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School of Computer Science

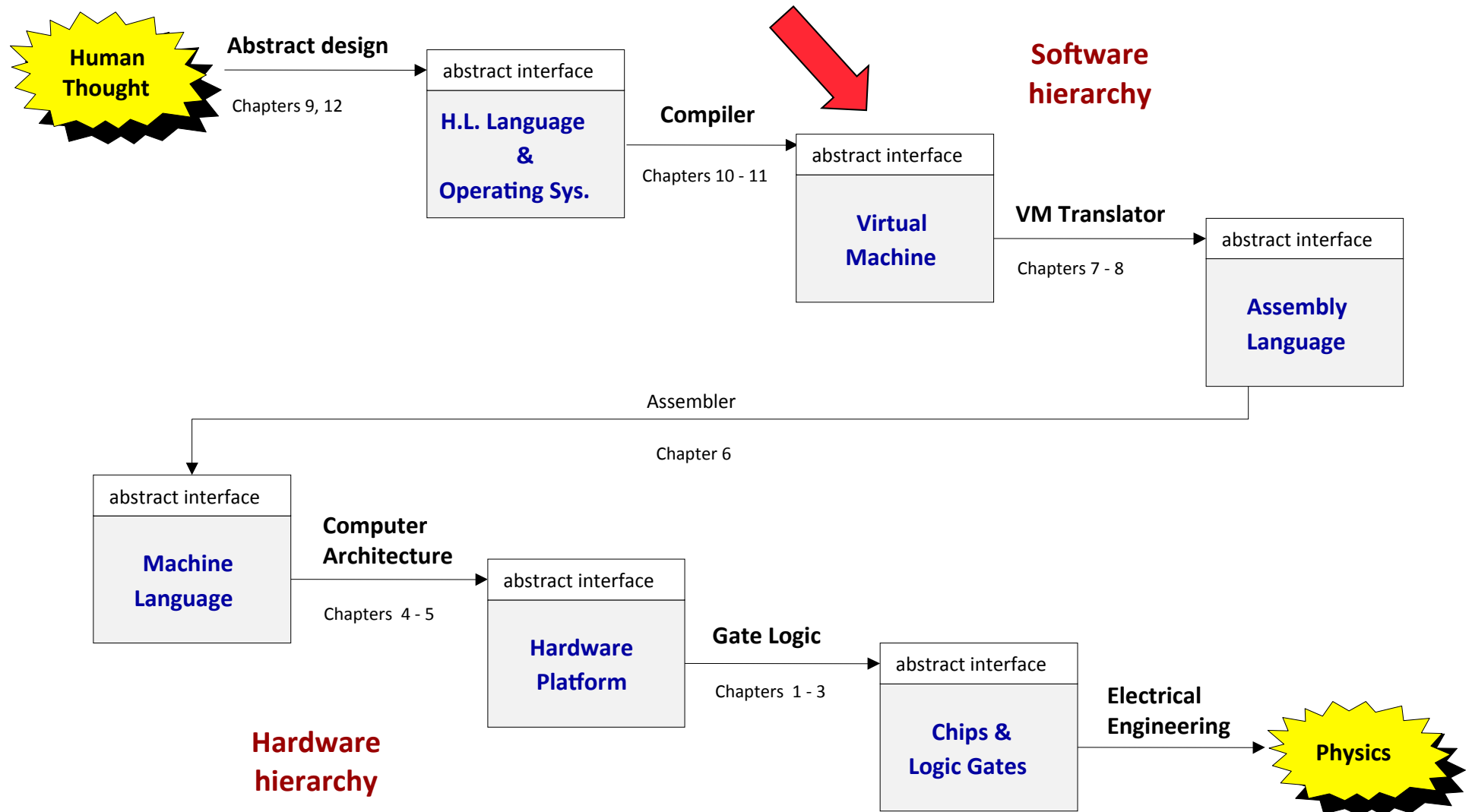
# COMP SCI 2000 Computer Systems

## Lecture 12

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*seek* LIGHT

# Review - where we are:





# Motivation

## Jack code (example)

```
class Main
{
    static int x;

    function void main()
    {
        // Inputs and multiplies two numbers
        var int a, b, x;
        let a = Keyboard.readInt("Enter a number");
        let b = Keyboard.readInt("Enter a number");
        let x = mult(a,b);
        return;
    }
}

// Multiplies two numbers.
function int mult(int x, int y)
{
    var int result, j;
    let result = 0; let j = y;
    while ~(j = 0)
    {
        let result = result + x;
        let j = j - 1;
    }
    return result;
}
```

Our ultimate goal:

Translate high-level  
programs into  
executable code.



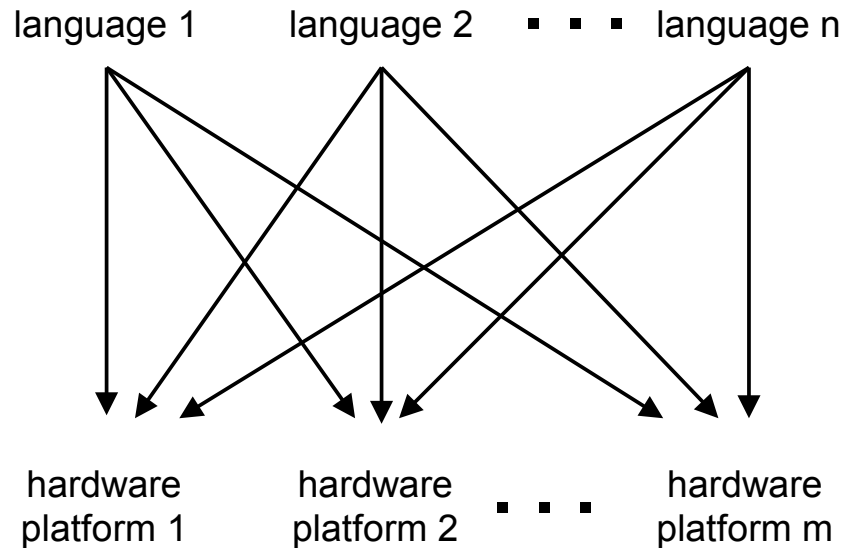
Compiler

## Hack code

```
0000000000010000
1110111111001000
0000000000010001
1110101010001000
0000000000010000
1111110000010000
0000000000000000
1111010011010000
0000000000010010
1110001100000001
0000000000010000
1111110000010000
0000000000010001
0000000000010000
1110111111001000
0000000000010001
1110101010001000
0000000000010000
1111110000010000
0000000000000000
1111010011010000
0000000000010010
1110001100000001
0000000000010000
1111110000010000
0000000000010001
...
```

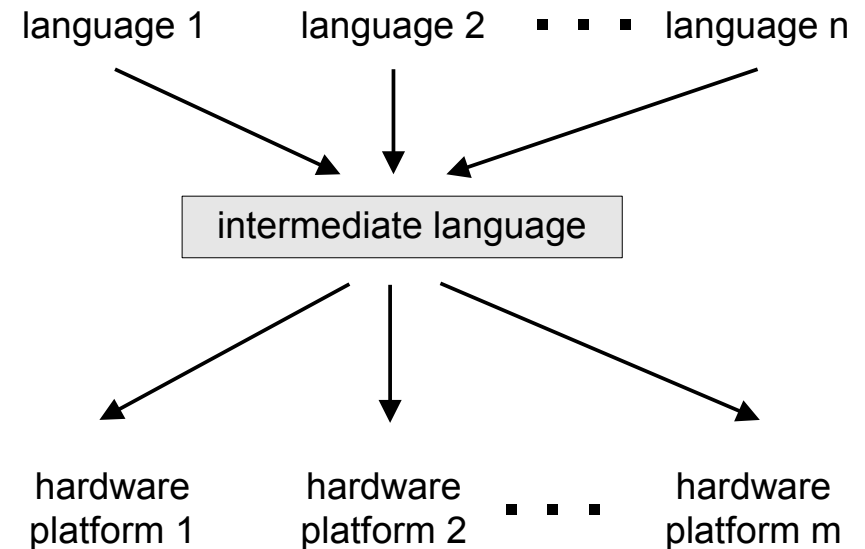
# Compilation models

## direct compilation:



requires  $n \cdot m$  translators

## 2-tier compilation:

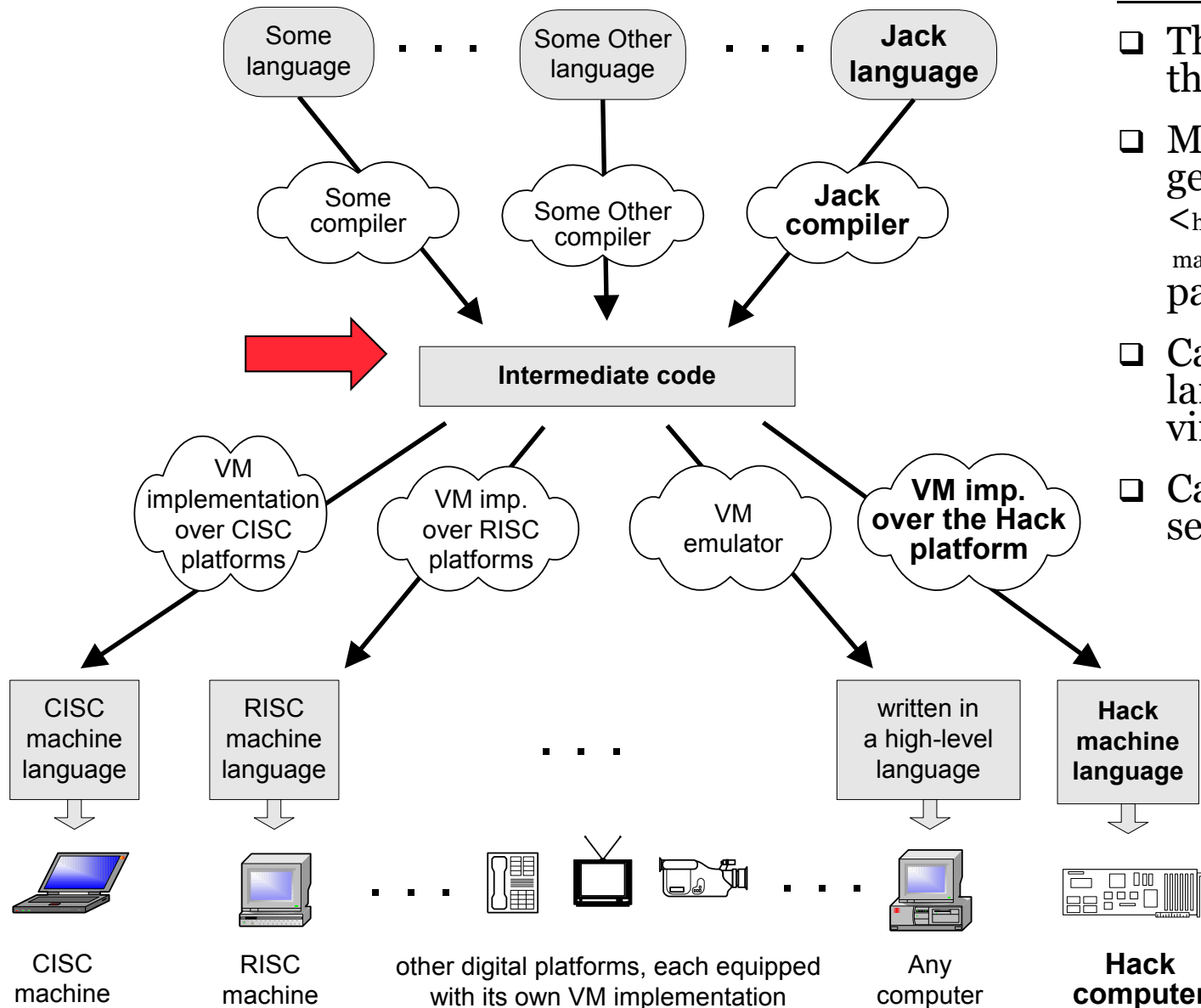


requires  $n + m$  translators

## Two-tier compilation:

- ❑ First compilation stage: depends only on the details of the source language
- ❑ Second compilation stage: depends only on the details of the target language.

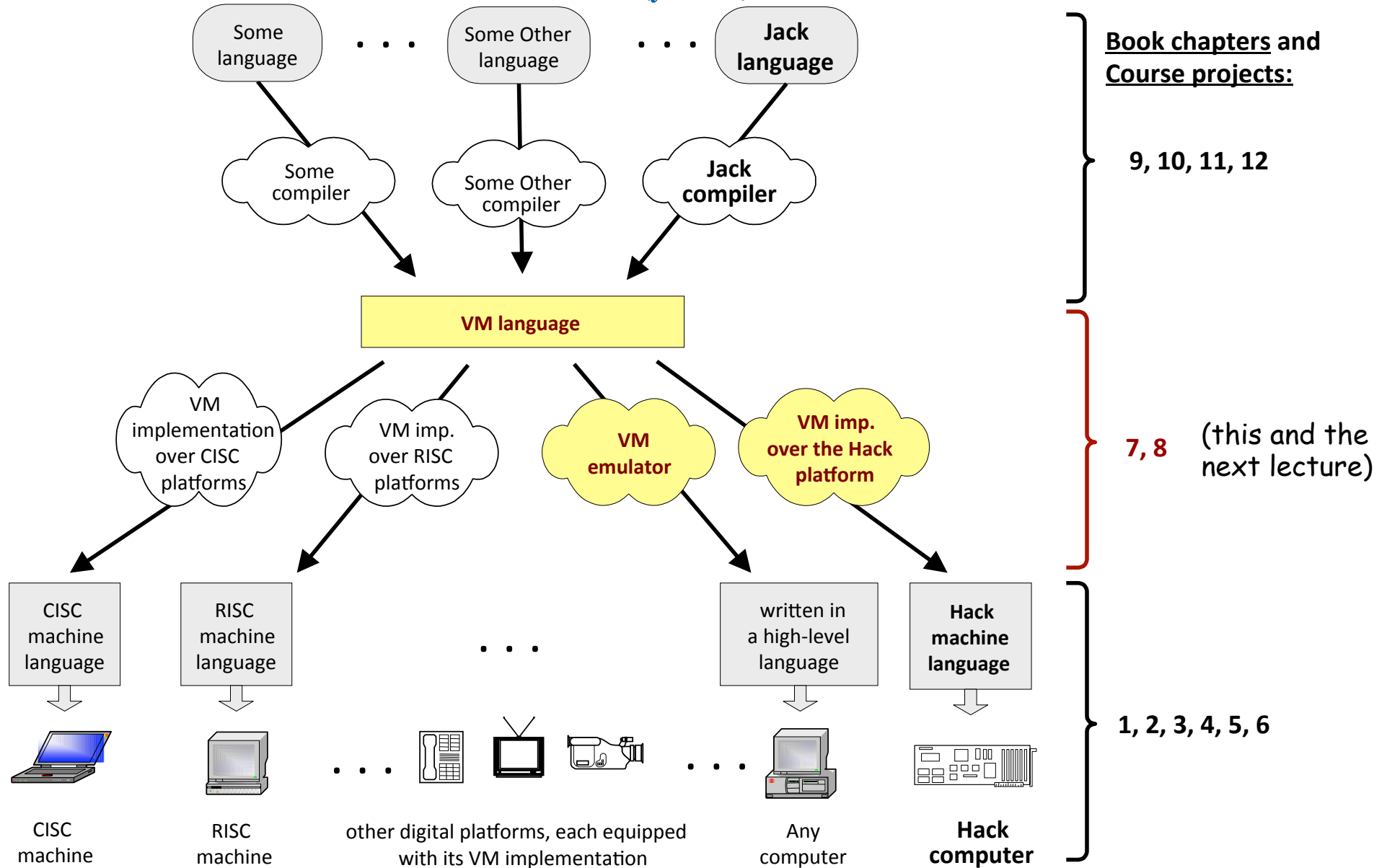
# The big picture



## The intermediate code:

- ❑ The interface between the 2 compilation stages
- ❑ Must be sufficiently general to support many <high-level language, machine-language> pairs
- ❑ Can be modeled as the language of an abstract virtual machine (VM)
- ❑ Can be implemented in several different ways.

# Focus of this lecture (yellow):



# Other Examples of Virtual Machines

- Worksheet Lecture 12, Question 1
  - Examples of virtual machines and their uses

# The VM model and language

## Perspective:

From here till the end of the next lecture we describe the VM model used in the Hack-Jack platform

Other VM models (like Java's JVM/JRE and .NET's IL/CLR) are similar in spirit but differ in scope and details.

## Several different ways to think about the notion of a virtual machine:

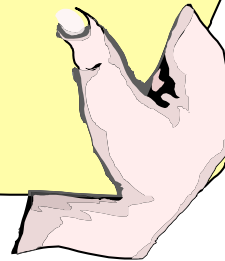
- ❑ **Abstract software engineering view:**  
the VM is an interesting abstraction that makes sense in its own right
- ❑ **Practical software engineering view:**  
the VM code layer enables “managed code” (e.g. enhanced security)
- ❑ **Pragmatic compiler writing view:**  
a VM architecture makes writing a compiler much easier  
(as we'll see later in the course)
- ❑ **Opportunistic empire builder view:**  
a VM architecture allows writing high-level code once and have it run on many target platforms with little or no modification.



## Yet another view (poetic)

"programmers are creators of universes for which they alone are responsible. Universes of virtually unlimited complexity can be created in the form of computer programs."

(Joseph Weizenbaum)



Our VM model + language are an example of one such universe.

# Lecture plan

Goal: Specify and implement a VM model and language:

## Arithmetic / Boolean commands

add  
sub  
neg  
eq  
gt  
lt  
and  
or  
not

This lecture

## Memory access commands

pop x (pop into x, which is a variable)  
push y (y being a variable or a constant)

## Program flow commands

label (declaration)  
goto (label)  
if-goto (label)

Next lecture

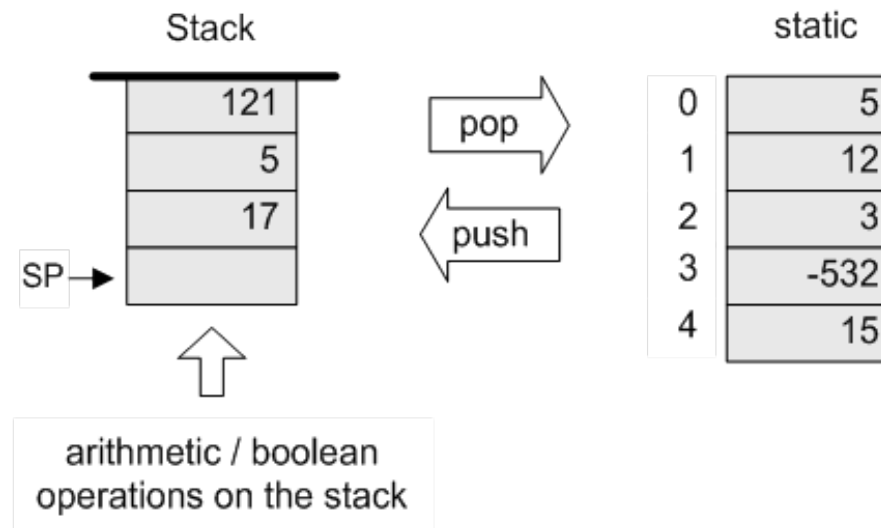
## Function calling commands

function (declaration)  
call (a function)  
return (from a function)

Our game plan: (a) describe the VM abstraction (above)  
(b) propose how to implement it over the Hack platform.

# Our VM model is *stack-oriented*

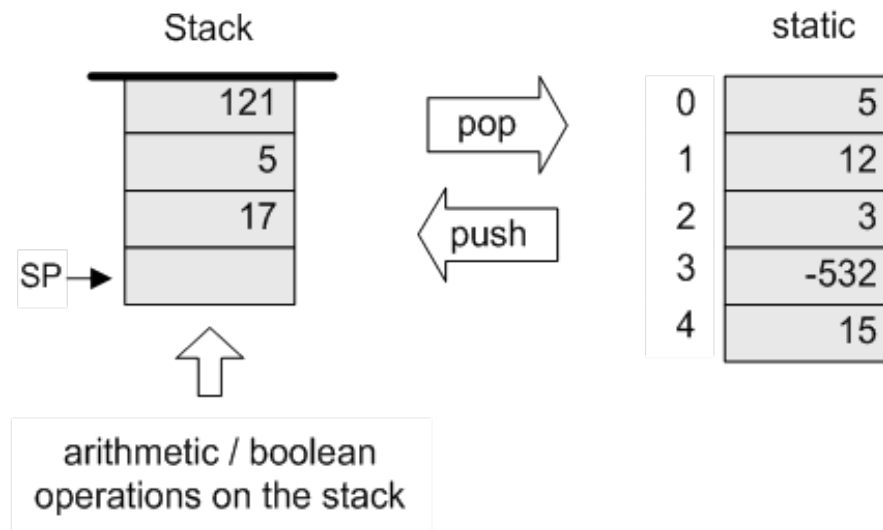
- All operations are done on a stack
- Data is saved in several separate *memory segments*
- All the memory segments behave the same
- One of the memory segments is called *static*, and we will use it (as an arbitrary example) in the following examples:



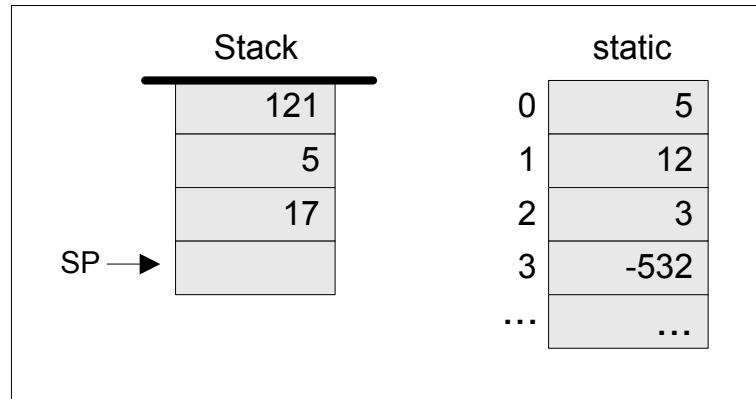
# Data types

Our VM model features a single 16-bit data type that can be used as:

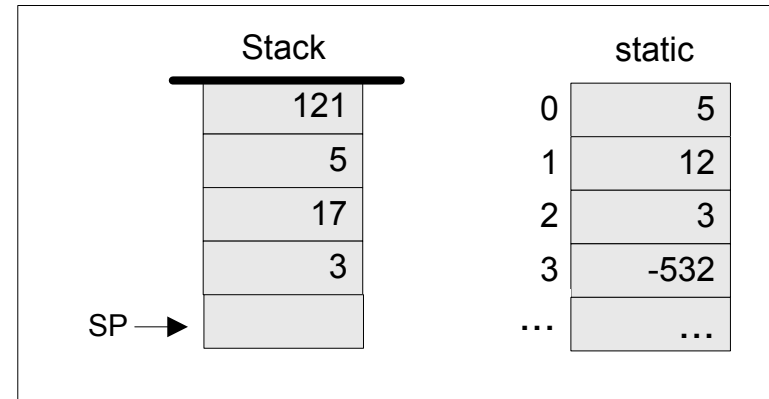
- ❑ an integer value (16-bit 2's complement: -32768, ... , 32767)
- ❑ a Boolean value (0 and -1, standing for true and false)
- ❑ a pointer (memory address)



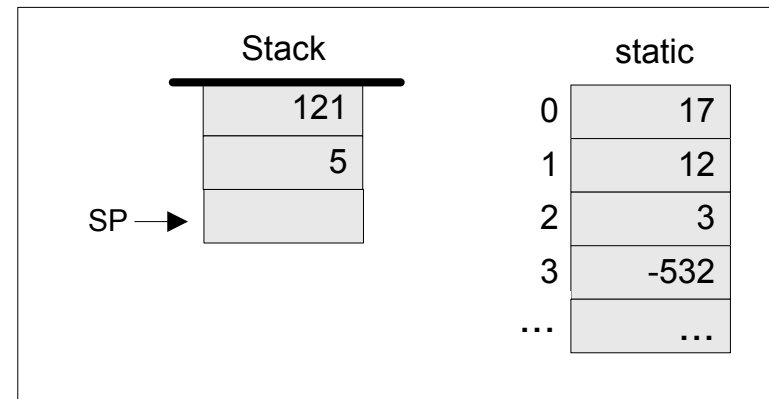
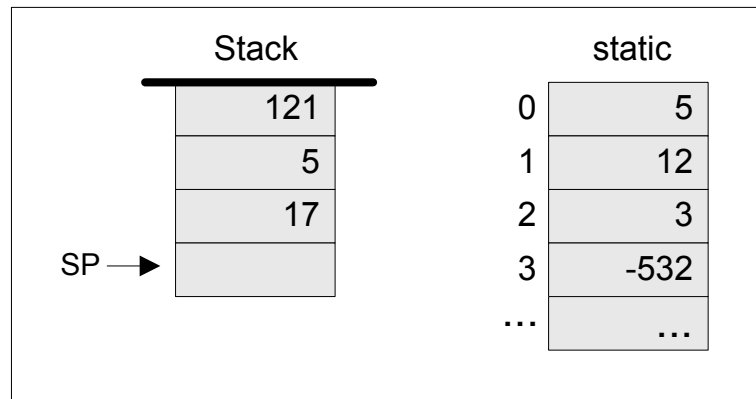
# Memory access operations



(before)



(after)



## The stack:

- A classical LIFO data structure
- Elegant and powerful
- Several hardware / software implementation options.

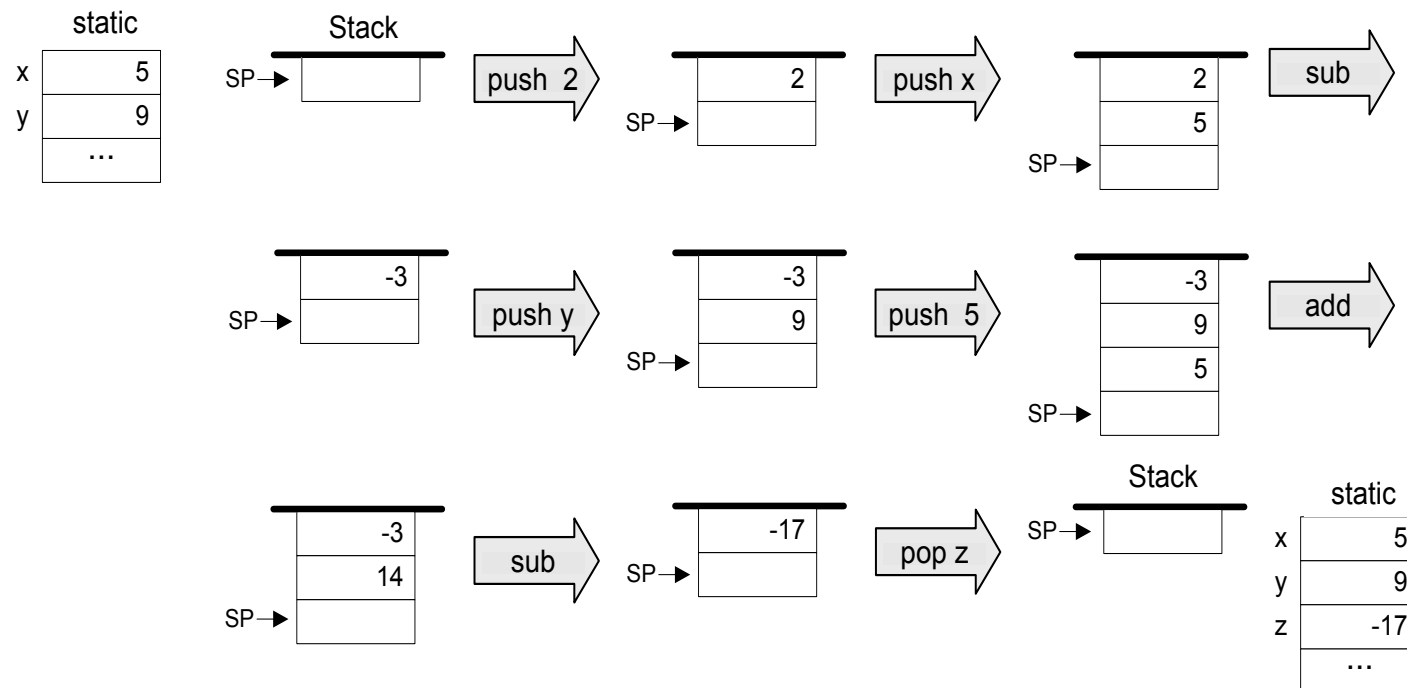


# Evaluation of arithmetic expressions

## VM code (example)

```
// z=(2-x)-(y+5)
push 2
push x
sub
push y
push 5
add
sub
pop z
```

(suppose that  
x refers to static 0,  
y refers to static 1, and  
z refers to static 2)

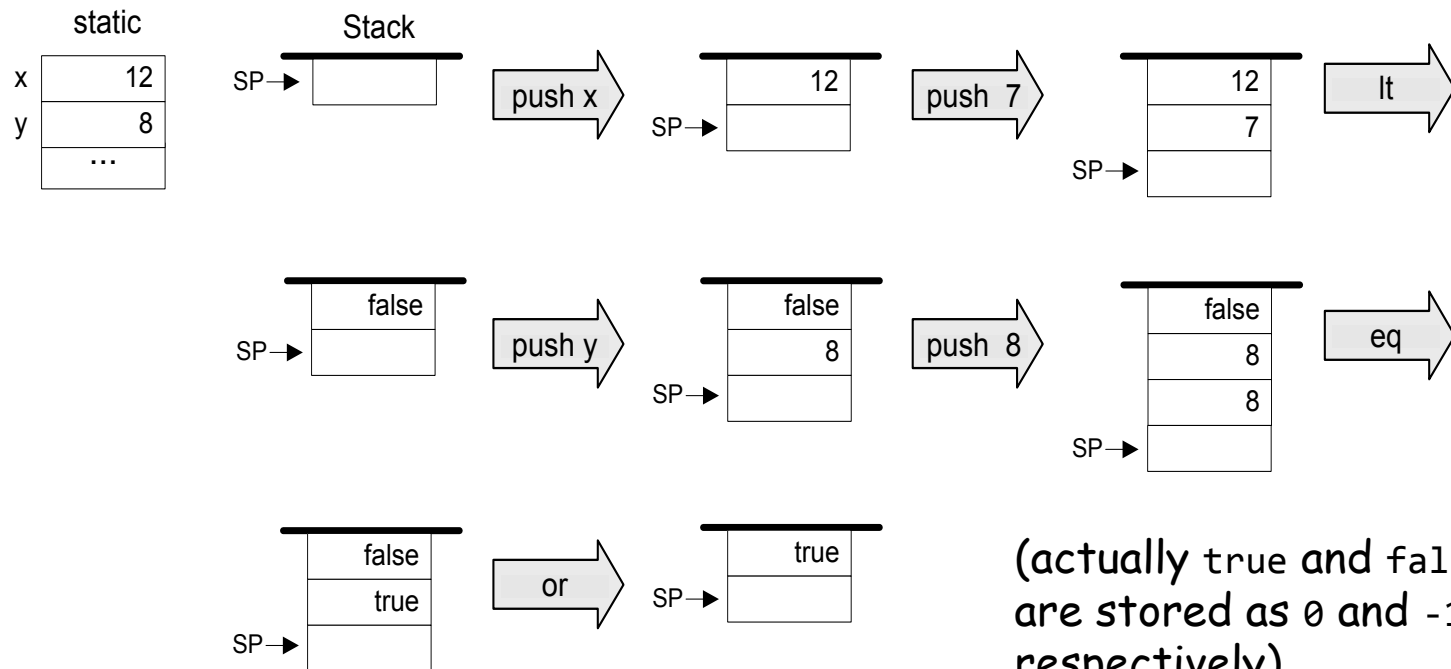


# Evaluation of Boolean expressions

## VM code (example)

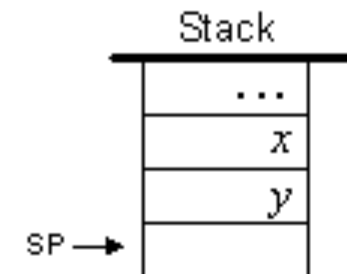
```
// (x<7) or (y=8)
push x
push 7
lt
push y
push 8
eq
or
```

(suppose that  
x refers to static 0, and  
y refers to static 1)



# Arithmetic and Boolean commands in the VM language (wrap-up)

Command	Return value (after popping the operand/s)	Comment
add	$x + y$	Integer addition (2's complement)
sub	$x - y$	Integer subtraction (2's complement)
neg	$-y$	Arithmetic negation (2's complement)
eq	true if $x = y$ and false otherwise	Equality
gt	true if $x > y$ and false otherwise	Greater than
lt	true if $x < y$ and false otherwise	Less than
and	$x \text{ And } y$	Bit-wise
or	$x \text{ Or } y$	Bit-wise
not	$\text{Not } y$	Bit-wise



# Stack Machine Examples

- Worksheet lecture 12, Question 2
  - Translating expressions into Hack Virtual Machine code

# Next lecture

- Virtual machine memory model
- Translating VM code into assembly language
- Keep working on your next assignment.
- Questions?