

## Problem

Problem Statement: Define f(n,2) to be the maximal size of a subset  $S \in K_n$ , which  $K_n = \{a | a \in \mathbb{N} \& a \leq n\}$  such that it follows the axiom that if  $b \in S$ , 2b is not in S. What is f(n,2)?

Cases:

- 1. f(1,2) = 1, only 1.
- 2. f(2,2) = 1, being 1 or 2.
- 3. f(3,2) = 2, either 1, 3 or 2, 3.
- 4. f(4,2) = 3, being 1, 3, 4.

Observation 1.1:  $f(n,2) \ge \left\lceil \frac{n}{2} \right\rceil$ 

*Proof:* Trivial because all odd positive integers less than or equal to n would work.  $\square$ 

Now, consider the power of 2 for each positive integers in the range, rewriting it in the form of  $k = 2^n p_i$  for which  $n = v_2(k)$ , then when n = 0 would always work. Now, we define

$$G_k = \{2^k a \mid a \nmid 2 \& 2^k a < n\}$$

For n, the sets are defined as  $G_0, G_1, ..., G_{\lfloor \log_2 n \rfloor}$ , and define a set  $S_G$  that encompasses all of the above subsets.

Then clearly, we cannot take  $2G_k$  if we take all the elements of  $G_k$ . Therefore, we either take  $G_0, G_2, G_4, ..., G_{2s}$  or  $G_1, G_3, ..., G_{2s+1}$ . Now, it suffices to find the value of  $G_s$ .

Claim 1.2:  $|G_s| = \left\lceil \frac{\left\lfloor \frac{n}{2^s} \right\rfloor}{2} \right\rceil$  which  $\lfloor a \rfloor$  is defined as the largest integer less than or equal to a, and  $\lceil a \rceil$  is defined as the smallest integer greater than or equal to a.

*Proof:* We divide this into two claims:

Subclaim 1: The number of positive integers k less than or equal to n such that  $2^s \mid k$  is equal to  $\left\lfloor \frac{n}{2^s} \right\rfloor$ .

*Proof:*  $2^s a \le n$  for some positive integer a, and from here,  $a \le \frac{n}{2^s}$ . If  $\frac{n}{2^s}$  is not an integer, then the maximum pos.int a that is less than or equal to  $\left|\frac{n}{2^s}\right|$  by definition.  $\square$ 

Subclaim 2: The number of odd positive integers less than n is  $\lceil \frac{n}{2} \rceil$ 



*Proof:* Followed by the prove of observation 1.1.  $\square$ 

The claim follows by combining these two subclaims.

Q.E.D.

To create a single expression of claim 1.2, we want to find a simpler form of  $G_s$ .

Conjecture 1.3:  $|G_s|$  can be written as  $\lceil \frac{m}{2^{s+1}} \rceil$  for some positive integer m.

Now we define 
$$p(x,y) = \left\lceil \frac{\left\lfloor \frac{x}{2^y} \right\rfloor}{2} \right\rceil$$
, then  $p(a_1 2^m + b_1, m) = \left\lceil \frac{a_1}{2} \right\rceil$  such that  $0 \le b_1 < 2^m$ 

From here, define  $n_s$  to be  $n \pmod{2^s}$ , then  $f(n,s) = \left\lceil \frac{n-n_s}{2^{s+1}} \right\rceil$ . Therefore,  $m = n - n_s$ . Then, we split into cases about the parity of  $a_1$ :

- 1. If  $a_1$  is even, then  $a_1 2^m + b_1 \equiv b_1 \pmod{2^{m+1}}$
- 2. If  $a_1$  is odd, then  $a_1 2^m + b_1 = (a_1 1)2^m + 2^m + b_1 \equiv 2^m + b_1 \pmod{2^{m+1}}$

Assume that n = 50, then

1. 
$$|G_0| = \left\lceil \frac{\lfloor \frac{50}{2^0} \rfloor}{2} \right\rceil = 25, 50 \equiv 0 \pmod{2}$$

2. 
$$|G_1| = \left\lceil \frac{\left\lfloor \frac{50}{2^1} \right\rfloor}{2} \right\rceil = 13, 50 \equiv 2 \pmod{4}$$

3. 
$$|G_2| = \left\lceil \frac{\lfloor \frac{50}{2^2} \rfloor}{2} \right\rceil = 6, 50 \equiv 2 \pmod{8}$$

4. 
$$|G_3| = \left\lceil \frac{\lfloor \frac{50}{2^3} \rfloor}{2} \right\rceil = 3, 50 \equiv 2 \pmod{16}$$

5. 
$$|G_4| = \left\lceil \frac{\left\lfloor \frac{50}{2^4} \right\rfloor}{2} \right\rceil = 2, 50 \equiv 18 \pmod{32}$$

6. 
$$|G_5| = \left\lceil \frac{\left\lfloor \frac{50}{25} \right\rfloor}{2} \right\rceil = 1, 50 \equiv 50 \pmod{64}$$

## Conjecture 1.4:

1. 
$$|G_s| = \left\lceil \frac{n}{2^{s+1}} \right\rceil$$
 if  $n \equiv a \pmod{2^{s+1}}$  and  $a \ge 2^s$ 

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Now, let 
$$g(m,s) = \left\lceil \frac{m}{2^{s+1}} \right\rceil$$
, then  $p(m,s) = g(m-2^s,s)$ 



Now let  $m = a_1 2^s + b_1$ , which by euclidean algorithm, this will always be possible. Then,  $f(m, s) = \left\lceil \frac{a_1}{2} \right\rceil$ .

$$g(m-2^s,s) = \left\lceil \frac{a_1 2^s + b_1}{2^{s+1}} \right\rceil$$

Which because of the bound that  $0 \le b_1 < 2^s$ ,  $\frac{b_1}{2^{s+1}} < \frac{2^s}{2^{s+1}} = \frac{1}{2}$ . Therefore,  $b_1$  will have no impact on the value of  $g(m-2^s,s)$ .  $\square$ 

Therefore,  $g(m-2^s,s) = \left\lceil \frac{a_1}{2} \right\rceil$ , and the prove is complete.

In conclusion,  $|G_m| = \lceil \frac{n-2^m}{2^{m+1}} \rceil$ , and this allows us to massively simplify our expression. Before that, we need to prove that  $|G_0| + |G_2| + ... > |G_1| + |G_3| + ...$  Consider the differences  $|G_0| - |G_1|, |G_2| - |G_3|, ..., |G_{2n}| - |G_{2n+1}|$ . For the remaining, consider  $u = \lfloor \log_2 n \rfloor \equiv 1 \pmod{2}$ . Now we separate the cases by parity:

1. if  $u \equiv 1 \pmod{2}$ , then we will be able to partition  $G_0, ..., G_u$  into  $\frac{u+1}{2}$  subsets of  $\{G_{2n}, G_{2n+1}\}$  such that  $0 \le n \le \frac{u-1}{2}$ , and clearly  $|G_{2n}| - |G_{2n+1}| > 0$ , thus, the inequality

$$\sum_{m=0}^{\frac{u-1}{2}} |G_{2m}| - |G_{2m+1}| > 0$$

holds.  $\square$ 

2. If  $u \equiv 0 \pmod{2}$ , then consider the set  $S \setminus \{G_u\}$ . Then,  $|S \setminus \{G_u\}| \equiv 1 \pmod{2}$ , which by subcase 1, the inequality holds. Now, this simply transforms to

$$\left(\sum_{m=0}^{\frac{u-2}{2}} |G_{2m}| - |G_{2m+1}|\right) + |G_u| > 0$$

Which is true since  $|G_u|$  is always larger than 0 unless if  $2^u > n$ .  $\square$ 

In conclusion, let 
$$a=\lfloor \log_2 n \rfloor$$
, then 
$$f(n,2)=\sum_{m=0}^{\left \lfloor \frac{a}{2} \right \rfloor} \left \lceil \frac{n-2^{2m}}{2^{2m+1}} \right \rceil$$

However, this will become much more complicated when k becomes large. Now, consider  $D_n$  to be the multiples of  $2^n$ . Then, we can express all of these as  $|D_{2n}| - |D_{2n-1}|$ , and we can just combine as a single summation.

$$f(n,2) = \sum_{i=0}^{\lfloor \log_2 n \rfloor} \lfloor \frac{n}{2^i} \rfloor$$

And the idea is the same for all k other than 2 because we just want to take  $G_{2n}$  by subtracting



all the multiples of  $k^{2n}$  by all the multiples of  $k^{2n+1}$ . Then,

$$f(n,k) = \sum_{i=0}^{\lfloor \log_2 n \rfloor} \left\lfloor \frac{n}{k^i} \right\rfloor$$

With simulation, it's evident that  $\lim_{n\to\infty}f(n,k)=\frac{kn}{k+1}$ 



The Extended Problem Statement: Define  $f(n, k_1, k_2)$  to be the maximal size of a subset  $S \in K_n$ , which  $K_n = \{a | a \in \mathbb{N} \& a \leq n\}$  such that it follows the axiom that if  $b \in S$ , kb is not in S for all  $k_1, k_2$ . What is  $f(n, k_1, k_2)$ ?

First of all, we can consider the multiples of  $k^0$  with the previous strategy. Notice that for all  $a \in K_n$ ,  $a \equiv b \pmod{n}$  such that  $b \in \frac{\mathbb{Z}}{k\mathbb{Z}}$ , and I will first solve the problem when  $\gcd(k_1, k_2) = 1$  and  $k_1, k_2 \mid n$ .

**Observation 2.1:** Assume that 
$$gcd(k_1, k_2) = 1$$
, then  $f(n, k) \ge (k_1 - 1)(k_2 - 1) \left| \frac{n}{k_1 k_2} \right|$ 

*Proof:* This is a result of the Chinese Remainder Theorem: We can take every positive integers a that aren't divisible by  $k_1$  or  $k_2$ , which by modular arithmetics,

- 1.  $a \equiv 1, ..., k_1 1 \pmod{k_1}$
- 2.  $a \equiv 1, ..., k_2 1 \pmod{k_2}$

Which since  $\gcd(k_1, k_2) = 1$ , then by Chinese Remainder Theorem, every pair of  $a_1 \pmod{k_1}$  and  $a_2 \pmod{k_2}$  will produce a unique residue modulo  $k_1k_2$ . Therefore, there are a total  $(k_1 - 1)(k_2 - 1)$  of residues in a cycle of length  $k_1k_2$  that is divisible by neither  $k_1$  nor  $k_2$ . From here, we simply want to find the number of such cycles, which there are  $\left|\frac{n}{k_1k_2}\right|$  of them.  $\square$ .

However, this is not the end. We will also have  $k_1^{2m}(k_1x_1+1), k_1^{2m}(k_1x_1+2)...$  and  $k_2^{2m}(k_2x_2+1), k_2^{2m}(k_2x_2+2),...$  Therefore, let  $G_{k_1,m}$  be the set of all positive integers a such that  $k_1^m a < n$  and  $a \in \frac{\mathbb{Z}}{k_1\mathbb{Z}} - \{0\}.$ 

Conjecture 2.2: Define  $G_{k_1,m}$  be the set of positive integers p less than or equal to n such that  $v_{k_1}(p) = m$ .  $G_{k_2,m}$  is defined analogously. then, the sets

$$S_{k_1,m}\{a \mid a \in k_1^{2m} G_{k_1 k_2,0} \text{ and } a \leq n\}$$

$$S_{k_2,m}\{b\mid b\in k_2^{2m}G_{k_1k_2,0}\text{ and }b\leq n\}$$

will work.

*Proof:* Follow directly from the Chinese Remainder Theorem as a result of *Observation 2.1* and the extended version of *Conjecture 1.4.*  $\square$ .

Now, we simply want to find the set of quotients when a positive integer  $a \in K_n$  is divided by powers of  $k_1$  and  $k_2$ . There yields  $p_1 = \left\lfloor \frac{n}{k_1^m} \right\rfloor$  such positive integers, and analogous defined for  $p_2$ , which the set will be in the form  $\{1, ..., a_{p_1}\} \leftrightarrow \{k_1 m^i, ..., k_1^m a_{p_1}\}$  and the former set is simply  $K_{p_1}$ .



From here, use the Euclid's Algorithm by rewriting  $p_1 = xk_1k_2 + r$ , which there will simply be  $x(k_1 - 1)(k_2 - 1) + r$  positive integers in the set  $S_{k_1,m}$ . From here, the answer is

$$x\left(\frac{p_1-r}{xk_1k_2}\right)\varphi(k_1k_2) + r = \frac{(p_1-r)\varphi(k_1k_2)}{k_1k_2} + r$$

To further simplify this, let's make clear of what  $p_1 - r$  is.  $p_1 - r$  is basically the largest multiple of  $k_1k_2$  less than  $p_1$ , and  $p_1$  is the largest quotient when a positive integer less than or equal to n divided by  $k_1^m$ , from our conjecture.  $p_1 = \left| \frac{n}{k_1^m} \right|$ 

$$p_1 - r$$
 is simply  $k_1 k_2 \left\lfloor \frac{n \choose k_1^m}{k_1 k_2} \right\rfloor$ . The conjecture so far is that this is equal to  $\left\lfloor \frac{n}{k_1^{m+1} k_2} \right\rfloor$ .

This conjecture is flawed because for f(6,2,3), you can take 1,4,5,6 which will yield a result of 4 instead of 2.



Salvaged Conjecture 2.3: Assume that  $k_1 < k_2$ , then

$$f(n, k_1, k_2) = \sum_{\zeta=0}^{\infty} \left( \left\lfloor \frac{n}{(k_1 k_2)^{\zeta}} \right\rfloor - \left\lfloor \frac{n}{k_1 (k_1 k_2)^{\zeta}} \right\rfloor \right)$$

Proof: It's always possible to partition  $K_n$  into various strings of  $\{p, ..., pk_1^a k_2^b\}$  such that  $\gcd(k_1, k_2, p) = 1$ . First, pick out all positive integers  $a \in (\frac{n}{k_1}, n]$ . It's easy to prove that this construction works because the function  $\frac{n}{x}$  is monotonically decreasing, and the maximum value of of such multiple chain will be extended from n, which is  $\frac{n}{k_1}$  and this is excluded from the set.  $\square$ 

Claim: that the next stage is between the set  $\left(\left\lfloor \frac{n}{k_1^2 k_2} \right\rfloor, \frac{n}{k_1 k_2} \right]$ 

*Proof:* For the sake of contradiction, assume that there exists an integer between  $\left[\frac{n}{k_1k_2}, \frac{n}{k_1}\right]$ . This is equivalent to the statement that there exist a chain of multiples  $a_1, k_1a_1$  or  $a_1, k_2a_1$  for any integer  $a_1 \in \left[\frac{n}{k_1k_2} + 1, \frac{n}{k_1} - 1\right]$ .

Proceed with bounding:  $\frac{n}{k_2} + k_1 \le a_1 k_1 \le n - k_1 < n$ .  $\square$ 

What if n is not divisible by  $k_1$ , nor  $k_2$ ? In this case, we simply just take the floor value.

**Example:** Consider f(1296, 2, 3), first we take all positive integers between [649, 1296], and then we take [109, 216], [19, 36], [4, 6], [1]. This yields  $648 + 108 + 18 + 3 + 1 = \boxed{778}$ 

For f(1296, 3, 5) instead, we first take [433, 1296]. Then, the maximum value of  $x \in [1, 432]$  is 86. Then, we take from 86 to  $\lfloor \frac{86}{3} \rfloor + 1$ , which is [29, 86]. Continue with this, we take [2, 5]. This gives  $864 + 58 + 4 = \lceil 926 \rceil$ 

However, we see a counterexample that if  $n=20, k_1=2, k_2=5$ , then we can choose [11, 20] and [2] by our algorithm, but notice that 5 can also be taken because  $5 \cdot 5 > n$ . Therefore, we have to also take positive integers  $a \in \left[\left|\frac{n}{k_1 k_2}\right|, \left|\frac{n}{k_1}\right|\right]$ , such that  $\frac{a}{k_1} > \frac{n}{k_1^2 k_2}$  and  $k_2 a > n$  and  $k_1 a < \frac{n}{k_1}$ .

Therefore, we will Salvage our current conjecture again:



Salvaged Conjecture 2.4: Assume that  $k_1 < k_2$ , then

$$f(n, k_1, k_2) = \sum_{\zeta=0}^{\infty} \left( \left\lfloor \frac{n}{(k_1 k_2)^{\zeta}} \right\rfloor - \left\lfloor \frac{n}{k_1 (k_1 k_2)^{\zeta}} \right\rfloor \right) + S_n$$

Such that  $S_n$  is a set that contains all positive integers a such that

$$\begin{cases} \frac{a}{k_1} > \frac{n}{k_1^2 k_2} \\ k_2 a > n \\ k_1 a < \frac{n}{k_1} \end{cases}$$

We bound the possible values a:  $\frac{n}{k_2} < a < \frac{n}{k_1^2}$ . We can ignore the first bound of  $a > \frac{n}{k_1 k_2}$  because  $\frac{n}{k_1 k_2} < \frac{n}{k_2}$ . From here, we deduce that  $k_1^2 < k_2$  must be true for  $|S_n| > 0$ , and this will produce about  $n \frac{k_2 - k_1^2}{k_1^2 k_2}$  values.

Now, we continue to approximate the maximal number of elements.

$$f(n, k_1, k_2) = \sum_{\zeta=0}^{\infty} \left( \left\lfloor \frac{n}{(k_1 k_2)^{\zeta}} \right\rfloor - \left\lfloor \frac{n}{k_1 (k_1 k_2)^{\zeta}} \right\rfloor \right) + S_n$$

Which when  $k_1^2 < k_2$ ,  $|S_n| \approx \frac{(k_2 - k_1^2)n}{k_1^2 k_2}$ . For the sake of approximation, assume that the floor values vanish. Then,

$$f(n, k_1, k_2) \approx \sum_{\zeta=0}^{\infty} \frac{n}{(k_1 k_2)^{\zeta}} \left( \frac{k_1 - 1}{k_1} \right) = \frac{(k_1 - 1)n}{k_1} \left( \frac{k_1 k_2}{k_1 k_2 - 1} \right) = n \left( \frac{k_1 k_2 - k_2}{k_1 k_2 - 1} \right)$$

## Conjecture 2.5:

1. If  $k_1^2 > k_2$ , then

$$f(n, k_1, k_2) \approx n \left( \frac{k_1 k_2 - k_2}{k_1 k_2 - 1} \right)$$

2. IF  $k_1^2 < k_2$ , then

$$f(n, k_1, k_2) \approx n \left( \frac{k_1 k_2 - k_2}{k_1 k_2 - 1} + \frac{k_2 - k_1^2}{k_1^2 k_2} \right)$$

Now, we will prove that this is indeed the maximal size.





Final Final final problem statement: Let  $P_n$  be the set of prime divisors of n(for example,  $P_{50} = \{2, 5\}$  and  $P_{30} = \{2, 3, 5\}$ ). Define  $K_n = \{a \in \mathbb{N} | a \leq n\}$  What is the maximal size of a subset  $S \in K_n$  such that it follows the axiom that if  $b \in S$ ,  $kb \notin S$  for all  $k \in P_n$ .