# Comp lecture #13

# Elements of MAL

- Directives start with '.':
- .data data segment,
- .text code segment.
- Labels end with ':'
- v:, main:, loop: and endloop:.
- Registers start with '\$'.
- Comments start with '#'.
- Instructions and macros (pseudo-instructions).
- Constants.

# Basic MAL Directives

.globl sym	Declares sym as global (accessible from other files)
.extern sym size	Declares sym as global with specified size
.ascii str	Saves str in memory
.asciiz str	Saves str and adds '\0' at the end
.byte b1, b2,, bn	Sequentially saves b1, b2,, bn allocating 8 bits for each
.half h1, h2,, hn	Sequentially saves h1, h2,, hn allocating 16 bits for each
.word w1, w2,, wn	Sequentially saves w1, w2,, wn allocating 32 bits for each
.dword dw1, dw2,, dwn	Sequentially saves dw1, dw2,, dwn allocating 64 bits for each
.float f1, f2,, fn	Sequentially saves floats f1, f2,, fn
double d1, d2,, dn	Sequentially saves double floats d1, d2,, dn
.space n	Allocates n bytes
.align n	Aligns data to 2 <sup>n</sup> bytes

# Labels in MAL

- in the .data segment, data labels define addresses of variables.
- in the .text segment instruction labels define addresses of instruction.

# **MIPS** Registers

- MIPS has 32 main register and 32 co processor registers
- Size of a register is 32 bits
- Registers saves address, data and instruction
- Max mem = 2<sup>32</sup> 4G bytes (doesnt make sense to use 64 bit if CPU is only 32-bit)
- All MIPS arithmetic operations use registers.
- Register's name starts with '\$' and has two

synonymic versions:

- Number
- symbolic name.

Register Number	Conventional Name	Usage
\$0	\$zero	Hard-wired to 0
\$1	\$at	Reserved for pseudo-instructions
\$2 - \$3	\$v0, \$v1	Return values from functions
\$4 - \$7	\$a0 - \$a3	Arguments to functions - not preserved by subprograms
\$8 - \$15	\$t0 - \$t7	Temporary data, not preserved by subprograms
\$16 - \$23	\$s0 - \$s7	Saved data registers, preserved by subprograms
\$24 - \$25	\$t8 - \$t9	More temporary registers, not preserved by subprograms
\$26 - \$27	\$k0 - \$k1	Reserved for kernel. Do not use.
\$28	\$gp	Global Area Pointer (base of global data segment)
\$29	\$sp	Stack Pointer
\$30	\$fp	Frame Pointer
\$31	\$ra	Return Address

### **MIPS Registers for Floating Point Operations**

Register Number	Conventional Name	Usage
\$f0 - \$f3	-	Floating point return values
\$f4 - \$f10	-	Temporary registers, not preserved by subprograms
\$f12 - \$f14	-	First two arguments to subprograms, not preserved by subprograms
\$f16 - \$f18	-	More temporary registers, not preserved by subprograms
\$f20 - \$f31	-	Saved data registers, preserved by subprograms

# **MAL Data Types**

Туре	Example	Size
Byte	a: .byte 0	8 bit
Halfword	b: .half 0	16 bit
Word	c: .word 0	32 bit
Doubleword	d: .dword 0	64 bit

- In MAL a type defines only the size of data.
- The interpretation of binary code depends on instruction.

# **Examples of Data Definitions**

- asciiz differe from .ascii in that .asciiz terminates

# **Source Code**

- Each line of code can be
  - Instruction (?:label)
  - Single directive
  - Empty line
  - Comment
- Comment starts:#, ends:newline

# **Special Instruction syscall**

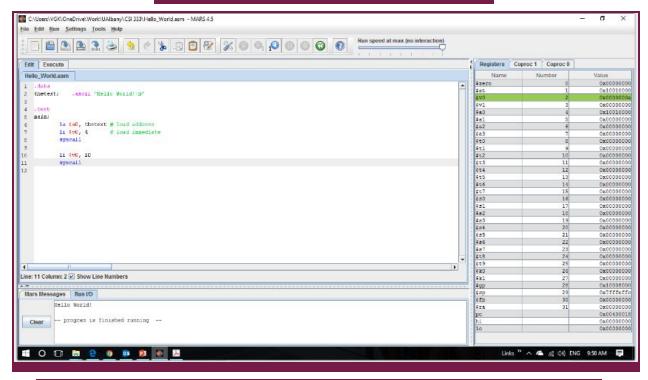
Service	Code in \$v0	Arguments
Print integer	1	\$a0 = integer to be printed
Print float	2	\$f12 = float to be printed
Print double	3	\$f12 = double to be printed
Print string	4	\$a0 = address of string in memory
Read integer to \$v0	5	
Read float to \$f0	6	
Read double to \$f0	7	
Read string	8	\$a0 = memory address of string input buffer \$a1 = length of string buffer (n)
Dynamic memory allocation	9	\$a0 = amount
Exit	10	
Print character	11	\$a0 = character to be printed
Read character in \$v0	12	

- Lets your code commmunicate with os
- If something goes wrong the OS usually tells you in
- If you mess up your processor, you can destroy your machine
- Use software (simulators) to run your mips code

# **MAL Simulators**

- Direct experiments with CPU can destroy it
- Use simulators such as SPIM and MARS to test your code

# **MAL Simulators**



- First part data that defines the strings
- The list to your right are 32 registers
- Before execution all registers are zero
- The values in the register change as they need to
- Change value of regsiter vo, to 4 becuase next instruction is syscall
- Goes to register vo to find what to do, print string at register (memory location) ao

- In MIPS you can execute your code step by step to see what is happening
- Exit called done by syscall, if code does not do it itself
- Values placed into registers are binary
- Use hexadecimals to save space
- ox tells you if a number is hexadecimal

### **SPIM Interface**

Make sure code runs well with SPIM simulator

### **MAL Instructions**

- Psuedo-instructions in the previous example
  - , <mark>la rdest, add</mark>r
- translates to
  - $_{\circ}$  lui at, hi(addr)
  - $_{\circ}$  ori rdest, \$at, lo(addr)

- Too hard to remember all instruction, but have to know
  - where code is
- Use psuedo-instruction

### **Types of Instructions**

- Type R(register)
  - $_{\odot}$  Three operands,
    - **Destination register (\$rd)**
    - **First argument (\$rs)**
    - Seconde argument (\$rt)
    - **add \$t2, \$t0, \$t1**
    - **\$t0 + \$t1 -> \$t2**
- **Type I (immediate)** 
  - $_{\odot}$  Two registers and constant
  - $_{\odot}$  addi \$t3, \$t2, 12
  - **\$12 + 12 -> \$13**

Type J (jump)
 one operand: new address for PC.
 E.g. j 128.

### TRANSlating C segments into MAL

- Assembly language can take a long time
- To introduce MAL instructions
- To convery level of detail involved

# Example 1

```
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A Segment in C:
                           Equivalent MAL Segment:
                           Register Assignment:
int f, g, h, i, j;
                           f: $15 g: $16 h: $17 i: $18 j: $19
/* ... */
f = (g + h) - (i + j);
                           add $8, $16, $17
                                                     #$8 has g
                           + h.
                           add $9, $18, $19
                                                      #$9 has i
                           + j.
                           sub $15, $8, $9
                                                      #$15 has
                           f.
```

### Example 2

```
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A Segment in C:
                         Equivalent MAL Segment:
                         Register Assignment:
int f, g, h, i, j;
                         f: $15 g: $16 h: $17 i: $18 j: $19
if (i == j)
        f -= i;
                                   beg $18, $19, then
                                   add $15, $16, $17 #f = g +
else
                         h.
         f = q + h;
                                   j exit
                         then: sub $15, $15, $18 #f -= i.
                         exit:
```

- Here the else statement comes right after the conditional and exits
- The then label comes to as representation of the if statement

# Example 3

#### A Segment in C:

Assume all variables are of type int.

```
if (a > 0) {
    x += y;
    p *= q;
}
```

#### **Equivalent MAL Segment:**

Register Assignment: a: \$15 x: \$16 y: \$17 p: \$18 q: \$19

```
bgt $15, $0, then
j end_if
then: add $16, $16, $17
mul $18, $18, $19
end_if:
```

### alt

#### A Segment in C:

Assume all variables are of type int.

```
if (a > 0) {
    x += y;
    p *= q;
}
```

#### Equivalent MAL Segment (alternative solution):

Register Assignment:

```
a: $15 x: $16 y: $17 p: $18 q: $19
```

```
ble $15, $0, end_if
add $16, $16, $17
mul $18, $18, $19
end if:
```

### **C and Assembly Conditional Operators**

C Conditional Operator	MIPS Assembly Instruction
a == b	beq \$t0, \$t1, then
a != b	bne \$t0, \$t1, then
a < b	blt \$t0, \$t1, then
a > b	bgt \$t0, \$t1, then
a <= b	ble \$t0, \$t1, then
a >= b	bge \$t0, \$t1, then
a == 0	beqz \$t0, then

# LOOPS

### TRANSLATING C SEGMENTS INTO MAL — FOR LOOP

#### A Segment in C: **Equivalent MAL Segment:** Assume all variables are of Register Assignment: type int. i: \$15 j: \$16 n: \$17 temporary: \$18. move \$16, \$0 #Set j = 0. j = 0;addi \$15, \$0, 1 #Set i = 1. for (i = 1; i <= n; i++) loop: bgt \$15, \$17, end\_loop #If i > n, loop is over. j += 2 \* i; mul \$18, \$15, 2 $\#Set\ temp = 2 * i.$ add \$16, \$16, \$18 #Set j += temp. addi \$15, \$15, 1 #i++ See Handout 13.1. j loop end loop:

### Important MAL Instructions: Load, Store and Move

#### General form:

```
lw R, label #Load to register from memory.
li R, constant #Load to register immediate.
sw R, label #Store in memory from register.
move Rd, Rs #Move from register to register.
```

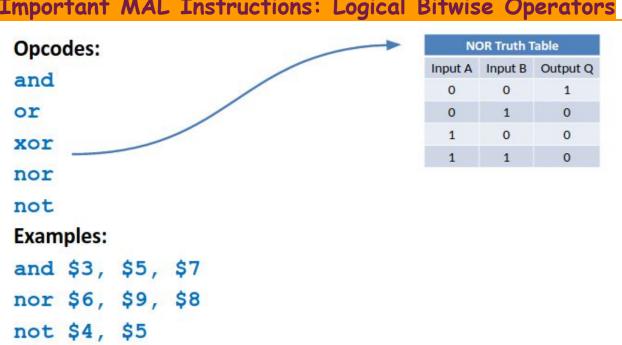
#### Example:

```
.data
val: .word 24
.text
lw $5, val
li $6, -21
move $7, $6
sw $7, val
```

# Important MAL Instructions: Arithmetic Instructions

```
General form:
add Rd, Rs1, Rs2
                     #Add
sub Rd, Rs1, Rs2
                     #Subtract
mul Rd, Rs1, Rs2
                    #Multiply
div Rd, Rs1, Rs2
                     #Divide
rem Rd, Rs1, Rs2
                     #Remainder
Examples:
sub $5, $3, $7
#subtracts the contents of $7 from that of $3 and
#stores the result in $5.
div $5, $3, $7
#divides the contents of $3 by the contents of $7 and
#stores the quotient in $5. (The remainder is ignored.)
rem $5, $3, $7
#divides the contents of $3 by the contents of $7 and
#stores the remainder in $5. (The quotient is ignored.)
```

### Important MAL Instructions: Logical Bitwise Operators



## **Important MAL Instructions:**

## **Shift Instructions**

#### General form:

```
sll Rd, Rs, const #Shift left logical
srl Rd, Rs, const #Shift right logical
sra Rd, Rs, const #Shift right arithmetic
```

The sll and srl instructions fill the vacated spots with zeros.

The sra instruction fills the vacated spots with the sign bit.

#### **Examples:**

```
sll $5, $5, 4
srl $6, $9, 3
sra $10, $10, 2
```

# Important MAL Instructions: Jump Instructions

- Unconditional Jump:
- j label #Jump using absolute addresses
  b label #Branch using PC-relative addresses
  Note: The difference is important for position-independent code.
- Conditional Jump:

```
beq R1, R2, label
beqz R1, label
```

Other Possible Jump Conditions:

```
bne blt bgt ble bge (see above)
bnez bltz bgtz blez bgez
```

# A Simple MAL Program

For 13.3 notes

### The syscall Instruction

- Used for I/O and for stopping the program
- The operation to be carried out is specified as a command (in register \$v0) to syscall.
- \$v0 is a synonym for \$2. This register must contain the command for syscall. This register also contains the return value (if any) produced by executing the command.
- \$a0 and \$a1 are synonyms for \$4 and \$5
  respectively. These registers must contain suitable
  values if such values are needed for the command

Using syscall to Read an Integer

MAL Code:

li \$v0, 5 syscall

Using syscall to Print an Integer

#### MAL Code:

Assume that we want to print the value in \$12.

```
move $a0, $12
li $v0, 1
syscall
```

### Using syscall to Print a String

#### MAL Code:

The following segment prints the string Hello followed by the newline character.

.data

hstr: .asciiz "Hello\n"

.text

la \$a0, hstr

li \$v0, 4

syscall

### Using syscall to Read a String, cont'd

Service	\$v0	Arguments	
Print integer	1	\$a0 = integer to be printed	
Print float	2	\$f12 = float to be printed	
Print double	3	\$f12 = double to be printed	
Print string	4	\$a0 = address of string in memory	
Read integer to \$v0	5		

Read float to \$f0	6	
Read double to \$f0	7	
Read string	8	\$a0 = memory address of string input buffer \$a1 = length of string buffer (n)
Dynamic memory allocation	9	\$a0 = amount
Exit	10	
Print character	11	\$a0 = character to be printed
Read character in \$v0	12	

#### MAL Code:

The following MAL segment prompts the user for a string and reads in a string consisting of at most 10 characters, including the newline character.

```
.data
hstr: .asciiz "Input: "
buf: .space 11

.text #Prompt the user.
la $a0, hstr
li $v0, 4
syscall

la $a0, buf #Read string.
li $a1, 11
li $v0, 8
syscall
```

# Using syscall to Stop a Program

### MAL Code:

li \$v0, 10 syscall

# Saving and Restoring Registers

```
Assume that we want to save and restore $2 and $4.

.data

s2: .word 0

s4: .word 0

.text

#

# <-- Code uses $2 and $4.

#

sw $2, s2 #Save registers.

sw $4, s4

#

# <-- Code uses $2 and $4 for I/O or other purposes

#

lw $2, s2 #Restore registers.

lw $4, s4

More common method - use the system stack (to be discussed later).
```