# Relational Model

**COURSE 2: Databases** 

## Relational model

#### Relational model

- Database = collection of RELATIONS
  - relation in relational model ≠ relationship in ERD.
- Relation Schema: A relation schema represents the name of the relation with its attributes.

Attribute domain – Each attribute has some pre-defined values.

#### Relational model

• Codd rules 1985 → Is DBMS relational? If yes, to what degree?

https://computing.derby.ac.uk/c/codds-twelve-rules/

Relational Integrity constraints

**RELATIONS** 

**OPERATORS** 

Relational Integrity constraints

**RELATIONS** 

**OPERATORS** 

- Domain constraints
  - the value of each attribute must be unique, specified data types integers, real numbers, characters, Booleans, variable length strings etc.
- Key constraint
  - Unique + not null PK
- Referential integrity constraints
  - the value of a FK is null or it coresponds to the value of a PK.

Relational Integrity constraints

**RELATIONS** 

**OPERATORS** 

• Relational shema  $R(A_1, A_2, ..., A_n)$ 

•  $R \subset D_1 \times D_2 \times \cdots \times D_n$ ,  $D_i$  domain

Relational Integrity constraints

**RELATIONS** 

**OPERATORS** 

• UNION, INTERSECT, PRODUCT, DIFFERENCE

- PROJECT
- SELECT
- JOIN
- DIVISION

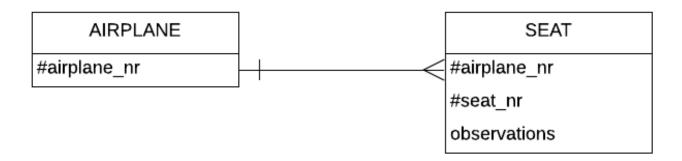
# Converting ER into RM

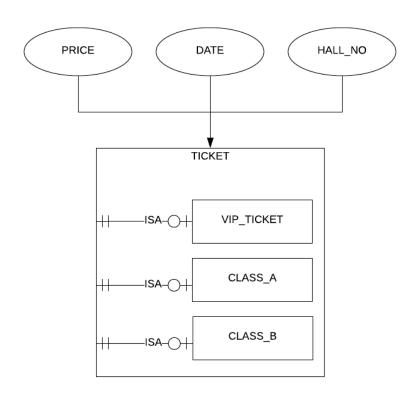
#### Rules for entities

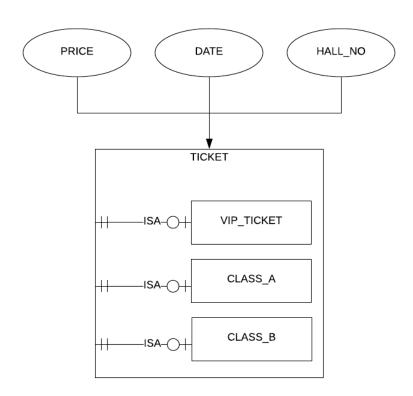
- Strong entities  $\rightarrow$  independent tables
  - PK doesn't contain foreign keys.
- Weak entities → table
  - PK contains the key of the related strong entity and or more key attributes.
- Sub-entities 

  one ore more tables, Boolean attribute,
  type\_attribute
  - PK may also represent a FK.

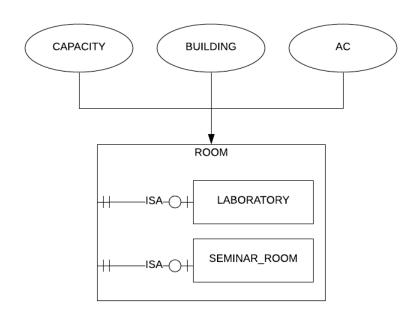
## Rules for entities strong – weak entity

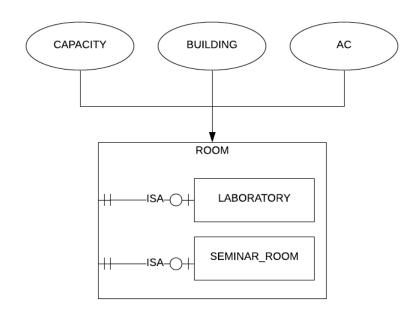




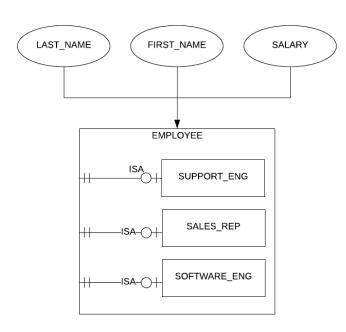


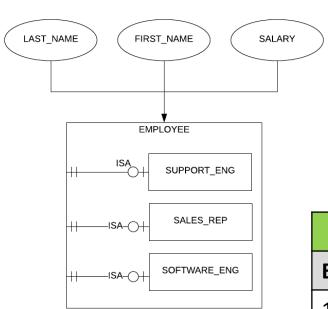
TICKET_ID	PRICE	HALL_NO	DATE	TYPE
1	200	Coliseum	08/03/20	VIP
2	150	Lyttelton	14/04/20	А
3	140	Olivier	01/05/20	А
4	90	Coliseum	04/06/20	В
5	220	Lyttelton	08/03/20	VIP
6	95	Olivier	14/04/20	В
7	210	Coliseum	20/03/20	VIP





ROOM_ID	CAPACITY	BUILDING	LAB	SEM
1	40	FMI	1	1
2	45	Magurele	1	0
3	30	Geografie	0	0
4	90	FMI	1	0
5	80	FMI	1	0
6	95	Drept	0	1
7	20	FMI	1	1



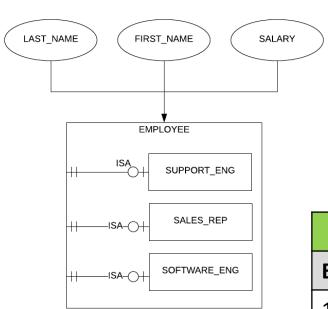


EMPLOYEES				
EMP_ID	LAST_NAME	FIRST_NAME	SALARY	
1	Smith	John	2500	
2	Grant	Anne	2700	
3	Brown	Gregory	2300	
•••				

SUPPORT_ENG			
EMP_ID	LEVEL		
1	3		
•••	•••		

SALES_REP			
EMP_ID	TARGET		
2	25		
	•••		

SOFTWARE_ENG			
EMP_ID TEEM			
3			

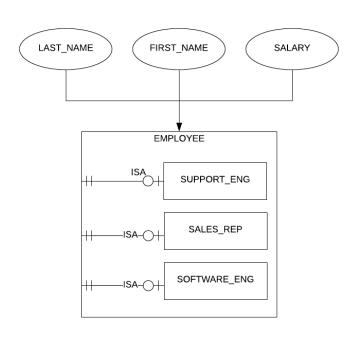


EMPLOYEES				
EMP_ID	LAST_NAME	FIRST_NAME	SALARY	
1	Smith	John	2500	
2	Grant	Anne	2700	
3	Brown	Gregory	2300	
•••				

SUPPORT_ENG			
EMP_ID	LEVEL		
1	3		
•••	•••		

SALES_REP			
EMP_ID	TARGET		
2	25		
	•••		

SOFTWARE_ENG			
EMP_ID TEEM			
3			



SUPPORT_ENG				
EMP_ID    LEVEL    LAST_NAME    FIRST_NAME    SALARY				
1	3	Smith	John	2500
	•••			

SALES_REP					
EMP_ID    TARGET    LAST_NAME    FIRST_NAME    SALARY					
2	25	Grant	Anee	2700	

SOFTWARE_ENG					
EMP_ID	TEEM	LAST_NAME	FIRST_NAME	SALARY	
3	3	Brown	Gregory	2300	
	•••				

## Rules for relationships

- 1 to 1 & 1 to M  $\rightarrow$  foreign keys.
  - 1 (PK) to M (FK)
  - Usually in 1 to 1 relationships the FK is placed in the tables with fewer rows.
- M to M  $\rightarrow$  associative table.
  - PK contains FKs and additional column.
- Ternary relationships  $\rightarrow$  associative table.
  - PK contains FKs and additional column.

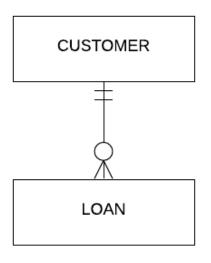
### One to One



ACCOUNT					
ACCOUNT_ID LAST_NAME FIRST_NAME DATE					
10	Snow	John	08/03/20		
22	Grant	Anee	14/04/20		
300	Brown	Gregory	01/05/20		
•••	•••		•••		

CARD					
CARD_ID	ACCOUNT_ID	CVN	DATE		
16897	10	125	18/04/21		
24789	22	987	14/04/22		
34597	300	875	03/05/21		

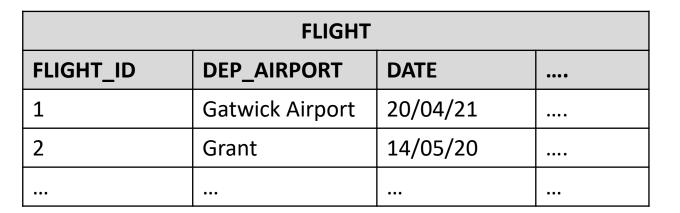
## One to Many

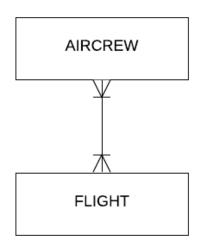


CUSTOMER					
CUSTOMER_ID	LAST_NAME	FIRST_NAME	••••		
10	Snow	John	••••		
22	Grant	Anee			
300	Brown	Gregory			

LOAN					
LOAN_ID CUSTOMER_ID VALUES DATE					
16897	10	125000	18/04/21		
24789	22	987000	14/04/22		
34597	300	87500	03/05/21		

## Many to Many

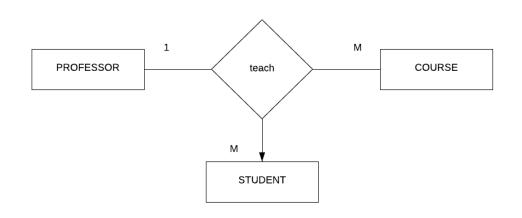




FLIGHT_CREW					
CREW_ID	FLIGHT_ID	OBSERVATIONS			
10	1				
22	1				
10	2				

AIRCREW					
CREW_ID	JOB_ID				
10	Snow	John	captain		
22	Grant Anee		first_officer		

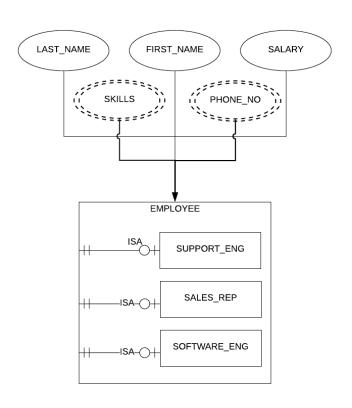
## Ternary Relationships



TEACH					
PROFESSOR_ID	COURSE_ID	STUDENT_ID GRAD			
1	BD	1001	9		
1	SGBD	1002	10		
1	BD	1002	8		
2	TAP	1001	8		
2	TAP	1002	10		
2	AG	1001	5		
			••••		

#### Rules for attributes

- Simple attribute → column
- Multivalued attributes  $\rightarrow$  weak entity  $\rightarrow$  table
  - → set of columns



EMPLOYEES						
EMP_ID	LAST_NAME	FIRST_NAME	SALARY	PHONE1	PHONE2	
1	Smith	John	2500	0745	0720	
2	Grant	Anne	2700	07497	NULL	
3	Brown	Gregory	2300	NULL	07458	

EMP_SKILL					
EMP_ID	SKILL	LEVEL			
1	Python	3			
1	C++	2			
1	NoSql	3			
2	SQL	1			

# Indexes

#### Indexes

Maps search key to data using specific data structures.

- Optimized search.
- Optimized joins (lookup in more than one table)
- Optimized order/group

- slower DML (insert and update operations).
- extra memory

**SELECT** 

Optimized search

Optimized joins

Optimized order/group

Index

slower DML

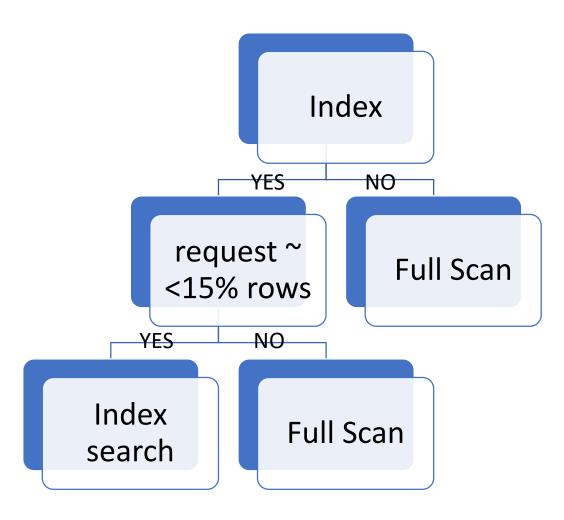
extra memory

extra load

INSERT, UDATE

Databases C2 Relational Model, indexes

## Sql Optimizer



## Autogenerated columns

- MySQL auto-generated index (key):
  - DB\_ROW\_ID increases monotonically as new rows are inserted.
  - DB\_ROLL\_PTR roll pointer, points to log record.
  - DB\_TRX\_ID last transaction that updated or inserted the row.

#### • Oracle rowid:

- Pseudo column 18 characters = 10 + 4 + 4 (block, row, file).
- Store and return row address in hexadecimal format (string).
- Unique identifier for each row.
- Immutable.

## Autogenerated columns

#### Oracle rowid:

Used in where clause to select/update/delete a row.

#### Oracle rownum:

- Sequential number in which oracle has fetched the row, before ordering the result
- Temporary generated along with a select statement.

#### Mongo

ObjectID (timestamp 4Bytes + random 5Bytes + Count 3Bytes.

#### Index

- Data structure that optimize search.
- Automatically created when a PK/unique constraint is defined.

#### **Primay key**

- Constraint imposed on insert/update behavior.
- NotNull & Unique.

```
MySQL SHOW EXTENDED INDEX FROM index_test;
```

```
Oracle select * from user_indexes where table_name = 'INDEX_TEST';
```

```
MySQL select * from information_schema.statistics where table_name = 'index_test1' and index_name = 'primary';
```

```
Oracle select * from user_constraints where table_name = 'INDEX_TEST';
```

# Index types

## Clustered index (SqlServer, MySql)

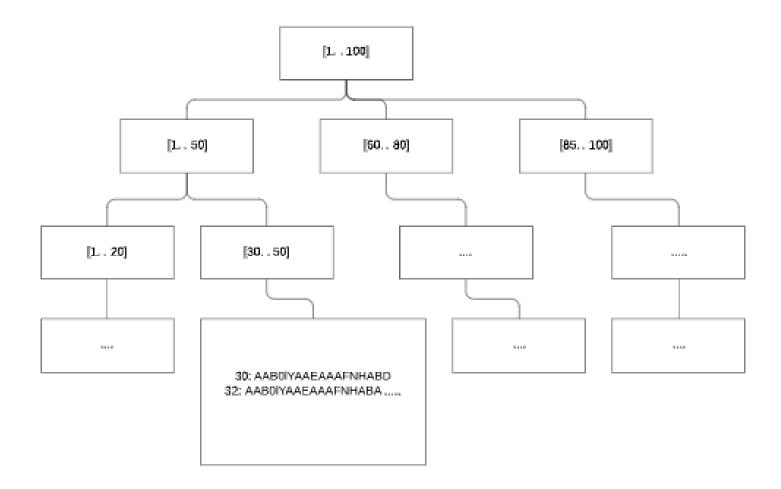
- Defines the order in which data is physically stored in a table. (index on column semester)
- Only one clustered index on a table (data can be stored in only one order)
- A cluster index is created automatically when a primary key is defined.
- No second data structure for the table
- Oracle: IOT index organized tables. Table is stored in a B-tree structure. (key and non-keys column are stored in leafs)

#### B – Tree

- B -- Balanced tree.
- Default index type in Oracle.
- Two types of nodes: branch blocks and leaf blocks.
- Branch blocks pointers to lower levels.
- Leaf blocks contain rowids/physical address.
- The number of blocks traversed in order to reach a leaf block is the same for each leaf block.

#### B – Tree

- create index idx\_emp\_id on employees(employee\_id).
  - Divide employee\_id values in sorted ranges.
  - Leaves nodes store rowid



#### Reverse index

- B tree where keys are in reverse order. Key 4573 is stored 3754.
- Optimized insert operations.
- Key 4573 will be stored in the same block with key 9573 while 4574 will be stored in a different block.

## Bitmap index

- Used for columns with limited number of distinct values.
- Example: language proficiency levels (en)

emp_id	en	fr
1	A1	B1
2	A2	B2
3	C1	A1
4	A1	B1
5	A1	

row_id	A1	A2	B1	B2	<b>C1</b>	<b>C2</b>
AABOIYAAEAAAFNHABD	1	0	0	0	0	0
AABOIYAAEAAAFNHABV	0	1	0	0	0	0
AABOIYAAEAAAFNHABX	0	0	0	0	1	0
AABOIYAAEAAAFNHAAv	1	0	0	0	0	0
AABOIYAAEAAAFNHAAV	1	0	0	0	0	0