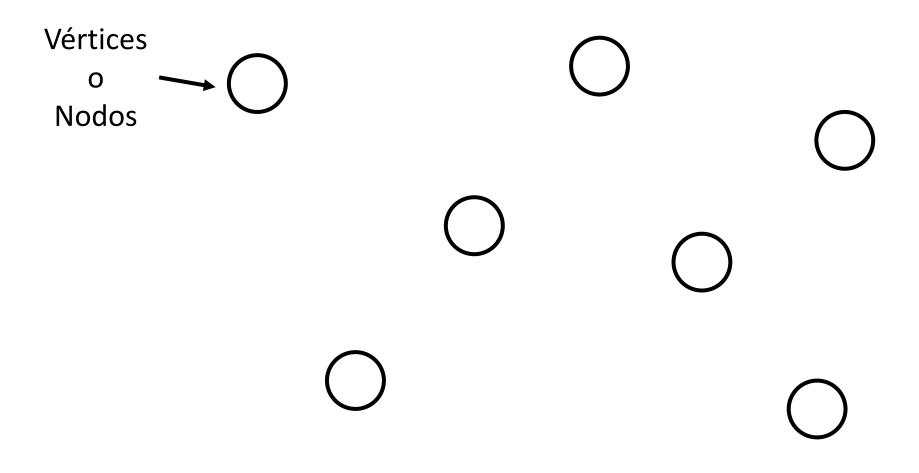
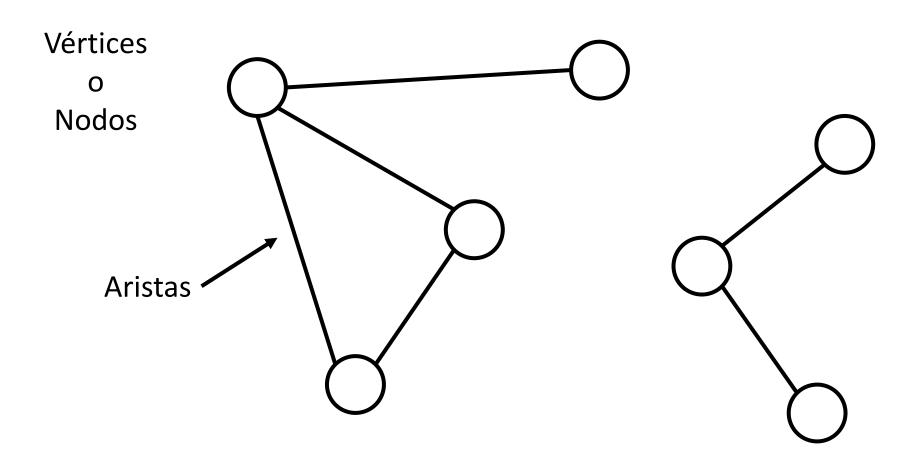
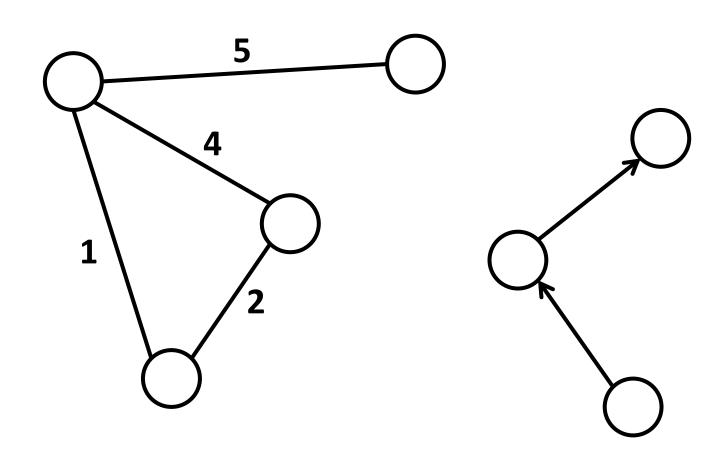
# Introducción a grafos, DFS, BFS y sus aplicaciones

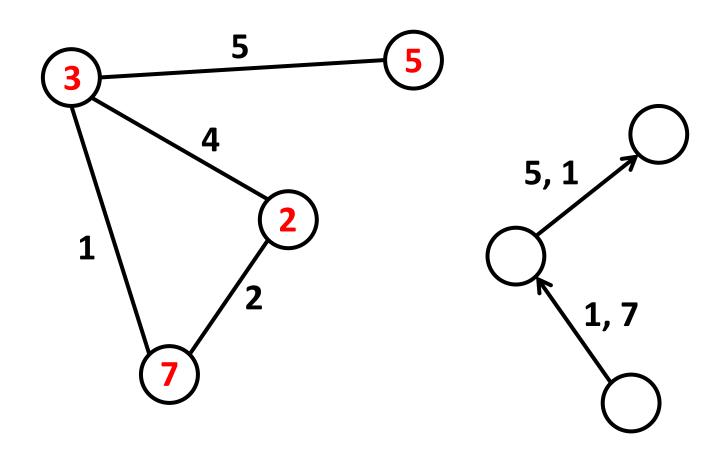
Miguel Ortiz

Clases para la Olimpiada Boliviana de Informática

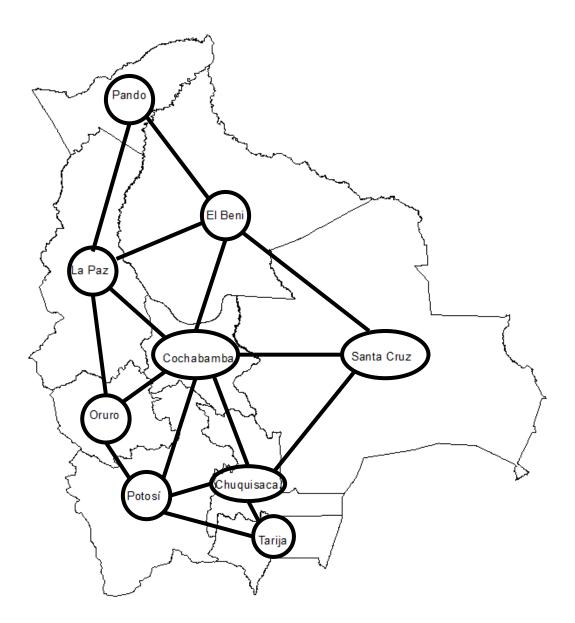




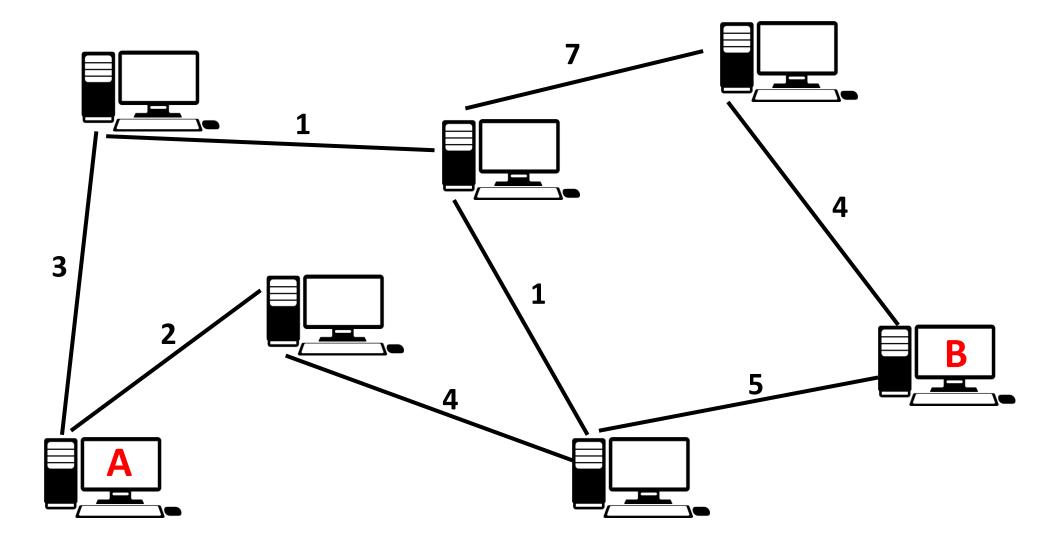




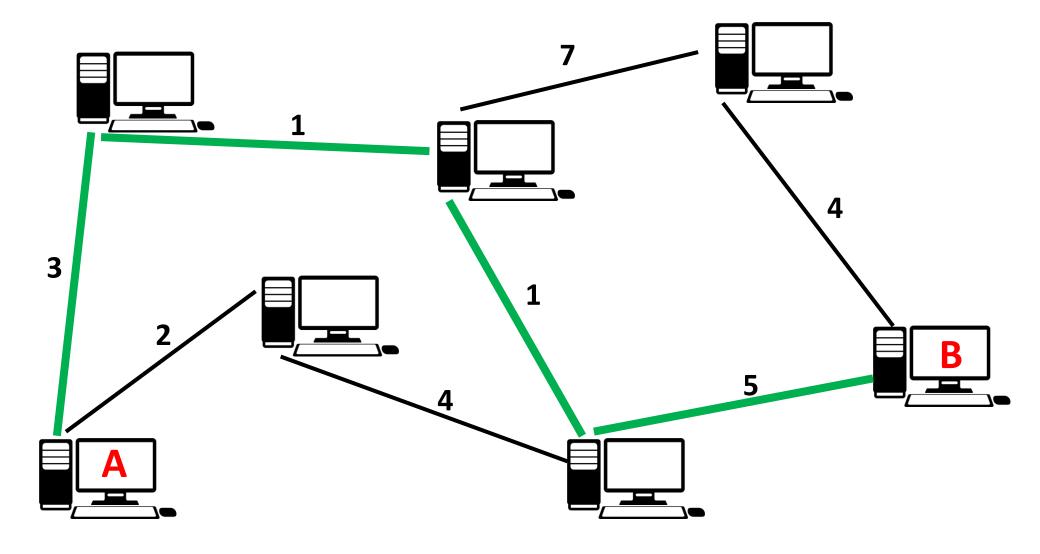
# Ejemplos



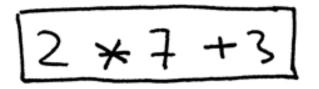
# Ejemplos



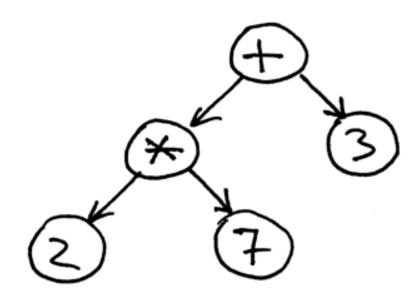
# Ejemplos



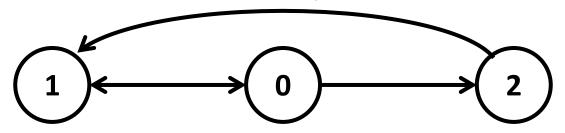
Ejemplos



# A ST



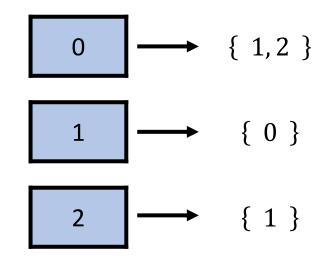
### ¿Cómo se representan?



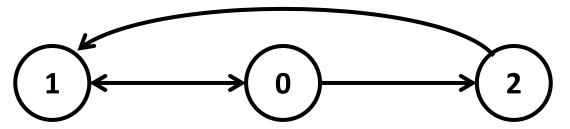
#### Matriz de adyacencia

#### 

#### Listas de adyacencia



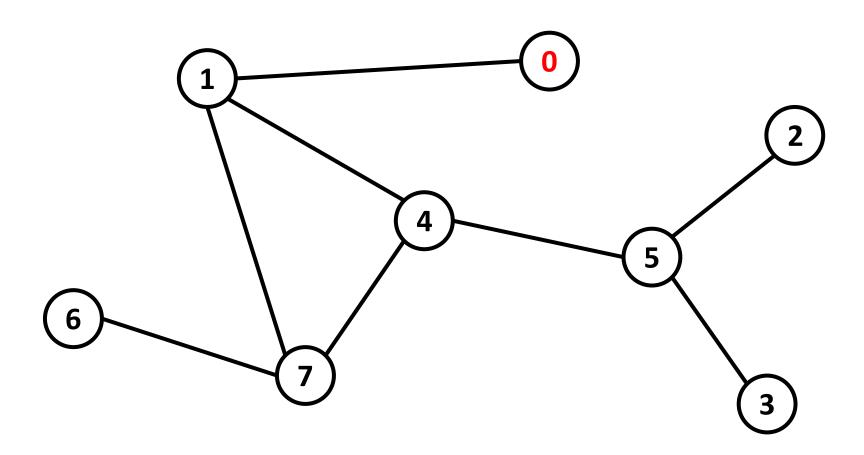
#### ¿Cómo se representan?

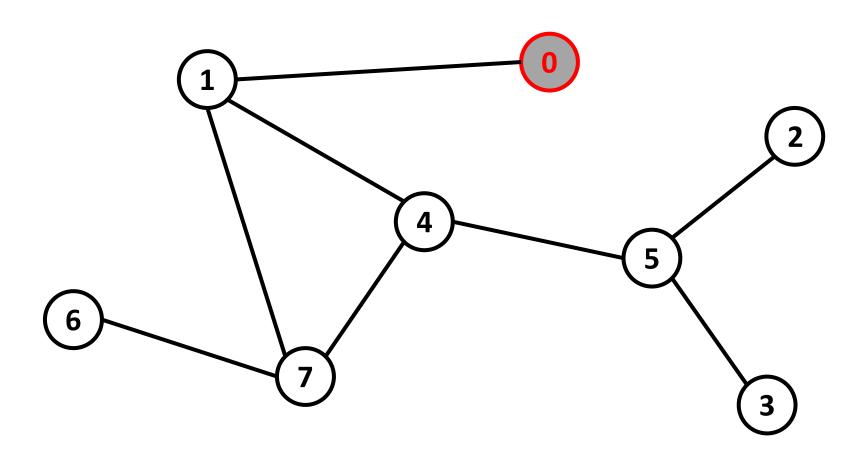


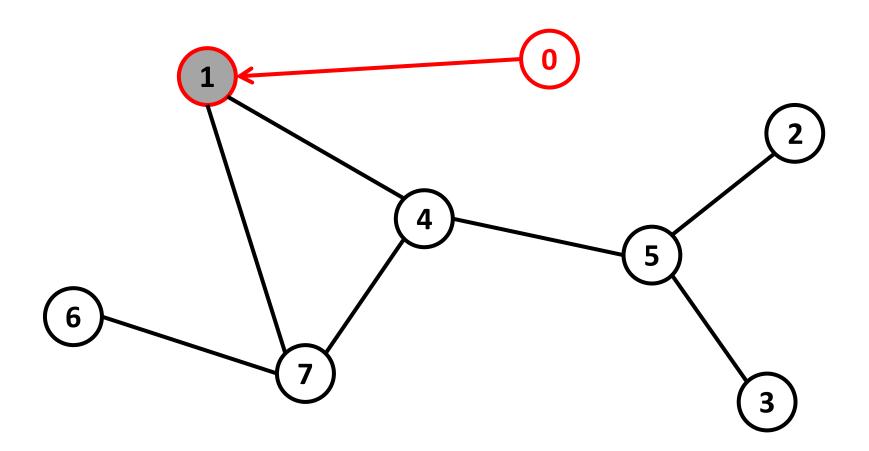
```
// Matriz de adyacencia
int g[n][n];

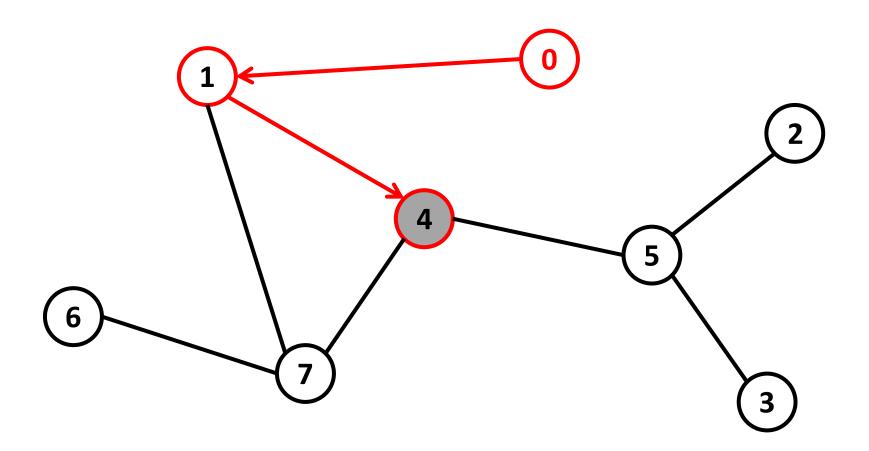
// Llenado
g[0][1] = 1;
g[0][2] = 1;
g[1][0] = 1;
g[2][1] = 1;
g[2][1] = 1;
//Lista de adyacencia
vector<int> g[n];
// vector<vector<int>> g;
// Llenado
g[0].push_back(1);
g[0].push_back(2);
g[2][1] = 1;
g[1].push_back(0);
g[2].push_back(1);
```

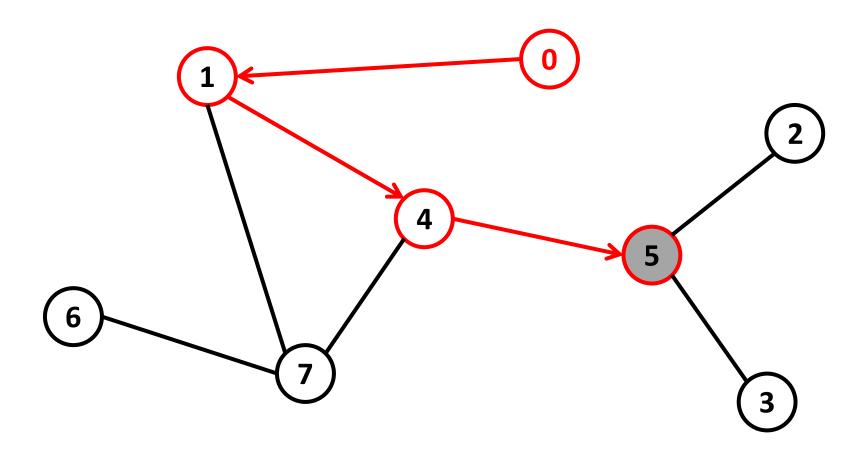
- Algoritmo de recorrido de grafos
- Empieza en un nodo, lo marca como visitado y se mueve a sus vecinos que aun no han sido visitados
  - Si hay varias opciones para moverse, se mueve a cualquiera
- Cuando no hay vecinos sin marca, vuelve por donde vino
- Hace el mismo procedimiento en cada paso hasta haber visitado todos los nodos posibles
- Sigue un solo camino por el grafo, como caminar por un laberinto con una mano siempre en la pared

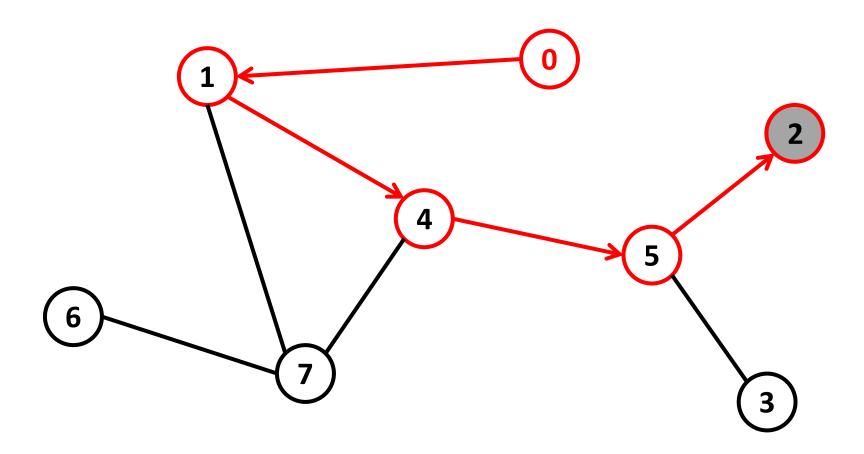


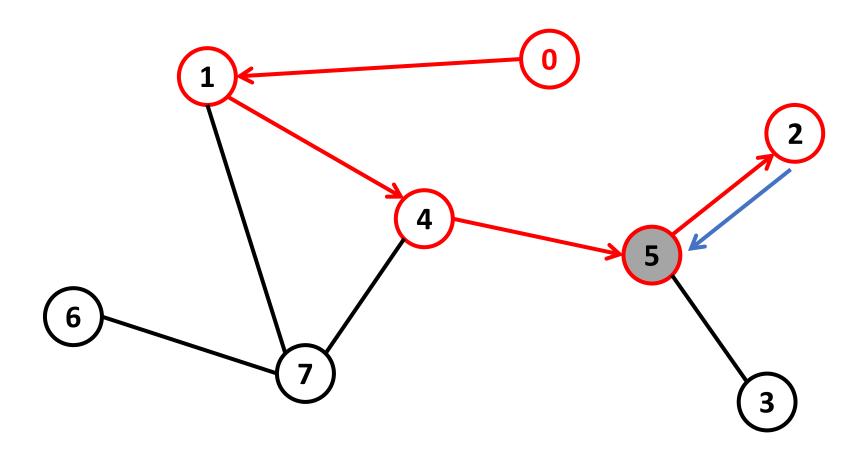


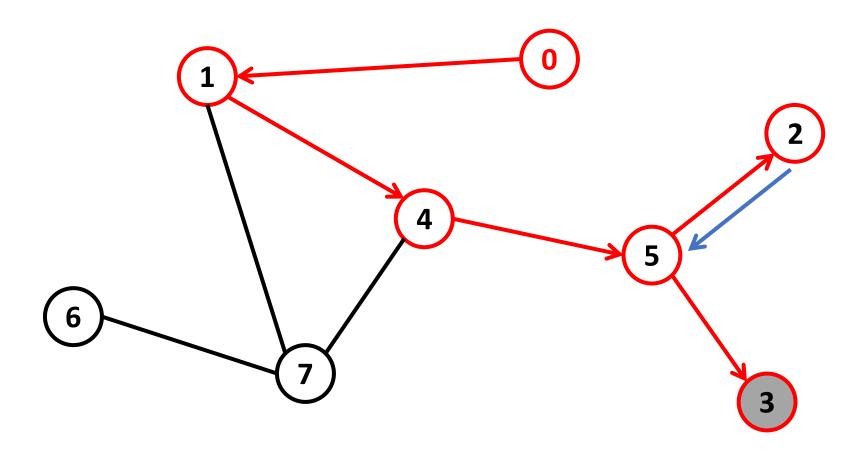


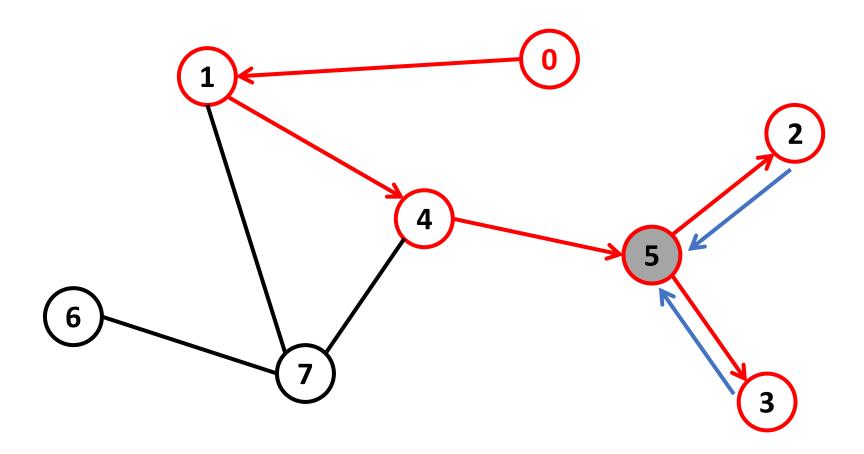


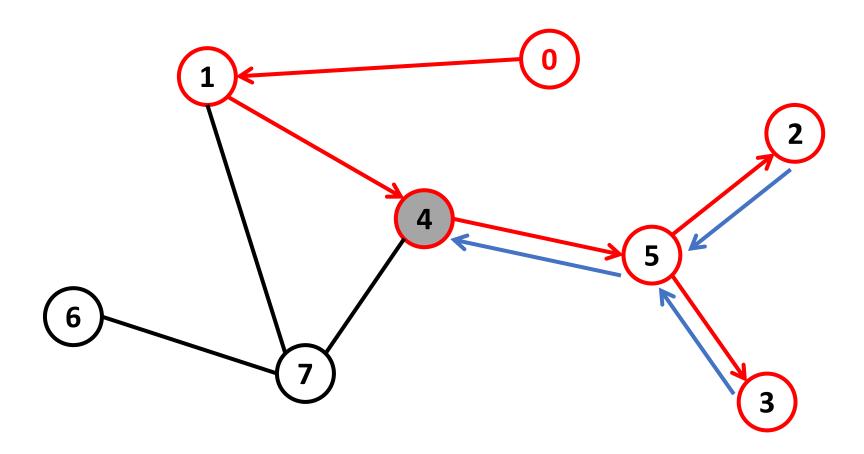


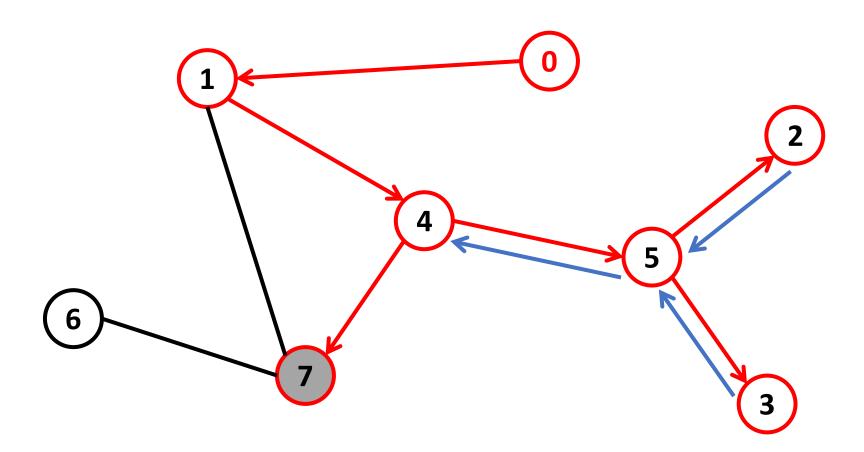


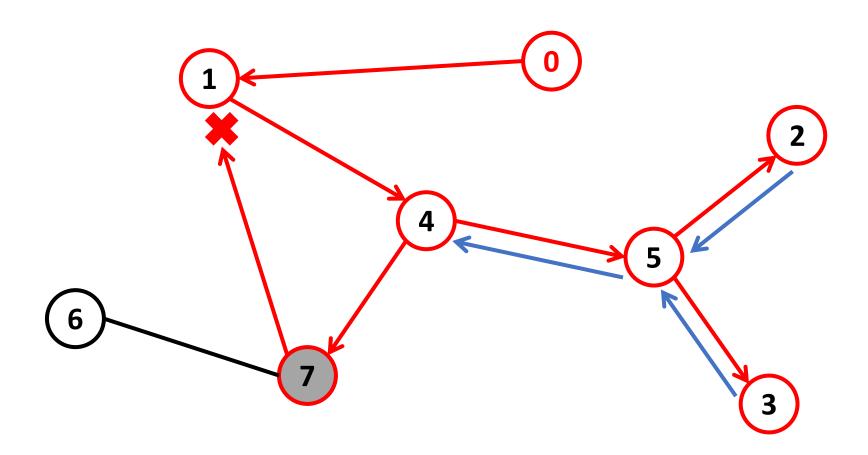


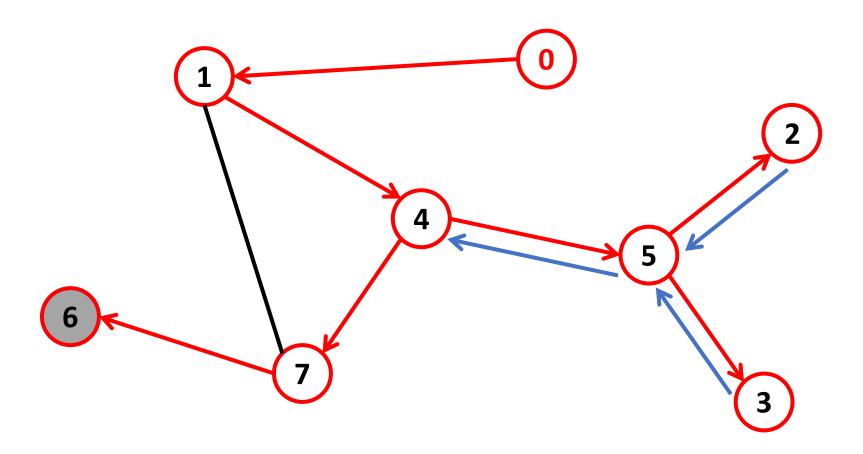


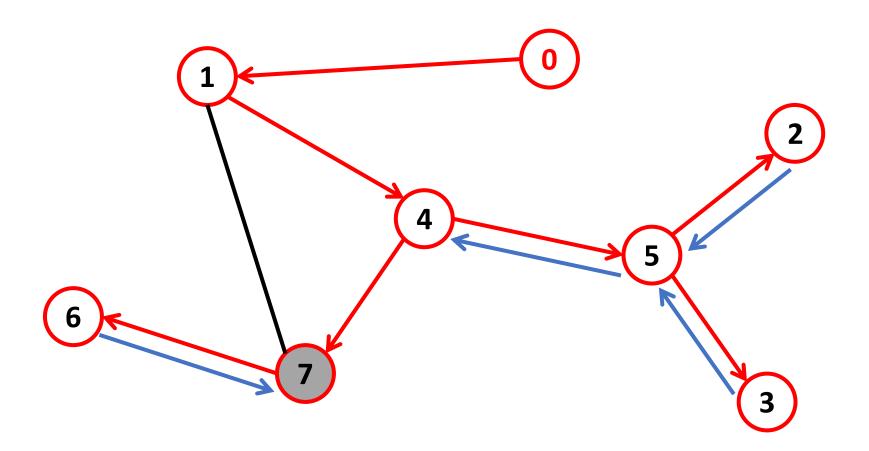


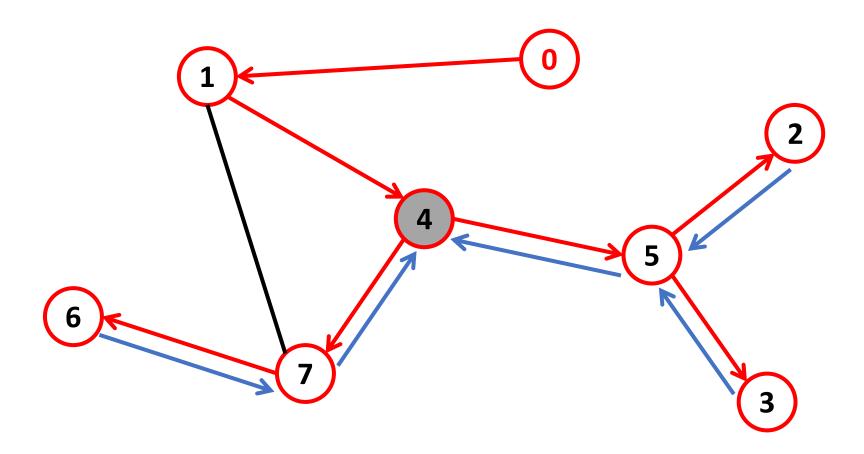


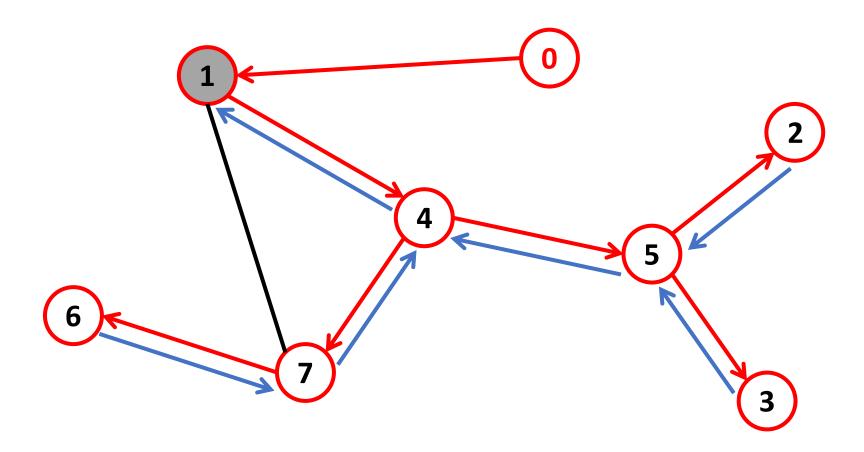


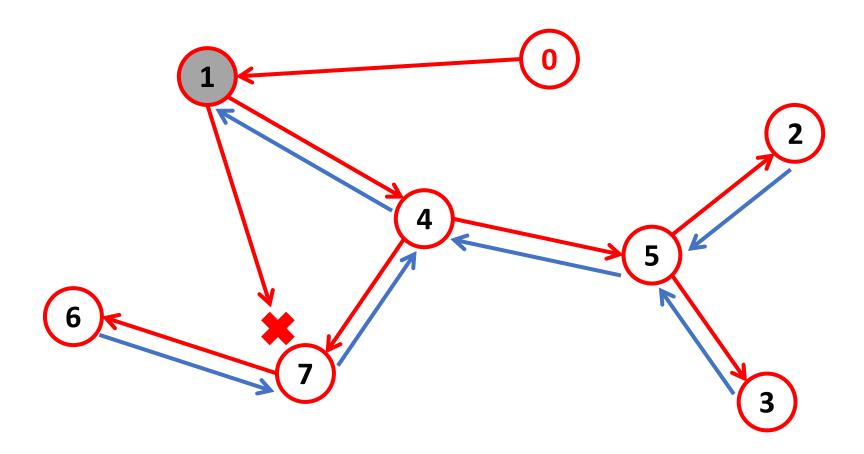


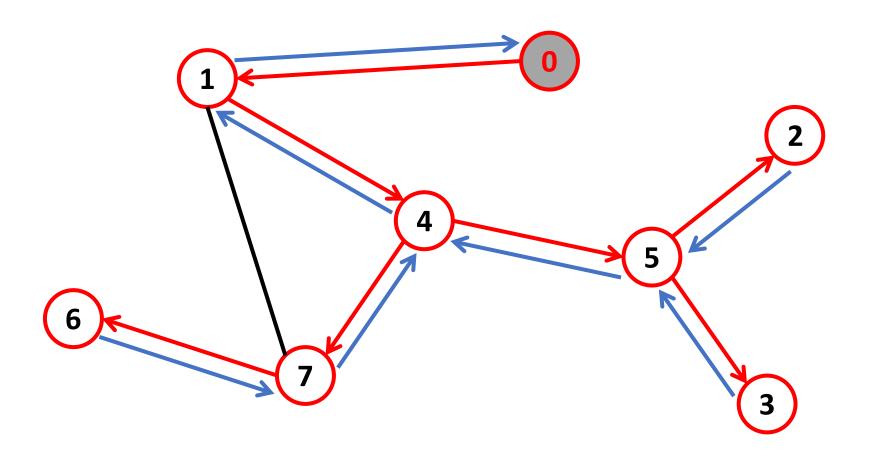












#### Necesitamos:

```
vector<int> g[n];
bool vis[n];
```

#### Necesitamos:

```
vector<int> g[n];
bool vis[n];
```

```
void dfs(int u) {
  vis[u] = true;
  for (int v : g[u]) {
    if (!vis[v]) {
       dfs(v);
    }
  }
}
```

Iterativo:

```
stack<int> st;
vector<int> vis(n, false);
st.push(u);
while (st.size() > 0) {
  int u = st.top();
  st.pop();
 vis[u] = true;
  for (int v : adj[u]) {
    if (!vis[v]) {
      st.push(v);
```

#### Recursivo:

```
void dfs(int u) {
  vis[u] = true;
  for (int v : g[u]) {
    if (!vis[v]) {
       dfs(v);
    }
  }
}
```

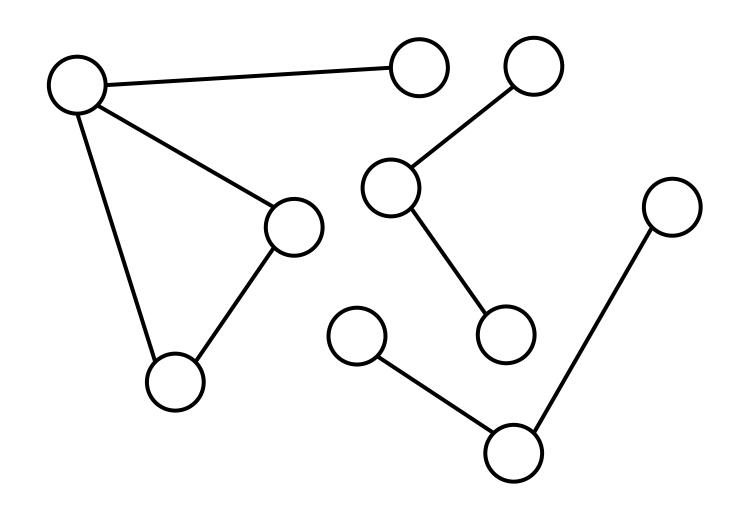
Iterativo:

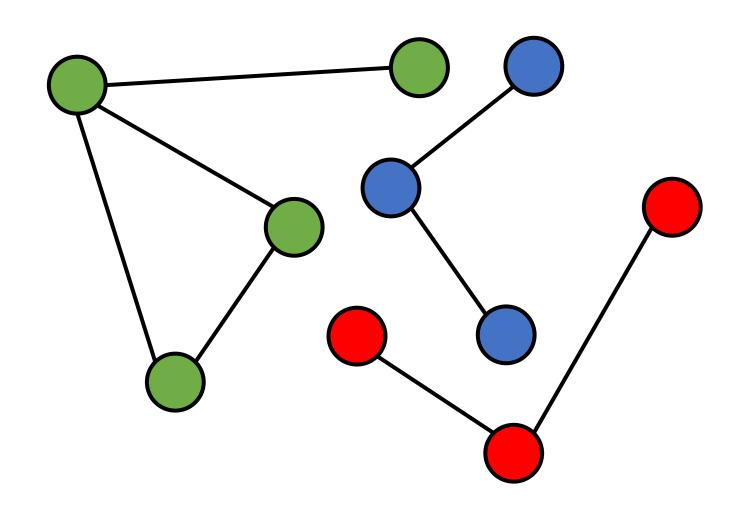
```
stack<int> st;
vector<int> vis(n, false);
st.push(u);
while (st.size() > 0) {
  int u = st.top();
  st.pop();
 vis[u] = true;
  for (int v : adj[u]) {
    if (!vis[v]) {
      st.push(v);
```

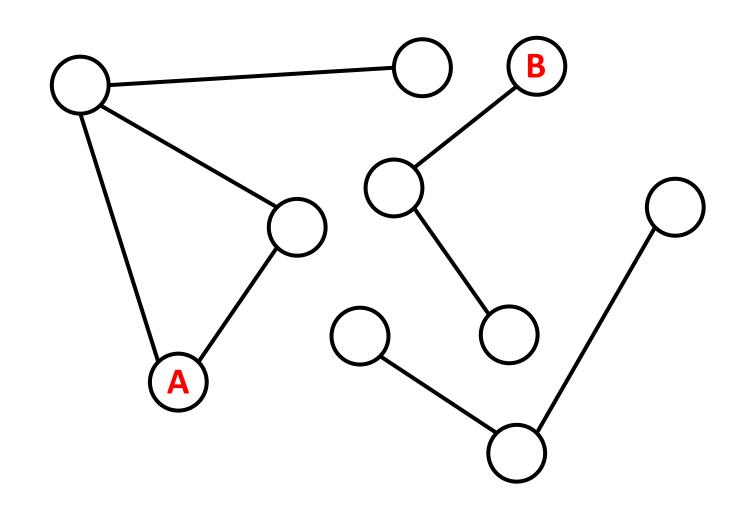
#### Recursivo:

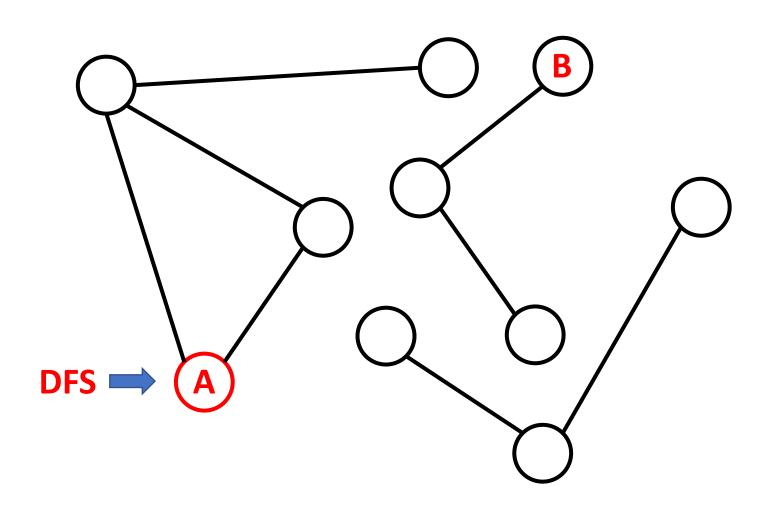
```
void dfs(int u) {
  vis[u] = true;
  for (int v : g[u]) {
    if (!vis[v]) {
       dfs(v);
     }
  }
}
```

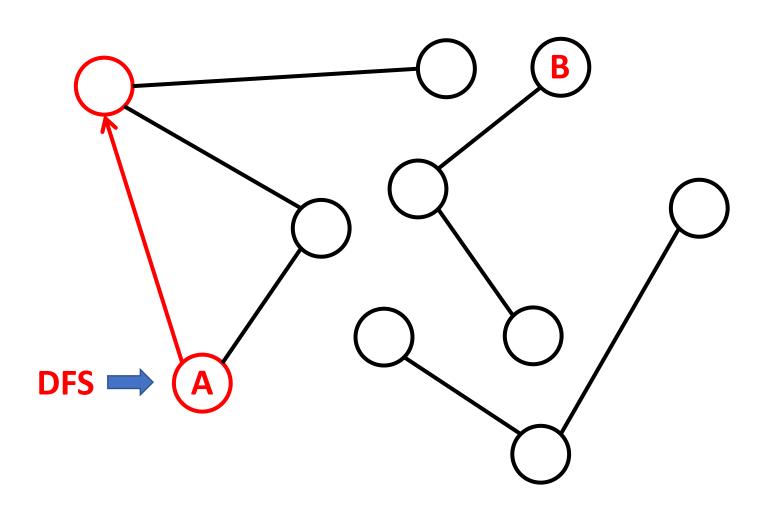
# Aplicación – Verificar conectividad Componentes Conexos

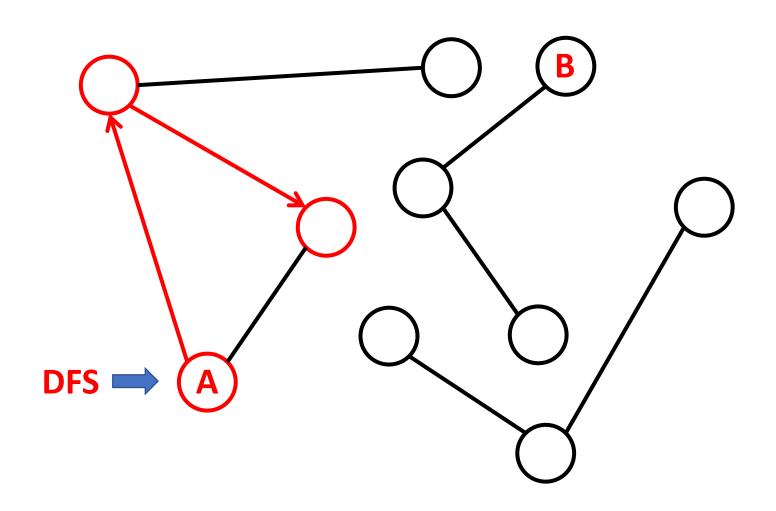


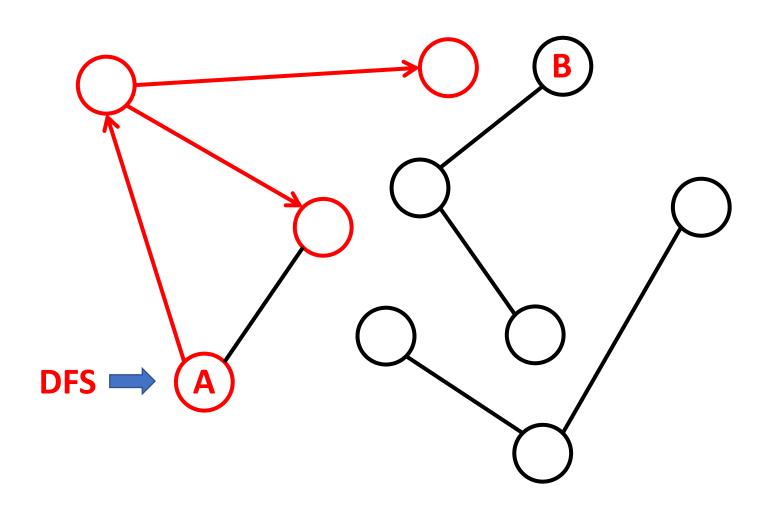


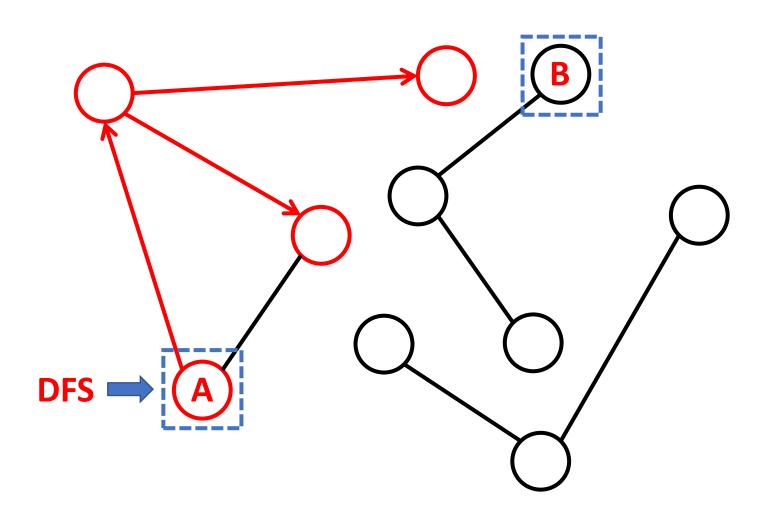


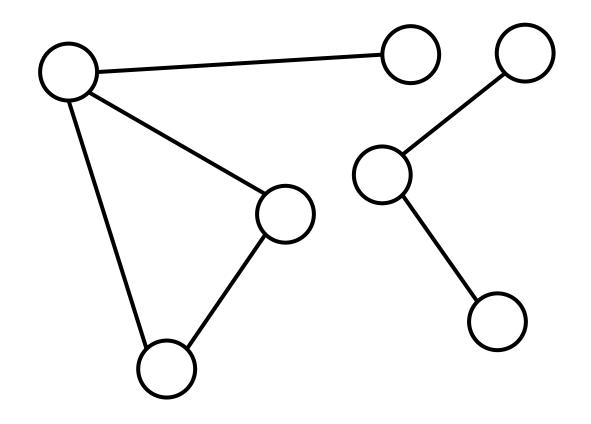


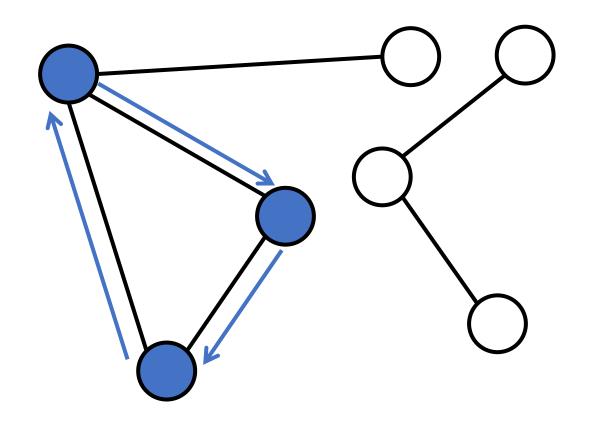


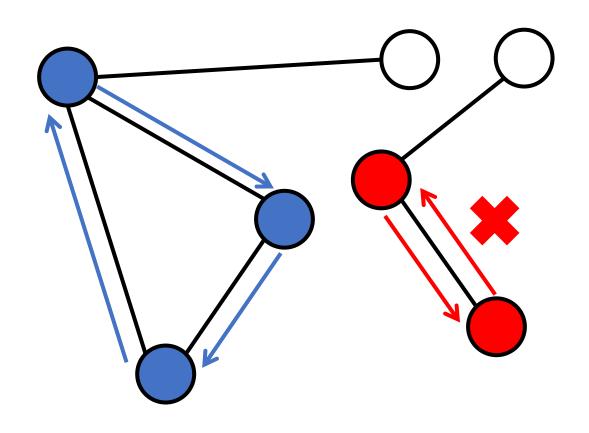


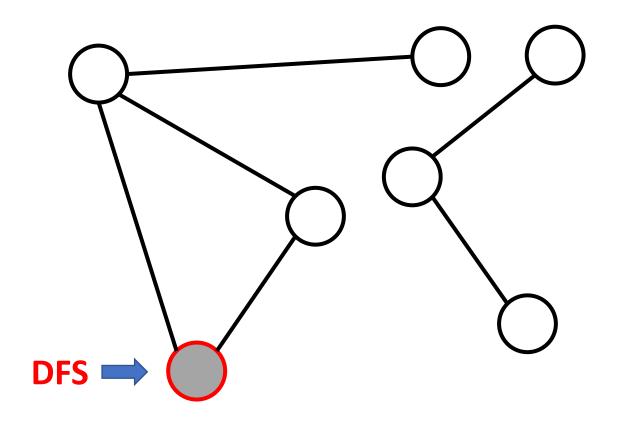


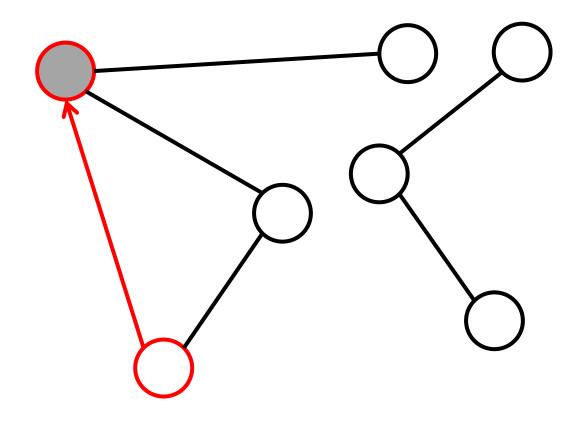


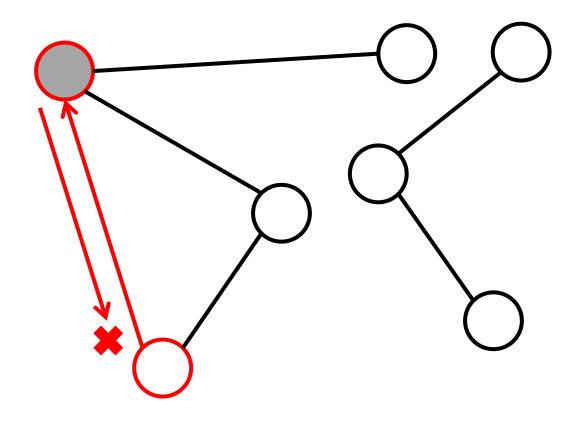


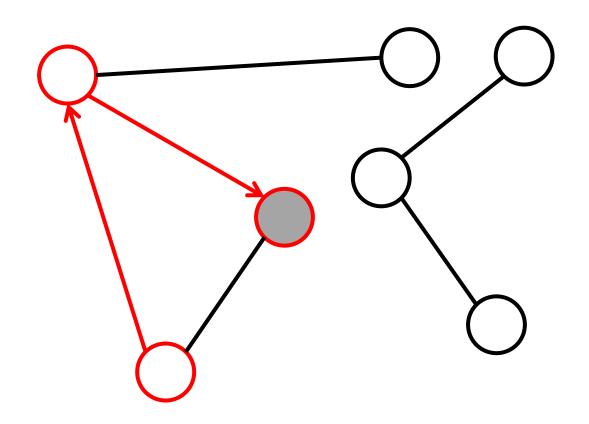


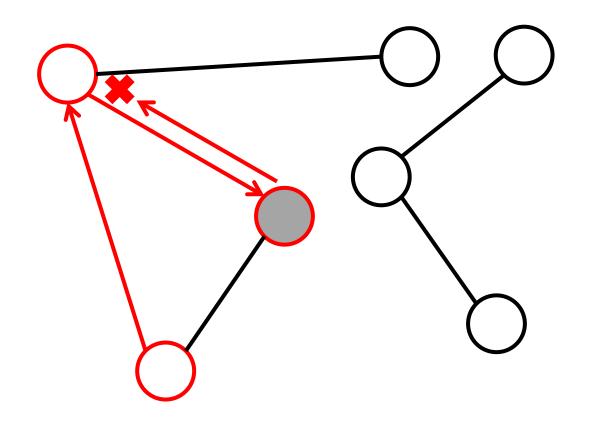


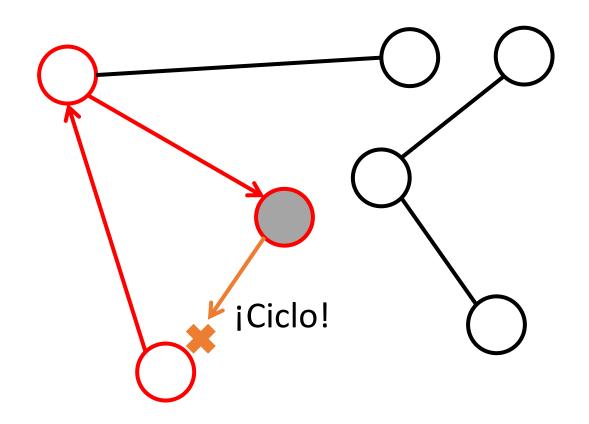


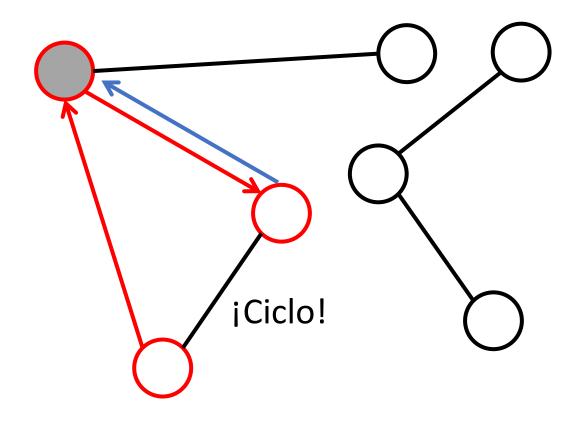


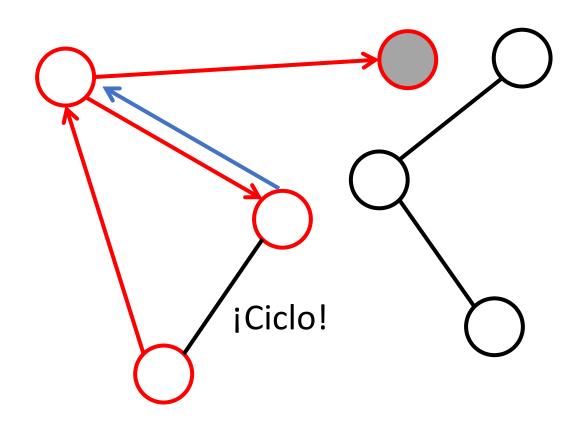


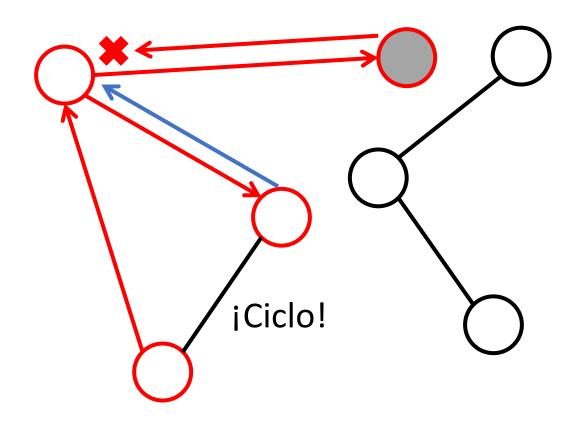


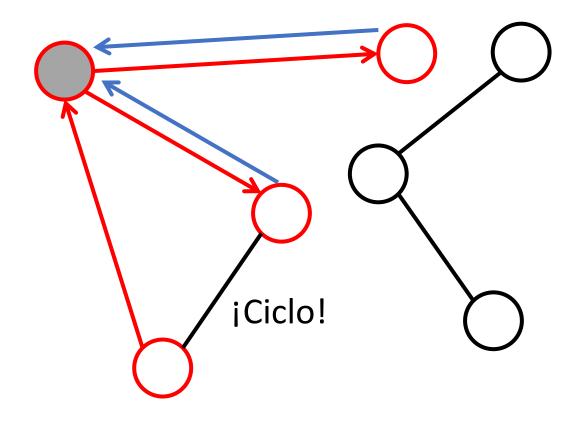


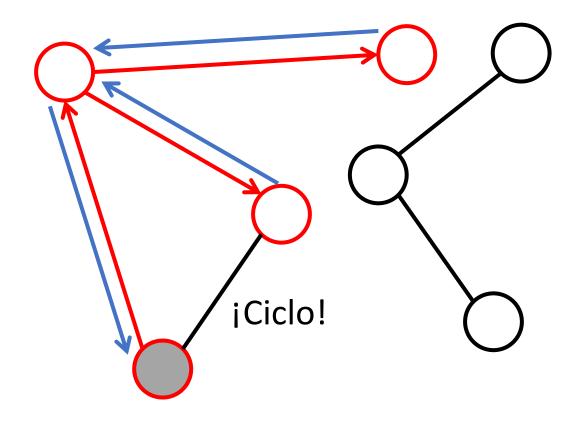


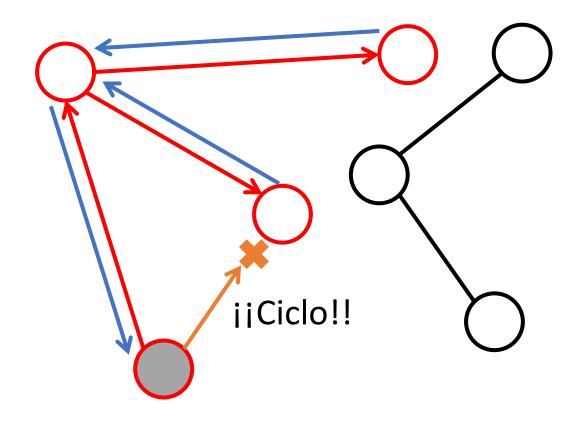


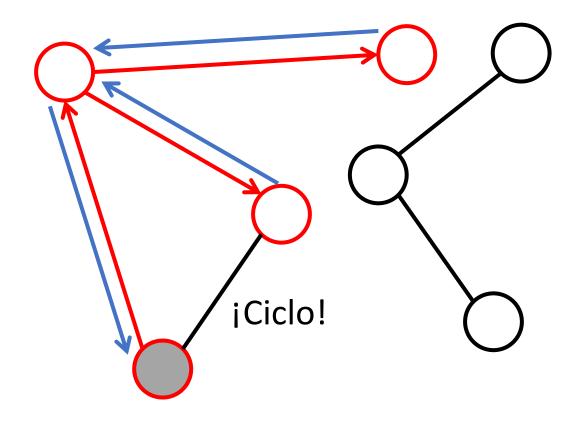


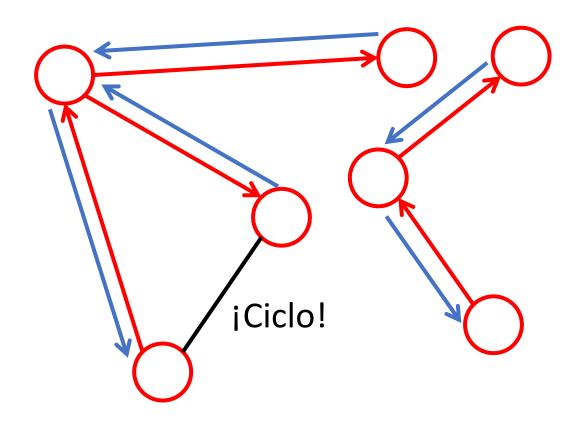


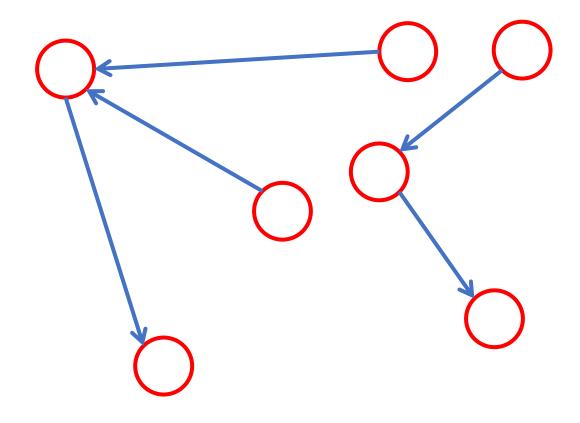






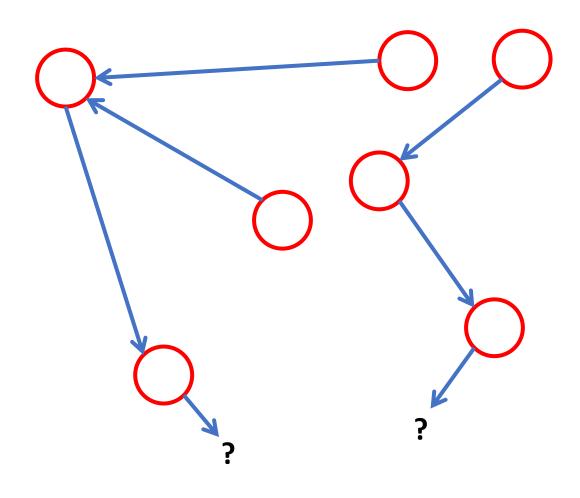


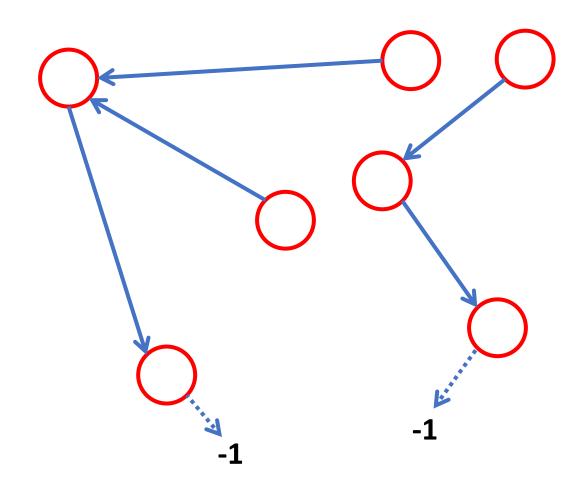


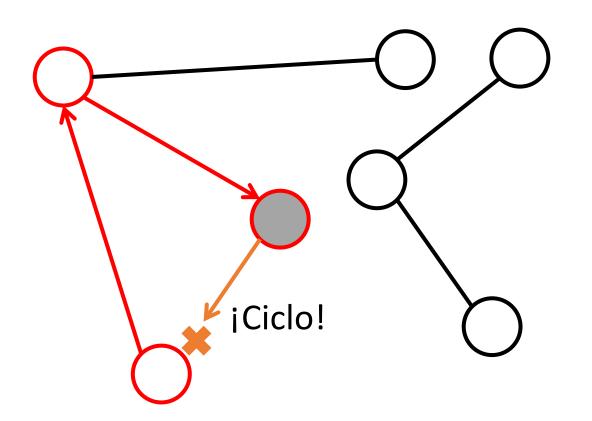


```
vector<int> g[n];
bool vis[n];
void dfs(int u) {
 vis[u] = true;
  for (int v : g[u]) {
    if (!vis[v]) {
     // u -> v
     // aqui se define el padre
      dfs(v);
```

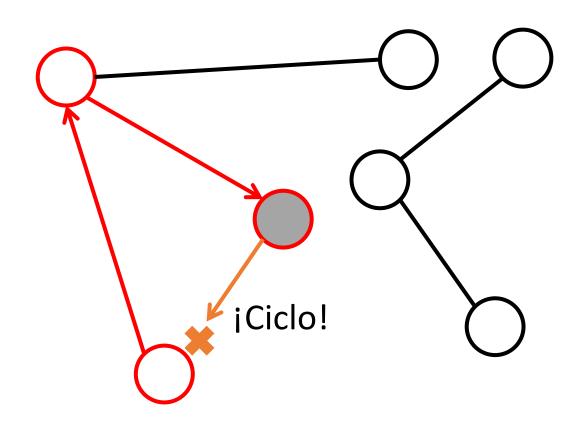
```
vector<int> g[n];
bool vis[n];
int padre[n];
void dfs(int u) {
 vis[u] = true;
  for (int v : g[u]) {
    if (!vis[v]) {
     // u -> v
      padre[v] = u;
      dfs(v);
```







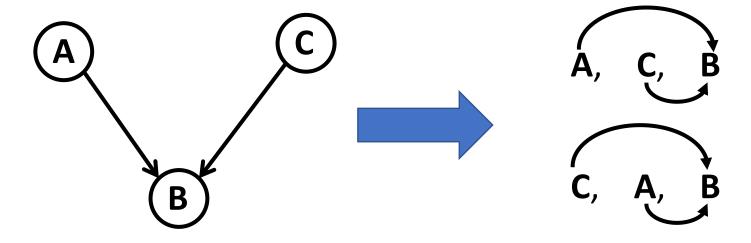
```
vector<int> g[n];
bool vis[n];
int padre[n];
void dfs(int u) {
 vis[u] = true;
  for (int v : g[u]) {
    if (!vis[v]) {
      // u -> v
      padre[v] = u;
      dfs(v);
```



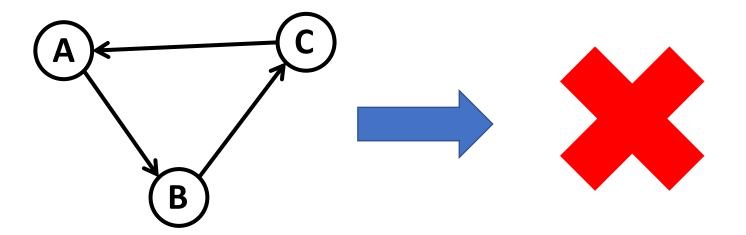
```
vector<int> g[n];
bool vis[n];
int padre[n];
bool hayCiclo = false;
void dfs(int u) {
 vis[u] = true;
  for (int v : g[u]) {
    if (!vis[v]) {
      // u -> v
      padre[v] = u;
      dfs(v);
    else if (v != padre[u]) {
      hayCiclo = true;
```

Es un concepto aplicado a grafos dirigidos que no contienen ciclos.

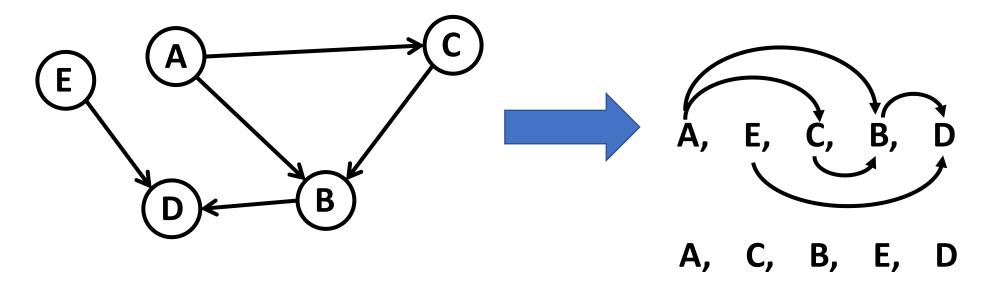
Es un concepto aplicado a grafos dirigidos que no contienen ciclos.



Es un concepto aplicado a grafos dirigidos **que no contienen ciclos**.



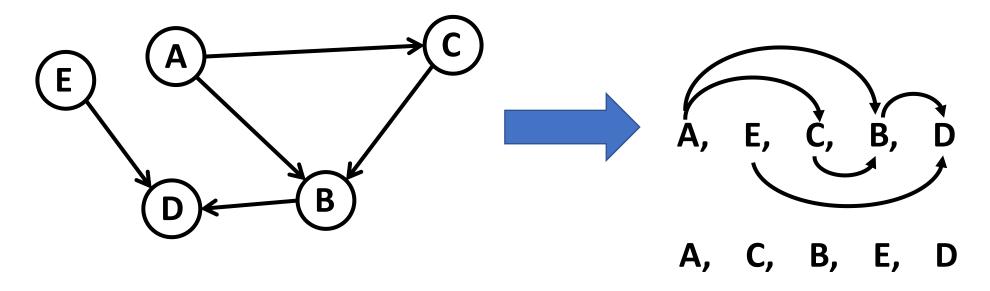
Es un concepto aplicado a grafos dirigidos que no contienen ciclos.

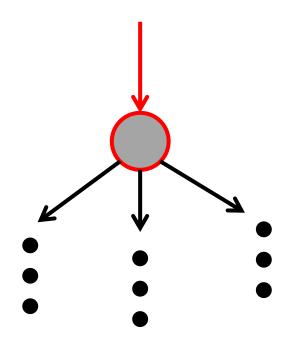


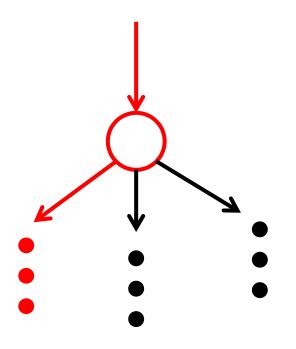
Es un concepto aplicado a grafos dirigidos que no contienen ciclos.

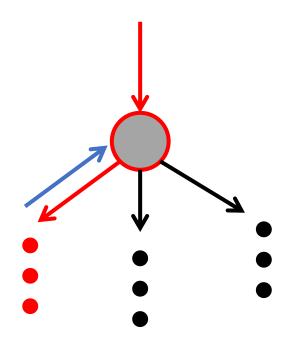
Orden lineal de los vértices del grafo de tal modo que por cada arista dirigida  $u \to v$ , u va antes que v en el orden.

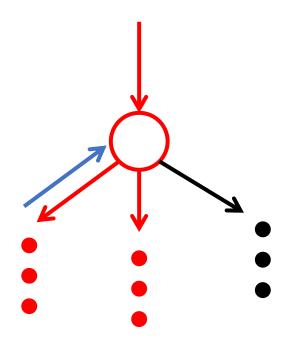
Aprovechamos una propiedad del recorrido DFS para calcular esto.

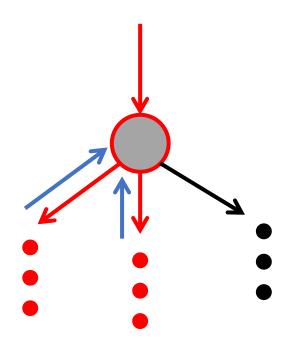


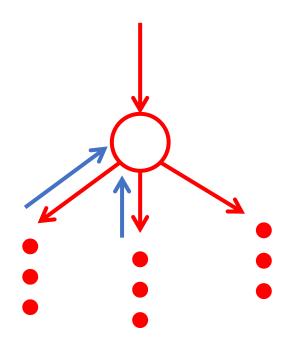


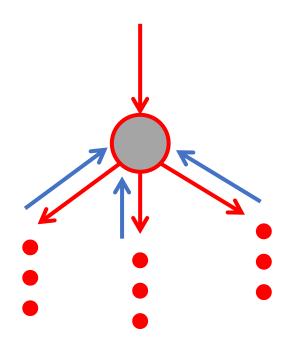


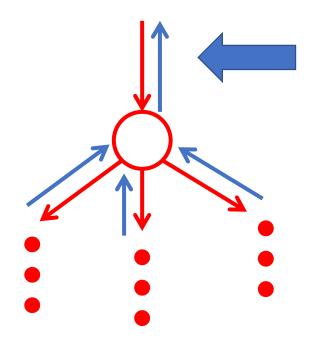




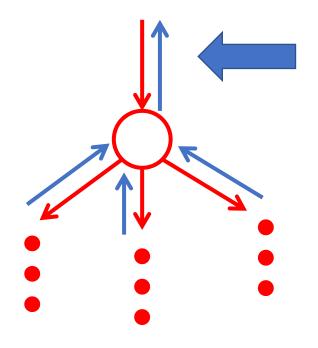






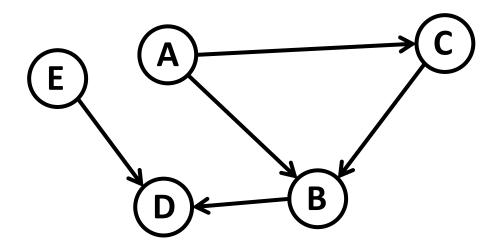


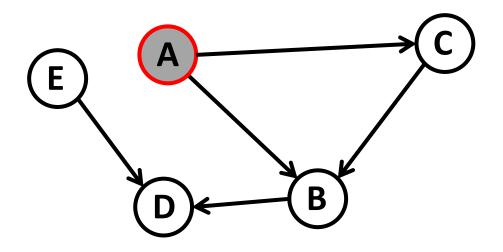
Regresar (backtracking) significa que ya se termino de procesar a **todos** sus vecinos.

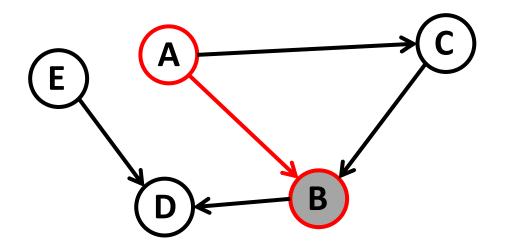


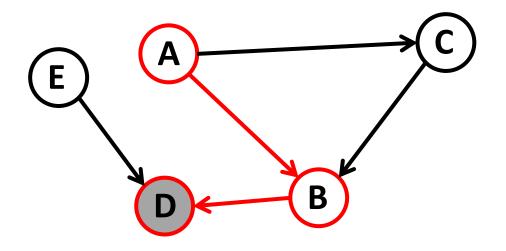
Regresar (backtracking) significa que ya se termino de procesar a **todos** sus vecinos.

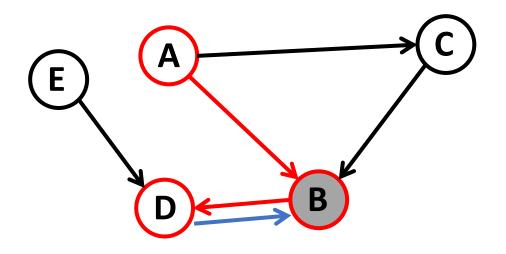
Entonces, podemos poner el nodo **antes** de sus vecinos en el orden topológico.

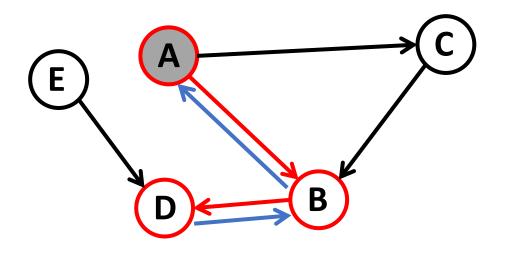




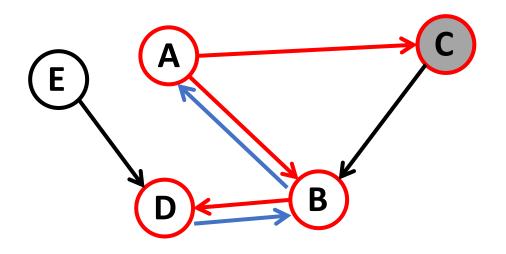




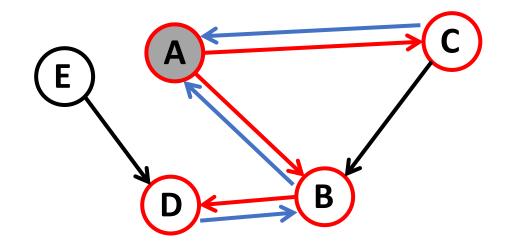




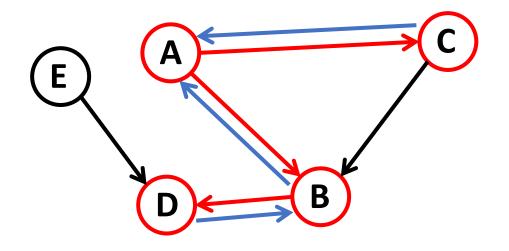
Orden topológico: **B D** 



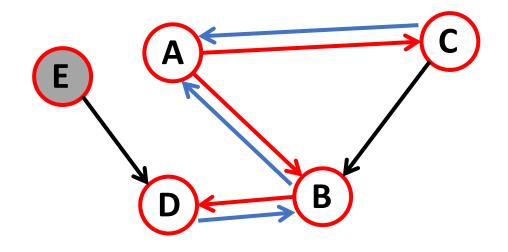
Orden topológico: **B D** 



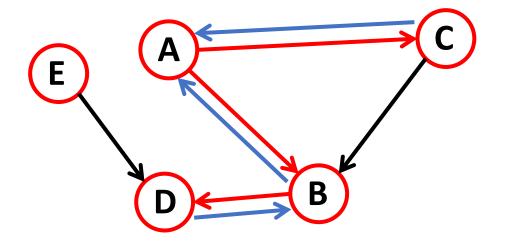
Orden topológico: C B D



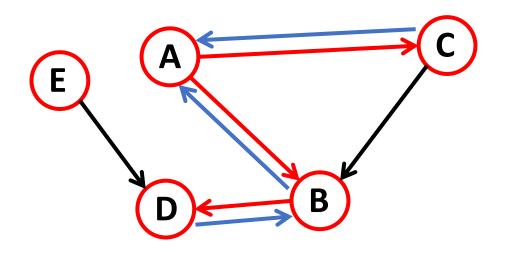
Orden topológico: A C B D



Orden topológico: A C B D

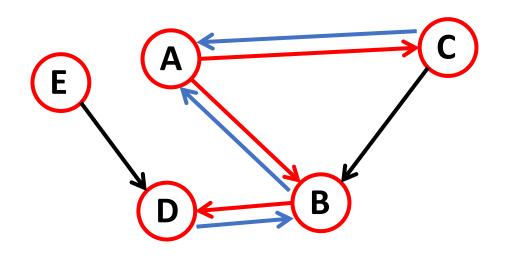


Orden topológico: E A C B D



```
vector<int> g[n];
bool vis[n];
void dfs(int u) {
  vis[u] = true;
  for (int v : g[u]) {
    if (!vis[v]) {
      // u -> v (ida)
      dfs(v);
      // v <- u (vuelta)
  // justo antes de salir del nodo (backtracking)
```

Orden topológico: E A C B D



Orden topológico: E A C B D

invTopSort: D B C A E

```
vector<int> g[n];
bool vis[n];
vector<int> invTopSort;
void dfs(int u) {
 vis[u] = true;
  for (int v : g[u]) {
    if (!vis[v]) {
      // u -> v (ida)
      dfs(v);
      // v <- u (vuelta)
  invTopSort.push_back(u);
  // justo antes de salir del nodo (backtracking)
```

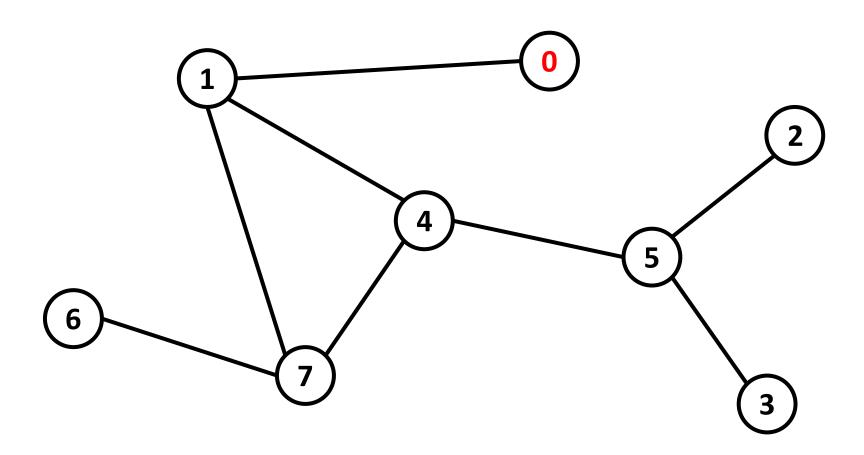
El nodo se pone al final de la lista en vez de al inicio, solo hace falta trabajar con la lista invertida.

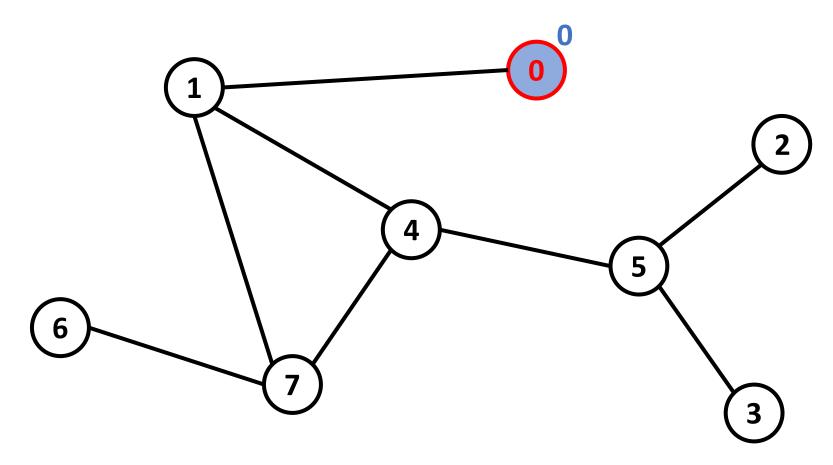
### Otros Algoritmos con DFS

- Flood Fill (DFS en matrices)
- Detectar ciclos en grafos dirigidos
- Puentes y Puntos de Articulación
- Componentes Fuertemente Conexos
  - Problema 2-SAT
- Heavy-Light Decomposition de arboles
- Euler Tour en arboles

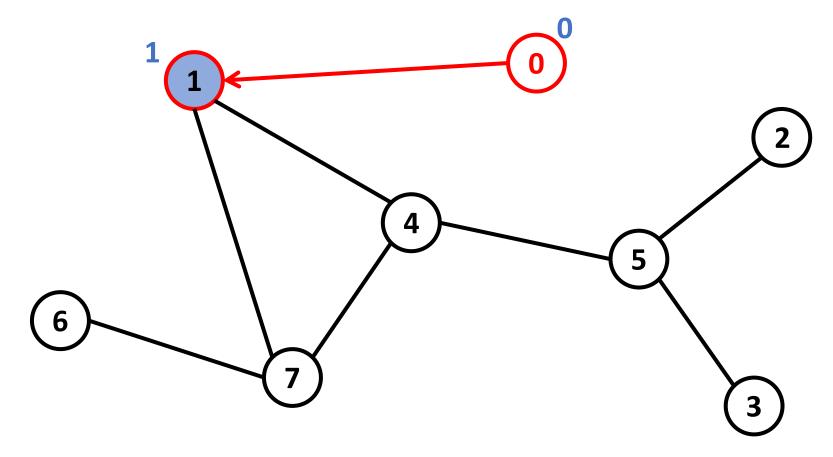
- Algoritmo de recorrido de grafos
- Empieza en un nodo y visita el resto en orden creciente segun la distancia al nodo inicial
- Es un poco mas dificil de implementar que DFS

BFS

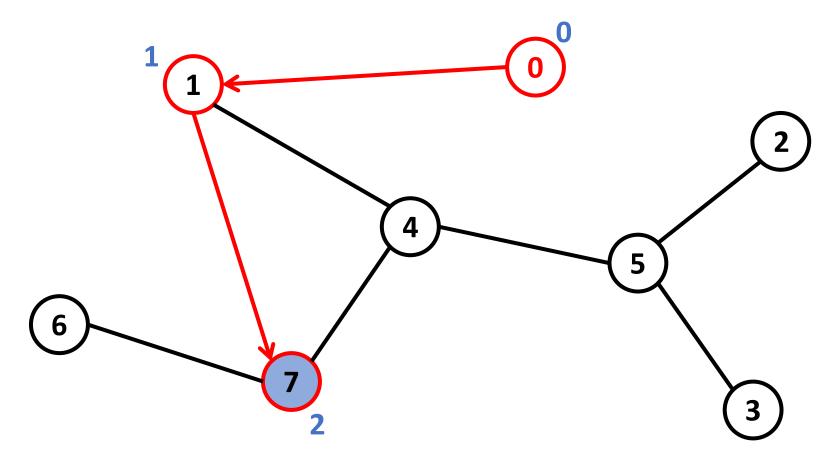




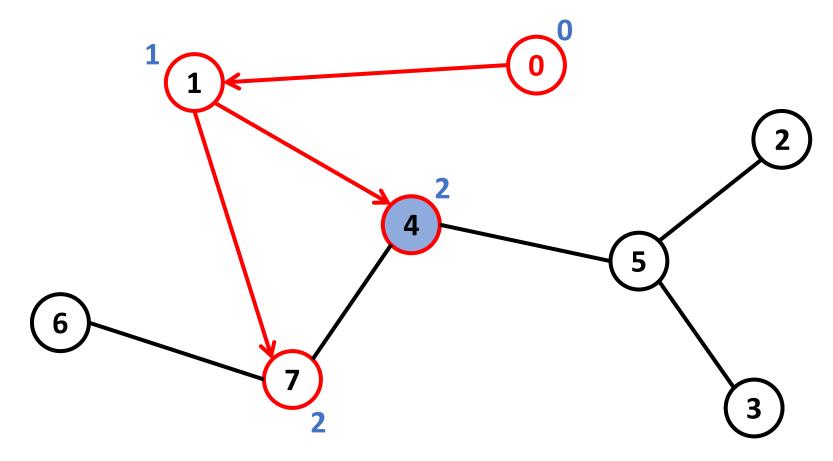
Orden de recorrido: 0



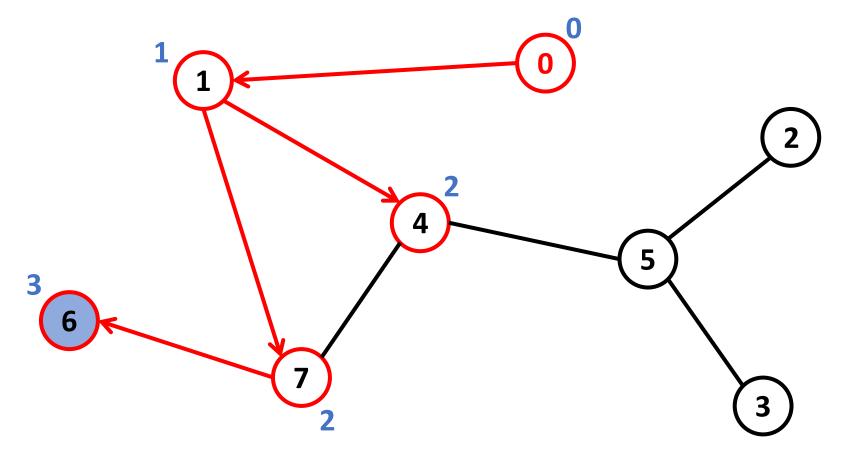
Orden de recorrido: 0, 1



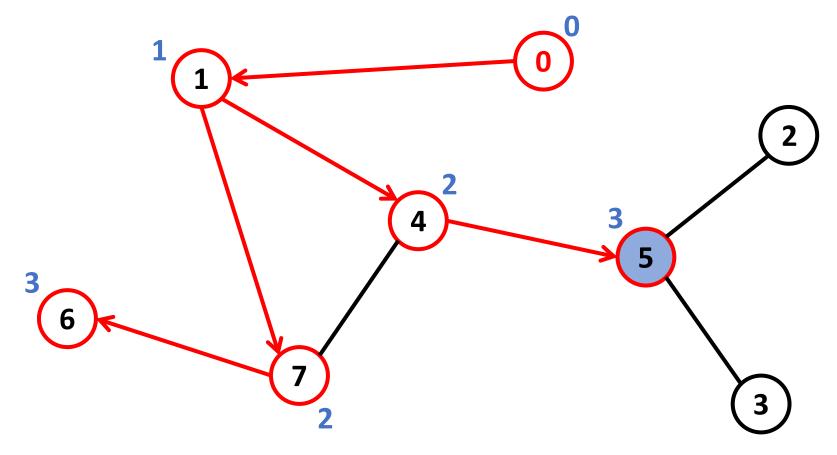
Orden de recorrido: 0, 1, 7



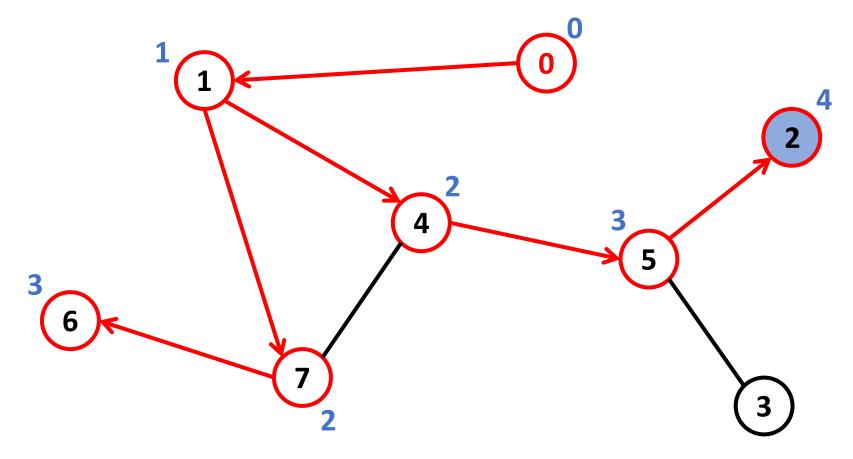
Orden de recorrido: 0, 1, 7, 4



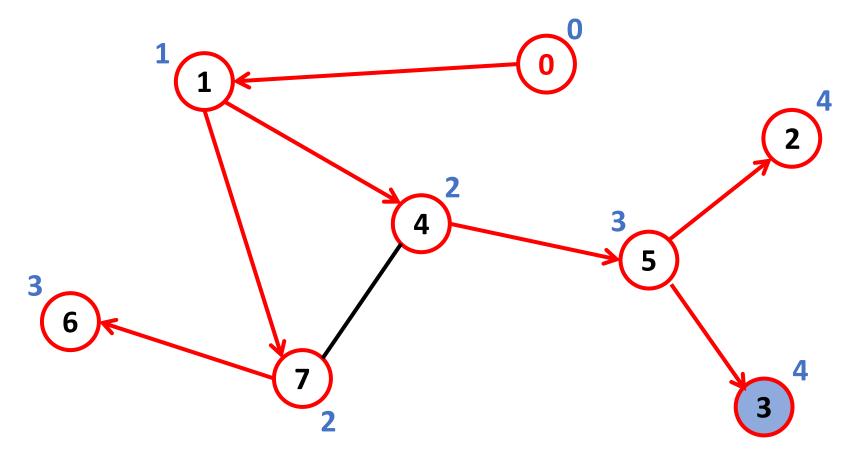
Orden de recorrido: 0, 1, 7, 4, 6



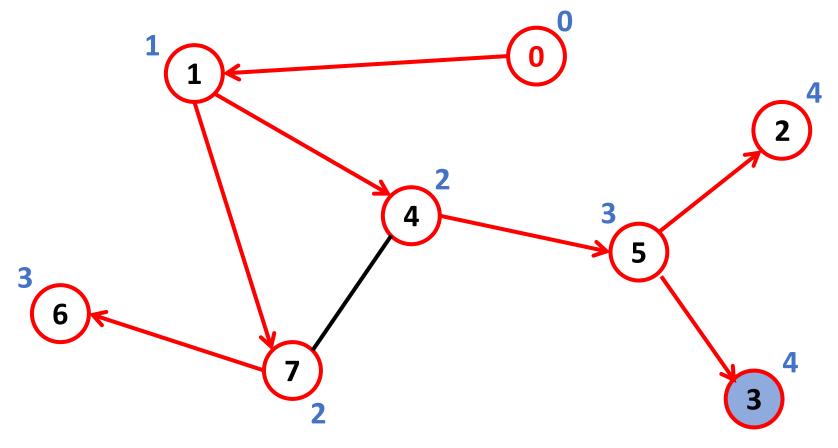
Orden de recorrido: 0, 1, 7, 4, 6, 5



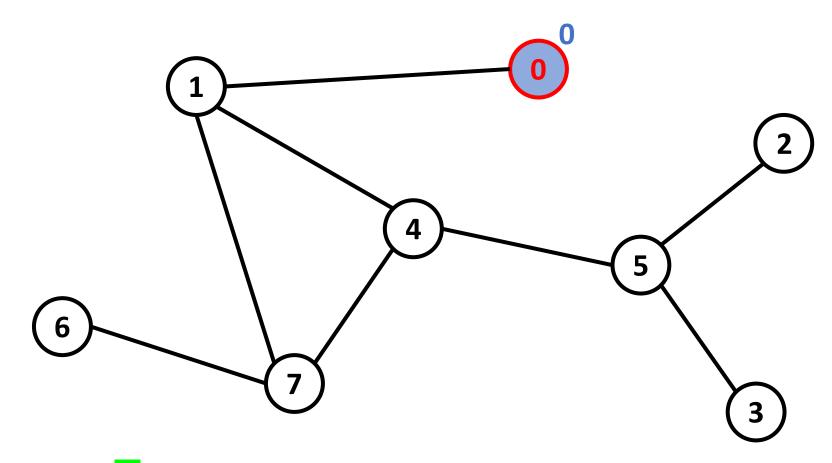
Orden de recorrido: 0, 1, 7, 4, 6, 5, 2



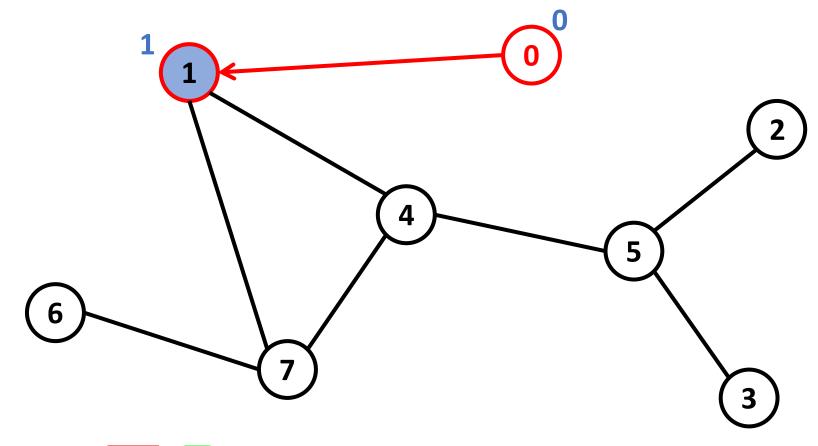
Orden de recorrido: 0, 1, 7, 4, 6, 5, 2, 3



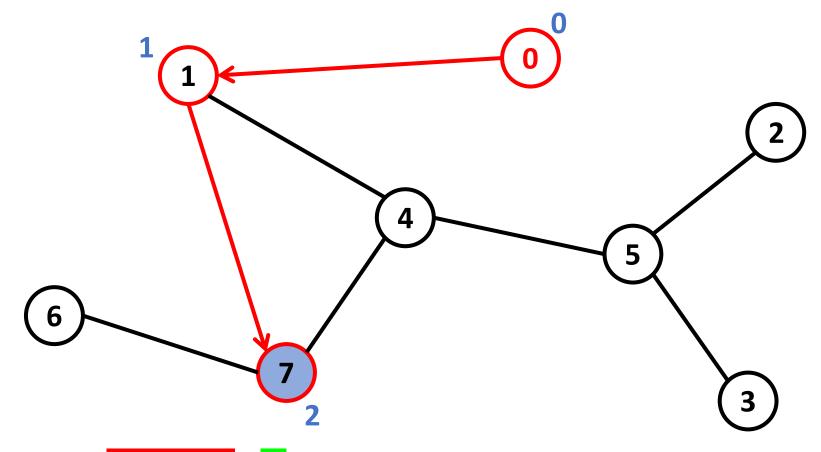
Orden de recorrido: 0, 1, 7, 4, 6, 5, 2, 3Distancia : 0 1 2 2 3 3 4 4



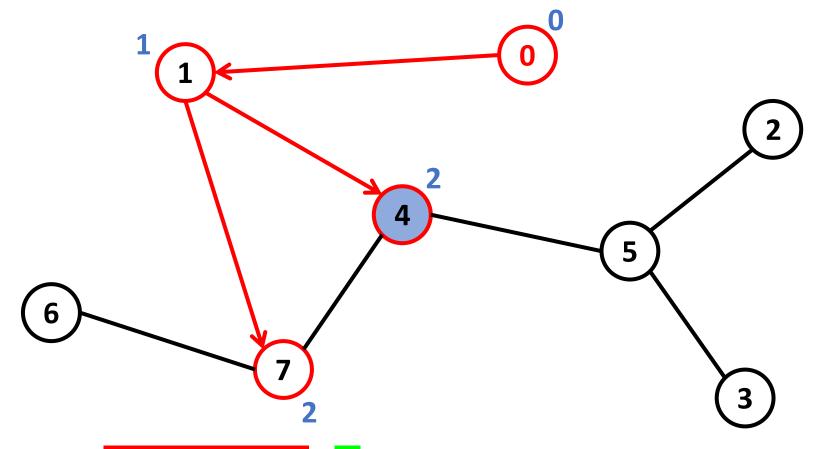
Cola: {1/1}
Nodo: 0



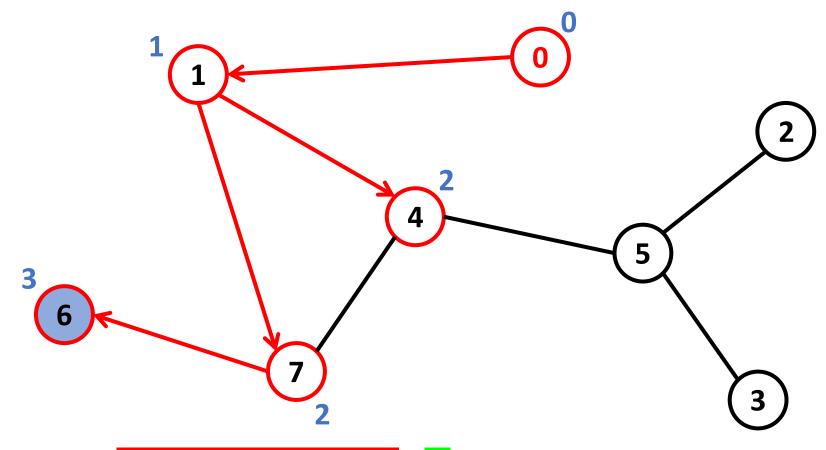
Cola: { 1, 7, 4}
Nodo: 1



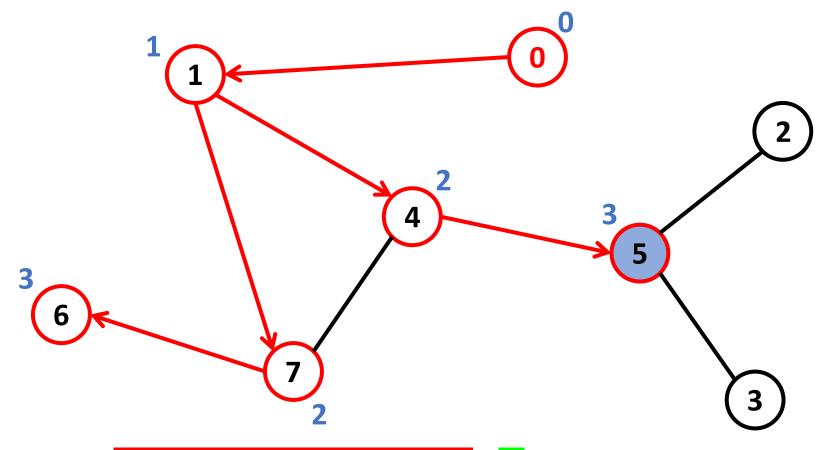
Cola: {1, 7, 4, 6}



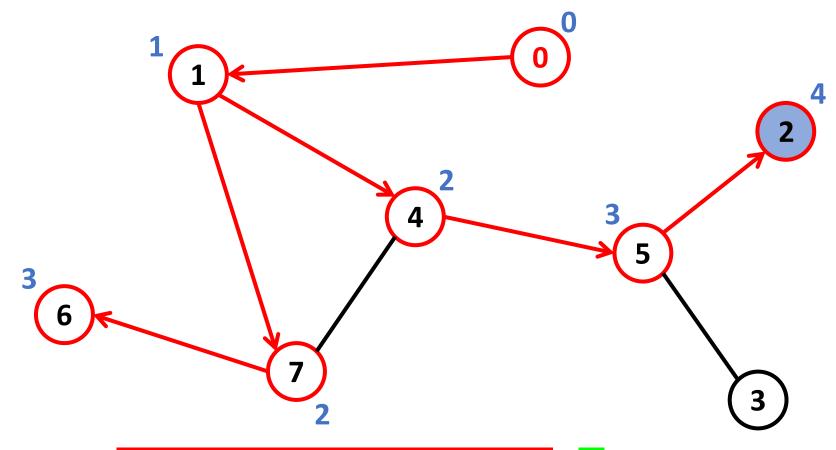
Cola: { 1, 7, 4, 6, 5}



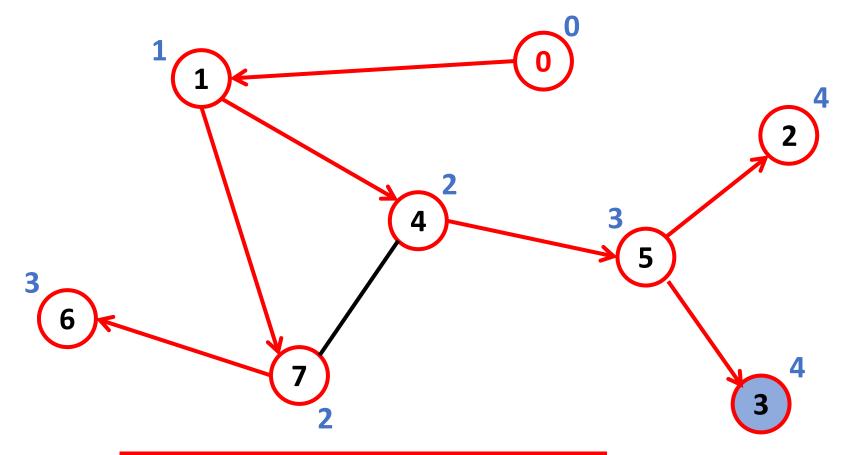
Cola: { 1, 7, 4, 6, 5}



Cola:  $\{\frac{1, 7, 4, 6, 5}{2}, 3\}$ 



Cola:  $\{\frac{1, 7, 4, 6, 5, 2}{3}\}$ 



Cola:  $\{\frac{1, 7, 4, 6, 5, 2, 3}{}\}$ 

### Código

#### **Necesitamos:**

```
vector<int> g[n];
vector<int> nivel(n, -1);
queue<int> q;
```

```
int inicio = 0;
nivel[inicio] = 0;
q.push(inicio);
while (!q.empty()) {
  int u = q.front();
  q.pop();
  for (int v : g[u]) {
    if (nivel[v] == -1) {
      nivel[v] = nivel[u] + 1;
      q.push(v);
```