



Branches

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Unconditional branch

- There are four instructions for unconditional branch:
 - branch `B <label>`
 - branch indirect `BX <Rn>`
 - branch and link `BL <label>`
 - branch indirect with link `BLX <Rn>`
- `BL` and `BLX` save the return address (i.e., the address of the next instruction) in `LR (r14)` and they are used to call subroutines.

Infinite loops

- A stand-alone program, without operating system, can not continue beyond the end, otherwise the behavior is unpredictable.
- An infinite loop is added as last instruction:

```
stop B stop
```

or

```
LDR r1, =stop
```

```
stop BX r1
```

LDR Vs ADR with BX

- Instructions are 16 or 32 bits long, so their address is always halfword aligned.
- BX requires that the last bit of the register is 1, otherwise a usage fault exception is raised.
- BX jumps to the address created by changing the last bit of the register to 0.
- LDR sets the last bit to 1 if the label is in a code area, to 0 if the label is in a data area.
- ADR and ADRL do not change the last bit.

LDR Vs ADR: example

1.	0x000000CC	LDR r1, =stop
	0x000000CE	stop BX r1

Ok: r1 = 0x000000CF

2.	0x000000CC	ADRL r1, stop
	0x000000D4	stop BX r1

No: r1 = 0x000000D4

3.	0x000000CC	ADRL r1, stop
	0x000000D4	ORR r1, r1, #1
	0x000000D8	stop BX r1

Ok: r1 = 0x000000D9

Branch range

- In the B instruction, the opcode is 8 bit and the immediate value is 24 bit.
- Since addresses are halfword-aligned, the immediate value specifies bit 24-1 of the relative address.
- The 25th is for the sign; so the relative address can be $\pm 2^{24}$ byte = ± 16 MB.
- BX can jump to any 32-bit value = 4 GB.

MOV for unconditional branch

- B and BX change the value of PC.
- Similarly, a jump can be implemented by changing the value of PC with MOV and LDR:
 - LDR <Rd>, =<label>
MOV PC, <Rd>
 - LDR PC, =<label>
- MOV and LDR force the last bit of PC to 0.
- MOV instead of BX is discouraged: the assembler generates a warning.

Conditional branch: B?? and BX??

??	Flags	Meaning	??	Flags	Meaning
EQ	$Z = 1$	equal	NE	$Z = 0$	not equal
CS HS	$C = 1$	unsigned \geq	CC LO	$C = 0$	unsigned $<$
MI	$N = 1$	negative	PL	$N = 0$	positive or 0
VS	$V = 1$	overflow	VC	$V = 0$	no overflow
HI	$C = 1 \ \& \ Z = 0$	unsigned $>$	LS	$C = 0 \ \& \ Z = 1$	unsigned \leq
GE	$N \geq V$	signed \geq	LT	$N \neq V$	signed $<$
GT	$Z = 0 \ \text{or} \ N = V$	signed $>$	LE	$Z = 1 \ \text{or} \ N \neq V$	signed \leq

Example: do you pass the exam?

```
; r0 contains the score of the exam
    CMP r0, #18
    BEQ refuse
    BLO reject
    BHI accept
    ...
refuse    ...    ; study more
reject    ...    ; study much more
accept    ...    ; go on holiday
```

Compare and branch

- Compare and branch if Zero:

`CBZ <Rn>, <label>`

jumps to `label` if `Rn = 0`

- Compare and branch if Nonzero:

`CBNZ <Rn>, <label>`

jumps to `label` if `Rn ≠ 0`

- `Rn` must be among `r0-r7`.
- Only forward branch is possible (4-130 byte).

CBZ-CBNZ Vs conditional branch

- These instructions are almost equivalent:

CBZ r0, myLabel	CMP r0, #0
	BEQ myLabel

CBNZ r0, myLabel	CMP r0, #0
	BNE myLabel

- Differences:
 - CMP sets the flags, while CBZ and CBNZ do not.
 - CBZ and CBNZ jump only forward, range is shorter
 - CBZ and CBNZ cannot be used within an IT block.

While loop

The pseudocode of the while loop is

```
while (r0 != N) {  
    ...          //do something  
}
```

While loop: implementation

```
1.          B test
            loop    ...      ; do something
            test    CMP r0, #N
                   BNE loop
2. test    CMP r0, #N
           BE  exit
           ...      ; do something
           B   test
exit
```

While loop with CBZ

If $N = 0$, an alternative implementation using CBZ is:

```
loop      CBZ  r0,  exit
          ...      ; do something
          B  loop
exit      ...
```

For loop

The pseudocode of the for loop is

```
for (i = 0; i < N; i++) {  
    ...           //do something  
}
```

For loop: naive implementation

```
                MOV  r0,  #0
loop            CMP  r0,  #N
                BHS  exit
                ...      ; do something
                ADD  r0,  r0,  #1
                B    loop
exit
```


For loop: optimization

```
        MOV  r0,  N
loop    ...      ; do something
        SUBS r0,  r0,  #1
        BNE  loop

exit
```

- `CBNZ r0, loop` can not be used instead of `BNE loop` **because the branch is backward.**

Do-While loop

The pseudocode

```
do {  
    ...           //do something  
} while (r0 != N);
```

can be implemented as:

```
loop    ...      ; do something  
test    CMP  r0, #N  
        BNE  loop
```

Example: absolute value of $N - M$

```
        MOV  r0,  #N
        MOV  r1,  #M
        CMP  r0,  r1
        BLT  neg
        SUB  r0,  r0,  r1
        B    exit
neg      SUB  r1,  r1,  r0
exit     ...      ; program continues
```

Branch penalty if $N < M$

[illegible]

Conditional execution

- The IT (If-Then) block avoids branch penalty because there is no change to program flow:

```
ITxyz <cond>
```

```
instr1<cond> <operands>
```

```
instr2<cond OR not cond> <operands>
```

```
instr3<cond OR not cond> <operands>
```

```
instr4<cond OR not cond> <operands>
```

Absolute value of $N - M$ with IT

```
MOV r0, #N
MOV r1, #M
CMP r0, r1
ITE GE
SUBGE r0, r0, r1
SUBLT r1, r1, r0
exit    ...    ; program continues
```

No branch penalty

	1	2	3	4	5	6	7	8	9
MOV	F	D	E						
MOV		F	D	E					
CMP			F	D	E				
ITE				F	D	E			
SUBGE					F	D	E		
SUBLT						F	D	E	
...							F	D	E
...								F	D

IT syntax

- first statement after `IT` must be the true case
- up to 4 instructions (true or false) are allowed
- the number of instructions in true and false cases must match the number of `T` and `E`
- false condition is inverse of true condition
- branches to `IT` instructions are not allowed
- an `IT` instruction can be a branch only if it is the last one.