WRF Scaling and Performance Assessment

Comparison of Compilers and MPI Libraries on Cheyenne

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Outline



- Background
 - WRF
 - Cheyenne
 - Benchmark Cases
- Compilers
- Message Passing Interface Libraries
- Run Time Scaling
- Computation Time Scaling
- MVAPICH scaling

Background

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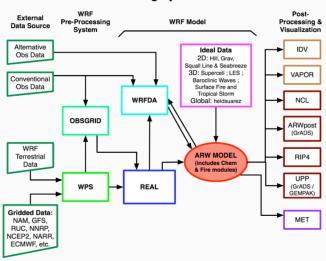


- The Weather Research and Forecast (WRF) model is a parallel mesoscale numerical weather forecasting application used in both operational and research environments.
- WRF is among the more commonly run codes by atmospheric scientists on NCAR's Cheyenne supercomputer.
 - Thus it is very important for WRF's users to know how to obtain the best performance of WRF on Cheyenne, especially as users scale their runs to larger core counts.

WRF System Flowchart



WRF Modeling System Flow Chart





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 - But cu_physics disabled and sf_sfclay_physics = 1 for both resolutions.

Summary of Benchmark Cases



Region	Resoultion	Horizontal	Vertical	Total	Time	Run
		Gridpoints	Gridpoints	Gridpoints	step	Hours
CONUS	12 km	425	300	127,500	72	6
CONUS	2.5 km	1901	1301	2,473,201	15	6
Maria	3 km	1396	1384	1,932,064	9	3
Maria	1 km	3665	2894	10,606,510	3	1

Compilers



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 - -fp-model fast=2 : similar to GNU's -Ofast optimization

Compiler Performance Results



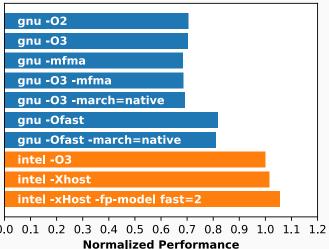


Fig. 1: Comparison of Intel 18.0.1 and Gnu 8.1.0 compilers with various compilation flags normalized to default Intel WRF compilation Runs made using CONUS 12 km Benchmark Case on 2 Nodes

Compiler Performance Summary



- Intel compiler is consistently 25-30% faster than the Gnu compiler across all flags tried.
- We also see that for both Intel and Gnu, the -Ofast (for Gnu) or -fp-model fast=2 (for Intel) are the only flags that make a significant difference in speed.
- Other flags tried such as -mfma or -march=native -Xhost made little to no difference in WRF's speed.

Message Passing Interface Libraries



• SGI's MPT version 2.18 (Default MPI for Cheyenne)



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MPI Comparison Results



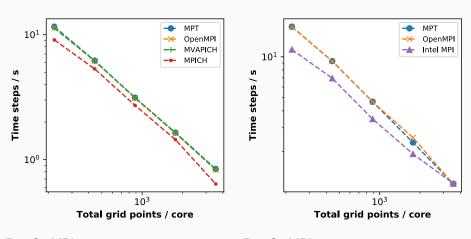


Fig. 2: MPI comparison using Gnu 8.1.0

Fig. 3: MPI comparison using Intel 18.0.1

Runs made using CONUS 12 km Benchmark Case

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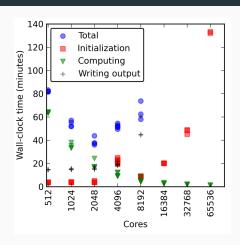


- MPT, MVAPICH and OpenMPI all have similar performance.
- MPICH has overall poor performance and the performance.
- Intel MPI does not scale well to large node counts.

Total Run Time Scaling

Run Time Scaling Comparison





Total 80 -Initialization Computing Wall-clock time (minutes) Writing output 60 40 -20 -0 -36864 -1152 2304 4608 18432 Cores

Fig. 4: Run Time Scaling on **Yellowstone**

Fig. 5: Run Time Scaling on Cheyenne

Runs made using Hurricane Maria 1 km Benchmark case.

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- In Fig 4 for Yellowstone, the initialization time scaled much poorer at large node counts, eventually leading to unfeasible long jobs.
 - This improvement in the scaling of the initialization time is likely due to the improvements made to WRF's initialization code with how the MPI calls are performed along with improvements in the MPI used in Cheyenne versus Yellowstone.

Computation Time Scaling

Computation Time Scaling Results



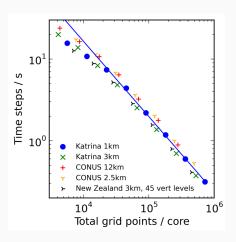


Fig. 6: Computation Scaling on Yellowstone

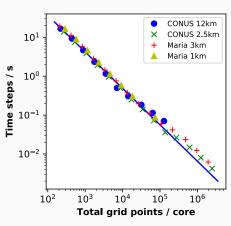


Fig. 7: Computation Scaling on **Cheyenne**



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 - Likely due to improvements in WRF's MPI code along and a better network interconnect on Cheyenne than Yellowstone



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 - Undersubscription of cores is an inefficient use of a user's core-hour allocation and users should in general not do this

MVAPICH Scaling



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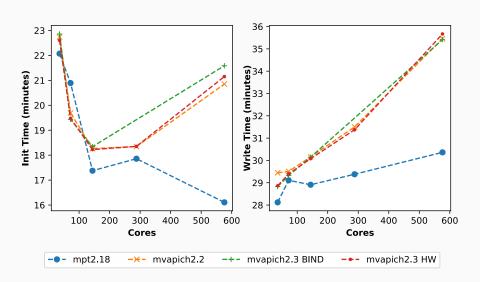


Fig. 8: MVAPICH CONUS 12 km Init and Write Scaling



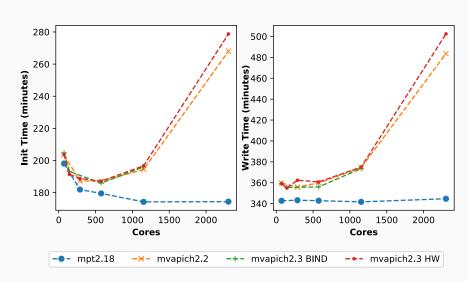


Fig. 9: MVAPICH Maria 1km Init and Write Scaling

Conclusion



- Intel compiler consistently faster than Gnu compiler
 - Users should use -fp-model fast=2 or -Ofast for a modest performance increase
- MPT, OpenMPI, and MVAPICH show similar performance while Intel MPI and MPICH have poorer performance
- WRF's initialization and writing time show improvements compared to previous results on Yellowstone with a previous WRF version
- WRF V4.0 scales well across entire run-able region
 - Will run out of memory on runs with too many of gridpoints per core
 - WRF will prevent runs with too few of gridpoints per core

Acknowledgments



- Davide Del Vento
- Brian Vanderwende
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