CARNEGIE MELLON UNIVERSITY COMPUTER SCIENCE DEPARTMENT 15-445/645 – DATABASE SYSTEMS (FALL 2021) PROF. LIN MA

Homework #2 (by Abi Kim)
Due: Sunday October 3, 2021 @ 11:59pm

IMPORTANT:

- Upload this PDF with your answers to Gradescope by 11:59pm on Sunday October 3, 2021.
- **Plagiarism**: Homework may be discussed with other students, but all homework is to be completed **individually**.
- You have to use this PDF for all of your answers.

For your information:

- Graded out of 100 points; 4 questions total
- Rough time estimate: \approx 1-4 hours (0.5-1 hours for each question)

Revision: 2021/09/19 21:22

Question	Points	Score
Cuckoo Hashing	20	
B+Tree	45	
Extendible Hashing	25	
B+Tree	10	
Total:	100	

Question 1: Cuckoo Hashing......[20 points]

Consider the following cuckoo hashing schema:

- 1. Both tables have a size of 4.
- 2. The hashing function of the first table returns the lowest two bits: $h_1(x) = x \& 0b11$.
- 3. The hashing function of the second table returns the next two bits: $h_2(x) = (x \gg 2) \& 0b11$.
- 4. When replacement is necessary, first select an element in the second table.
- 5. The original entries in the table are shown in the figure below.

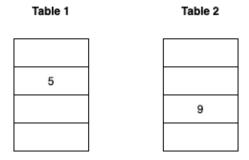


Figure 1: Initial contents of the hash tables.

- (a) [2 points] Select the sequence of insert operations that results in the initial state.
 - \square Insert 5, insert 9 \square Insert 9, insert 5 \square None of the above
- (b) [4 points] Insert keys 2 and 1. Select the resulting two tables.

□ A)	
Table 1	Table 2
1	5
2	9

□ B)	
Table 1	Table 2
	1
9	5
2	

□ C)	
Table 1	Table 2
	2
1	5
	9

Table 1	Table 2
	1
5	
2	9

????,没有正确答案,文字版在下面

(c) [4 points] Then insert 6, and delete 5. Select the resulting two tables.

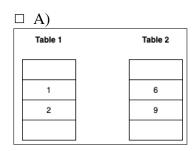


Table 1	Table 2
	2
1	
6	9

□ B)	
Table 1	Table 2
	2
1	6
	9

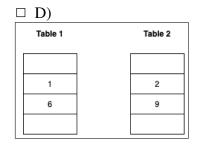


Table1: NULL, NULL, 2, NULL Table2: 1, 6, 9, NULL

(d) [4 points] Finally, insert 25. Select the resulting two tables.

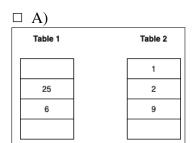
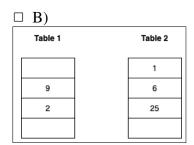


Table 1	Table 2
	2
1	6
25	9



□ D)	
Table 1	Table 2
	1
25	6
2	9

(e) [6 points] What is the smallest key that potentially causes an infinite loop given the tables in (d)?

 \square 0 \square 3 \square 4 \square 5 \square 9 \square 10 \square None of the above

加上5之后,共有6个数,他们能在的位置的集合共有5个空。

Question 2: B+Tree.....[45 points] Consider the following B+tree.

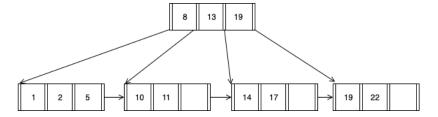
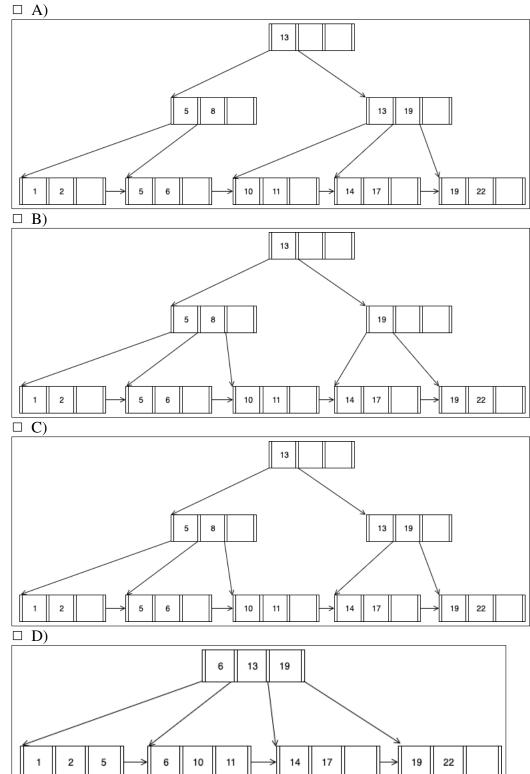


Figure 2: B+ Tree of order d=4 and height h=2.

When answering the following questions, be sure to follow the procedures described in class and in your textbook. You can make the following assumptions:

- A left pointer in an internal node guides towards keys < than its corresponding key, while a right pointer guides towards keys ≥.
- A leaf node underflows when the number of **keys** goes below $\lceil \frac{d-1}{2} \rceil$.
- An internal node underflows when the number of **pointers** goes below $\lceil \frac{d}{2} \rceil$.

(a) [15 points] Insert 6* into the B+tree. Select the resulting tree.



(b) [5 points] How many pointers (parent-to-child and sibling-to-sibling) do you chase to find all keys between 6^* and 17^* ?

 $\square \ 2 \quad \square \ 3 \quad \square \ 4 \quad \square \ 5 \quad \square \ 6 \quad \square \ 7$

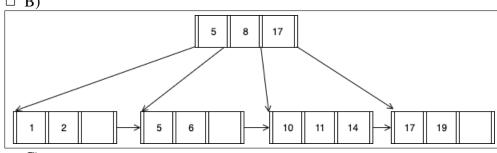
(c) [15 points] Then delete 22*. Select the resulting tree.

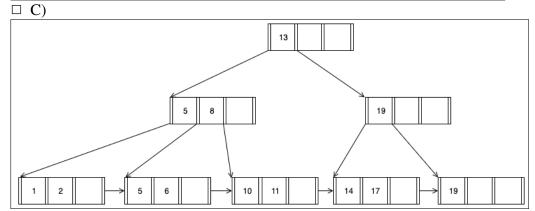
□ A)

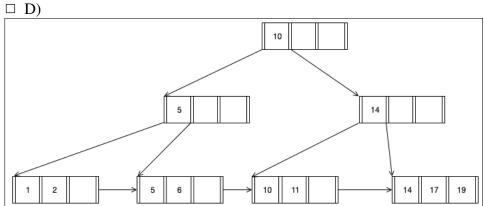
5 8 13

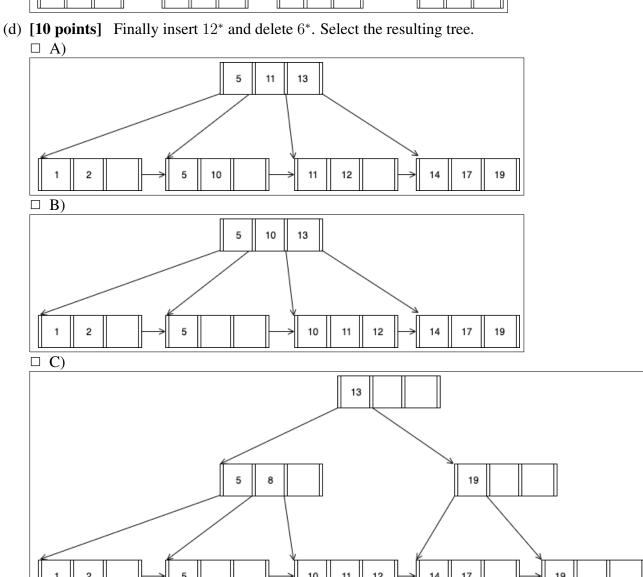
1 2 5 6 10 11 14 17 19

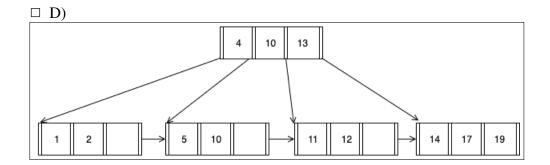
□ B)











Question 3: Extendible Hashing......[25 points]

Consider an extendible hashing structure such that:

- Each bucket can hold up to two records.
- The hashing function uses the lowest g bits, where g is the global depth.
- (a) Starting from an empty table, insert keys 6, 15, 34, 18.
 - i. [3 points] What is the global depth of the resulting table? \Box 0 \Box 1 \Box 2 \Box 3 \Box 4 \Box None of the above
 - ii. [3 points] What is the local depth the bucket containing 34?
 - \square 0 \square 1 \square 2 \square 3 \square 4 \square None of the above
 - iii. [3 points] What is the local depth of the bucket containing 15?
 - \square 0 \square 1 \square 2 \square 3 \square 4 \square None of the above
- (b) Starting from the result in (a), you insert keys 16, 7, 10, 20, 9.
 - i. [4 points] Which key will first cause a split (without doubling the size of the table)?
 - \square 16 \square 7 \square 10 \square 20 \square 9 \square None of the above
 - ii. [4 points] Which key will first make the table double in size?
 - \square 16 \square 7 \square 10 \square 20 \square 9 \square None of the above
- (c) Now consider the table below, along with the following deletion rules:
 - 1. If two buckets have the same local depth d, and share the first d-1 bits of their indexes (e.g. 010 and 110 share the first 2 bits), then they can be merged if the total capacity fits in a single bucket. The resulting local depth is d-1.
 - 2. If the global depth g becomes strictly greater than all local depths, then the table can be halved in size. The resulting global depth is g 1.

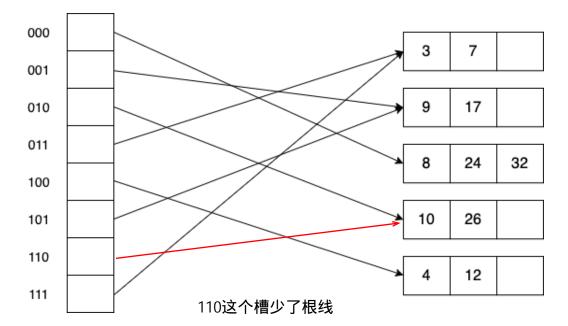


Figure 3: Extendible Hash Table along with the indexes of each bucket

Star	Starting from the table above, delete keys 10, 12, 7, 24, 8.										, 7, 24, 8.
i. [4 points] Which deletion first causes a reduction in a local d								duction in a local depth			
		10		12		7 🗆	24		8		None of the above
ii.	[4	poir	ıts]	Wh	ich de	eletion	first	cause	es a	rec	duction in global depth.
		10		12		7 🗆	24		8		None of the above

Consider the following B+trees shown below. Assume that threads use binary search to find matching keys in each node.

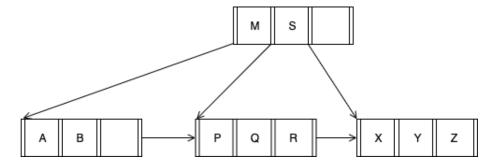


Figure 4: Figure 1

Consider the B+Tree shown in Figure 1. Answer the following questions for the resulting tree after deleting key B from the tree. If more than one solution exists, choose the tree that results in the most packed left-most leaf node.

(a)	[1 point	t] How	many nod	les will th	ne resultin	ig tree ha	ve?			
	□ 1	□ 2 □	3 🗆 4	4 □ N	ot possibl	le to deter	rmine			
(b)	[2 point	ts] Whi	ch key(s)	will be in	the left-	most leaf	node? M	lark all th	at apply.	
	\Box A	\Box B	\square M	\Box P	\Box Q	\Box R	\Box S	$\Box X$	\Box Y	\Box Z
	□ Not p	possible	to determ	ine						
(c)	[2 point	ts] Whi	ch key(s)	will be in	the root	node? M	ark all th	at apply.		
	\Box A	\Box B	\square M	\Box P	\Box Q	\Box R	\Box S	$\Box X$	\Box Y	\Box Z
	□ Not p	possible	to determ	ine						

(d) [5 points] The B+Tree shown in Figure 2 may be invalid. That is, it may or may not violate the correctness properties of B+Trees that we discussed in class. If the tree is invalid, select all the properties that are violated for each of the three nodes in the tree (i.e., Root, Leaf1, and Leaf2). If the tree is valid, then select 'None'. There will be no partial credit for missing violations.

Homework #2

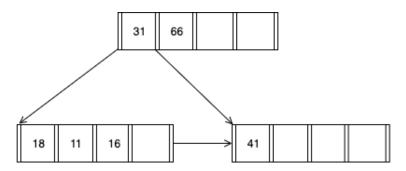


Figure 5: Figure 2

- \square Key order property is violated in **Root**.
- \Box Key-order property is violated in **Leaf1**.
- \Box Key-order property is violated in **Leaf2**.
- ☐ Half-full property is violated in **Root**.
- ☐ Half-full property is violated in **Leaf1**.
- ☐ Half-full property is violated in **Leaf2**.
- □ Balance property is violated in **Root**.
- □ Balance property is violated in **Leaf1**.
- □ Balance property is violated in **Leaf2**.
- \square Separator key violation in **Root**.
- ☐ Separator key violation in **Leaf1**.
- □ Separator key violation in **Leaf2**.