



ALU instructions

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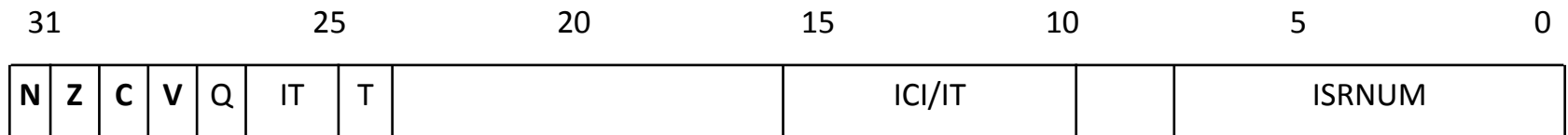
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Flags

- The uppermost nibble of the program status register contains the flags.



- Flags are set and cleared by means of:
 - comparison instructions, e.g., CMP and TST
 - ALU instructions if they have the suffix 'S'
 - a direct write to the program status register.

Carry flag C

- After a sum, C = 1 if the result size is 33 bit.

$$\begin{array}{r} 1010\ 0001\ 0001\ 0001\ 0001\ 0001\ 0001\ 0001\ + \\ 1011\ 0101\ 0101\ 0101\ 0101\ 0101\ 0101\ 0101\ = \\ \hline \text{C} = 1\ 0101\ 0110\ 0110\ 0110\ 0110\ 0110\ 0110\ 0110 \end{array}$$

- After a subtraction, C is inverted. Here C = 1:

$$\begin{array}{r} 1101\ 0000\ 0000\ 0000\ 0000\ 0000\ 0000\ 0000\ - \\ 0101\ 0000\ 0000\ 0000\ 0000\ 0000\ 0000\ 0000\ = \\ \hline 1000\ 0000\ 0000\ 0000\ 0000\ 0000\ 0000\ 0000 \end{array}$$

- After a move or logical instruction, C is the result of an inline barrel shifter operation.

Negative flag N

- It corresponds to the first bit of the result.
- If $N = 1$, a 2's complement number is negative.
- Example: $-4 + (-3) = -7 \rightarrow N = 1$

$$\begin{array}{r} 1111\ 1111\ 1111\ 1111\ 1111\ 1111\ 1111\ 1100\ + \\ 1111\ 1111\ 1111\ 1111\ 1111\ 1111\ 1111\ 1101\ = \\ \hline N = 1111\ 1111\ 1111\ 1111\ 1111\ 1111\ 1111\ 1001 \end{array}$$

- Why $N = 1$ in the following sum?

$$\begin{array}{r} 0111\ 0001\ 0001\ 0001\ 0001\ 0001\ 0001\ 0001\ + \\ 0011\ 0101\ 0101\ 0101\ 0101\ 0101\ 0101\ 0101\ = \\ \hline N = 1010\ 0110\ 0110\ 0110\ 0110\ 0110\ 0110\ 0110 \end{array}$$

Overflow flag V

- It is set if, in a sum of values with the same sign, there is a change in the MSB

This is a carry

$$\begin{array}{r} 1010\ 0001\ 0001\ 0001\ 0001\ 0001\ 0001\ 0001\ 0001\ 0001\ + \\ 1011\ 0101\ 0101\ 0101\ 0101\ 0101\ 0101\ 0101\ 0101\ 0101\ = \\ \hline (V = 1)\ 1\ 0101\ 0110\ 0110\ 0110\ 0110\ 0110\ 0110\ 0110\ 0110 \end{array}$$

- $V = 0$ if both conditions occur: the result is right

$$\begin{array}{r} 1110\ 0001\ 0001\ 0001\ 0001\ 0001\ 0001\ 0001\ 0001\ 0001\ + \\ 1011\ 0101\ 0101\ 0101\ 0101\ 0101\ 0101\ 0101\ 0101\ 0101\ = \\ \hline (V = 0)\ 1\ 1001\ 0110\ 0110\ 0110\ 0110\ 0110\ 0110\ 0110\ 0110 \end{array}$$

Zero flag Z

- It is set if the result is zero.

1111 1111 1111 1111 1111 1111 1111 1101 +
0000 0000 0000 0000 0000 0000 0000 0011 =

Z = 0000 0000 0000 0000 0000 0000 0000 0000

Comparison instructions

- They compare the value or test some bits:

`compare/test <Rd>, <operand2>`

- They set the flags without updating Rd.
- The second operand can be:
 - a register with an optional shift
 - a constant obtained by shifting left an 8-bit value
 - a constant of the form 0x00XY00XY
 - a constant of the form 0xXY00XY00
 - a constant of the form 0xXYXYXYXY.

CMP and CMN

- **CMP (compare)** subtracts `operand2` from `Rd` and updates the flags.
- **CMN (compare negative)** adds `operand2` to `Rd` and updates the flags.
 - The execution of the arithmetic operations at the base of the `CMP` and `CMN` is not modifying the content of the operands.
- **Examples with `r0 = #12`**
 - `CMP r0, #10` -> `N = 0, Z = 0, C = 1, V = 0`
 - `CMP r0, #12` -> `N = 0, Z = 1, C = 1, V = 0`
 - `CMP r0, #14` -> `N = 1, Z = 0, C = 0, V = 0`
- **The assembler can change `CMP` and `CMN`:**
 - `CMP r0, #-8` becomes `CMN r0, #8`

TST and TEQ

- TST (test) computes the logical AND between `operand2` and `Rd`; then updates all the flags except `V`.
- TEQ (test equivalence) computes the logical EOR between `operand2` and `Rd`; then updates all the flags except `V`.
- Examples with `r0 = #1`
 - TST `r1, r0, LSL#4` -> is the 4th bit of `r1` set?
 - TEQ `r2, r3` -> are `r2` and `r3` equal?

Accessing Program Status Register

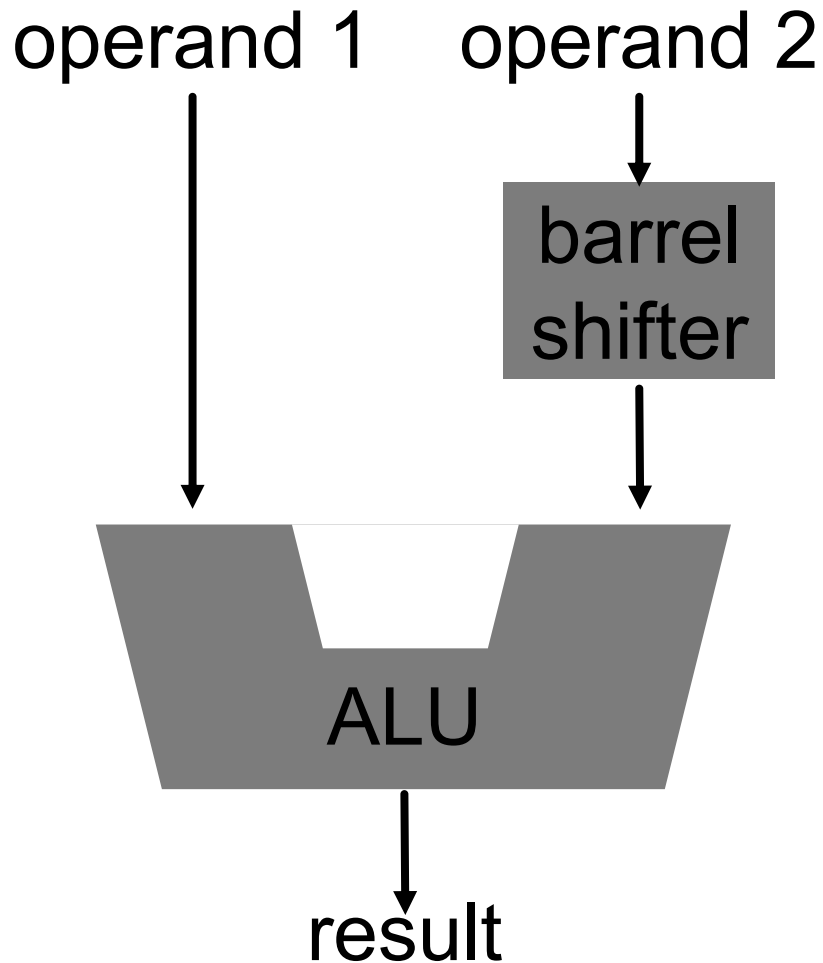
- `MRS <Rn>, <Sreg>` copies a special register into a register.
- `MSR <Sreg>, <Rn>` copies a general purpose register into a special register.
- `Sreg` can be APSR, EPSR, IPSR, and PSR.
- `MRS r0, APSR` reads the flags and copies them to the uppermost nibble of `r0`.

ALU instructions and flags

- ALU can perform:
 - arithmetic operations
 - logic operations
 - shift and rotate operations.
- The flags are affected only if the suffix '**S**' is appended to the instruction.
- ```
LDR r0, =0xFFFFFFFF9
ADDS r1, r0, #7
```

Flags are N = 0, Z = 1, C = 1, V = 0.

# Internal data path to ALU



Operand 1 is a register.

Operand 2 can be:

- shifted register
- shifted 8-bit constant
- 0x00XY00XY
- 0xXY00XY00
- 0xXYXYXYXY
- 12-bit constant (ADD and SUB only).

# Arithmetic instructions

- Addition
  - addition with carry
- Subtraction
  - subtraction with carry
  - reverse subtraction
- Multiplication
  - multiplication with accumulation
  - multiplication with subtraction
- Division

# Addition

- `ADD <Rd>, <Rn>, <op2>`     $Rd = Rn + op2$
- `ADC <Rd>, <Rn>, <op2>`  
     $Rd = Rn + op2 + C$
- **ADDW is like ADD, but it takes only a 12-bit value and it can not update flags.**
- **With ADC it is possible to add 64-bit values:**

`ADDS r4, r0, r2`

`ADC r5, r1, r3`

**In the example** `r5, r4 = r1, r0 + r3, r2`

# Subtraction

- SUB <Rd>, <Rn>, <op2>     $Rd = Rn - op2$

- SBC <Rd>, <Rn>, <op2>

$$Rd = Rn - op2 + C - 1$$

- SUBW is like SUB, but it takes only a 12-bit value and it can not update flags.

- With SBC it is possible to subtract 64-bit values

SUBS r4, r0, r2

SBC r5, r1, r3

In the example  $r5, r4 = r1, r0 - r3, r2$

# Reverse subtraction

- RSB <Rd>, <Rn>, <op2>

$$Rd = op2 - Rn + C - 1$$

- Advantages:

- either one or the other operand can be shifted before the subtraction

SUB r0, r1, r2, LSL #2 ; r0 = r1 - r2\*4

RSB r0, r2, r1, LSL #2 ; r0 = r1\*4 - r2

- a register can be subtracted from a constant.



# Multiplication

- multiplication with 32-bit result

MUL <Rd>, <Rn>, <Rm>

- unsigned multiplication with 64-bit result

UMULL <Rd1>, <Rd2>, <Rn>, <Rm>

- signed multiplication with 64-bit result

SMULL <Rd1>, <Rd2>, <Rn>, <Rm>

- Note 1: there is no distinction between signed and unsigned multiplication with 32-bit result.
- Note 2: all operands must be registers.

# Multiplication with accumulation

- MLA <Rd>, <Rn>, <Rm>, <Ra>

$$Rd = Rn * Rm + Ra$$

- MLS <Rd>, <Rn>, <Rm>, <Ra>

$$Rd = Rn * Rm - Ra$$

- UMLAL <Rd1>, <Rd2>, <Rn>, <Rm>

$$Rd1, Rd2 = Rn * Rm + Rd1, Rd2$$

- SMLAL <Rd1>, <Rd2>, <Rn>, <Rm>

same as UMLAL, but with signed values.

# Division

- unsigned division

UDIV <Rd>, <Rn>, <Rm>

- signed division

SDIV <Rd>, <Rn>, <Rm>

- If  $R_n$  is not exactly divisible by  $R_m$ , the result is rounded toward zero.
- UDIV and SDIV do not change the flags (the suffix 'S' can not be added).

# Logic instructions

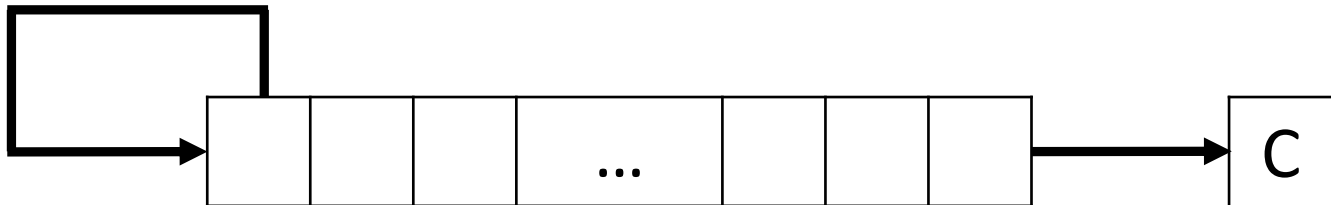
- AND <Rd>, <Rn>, op2 ; Rn **AND** op2
- BIC <Rd>, <Rn>, op2 ; Rn **AND NOT** op2
- ORR <Rd>, <Rn>, op2 ; Rn **OR** op2
- EOR <Rd>, <Rn>, op2 ; Rn **XOR** op2
- ORN <Rd>, <Rn>, op2 ; Rn **OR NOT** op2
- MVN <Rd>, <Rn> ; **NOT** Rn

# Shift instructions

- LSL <Rd>, <Rn>, <op2>

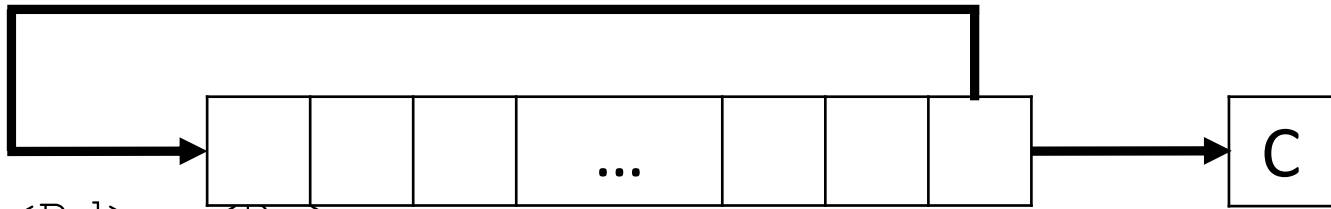


- LSR <Rd>, <Rn>, <op2>

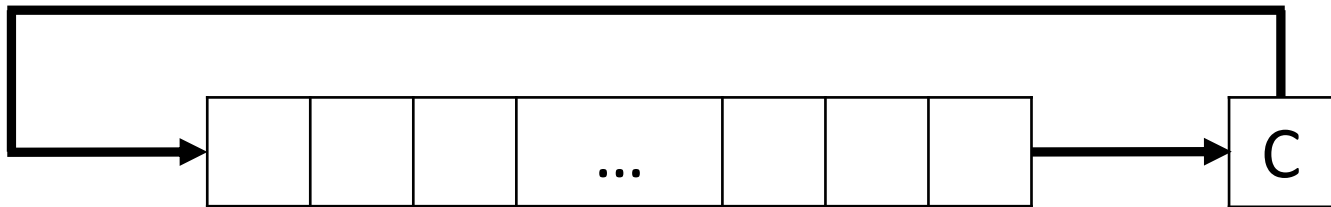


# Rotate instructions

- ROR <Rd>, <Rn>, <op2>



- RRX <Rd>, <Rn>



## Exercise: 64-bit routines for shift

- shift left  $r1, r0$  by one bit and save the result in  $r3, r2$ .
- logical shift right  $r1, r0$  by one bit and save the result in  $r5, r4$ .
- arithmetic shift right  $r1, r0$  by one bit and save the result in  $r7, r6$ .
- rotate right  $r1, r0$  by one bit and save the result in  $r9, r8$ .
- rotate right through carry flag  $r1, r0$  by one bit and save the result in  $r11, r10$ .