Algorithms and Data Structures 2

Sorting Algorithms

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Outline

- Overview of Sorting
 - Measuring Performance
- Sorting Algorithms
 - Bubble Sort
 - Insertion Sort
 - Merge Sort
 - Selection Sort
 - Quick Sort
- Visualizing Sorting Performance
 - Sample Source Code and Demo



Sorting Overview

- Arranging data in a collection based on a comparison algorithm
 - E.g., Any object with a notion of greater-than/less-than/equality
- Two general families of sorting algorithms
 - Linear Sorting
 - Divide and Conquer
- Linear algorithms treat the problem of sorting as a single large operation.
- Divide and Conquer algorithms partition the data to be sorted into smaller sets that can be independently sorted.



Measuring Performance

Comparisons

- When two values of the input array are compared for relative equality
 - Equal to, Greater then, Less then

Swaps

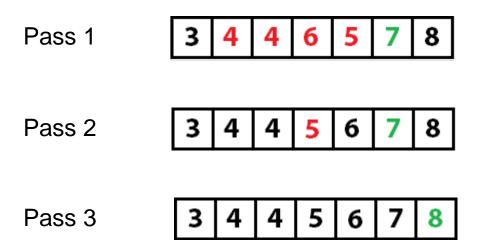
- When two values stored in the input array are swapped
- \Box E.g., [1, 0] => [0, 1]

Performance Considerations

- Comparisons and Swaps both have a cost
- Reducing either or both can improve performance
- The cost of both operations depends on many factors.

Bubble Sort

- Simplest sorting algorithm
- On Each Pass
 - Compare each array item to it's right neighbor
 - If the right neighbor is smaller then Swap right and left
 - Repeat for the remaining array items
- Perform subsequent passes until no swaps are performed





Bubble Sort Performance

- Worst case performance: O(n²)
 - Not appropriate for large unsorted data sets.
- Average case performance: O(n²)
 - Not appropriate for large unsorted data sets.
- Best case performance: O(n)
 - Very good best case performance and can efficiently sort small and nearly sorted data sets
- Space required: O(n)
 - Bubble sort operates directly on the input array meaning it is a candidate algorithm when minimizing space is paramount.



Insertion Sort

- Sorts each item in the array as they are encountered
- As the current item works from left to right
 - Everything left of the item is known to be sorted
 - Everything to the right is unsorted
- The current item is "inserted" into place within the sorted section





Insertion Sort Performance

- Worst case performance: O(n²)
 - Not appropriate for large unsorted data sets.
- Average case performance: O(n²)
 - Not appropriate for large unsorted data sets.
- Best case performance: O(n)
 - Very good best case performance and can efficiently sort small and nearly sorted data sets
- Space required: O(n)
 - Insertion sort operates directly on the input array meaning it is a candidate algorithm when minimizing space is paramount.



Merge Sort

- The array is recursively split in half recursively
- When the array is in groups of 1, it is reconstructed in sort order
- Each reconstructed array is merged with the other half

3 8 2 1 5 4 6 7



Merge Sort Performance

- Worst case performance: O(n log n)
 - Appropriate for large data sets
 - Data splitting means that the algorithm can be made parallel.
- Average case performance: O(n log n)
 - Appropriate for large data sets
- Best case performance: O(n log n)
 - Appropriate for large data sets
- Space required: O(n)
 - Merge can be, but is often not, performed in-place. These extra allocations increase the memory footprint required to sort data.



Selection Sort

 Sorts the data by finding the smallest item and swapping it into the array in the first unsorted location.

• Algorithm:

- Enumerate the array from the first unsorted item to the end
- Identify the smallest item
- Swap the smallest item with the first unsorted item





Selection Sort Performance

- Worst case performance:O(n²)
 - Not appropriate for large unsorted data sets.
- Average case performance:O(n²)
 - Not appropriate for large unsorted data sets.
 - Typically performs better than bubble but worse than insertion sort
- Best case performance:O(n²)
 - Not appropriate for large unsorted data sets.
- Space required: O(n)
 - Selection sort operates directly on the input array meaning it is a candidate algorithm when minimizing space is paramount.



Quick Sort

- Divide and Conquer algorithm
- Pick a pivot value and partition the array
- Put all values before the pivot to the left and above to the right
 - The pivot point is now sorted everything right is larger, everything left is smaller.
- Perform pivot and partition algorithm on the left and right partitions
- Repeat until sorted





Quick Sort Performance

- Worst case performance:O(n²)
 - Not appropriate for large pathologically sorted (inverse sorted) data sets.
- Average case performance:O(n log n)
 - Appropriate for large data sets
- Best case performance:O(n log n)
 - Very good best case performance and can efficiently sort small and nearly sorted data sets
- Space required: O(n)
 - As a recursive algorithm the array space as well as the stack space must be considered. There exist optimizations to reduce space usage further.

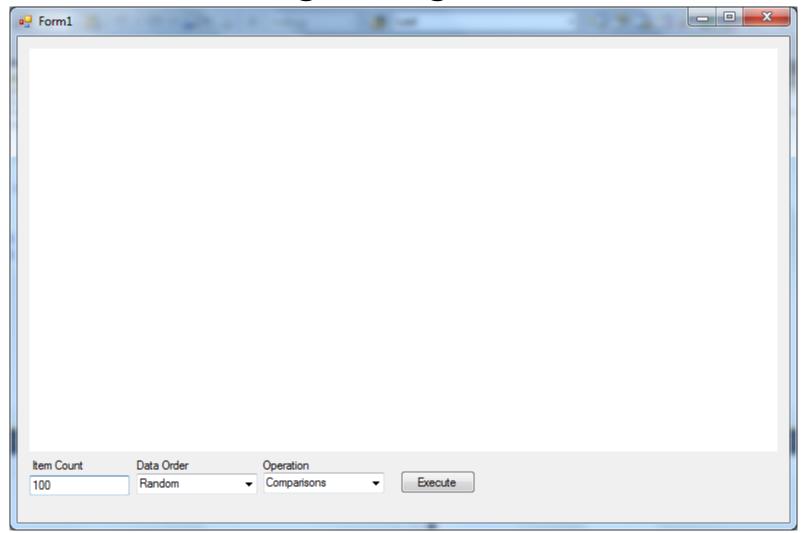


Visualizing Sorting Performance

- Graphically comparing the cost of sorting various data sets
- Variable Item Counts
 - Ability to set the item count to any range from 0 to billions.
- Variable data sorting
 - Pre-Sorted
 - Reversed Sorted
 - Random Data
- Measurement type
 - Comparisons
 - Swaps (or merge sort assignments)

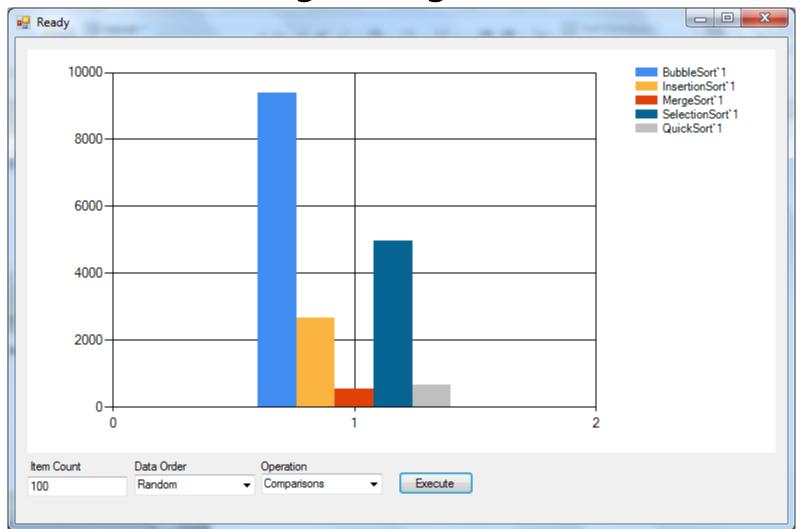


Visualizing Sorting Performance





Visualizing Sorting Performance





Summary

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 - Quick Sort
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References

Wikipedia.org Sorting Articles

- http://en.wikipedia.org/wiki/Bubble_sort
- http://en.wikipedia.org/wiki/Insertion_sort
- http://en.wikipedia.org/wiki/Merge_sort
- http://en.wikipedia.org/wiki/Selection_sort
- http://en.wikipedia.org/wiki/Quicksort

MSDN Charting API (Sample Application)

- System.Windows.Forms.DataVisualization.Charting.Chart
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