

COMMONWEALTH OF AUSTRALIA

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Data structures and Algorithms

Lecture 2b: stacks and queues [GT 2.1]

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Stacks and queues

These ADTs are restricted forms of List, where insertions and removals happen only in particular locations:

- stacks follow last-in-first-out (LIFO)
- queues follows first-in-first-out (FIFO)

So why should we care about a less general ADT?

- operations names are part of computing culture
- numerous applications
- simpler/more efficient implementations than Lists

Stack ADT



Main stack operations:

- **push**(e): inserts an element, e
- **pop**(): removes and returns the last inserted element

Auxiliary stack operations:

- **top**(): returns the last inserted element without removing it
- **size**(): returns the number of elements stored
- **isEmpty**(): indicates whether no elements are stored

Stack Example

operation	returns	stack
push(5)	-	[5]
push(3)	-	[5, 3]
size()	2	[5, 3]
pop()	3	[5]
isEmpty()	False	[5]
pop()	5	[]
isEmpty()	True	[]
push(7)	-	[7]
push(9)	-	[7, 9]
top()	9	[7, 9]
push(4)	-	[7, 9, 4]
pop()	4	[7, 9]

Stack Applications

Direct applications

- Keep track of a history that allows undoing such as Web browser history or undo sequence in a text editor
- Chain of method calls in a language supporting recursion
- Context-free grammars

Indirect applications

- Auxiliary data structure for algorithms
- Component of other data structures

Method Stacks

The runtime environment keeps track of the chain of active methods with a stack, thus allowing **recursion**

When a method is called, the system pushes on the stack a frame containing

- Local variables and return value
- Program counter

When a method ends, we pop its frame and pass control to the method on top

```
def main()  
    i = 5;  
    foo(i);
```

```
def foo(j)  
    k = j+1;  
    bar(k);
```

```
def bar(m)
```

...

bar
PC = 1
m = 6

foo
PC = 2
j = 5
k = 6

main
PC = 2
i = 5

Parentheses Matching

Each “(”, “{”, or “[” must be paired with a matching “)”, “}”, or “]”

- correct: ()(()){([()])}
- correct: ((()(()){([()])}))
- incorrect:)(()){([()])}
- incorrect: ({[]})
- incorrect: (

Scan input string from left to right:

- If we see an opening character, push it to a stack
- If we see a closing character, pop character on stack and check that they match

Stack implementation based on arrays

A simple way of implementing the Stack ADT uses an array:

- Array has capacity **N**
- Add elements from left to right
- A variable **t** keeps track of the index of the top element

```
def size()  
    return t + 1
```

```
def pop()  
    if isEmpty() then  
        return null  
    else  
        t ← t - 1  
        return S[t + 1]
```



Stack implementation based on arrays

- The array storing the stack elements may become full
- A push operation will then either grow the array or signal a “stack overflow” error.

```
def push(e)
  if  $t = N - 1$  then
    return “stack overflow”
  else
     $t \leftarrow t + 1$ 
     $S[t] \leftarrow e$ 
```



Stack implementation based on arrays

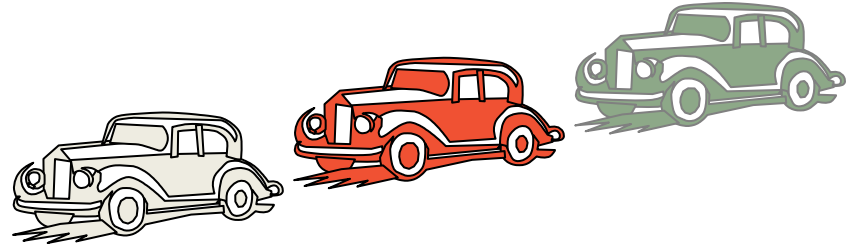
Performance

- The space used is $O(N)$
- Each operation runs in time $O(1)$

Qualifications

- Trying to push a new element into a full stack causes an implementation-specific exception or
- Pushing an item on a full stack causes the underlying array to double in size, which implies each operation runs in $O(1)$ amortized time.

Queue ADT



Main queue operations:

- **enqueue(e)**: inserts an element, e, at the end of the queue
- **dequeue()**: removes and returns element at the front of the queue

Auxiliary queue operations:

- **first()**: returns the element at the front without removing it
- **size()**: returns the number of elements stored
- **isEmpty()**: indicates whether no elements are stored

Boundary cases:

- Attempting the execution of dequeue or first on an empty queue signals an error or returns null

Queue Example

Operation		Output	Q
enqueue(5)	—	(5)	
enqueue(3)	—	(5, 3)	
dequeue()	5	(3)	
enqueue(7)	—	(3, 7)	
dequeue()	3	(7)	
first()	7	(7)	
dequeue()	7	()	
dequeue()	<i>null</i>	()	
isEmpty()	<i>true</i>	()	
enqueue(9)	—	(9)	
enqueue(7)	—	(9, 7)	
size()	2	(9, 7)	
enqueue(3)	—	(9, 7, 3)	
enqueue(5)	—	(9, 7, 3, 5)	
dequeue()	9	(7, 3, 5)	

Queue applications

Buffering packets in streams, e.g., video or audio

Direct applications

- Waiting lists, bureaucracy
- Access to shared resources (e.g., printer)
- Multiprogramming

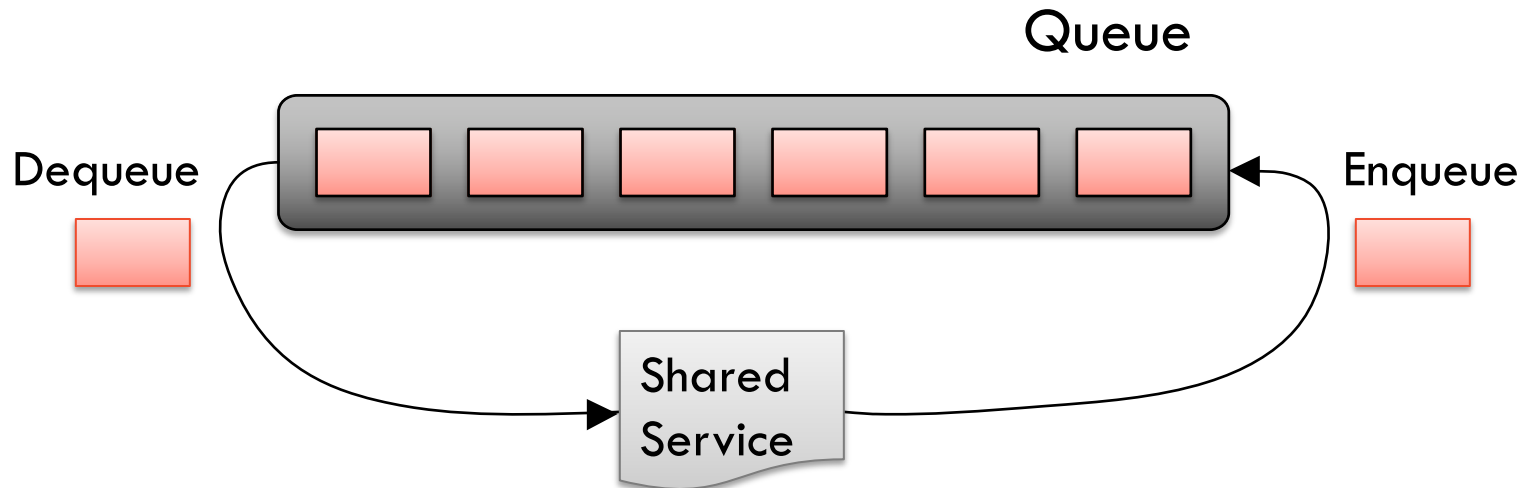
Indirect applications

- Auxiliary data structure for algorithms
- Component of other data structures

Queue application: Round Robin Schedulers

Implement a round robin scheduler using a queue Q by repeatedly performing the following steps:

1. $e = Q.dequeue()$
2. Service element e
3. $Q.enqueue(e)$



Queue implementation based on arrays

Use an array of size N in a circular fashion

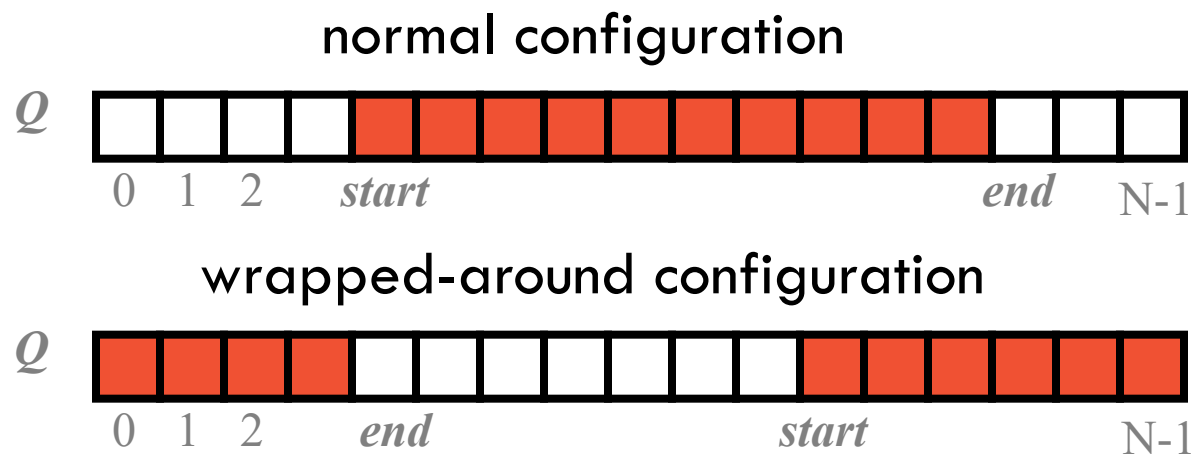
Two variables keep track of the front and size

start : index of the front element

end : index past the last element

size : number of stored elements

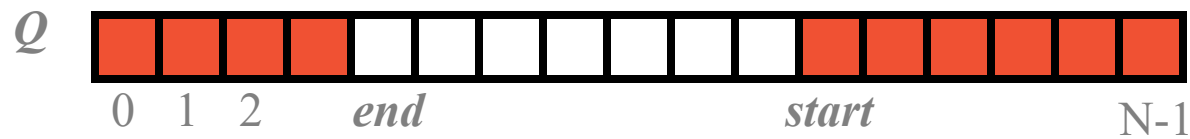
These are related as follows $\text{end} = (\text{start} + \text{size}) \bmod N$,
so we only need two, **start** and **size**



Queue Operations: Enqueue

Return an error if the array is full. Alternatively, we could grow the underlying array as dynamic arrays do

```
def enqueue(e)
  if size = N then
    return "queue full"
  else
    last ← (first + size) mod N
    Q[last] ← e
    size ← size + 1
```

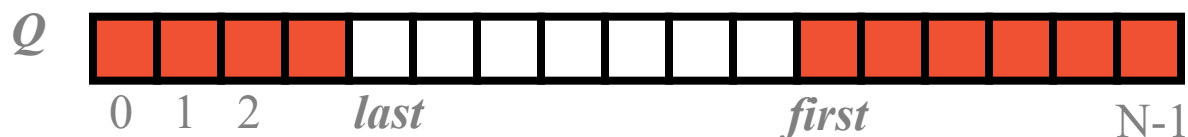


Queue Operations: Dequeue

Note that operation dequeue returns error if the queue is empty

One could alternatively signal an error

```
def dequeue()  
    if isEmpty() then  
        return "queue empty"  
    else  
        e ← Q[first]  
        first ← (first + 1) mod N  
        size ← (size - 1)  
        return e
```



Double-ended queues: Deques

- A linear structure that allows insertions and deletions at both ends

Method	Time
size, isEmpty	$O(1)$
getFirst, getLast	$O(1)$
addFirst, addLast	$O(1)$
removeFirst, removeLast	$O(1)$

Table 5.4: Performance of a deque realized by a doubly linked list.

Double-ended queue operations

The deque abstract data type is richer than both the stack and the queue ADTs. The fundamental methods of the deque ADT are as follows:

- `addFirst(e)`: Insert a new element e at the head of the deque.
- `addLast(e)`: Insert a new element e at the tail of the deque.
- `removeFirst()`: Remove and return the first element of the deque; an error occurs if the deque is empty.
- `removeLast()`: Remove and return the last element of the deque; an error occurs if the deque is empty.

Additionally, the deque ADT may also include the following support methods:

- `getFirst()`: Return the first element of the deque; an error occurs if the deque is empty.
- `getLast()`: Return the last element of the deque; an error occurs if the deque is empty.
- `size()`: Return the number of elements of the deque.
- `isEmpty()`: Determine if the deque is empty.

This week

Tutorial Sheets 2: Available on Ed

Quiz 1: Available on Canvas at the end of this lecture

Assignment 1: Available on Ed and Gradescope