ . Similarly, left-multiplying  $A$  by  $x_{ij}(t)$  corresponds to the row operation of  $R_i \mapsto R_i + tR_j$ , where  $R_k$  is row  $k$  of  $A$ . Right-multiplying  $A$  by  $x_{ii}(\lambda - 1)$  corresponds to the column operation  $C_i \mapsto \lambda C_i$ , for some scalar  $\lambda$ . Similarly, left-multiplying  $A$  by  $x_{ii}(\lambda - 1)$  corresponds to the row operation  $R_i \mapsto \lambda R_i$ . We see that we can perform the usual row and column operations by right- and left-multiplying by elements of  $B$ .

*Proof of Lemma 16.* We begin by proving that  $\mathrm{GL}_n(\mathbb{F}_q) = \bigcup_{w \in W} BwB$  and will prove disjointness of the union later. We proceed by induction. The  $n = 1$  case is clearly true since all matrices in  $\mathrm{GL}_1(\mathbb{F}_q)$  are upper-triangular. Now let  $n > 1$  and  $g \in \mathrm{GL}_n(\mathbb{F}_q)$ . We wish to find a permutation matrix  $w$  in  $BgB$ . We have two cases:  $g_{n,1} \neq 0$  and  $g_{n,1} = 0$ .

In the first case, the previous discussion tells us that we can multiply  $g$  on the left and the right by appropriate elements of  $B$  so that the resulting matrix has zeros in the left column and bottom row, except for the bottom left entry, which is  $g_{n,1}$ . This is non-zero so we can normalise this resulting matrix by  $g_{n,1}$  to yield  $\begin{pmatrix} 0 & g' \\ 1 & 0 \end{pmatrix}$ . Here  $g'$  lies in  $\mathrm{GL}_{n-1}(\mathbb{F}_q)$ . The inductive hypothesis that the  $n - 1$  is true tells us that  $g'$  lies in a double coset  $Bw'B$  for some  $(n - 1) \times (n - 1)$  permutation matrix  $w'$ . Then the desired  $w$  is obtained by setting  $w = \begin{pmatrix} 0 & w' \\ 1 & 0 \end{pmatrix}$ .

In the second case, choose  $g_{i1} \neq 0$  and  $g_{nj} \neq 0$  so that  $i$  is as large as possible and  $j$  is as small as possible. This amounts to choosing the two non-zero entries in the left column and bottom row that are closest to the bottom left entry. Left- and right-multiplication by appropriate elements of  $B$  yields a matrix whose first and  $j^{\text{th}}$  columns and  $i^{\text{th}}$  and last rows are empty, except the entries  $g_{i1}$  and  $g_{nj}$ . Since these entries are non-zero, we can normalise these to 1 as well. Now we apply the inductive hypothesis to the matrix obtained by removing these two rows and two columns. We are left with a permutation matrix and this completes the induction.

We verify that the union is disjoint. Let  $w_1, w_2 \in W$  be representatives for the same double coset. Then  $Bw_1B = Bw_2B$  and, given any  $b \in B$ , there exists  $b' \in B$  with  $w_1bw_2^{-1} = b'$ . In particular,  $w_1w_2^{-1} \in B \cap W = \{1\}$ . Thus  $w_1 = w_2$ .

We now prove the modified decomposition. Consider the subgroup  $T$  of diagonal matrices in  $\mathrm{GL}_n(\mathbb{F}_q)$ . Notice that  $B = TU = UT$  and  $M = TW = WT$  so the result follows from the regular Bruhat decomposition. Disjointness is proven as before.  $\square$

*Proof of Theorem 15.* Consider the involutive anti-automorphism  $\varphi: G \rightarrow G$  defined by

$$\varphi(g) := w_0 g^t w_0, \quad \text{where } w_0 = \begin{pmatrix} & & 1 \\ & \ddots & \\ 1 & & \end{pmatrix}.$$

We verify that  $U\varphi(g)U = UgU$  for all  $g \in G$ . For each double coset  $UgU$ , we will show that  $UgU$  has a certain coset representative  $g'$  with  $\varphi(g') = g'$ , or  $f(g) = 0$  for all  $f \in \mathcal{H}$ .

The modification of the Bruhat decomposition in Lemma 16 tells us that  $UgU = UmU$  for some monomial matrix  $m$ . Let  $f \in \mathcal{H}$  be non-vanishing on  $UmU$ . That is,  $f(m) \neq 0$ . We show that  $m$  has the form

$$m = \begin{pmatrix} & & D_1 \\ & & \\ & D_2 & \\ & \ddots & \\ D_r & & \end{pmatrix},$$

for some diagonal matrices  $D_1, \dots, D_r$ . Equivalently, we show that if  $m_{ij}$  and  $m_{i+1,k}$  are non-zero, then we must have  $k \leq j + 1$ .

To see this, assume that  $m_{ij}, m_{i+1,k} \neq 0$  and  $k > j + 1$ . Then define  $x := I_n + m_{ij}E_{i,i+1} \in U$  and  $y := I_n + m_{i+1,k}E_{jk} \in U$ . Simple computations tell us that  $xm = m + m_{ij}m_{i+1,k}e_{ik} = my$ ,  $\pi(x) = \psi(m_{ij}) \neq 1$  and  $\pi(y) = \pi(0) = 1$ . Then, since  $f \in \mathcal{H}$ , there holds  $\pi(x)f(m) = f(xm) = f(my) = f(m)\pi(y)$ . Thus  $(\pi(x) - \pi(y))f(m) = 0$ , so  $f(m) = 0$  since  $\pi(x) \neq \pi(y)$ .

Now we show that each diagonal matrix  $D_i$  is actually a matrix of scalars. In particular, we show that if  $m_{i,j}$  and  $m_{i+1,j+1}$  are non-zero then they are equal. Consider  $x$  and  $y$  as given above, with  $k = j + 1$ . Then  $xm = my$ ,  $\pi(x) = \psi(m_{ij})$ ,  $\pi(y) = \psi(m_{i+1,j+1})$  and  $(\pi(x) - \pi(y))f(m) = 0$ . Recall that  $f$  doesn't vanish on  $UmU$  so  $f(m) \neq 0$ . Thus  $\pi(x) = \pi(y)$ , which tells us that  $\psi(m_{ij}) = \psi(m_{i+1,j+1})$  and  $m_{ij} = m_{i+1,j+1}$  by injectivity of  $\psi$ .

Finally, notice  $\varphi(m) = m$ . This is easy to see, since  $m^t$  is simply  $m$  with the elements on the opposite diagonal reversed, and left- and right-multiplying by  $w_0$  also reverses the opposite diagonal. This completes the proof.  $\square$



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