

1. (a)

$$\begin{aligned} S &\Rightarrow X \\ &\Rightarrow bX \\ &\Rightarrow bbX \\ &\Rightarrow bbbX \\ &\Rightarrow bbb \end{aligned}$$

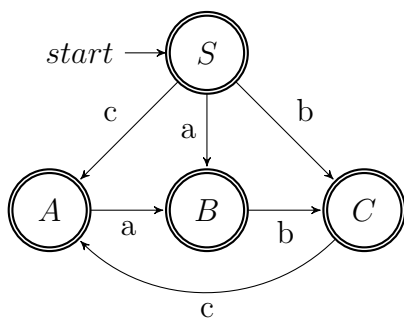
(b)

$$\begin{aligned} S &\Rightarrow aS \\ &\Rightarrow aaS \\ &\Rightarrow aaX \\ &\Rightarrow aabX \\ &\Rightarrow aabbX \\ &\Rightarrow aabb \end{aligned}$$

(c)

$$\begin{aligned} S &\Rightarrow X \\ &\Rightarrow \epsilon \end{aligned}$$

2.



3. (a)

$$\begin{aligned} S &\rightarrow aX \mid \epsilon \\ X &\rightarrow bS \end{aligned}$$

(b)

$$\begin{aligned} S &\rightarrow aX \mid X \\ X &\rightarrow bS \mid \epsilon \end{aligned}$$

(c)

$$\begin{aligned} S &\rightarrow aX \mid X \\ X &\rightarrow bS \mid \epsilon \end{aligned}$$

(d)

$$\begin{aligned} S &\rightarrow aX \mid X \\ X &\rightarrow bS \mid \epsilon \end{aligned}$$

4. Prove  $\{a^n b^n c^n\}$  is not regular.

**Proof.** Suppose not. That is, let  $L_1 = \{a^n b^n c^n\}$  be regular such that the Pumping Lemma gives us  $k$ . Choose some  $x = a^k, y = b^k, z = c^k$  such that  $xyz = a^k b^k c^k \in L_1$  and  $|y| \geq k$ . By the Pumping Lemma, let  $y = uvw$  with  $|v| > 0$  such that  $xuv^i wz \in L_1$  for all  $i \geq 0$ . Since  $v$  is a sequence of one or more  $bs$ ,  $uv^2w$  has more  $bs$  than  $uvw$ , and  $xuv^2wz$  has more  $bs$  than  $as$  and  $cs$ . This contradicts our initial assumption that  $|a| = |b| = |c|$ . Thus,  $\{a^n b^n c^n\}$  is not regular.  $\square$

5. (a)

$$\begin{aligned} S &\rightarrow 0A \mid 1A \mid \epsilon \\ A &\rightarrow 0B \mid 1B \\ B &\rightarrow 0S \mid 1S \end{aligned}$$

(b)

$$S \rightarrow aaSbb \mid \epsilon$$

(c)

$$\begin{aligned} S &\rightarrow aSb \mid A \mid \epsilon \\ A &\rightarrow bAa \mid S \mid \epsilon \end{aligned}$$

(d)

$$S \rightarrow aSa \mid bSb \mid \epsilon$$

6. Two different parse trees for the string  $() []$ .

 $\langle s \rangle$  $\langle \text{round} \rangle \langle \text{square} \rangle$  $() []$

Second parse tree for the same string.

