

Chapter 2, Section 2.5: Architecture and Characteristics of a RTOS kernel (μ C/OS-II)

1. Introduction
2. Interrupt Management
3. Task Management
4. Time Management
5. *Event Management*
6. Memory Management
7. Porting μ COS-II

Chap 2, Section 2.5, Page 74

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5. Event Management

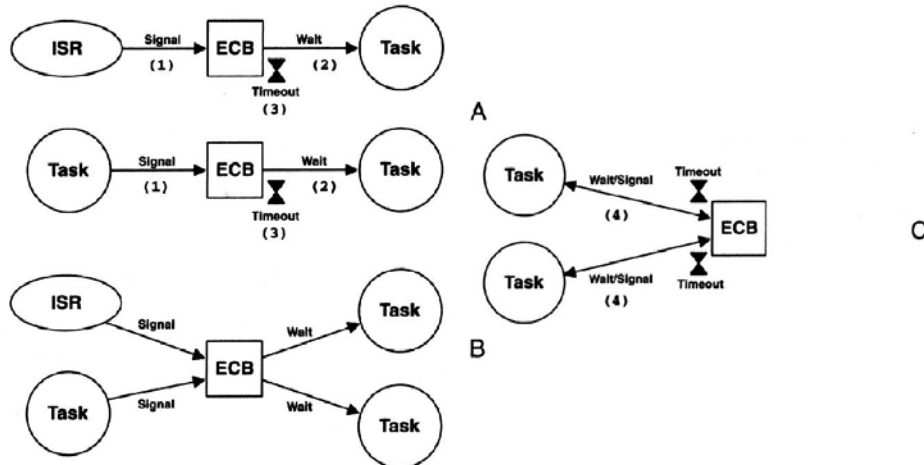
- 5.1 Definition of Event Control Block (ECB)
- 5.2 ECB Management
- 5.3 TCB vs ECB
- 5.4 Semaphore Management
- 5.5 Mutex Management
- 5.6 Mailbox Management
- 5.7 Message Queues Management

Chap 2, Section 2.5, Page 75

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5.1 Definition of ECB



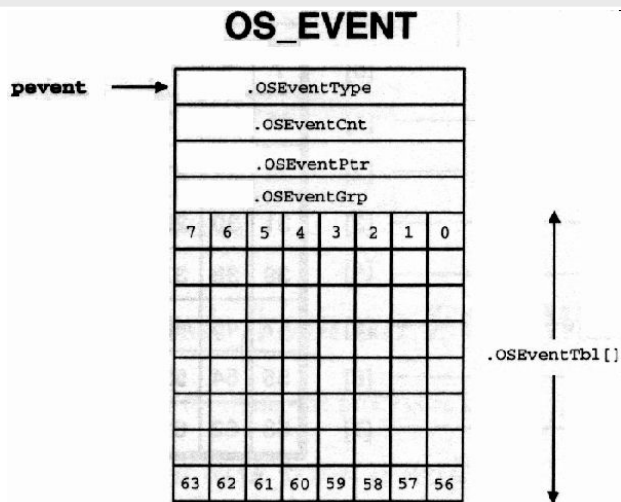
Chap 2, Section 2.5, Page 76

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5.1 Definition of ECB



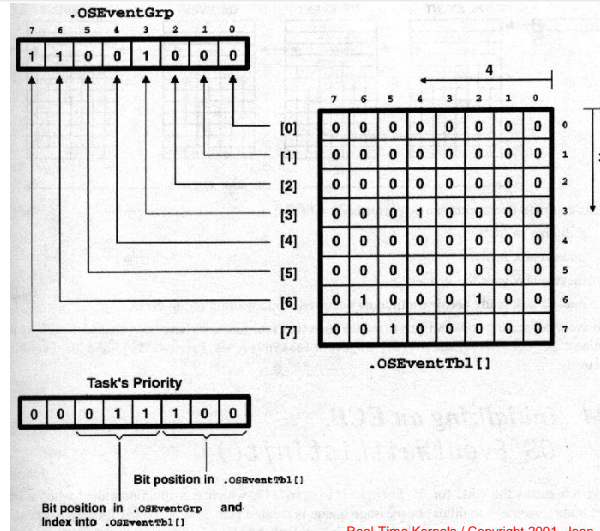
Chap 2, Section 2.5, Page 77

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5.1 Definition of ECB



Chap 2, Section 2.5, Page 78

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5.2 ECB Management

When *PEND* occurs, *OSEventTaskWait* will:

1. Remove the task from OSRdyTbl and OSRdyGr
2. Store the task in OSEventTbl and OSEventGrp

Chap 2, Section 2.5, Page 79

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5.2 ECB Management

When *POST* occurs, *OSEventTaskReady* will:

1. Elect the next task ready for the critical section (see next slide)
2. Remove the task from *OSEventTbl* and *OSEventGrp*
3. Reactivate the task in *OSRdyTbl* and *OSRdyGrp*

5.2 ECB Management

When *PEND after timeout* occurs, *OSEventTO*:

1. Remove the task from *OSEventTbl* and *OSEventGrp* lists

N.B. The task does not need to be reactivated into *OSRdyTbl* and *OSRdyGrp*. *OSTimeTick* will update *OSTCBDly* and the task will be reactivated after countdown reaches 0.

Diapositive 8

LF1 Landry Frechette; 2008-04-12

5.3 TCB vs ECB

• Operations that toggle a bit from 0 to 1 in OSxxxTbl and OSxxxGrp vector

Making a task ready to run

- consist of activating a task in OSRdyTbl and OSRdyGrp
- Some examples: OSTimeTick and **OS_EventTaskReady***
- See slides 83 and 84

≡

Making a task wait for an ECB

- consist in making a task wait for a semaphore, mutex, mailbox or queue in OSEventTbl and OSEventGrp
- some examples: **OS_EventTaskWait*** called by OSSemPend, OSMutexPend, OSMboxPend and OSQPend
- See slides 83 and 85

* Functions marked in red can be observed in the tutorial code

Chap 2, Section 2.5, Page 82

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5.3 TCB vs ECB

• In both cases, OSxxxTbl and OSxxxGrp are updated as followed:

- OSTCBY => ptcb->OSTCBY = priority >> 3;
- OSTCBBitY => ptcb->OSTCBBitY = OSMaTbl[ptcb->OSTCBY]
- OSTCBX => ptcb->OSTCBX = priority & 0x07;
- OSTCBBitX => ptcb->OSTCBBitX = OSMaTbl[ptcb->OSTCBX];

Index	Bit Mask
0	00000001
1	00000010
2	00000100
3	00001000
4	00010000
5	00100000
6	01000000
7	10000000

Chap 2, Section 2.5, Page 83

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5.3 TCB vs ECB

Code to change the state of a task from the waiting state to ready to execute state:

```
OSRdyGrp           |=      ptcb->OSTCBBitY;  
OSRdyTbl[ptcb->OSTCBY] |=    ptcb->OSTCBBitX;
```

5.3 TCB vs ECB

• Making a task wait for an ECB:

```
pevent->OSEventGrp           |=      ptcb->OSTCBBitY;  
pevent->OSEventTbl[ptcb->OSTCBY] |=    ptcb->OSTCBBitX;
```

5.3 TCB vs ECB

• Operations that toggle a bit from 1 to 0 in OSxxxTbl and OSxxxGrp vector

Removing a task from the ready list

- Some examples: OSTimeDly, OSTaskSuspend, OSChangePrio, and **OS_EventTaskWait**
- See slides 83 and 87

Removing a task from a waiting list

- **OS_EventTaskReady** called by SSemPost, OSMutexPost, OSMboxPost, OSQPost
- **OS_EventTO** called by OSSemPend, OSMutexPend, OSMboxPend, OSQPend
- See slides 83 and 88

5.3 TCB vs ECB

• Removing a task from the ready list :

```
if ((OSRdyTbl[ptcb->OSTCBY] &= ~ptcb->OSTCBBitX) == 0)
{OSRdyGrp &= ~ptcb->OSTCBBitY;
}
```


5.3 TCB vs ECB

- Removing a task from a waiting list

```
if ((pevent->OSEventTbl[ptcb->OSTCBBY]&= ~ ptcb->OSTCBBitX) ==0)
{
    pevent->OSEventGrp &= ~ ptcb->OSTCBBitY;
}
```

5.3 TCB vs ECB

- Two operations to determine the top priority task from OSxxxTbl and OSxxxGrp vector:

Finding the highest priority task ready to run

- For example: e.g. OSSched, OSIntExit
- See slide 90

≡

Finding the highest priority task waiting for the event

- For example: **OS_EventTaskReady**
- See slide 91

5.3 TCB vs ECB

- **Code to determine the top priority task (from the ready list):**

```
Y      =   OSUnMapTbl[OSRdyGrp] ;  
X      =   OSUnMapTbl[OSRdyTbl[Y]] ;  
prio   =   [Y << 3] + X ;
```

5.3 TCB vs ECB

- **Code to determine the top priority task waiting for an event:**

```
Y      =   OSUnMapTbl[pevent->OSEventGrp] ;  
X      =   OSUnMapTbl[pevent -> OSEventTbl [Y]] ;  
prio   =   [Y << 3] + X ;
```

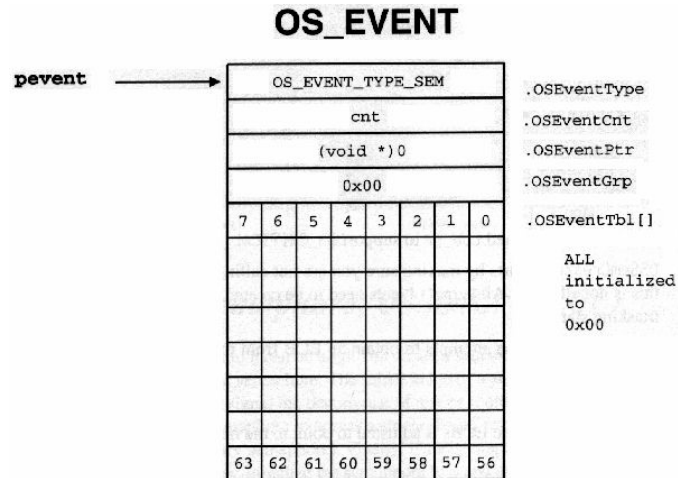
5.3 TCB vs ECB

```
*****
* Note(s): 1) Index into table is bit pattern to resolve highest priority.
*           2) Indexed value corresponds to highest priority bit position
* (i.e. 0..7)
*****
INT8U const OSUnMapTbl[] = {
    0, 0, 1, 0, 2, 0, 1, 0, 3, 0, 1, 0, 2, 0, 1, 0, // 0x00-0x0F
    4, 0, 1, 0, 2, 0, 1, 0, 3, 0, 1, 0, 2, 0, 1, 0, // 0x10-0x1F
    5, 0, 1, 0, 2, 0, 1, 0, 3, 0, 1, 0, 2, 0, 1, 0, // 0x20-0x2F
    4, 0, 1, 0, 2, 0, 1, 0, 3, 0, 1, 0, 2, 0, 1, 0, // 0x30-0x3F
    6, 0, 1, 0, 2, 0, 1, 0, 3, 0, 1, 0, 2, 0, 1, 0, // 0x40-0x4F
    4, 0, 1, 0, 2, 0, 1, 0, 3, 0, 1, 0, 2, 0, 1, 0, // 0x50-0x5F
    5, 0, 1, 0, 2, 0, 1, 0, 3, 0, 1, 0, 2, 0, 1, 0, // 0x60-0x6F
    4, 0, 1, 0, 2, 0, 1, 0, 3, 0, 1, 0, 2, 0, 1, 0, // 0x70-0x7F
    7, 0, 1, 0, 2, 0, 1, 0, 3, 0, 1, 0, 2, 0, 1, 0, // 0x80-0x8F
    4, 0, 1, 0, 2, 0, 1, 0, 3, 0, 1, 0, 2, 0, 1, 0, // 0x90-0x9F
    5, 0, 1, 0, 2, 0, 1, 0, 3, 0, 1, 0, 2, 0, 1, 0, // 0xA0-0xAF
    4, 0, 1, 0, 2, 0, 1, 0, 3, 0, 1, 0, 2, 0, 1, 0, // 0xB0-0xBF
    6, 0, 1, 0, 2, 0, 1, 0, 3, 0, 1, 0, 2, 0, 1, 0, // 0xC0-0xCF
    4, 0, 1, 0, 2, 0, 1, 0, 3, 0, 1, 0, 2, 0, 1, 0, // 0xD0-0xDF
    5, 0, 1, 0, 2, 0, 1, 0, 3, 0, 1, 0, 2, 0, 1, 0, // 0xE0-0xEF
    4, 0, 1, 0, 2, 0, 1, 0, 3, 0, 1, 0, 2, 0, 1, 0 // 0xF0-0xFF };
```

5.4 Semaphore Management

- To create a semaphore: *OSSemCreate()*
- To retrieve semaphore's info: *OSSemQuery()*
- To delete a semaphore: *OSSemDel()*
- To wait on a semaphore (W):
 - blocking: *OSSemPend()*
 - non-blocking *OSSemAccept()*
- To release (S) a semaphore: *OSSemPost()*

5.4 Semaphore Management



Chap 2, Section 2.5, Page 94

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5.5 Mutex Management

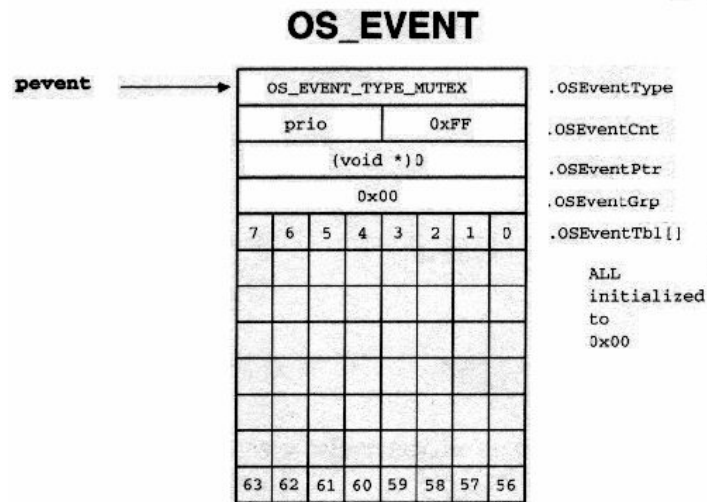
- To create a Mutex: *OSMutexCreate()*
- To retrieve mutexes' info: *OSMutexQuery()*
- To delete a Mutex *OSMutexDel()*
- To wait (W) on a mutex:
 - blocking: *OSMutexPend()*
 - non-blocking: *OSMutexAccept()*
- To release (S) a mutex: *OSMutexPost()*

Chap 2, Section 2.5, Page 95

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5.5 Mutex Management



Chap 2, Section 2.5, Page 96

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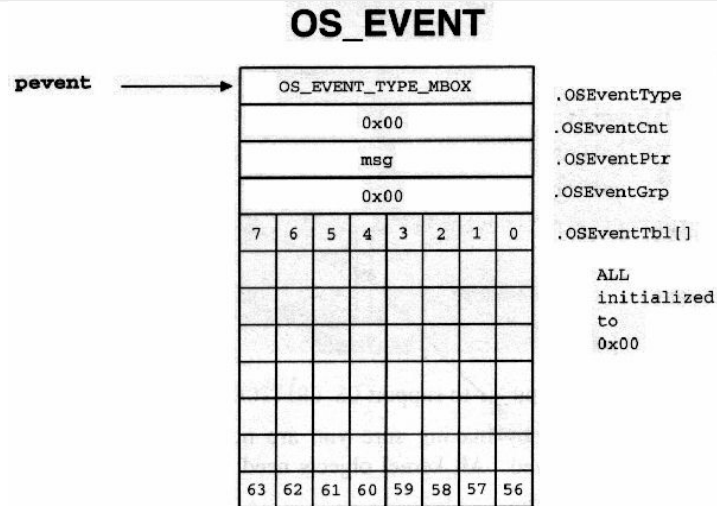
5.6 Mailbox Management

- To create a mailbox: *OSMboxCreate()*
- To retrieve mailbox' info: *OSMboxQuery()*
(*e.g. is there data in the mailbox?*)
- To delete a mailbox: *OSMboxDel()*

Chap 2, Section 2.5, Page 97

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5.6 Mailbox Management



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Chap 2, Section 2.5, Page 98

5.6 Mailbox Management

•To send a **non-blocking** message...

- in a single mailbox: **OSMboxPost()**
- to all waiting tasks: **OSMboxPostOpt()**

(we can send to the top priority waiting task, or to all tasks, e.g. broadcast)

•To read from a mailbox (take a message):

- blocking: **OSMboxPend()**
- non-blocking: **OSMboxAccept()**

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Chap 2, Section 2.5, Page 99

5.7 Message Queues Management

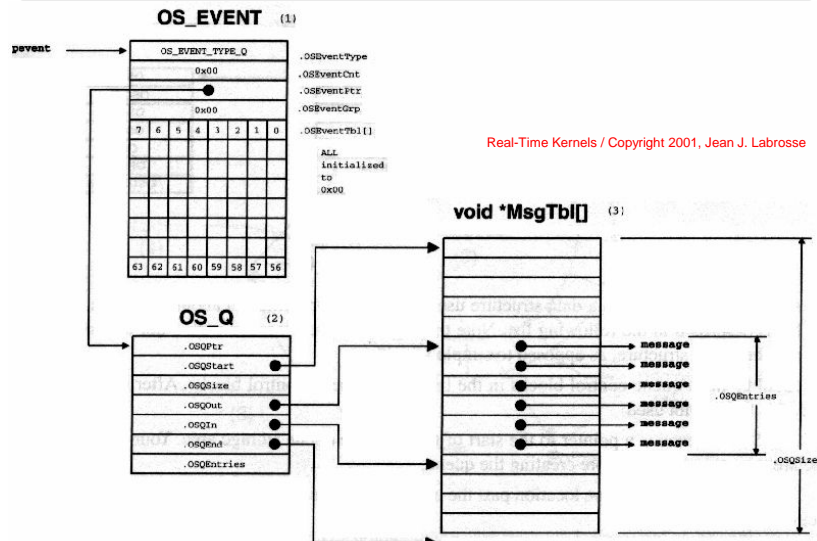
- To create a message queue: *OSQCreate()*
- To retrieve queue's info: *OSQQuery()*
- To delete a queue: *OSQDel()*
- To clear a queue: *OSQFlush()*

Chap 2, Section 2.5, Page 100

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5.7 Message Queues Management

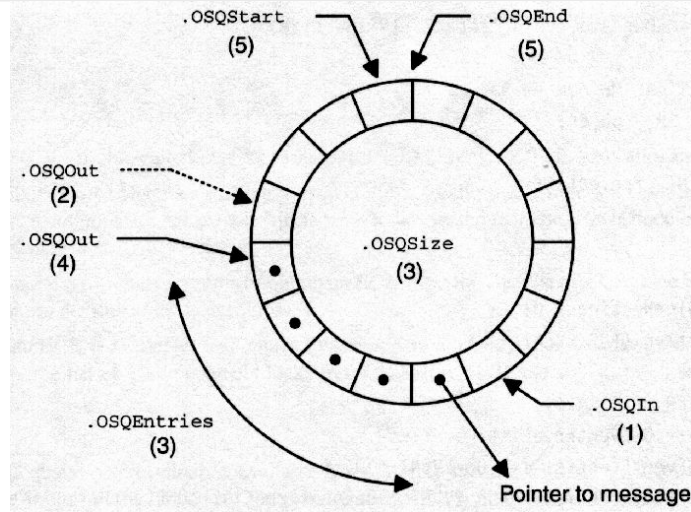


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5.7 Message Queues Management



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Chap 2, Section 2.5, Page 102

5.7 Message Queues Management

- To send a **non-blocking** message...
 - in a FIFO queue: *OSQPost()*
 - in a LIFO queue: *OSQPostFront()*
 - to all waiting tasks: *OSQPostOpt()*
- To read from a queue:
 - blocking: *OSQPend()*
 - non-blocking: *OSQAccept()*

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Chap 2, Section 2.5, Page 103